

Structured Programming in Assembler Session 17871

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Who am I?

- Richard Cebula HLASM, IBM Hursley, UK
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- Develop and support the following products:
 - HLASM
 - SuperC
 - XREF
 - IDF
 - DISASM
 - Structured Programming Macros



- Structured Programming for Assembler!?
- Control sections C-, R-, DSECTs, COM and DXDs
- Subroutines, EXTRN / WXTRN and Conditional-Sequential RLDs
- The location counter and USING statements
- Type-checking in HLASM
- Program Objects
- Controlling Listings
- Structured Programming Macros (SPMs)



- Structured Programming for Assembler!?
 - Yes!
 - HLASM is the High Level Assembler providing an extensive list of features for assembler programmers
 - HLASM can help programmers to:
 - Organise their assembler code better
 - Maintain their code better
 - Increase code reuse
- Remember, HLASM is available for all z Systems operating systems: z/OS, z/VM, z/VSE, z/Linux (inc. z/TPF)



Control sections

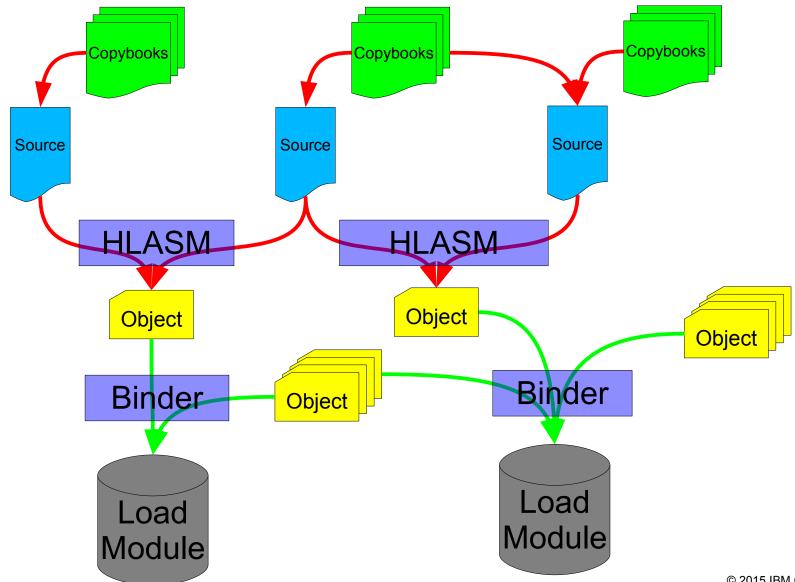


HLASM Structured Programming – Source and Object Modules

- An HLASM program consists of a number of sections of varying types
- The programmer should distinguish between:
 - Source Module
 - The source module is the division of code during assembly time
 - Each source module is assembled into a separate object module
 - Note that there is not always a 1-to-1 relationship between a source file and a source module
 - Object Module
 - The object module is the produced output from the assembler
 - The layout of the object module is determined by the type assembler options used
 - HLASM supports OBJ, GOFF and ELF object file formats



HLASM Structured Programming – Source and Object Modules



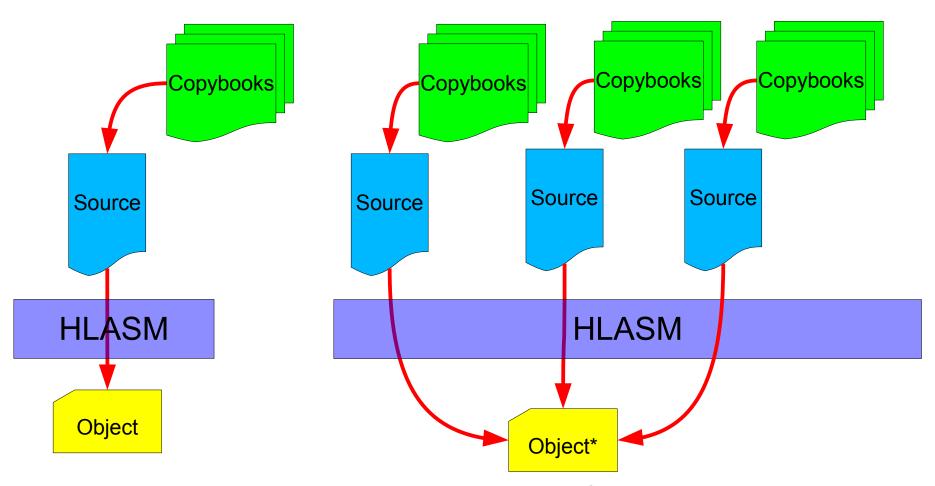


HLASM Structured Programming – Source Module – Sections

- Source module comprises 1 or more assembler statements
- Starts with any assembler statement except for MEXIT and MEND
- Ends with an END statement
- The BATCH option allows for more than a single source module to be specified in the same input stream
 - HLASM can take many source files on the same input stream, e.g. via use of the COPY statement
 - If HLASM encounters an END statement during processing then it completes assembling the source module.
 - If further input is found and the BATCH option is on (the default) then HLASM begins the assembly of a new source module



HLASM Structured Programming – Single VS Batch Assembly



In batch mode, all output is written to a single output object file.

*On z/Linux, use RPM asma90-1.6.0-27 or higher with the -R option to output to an object archive



HLASM Structured Programming – Object Module – Sections

- In the load module model, a control section is the smallest subdivision of a program that can be located as a unit
- Each assembled control section contains the object code for machine instructions and data
- Each source module (consisting of 1 or more source files) is assembled into 1 relocatable object module.
- The binder combines the object code of 1 or more sections into a load module (or program object)
 - The binder also calculates any addresses and space needed for common sections and external dummy sections from different object module
- The program loader loads the load module into virtual storage and converts any relocatable addresses into fixed locations



HLASM Structured Programming – Creating Sections

- The first section of a program may be initiated via the START, CSECT or RSECT assembler instructions
 - START initiates the first or only section of a source module
 - CSECT can be used anywhere in the source module to initiate or continue a control section
 - RSECT similar to CSECT but causes the assembler to check for possible violations of reenterability
- Unnamed sections are those that do not have a name although this is valid, it is not recommended.



- Reference control sections are used for referencing storage areas or to describe data and are not assembled into object code
 - Started by the DSECT, COM and DXD statements
 - As with other forms of control sections, they continue until interrupted by either another control section or an END statement
- DSECT Dummy control section
 - Reference control section that describes the layout of data in storage without reserving any virtual storage
 - There is no object code nor space in the object module reserved. A DSECT will cause the assembler to assign location values for the DSECT's symbols relative to its beginning
 - Data can be referred to symbolically by using the symbols defined in the dummy section



0000004DE2D4C9E3C8 404040404040404040 404040404040C2D6C2 404040404040404040 4040404040404040D4 D94040F2F0404040C8 C9C7C840E2E3 Etc.

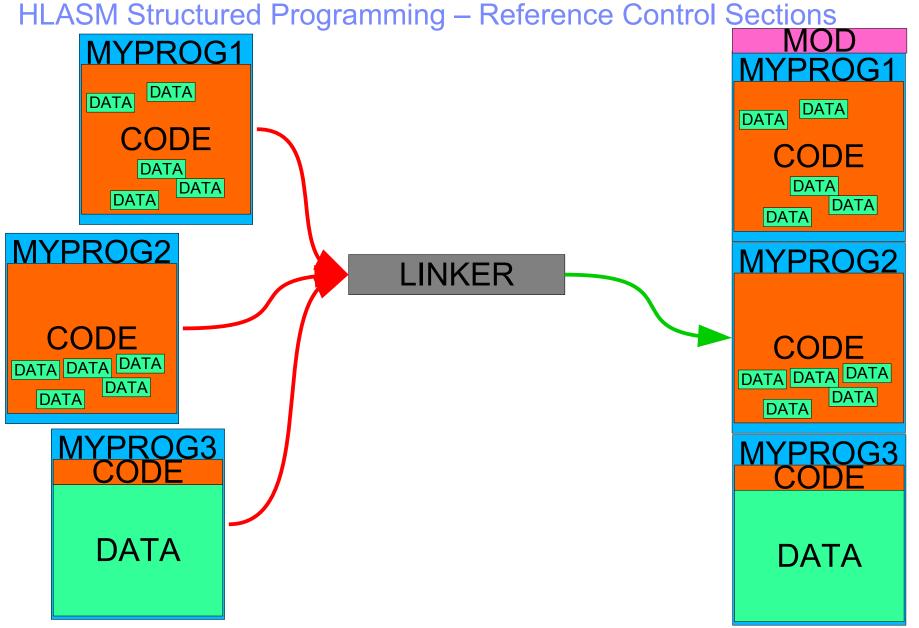


| | CUSTOMER | DSECT | | |
|---------------------------------|----------------------|-------|-------|---|
| 0000004DE2D4C9E3C8 | cust_number | dc | f'0' | |
| 4040404040404040 | cust_name | dc | c120' | 7 |
| 4040404040C2D6C2 | ←cust_cname | dc | c120' | 7 |
| 404040404040404 | cust_title | dc | C14' | 7 |
| 40404040404 04040 D4 | house_num | dc | C15' | 7 |
| D94040F2F0404040C8 | <pre>ddr_line1</pre> | dc | c140' | 7 |
| C9C7C840E2E3 | addr_line2 | dc | c140' | 7 |
| Etc. | city | dc | c140' | • |

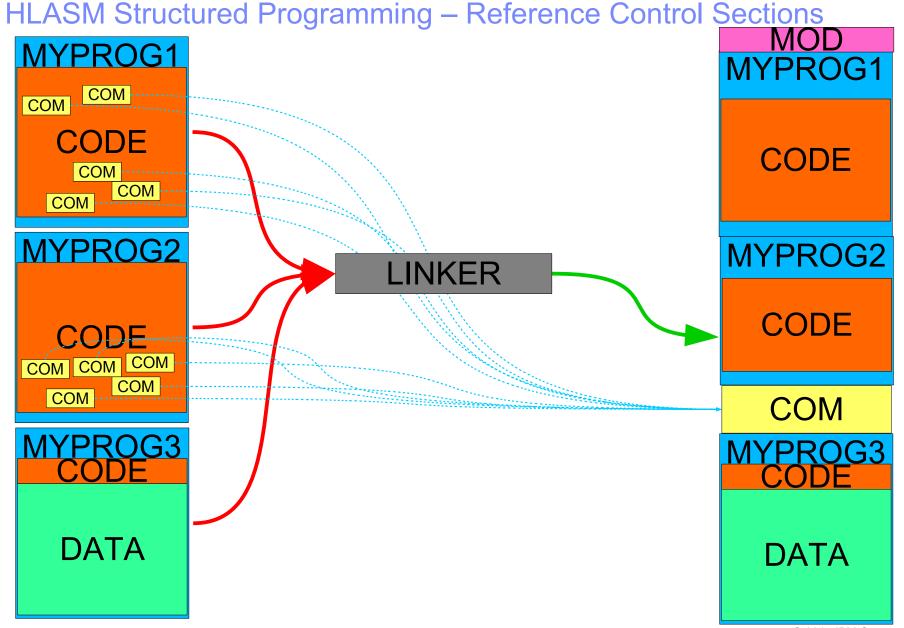


- COM Common control section
 - Allows for the definition of a common storage area in one or more source modules
 - At link time, only one copy of a COM section is created for all the object modules being linked as a single program
 - Only the storage area is provided the data must be provided at execution time, i.e.
 COM is uninitialised!
- External dummy sections
 - Created via DXD, DSECT and CXD instructions or via the Q-type address constant
 - To use:
 - Use a DXD instruction to define the external dummy section
 - Provide a Q-type constant to address the external dummy section
 - Use the CXD instruction to obtain the total length of all external dummy sections
 - Allocate the storage required as calculated by the CXD
 - Address the allocated storage (plus any offset into the areas as required)











- What happens with my DXDs at link time?
 - All DXDs are not allocated as part of the module instead just the references to the DXD items in the linked module are retained.
 - The linker collects all the DXDs in all the modules it is linking together and creates the complete External Dummy Section.
 - All DXD Q-type references throughout the various modules are replaced by offsets from the start of the completed External Dummy Section.
 - The total length of the External Dummy Section is placed into every CXD.
- To make use of the DXDs, the program has to allocate the correct amount of storage as per the CXD and then reference each member of the DXD via Q-type offsets, e.g.:

```
my_prog3:
...code goes here...
L r0,DXD_TOTAL
GETMAIN R,LV=(0)
...code goes here...
DXD_TOTAL CXD
```

```
my_prog1:
L r7,=Q(MY_DXD_SYM)
AR r7,r1
L r7,0(,r7)
```

```
my_prog2:
    L r7,=Q(MY_OTHER_DXD)
    AR r7,r1
    L r7,0(,r7)
```



HLASM Structured Programming – Reference Control Sections DXD MYPROG2 LINKER DXD DXD DXD DXD DXD DATA **DATA** DXD DXD



Subroutines



HLASM Structured Programming – Subroutines

- Breaking sections of code out into subroutines is one of the easiest ways to structure programs.
- If the subroutine is close to the code that calls it then the following code should suffice:

```
bras r14, mysubroutine
```

However, if the subroutine is beyond the range of an active using then an A-type address constant should be used instead:

```
1     r15,=A(mysubroutine)
basr     r14,r15
```

- This technique only works for addresses which can be resolved during assembly time.
- It is much more common to use BRAS/BRASL. BRASL also works with for external entry points



HLASM Structured Programming – External Subroutines

- An external reference is a symbol that is unresolvable at assembly time. Instead, the symbol is resolved by the binder when the object files are linked together.
- To make a symbol externally available, the module in which the symbol is declared must either define it as the name of a control section or as an ENTRY, e.g.:

MYSUB ENTRY

- Note that START, CSECT, RSECT are automatically considered as ENTRY symbols
- In the module that references the external symbol, the external symbol must be declared as EXTRN before being used:

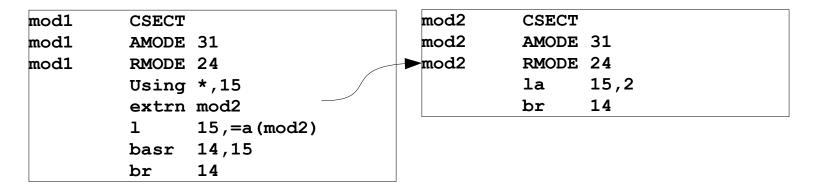
```
EXTRN MYSUB
...some code...
subaddr dc a(MYSUB)
```

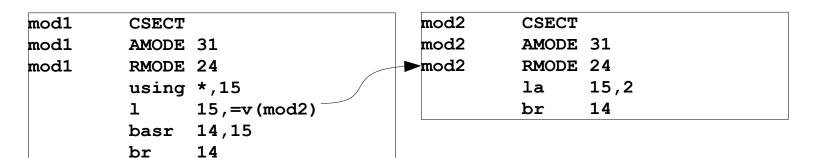
- The WXTRN instruction can also be used instead of EXTRN. However
 - For EXTRN, if the symbol is not found from the list of modules being linked, then a library search is performed to try and find the symbol
 - For WXTRN, no library search is performed



HLASM Structured Programming – External Subroutines

• Alternatively, the V-type constant automatically declares a symbol as EXTRN:







HLASM Structured Programming – Using CSRs – z/OS only

- For z/OS 1.13 and above, with HLASM APAR PM74898 applied and for GOFF only...
- Conditional Sequential RLDs, (CSRs) allow an assembler programmer to specify multiple external references in a single V-type address constant separated by colons, e.g.:

```
my modules     dc v(mod1:mod2:mod3)
```

- The assembler program continues to use the V-type as normal but:
 - When the binder goes to resolve the external symbol, it will resolve it to the first available external symbol
 - If the first symbol isn't present, then the 2nd available symbol is chosen and so on...
- For a more detailed look at how to use CSRs see:

http://www.ibm.com/support/docview.wss?uid=swg21598283



The location counter and USINGs



HLASM Structured Programming – Using the location counter

- The location counter is an internal count from the start of the source module in bytes and is represented by the symbol *
- The LOCTR instruction may be used to specify multiple location counters within a control section.
- Each location counter is given consecutive addresses and the order of location counters produced is dependent on the order in which they are defined
- A location counter continues until it is interrupted by the START, CSECT, DSECT, RSECT, CATTR or LOCTR instructions
- Specifying LOCTR to an already defined location counter will resume that counter



HLASM Structured Programming – The USING statement

- The USING specifies a base register for a series of location counter displacements to a particular symbol within a control section
- There have been improvements to HLASM's USING statement including 'Labeled USINGs' and 'Dependent USINGs'
 - Labeled USINGs
 - Allow simultaneous references to multiple instances of an object
 - Allow one object to be referenced per register
 - Dependent USINGS
 - Address multiple objects with a single register
 - Allow for a program to require fewer base registers
 - Allow for dynamic structure remapping during execution



HLASM Structured Programming – Labeled USINGs

- The labeled USING statement has the syntax in the form:
 - qualifier USING base,register
- A symbol can be referenced by prefixing it with a qualifier and therefore the same symbol may be addressed by 2 or more base registers at the same time, e.g.

```
* COPY THE CUSTOMER DETAILS TO THE NEW CUSTOMER RECORD

CSTMR1 USING CUST_DATA,R4

CSTMR2 USING CUST_DATA,R5

MVC CSTMR2.CUST_DETAILS,CSTMR1.CUST_DETAILS

L R3,NEW_CUST_NUMBER

ST R3,CSTMR2.CUST_NUM
```

 Without labeled USINGs, the above MVC would have to be written using manually calculated offsets, e.g.

```
* COPY THE CUSTOMER DETAILS TO THE NEW CUSTOMER RECORD
USING CUST_DATA,R5

MVC CUST_DETAILS,CUST_DETAILS-CUST_DATA(R4)
```



HLASM Structured Programming – Dependent USINGs

- A dependent USING is one that specifies an anchor location rather than a register as its operand
- Allows the programmer to address more than one DSECT at the same time using the same base register

```
* ADDRESS THE DATA IN BOTH XPL_DATA AND REQI_DATA
USING XPL_DATA,R10
USING REQI_DATA,L_XPL_DATA+XPL
LA R1,XPL_INPUT
LA R2,REQI_INPUT
```

- HLASM will add the offset from REQI_DATA to XPL_DATA in order to generate the offsets for the fields in the REQI_DATA DSECT based off register 10
- The displacement limit for a dependent using is still limited to 4095 (as usual)



HLASM Structured Programming – Labeled Dependent USINGs

 Dependent USINGs can also have a label allowing for some complex USING issues to be easily resolved using a single a base register

| | USING E,7 | | 1 Top level |
|------|----------------|---|------------------------|
| * | | | |
| D1E | USING D,D1 | 1 | 2 Map D1 into E at D1 |
| D1F1 | USING F,D1E.F1 | 2 | 3 Map F1 into D1 at F1 |
| D1F2 | USING F,D1E.F2 | 2 | 3 Map F2 into D1 at F2 |
| D1F3 | USING F,D1E.F3 | 2 | 3 Map F3 into D1 at F3 |
| * | | | 2 Middle level |
| D2E | USING D,D2 | 1 | Map D2 into E at D2 |
| D2F1 | USING F,D2E.F1 | 3 | 3 Map F1 into D2 at F1 |
| D2F2 | USING F,D2E.F2 | 3 | 3 Map F2 into D2 at F2 |
| D2F3 | USING F,D2E.F3 | 3 | 3 Map F3 into D2 at F3 |
| * | | | 2 Middle level |
| D3E | USING D,D3 | 1 | Map D3 into E at D3 |
| D3F1 | USING F,D3E.F1 | 4 | 3 Map F1 into D3 at F1 |
| D3F2 | USING F,D3E.F2 | 4 | 3 Map F2 into D3 at F2 |
| D3F3 | USING F,D3E.F3 | 4 | 3 Map F3 into D3 at F3 |



HLASM Structured Programming – Labeled Dependent USINGs

```
* Move fields named X within DSECTs described by F
MVC D1F1.X1,D1F1.X2 Within bottom-level DSECT D1F1
MVC D1F3.X2,D1F1.X1 Across bottom-level DSECTs in D1
MVC D3F2.X2,D3F3.X2 Across bottom-level DSECTs in D3
MVC D2F1.X1,D3F2.X2 Across bottom-level DSECTs in D2 & D3
```

- * Move DSECTs named F within DSECTs described by D

 MVC D3E.F1,D3E.F3 Within mid—level DSECT D3E

 MVC D1E.F3,D2E.F1 Across mid—level DSECTs D1E, D2E
- * Move DSECTs named D within E

 MVC D1,D2 Across top—level DSECTs D1, D2

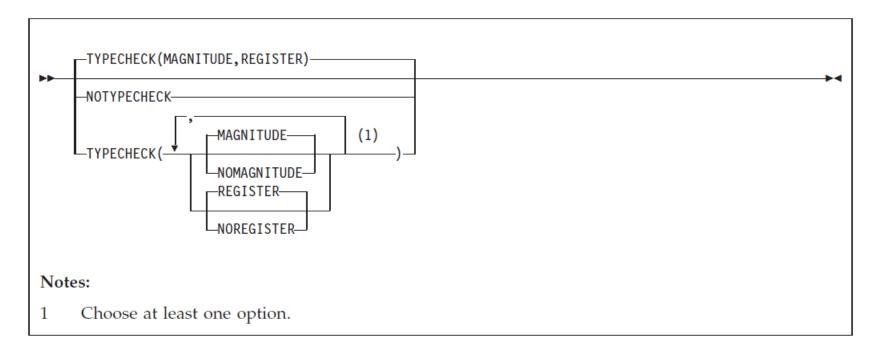


HLASM Type Checking



HLASM Structured Programming – Type checking - TYPECHECK

- HLASM provides a number of type-checking facilities to assist make programs safer to change
- HLASM type-checking is provided by the TYPECHECK option:





HLASM Structured Programming – Type checking - TYPECHECK

MAGNITUDE

 Causes HLASM to perform magnitude validation of signed immediate data fields, e.g.:

```
000000 A718 FFFF 0FFFF 7 lhi 1,x'ffff'

** ASMA320W Immediate field operand may have incorrect sign or magnitude
```

REGISTER

 Causes HLASM to perform type checking of register fields of machine instruction operands, e.g.:

```
00001 7 fprl equ 1,,,,fpr
000000 A718 FFFF 0FFFF 8 lhi fprl,x'ffff'
** ASMA323W Symbol fprl has incompatible type with general register field
```

NOTYPECHECK

Turns off all type-checking



HLASM Structured Programming – Type checking – User types

• The DC and DS statements allow users to create their own "Program Types" using the P parameter, e.g.:

```
my_data dc cap(c'read')13'this is my data'
```

- The value of the program type is returned via the SYSATTRP macro function
- The value of the assembler type is returned via the SYSATTRA macro function
- The value of the type *attribute* is returned via the macro T' operator



HLASM Structured Programming – Type checking – EQU

- The EQU statement can be used to:
 - Specify a 32-bit value used as the program type (4th operand of EQU)
 - Specify a register type (5th operand of EQU)
- The program type for an EQU symbol will be returned via the SYSATTRP function as for macro symbols
- The register types that are available to use are:
 - AR Access registers
 - CR, CR32, CR64 Control registers
 - FPR Floating point registers
 - GR, GR32, GR64 General purpose registers
 - VR z13 vector registers
- Note that there is no vector register type-checking on pre-z/Architecture vector instructions
 - The hardware hasn't been available to use for a while now...don't use them!!



HLASM Structured Programming

Program Objects



- Program Objects are a newer form of the z/OS load module which reside in PDSEs.
- There are 3 ways of creating program objects:
 - Use the GOFF assembler option. This will generate default classes causing the binder to store the linked program object into a PDSE.
 - Use the GOFF assembler option and create your own classes.
 - Store the output of binding a program into a PDSE. The binder will create default classes and create the program object for you.
- Program objects have advantages over load modules such as:
 - Can be stored in PDSE and zFS / HFS
 - Multiple dimensions and multiple A/RMODEs in the same module
 - Limited to 1GB in size load modules are limited to 16MB in size
 - External symbols can have up to 32K characters rather than 8
 - Open-ended architecture use the Binder interfaces to get to the data



- A load module consists of a number of control sections whereas a PO is a collection of classes and sections.
 - A *class* is an independently loadable module of data. The class itself also has various attributes to determine whether the data is to be loaded, reusable, etc.
 - A section is a collection of data from one or more classes.
 - An *element* is an intersection between a class and a section. The attributes of an element are defined by the class for that element's particular section.

| MYPROG | Class X | Class Y | Class Z |
|-----------|---------|---------|---------|
| Section A | Element | Element | Element |
| Section B | Element | Element | Element |
| Section C | Element | Element | Element |



- The CSECT instruction is now used to create a section to which a number of (default) classes belong. The classes which are created by HLASM are:
 - B_IDRL used to contain IDR (identification data) for each section
 - B_PRV used to contain any external dummy section data for a section
 - B_TEXT used to contain a section's object code
- Each of these classes has an element definition (ESD symbol type ED) whose owning ID is the ESDID of the section.
- The CSECT instruction will also generate a *label definition* (ESD symbol LD) to be an entry point for the section and will refer to the B_TEXT class for that section.



External Symbol Dictionary

| Symbol | Туре | e Id | Address | Length | Owner | Id | Flags | Alias-of |
|--------------|------|---------|---------|----------|--------|----|-------|----------|
| LISTINGB | SD | 0000001 | | | | | | |
| B_IDRL | ED | 0000002 | | | 000000 | 01 | | |
| B_PRV | ED | 0000003 | | | 000000 | 01 | | |
| B_TEXT | ED | 0000004 | 0000000 | 00000084 | 000000 | 01 | 08 | |
| LISTINGB | LD | 0000005 | 0000000 | | 000000 | 04 | 08 | |
| EXTERNAL | FUNC | CTION | | | | | | |
| | ER | 0000006 | | | 000000 | 01 | | |
| FUNCY | ER | 0000007 | | | 000000 | 01 | | |
| listme | ER | 8000000 | | | 000000 | 01 | | LISTINGZ |
| COMMON_DA | ATA | | | | | | | |
| _ | SD | 0000009 | | | | | | |
| B_IDRL | ED | A000000 | | | 000000 | 09 | | |
| B_PRV | ED | 000000B | | | 000000 | 09 | | |
| B_TEXT | ED | 000000C | 0000000 | 0000018 | 000000 | 09 | 00 | |
| COMMON_DA | ATA | | | | | | | |
| _ | CM | 000000D | 0000000 | | 000000 | 0C | 00 | |

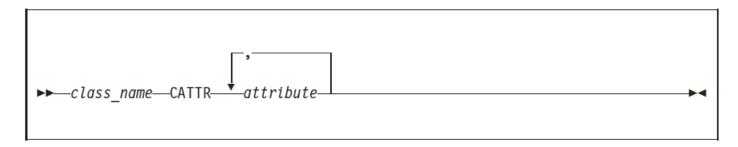


External Symbol Dictionary

| Symbol | Туре | e Id | Address | Length | Owner | Id | Flags | Alias-of |
|-----------|-------|--------------------------|---------|----------|--------|-----|-------|----------|
| LISTINGB | SD | 00000001 | | | | | | |
| B_IDRL | ED | 0000002 | | | 000000 | 01 | | |
| B_PRV | ED | 0000003 | | | 000000 | 01 | | |
| B_TEXT | ED | 0000000 | 0000000 | 00000084 | 000000 | 01 | 08 | |
| LISTINGB | LD | 00000005 | 0000000 | | 000000 | 04 | 08 | |
| EXTERNAL | _FUNC | CTION | | | | | | |
| | ER | 0000006 | | | 000000 | 01 | | |
| FUNCY | ER | 0000007 | | | 000000 | 01 | | |
| listme | ER | 8000000 | | | 000000 | 01 | | LISTINGZ |
| COMMON_D | ATA | | | | | | | |
| | SD | 0000009 | | | | | | |
| B_IDRL | ED | A 000000 A | | | 000000 | 009 | | |
| B_PRV | ED | 000000B | | | 000000 | 009 | | |
| B_TEXT | ED | 000000C | 0000000 | 0000018 | 000000 | 009 | 00 | |
| COMMON_DA | ATA | | | | | | | |
| | CM | 000000D | 0000000 | | 000000 | 00C | 00 | |



Classes can be created by using the CATTR instruction.



- Some of the attributes that can be specified are:
 - ALIGN(n) align the class on a 2^n boundary (values of n are 0, 1, 2, 3, 4, 12)
 - READONLY text is storage protected
 - REFR, RENT, REUS same as for load modules but for just a class
 - RMODE(24|31|ANY)
 - PART(name) a named subdivision of an element that can have an internal structure. The ESD type for a part is PD.
- The default attributes for a class are:
 - ALIGN(3), EXECUTABLE, NOTREUS, RMODE(24)



HLASM Structured Programming

Controlling Listings



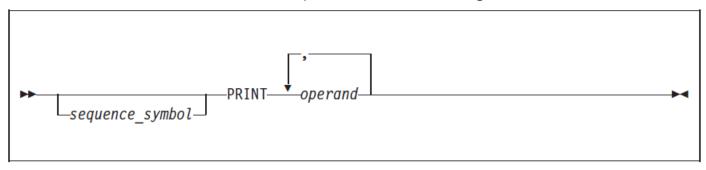
HLASM Structured Programming – Controlling Listings

- HLASM listings contain a large amount of very useful information but may easily become daunting to look at. Fortunately, HLASM has some useful instructions to help manage the data is presented in the listing.
- TITLE Specifies headings for each page of the source and object section of the listing.
 Using the TITLE instruction also causes the page header to be printed.
- EJECT Causes the next line of assembler listing to be printed at the top of a new page.
- SPACE Inserts one or more blank lines into the listing on the current page.



HLASM Structured Programming – Controlling Listings

PRINT – Controls the amount of detail printed in the listing.



- ONIOFF Print or stop printing the source and object section of the listing
- GEN|NOGEN Print or stop printing macro expansion. Note that the MNOTE instruction always causes a message to be printed.
- NODATA|DATA Print only the first 8-bytes of a data constant in the listing. If PRINT DATA is used then the entire constant is printed in the object code.
- NOMCALL|MCALL Control whether or not nested macro calls are printed
- MSOURCE|NOMSOURCE Control whether or not the printing of source statements produced from macro processing are printed
- UHEAD|NOUHEAD Control whether or not to print the ACTIVE USINGS at the start of each page of listing.



HLASM Structured Programming

Structured Programming Macros



- HLASM provides a set of Structured Programming Macros (SPMs) as part of the HLASM Toolkit feature
- The SPMs provide
 - Branching structures (if) IF, ELSEIF, ELSE, ENDIF
 - Looping structures (Do) DO, ITERATE, DOEXIT, ASMLEAVE, ENDDO
 - Searching STRTSRCH, ORELSE, ENDLOOP, ENDSRCH, EXITIF
 - N-way branching CASENTRY, CASE, ENDCASE
 - Selection on general sequential cases SELECT, WHEN, NEXTWHEN, OTHERWISE, ENDSEL
- Why use SPMs?
 - Improve code readability, maintainability, understandability
 - Faster application development
 - Cleaner code
 - Eliminates extraneous labels makes it easier to revise



• Ever seen some code like this?

| CHK1 | ICM | 2,15,RETCODE | RETURN CODE 0? |
|----------|------|--------------|----------------------------------|
| | BC | 8, RETOK | BRANCH TO RETURN OK |
| | SLL | 2,2 | MAKE BRANCH OFFSET |
| | ICM | 3,15,RESN | RETURN NOT 0 AND REASON PRESENT? |
| | BC | 7, RESTAB(2) | BRANCH INTO REASON TABLE |
| CHK2 | L | 15, CHKRES | CALL CHECKRES SUBROUTINE |
| | BALR | 14,15 | |
| | LTR | 15,15 | RETURN OK? |
| | BC | 8,CHK1 | GOTO REPEAT CHECKS |
| | BC | 1,RETOK | RETURN WILL DO |
| | BC | 4,CHK1 | LOOKUP REASON |
| RETOK | С | 2,=F'3' | WAS RETURN LESS THAN 3? |
| | BC | 4,CHKERRL4 | YES |
| | LA | 3,0 | |
| | BC | 15,CONT242 | CONTINUE CODE |
| CHKERRL4 | L | 15,0MINERR | OUTPUT MINOR ERROR |
| CONT242 | • • | • | |



• Ever seen someone make a fix like this?

| CHK1 | ICM | 2,15,RETCODE | RETURN CODE 0? |
|----------|------|-----------------|----------------------------------|
| | BC | 8, RETOK | BRANCH TO RETURN OK |
| | SLL | 2,2 | MAKE BRANCH OFFSET |
| CHK1B | ICM | 3,15,RESN | RETURN NOT 0 AND REASON PRESENT? |
| | BC | 7, RESTAB(2) | BRANCH INTO REASON TABLE |
| CHK2 | L | 15,CHKRES | CALL CHECKRES SUBROUTINE |
| | BALR | 14,15 | |
| | LTR | 15,15 | RETURN OK? |
| | BC | 8,CHK1 | GOTO REPEAT CHECKS |
| | BC | 1, RETOK | RETURN WILL DO |
| | BC | 4,CHK1 B | LOOKUP REASON |
| RETOK | С | 2,=F'3' | WAS RETURN LESS THAN 3? |
| | BC | 4,CHKERRL4 | YES |
| | LA | 3,0 | |
| | BC | 15, CONT242 | CONTINUE CODE |
| CHKERRL4 | L | 15,0MINERR | OUTPUT MINOR ERROR |
| CONT242 | | | |



- Maybe the previous slides were a contrived example...you'd be surprised...
 - CHK1B might be an easy fix for now but what will the code be like when we reach CHK9C ... We've all seen it happen...
- Why does assembler code have a reputation for always being an unmaintainable mess?
 - Code gets maintained rather than developed...
 - Documentation gets lost...
 - Programmer's make bad choices...
- What can we do to fix this?
 - Structure your code
 - Don't call everything tmp1, tmp2, temp1, temp2, etc.
 - Use HLASM 's facilities to help you



- To use just copy in the ASMMSP copybook
 - By default, the SPMs produce based branch on condition instructions. To cause the SPMs to produce relative branch instructions, use the ASMMREL macro:

ASMMREL ON

- Global variables used by the macros begin with &ASMA_
- User-visible macros have meaningful mnemonics
 - Internal non-user macros begin with ASMM



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF (condition)
 - IF (instruction, parm1,parm2,condition)
 - IF (compare instruction,parm1,condition,parm2)
 - IF CC=condition_code
- Conditions may be joined together with Boolean operators
 - AND
 - OR
 - ANDIF
 - A OR B AND C \rightarrow A OR (B AND C) but A OR B ANDIF C \rightarrow (A OR B) AND C
 - ORIF
 - A AND B ORIF C OR D → (A AND B) OR (C OR D)
 - NOT



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF (condition)

```
LT R1,CUSTOMER_NUMBER

IF (Z) THEN

Code to assign a new customer number

ELSEIF (N) THEN

Code for negative customer number - error

ELSE

Look up the customer and do something useful

ENDIF
```



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF (condition)

```
LT R1,CUSTOMER_NUMBER

IF (Z) THEN

Code to assign a new customer number

ELSEIF (N) THEN

Code for negative customer number - error

ELSE

Look up the customer and do something useful

ENDIF
```

THEN is actually a comment – not needed



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF (instruction, parm1,parm2,condition)

```
IF (LT,R1,CUSTOMER_NUMBER,Z) THEN

Code to assign a new customer number

ELSEIF (N) THEN

Code for negative customer number - error

ELSE

Look up the customer and do something useful

ENDIF
```



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF (compare instruction, parm1, condition, parm2)

```
if (cli,x_val,eq,X'55'),OR, Value is x55? X
        (cli,x_val,eq,X'5D'),OR, Value is x5D? X
        (cli,x_val,eq,X'51'),OR, Value is x51? X
        (cli,x_val,eq,X'59') Value is x59?
        Code goes here...
endif
...
```



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF (compare instruction,parm1,condition,parm2)

```
if (cli,x_val,eq,X'55'),OR, Value is x55?
          (cli,x_val,eq,X'5D'),OR, Value is x5D?
          (cli,x_val,eq,X'51'),OR, Value is x51?
          (cli,x_val,eq,X'59') Value is x59?
          Code goes here...
endif
...
```

Remember to continue statements

X

X



The following are the condition mnemonics permitted by the SPMs

| Case | Condition Mnemonics | Meaning | Complements |
|------------------------------------|------------------------|-----------------|-------------|
| After compare instructions | H, GT | Higher, greater | NH, LE |
| | L, LT | Less than | NL, GE |
| | E, EQ | Equal | NE |
| After arithmetic instructions | P | Plus | NP |
| | M | Minus | NM |
| | Z | Zero | NZ |
| | O | Overflow | NO |
| After test under mask instructions | O | Ones | NO |
| | M | Mixed | NM |
| | Z | Zero | NZ |



- Provides simple selection for a given condition or instruction
- Multiple forms of IF statement:
 - IF CC=condition code

```
L R15,TRREME

TRT 0(1,R6),0(R15)

IF CC=0,OR,CC=2 Did the operation complete?

Yes - let's examine register 2 for TRT result...

ELSE

Data was examined but there is still more to go...

Examine the result and re-drive the instruction

ENDIF
```



- Provides iterative execution
- Multiple forms of DO statement:
 - DO ONCE/INF...
 - DO FROM=(Rx,i),TO=(Ry+1,j),BY=(Ry,k)...
 - DO WHILE=(condition)...
 - DO UNTIL=(condition)...
- Loop control keywords
 - DOEXIT uses IF-macro style syntax to exit the current DO or any containing labeled DO
 - ASMLEAVE unconditionally exit the current DO or any containing labeled DO
 - ITERATE requests immediate execution of the next loop iteration for the current or any containing labeled DO



HLASM Structured Programming - SPMs - DO...INF

- Provides iterative execution
- Multiple forms of DO statement:
 - DO INF → infinite loop

```
* Count the number of customers

XGR R2,R2 Clear the counter

DO INF

LT R3,CUSTOMER_NUMBER(R2) Get the next customer

DOEXIT (Z) Exit if at end

ALGFI R2,L'CUSTOMER_REC Increment counter

ENDDO
```



- Provides iterative execution
- Multiple forms of DO statement:
 - DO FROM → Counted loop note that this generates a BCT

```
* Output only top 10 customers

DO FROM=(R1,10)

L R3,CUSTOMER_NUMBER(R1) Get customer

OUTPUT_CUSTOMER CUST=R3 Call subroutine
ENDDO
```



- Provides iterative execution
- Multiple forms of DO statement:
 - DO FROM → Counted loop note that this generates a BXH

```
Clear process flag
XC FLAG, FLAG
IF
    (CLI, INPUT RECORD, NE, C'*')
                                            Not a comment?
    DO
            FROM = (R1, 1), TO = (R5, 72), BY = (R4, 1)
        LA R2, INPUT RECORD (R1)
                                   Get next character
        IF
           (CLC, 0(5,R2), EQ, =C'END'), AND,
                                                                X
            (CLI, FLAG, EQ, 2)
                                        Process the END record
            ASMLEAVE
                                        Leave the process loop
        ENDIF
    ENDDO
ENDIF
```



HLASM Structured Programming – SPMs – DO...Backwards...

- Provides iterative execution
- Multiple forms of DO statement:
 - DO FROM → Counted loop note that this generates a BXH

```
Clear process flag
XC FLAG, FLAG
IF
    (CLI, INPUT RECORD, NE, C'*')
                                             Not a comment?
            FROM= (R1,1), TO= (R5,72), BY= (R4,-1)
    DO
        LA R2, INPUT RECORD (R1)
                                             Get next character
            (CLC, 0(5,R2), EQ, =C'END'), AND,
        IF
                                                                 X
            (CLI, FLAG, EQ, 2)
                                         Process the END record
            ASMLEAVE
                                         Leave the process loop
        ENDIF
    ENDDO
ENDIF
```



HLASM Structured Programming – SPMs – SELECT...

Provides selective execution similar to an IF

```
WHEN (CLI, PREFIX, EQ, C'+')

MVC OUT_BUFFER (L'STMT_MACROEXP), STMT_MACROEXP

WHEN (CLI, PREFIX, EQ, C'=')

MVC OUT_BUFFER (L'STMT_COPYBOOK), STMT_COPYBOOK

WHEN (CLI, PREFXI, EQ, C' ')

MVC OUT_BUFFER (L'STMT_NORMAL), STMT_NORMAL

OTHERWISE

BAL R14, ERROR_ROUTINE Listing isn't correct...

J EXIT_ROUTINE

ENDSEL

* Output the type of statement encountered
```



Summary

- Control sections C-, R-, DSECTs, External dummy sections
 - CSECT, RSECT, DSECT, DXD, COM
- Subroutines
 - EXTRN / WXTRN
 - Conditional Sequential RLDs (CSRs)
- The location counter and USING statements
 - LOCTR
 - Labeled USINGs
 - Dependent USINGs
- Program Objects
- Controlling Listings
- Structured Programming Macros (SPMs)
 - IF...ELSEIF...ELSE...ENDIF
 - DO WHILE / UNTIL / FROM...ENDDO
 - SELECT...WHEN...OTHERWISE...ENDSEL
 - ASMLEAVE, DOEXIT



Where can I get help?

- z/OS V2R1 Elements and Features
 http://www.ibm.com/systems/z/os/zos/bkserv/v2r1pdf/#IEA
- HLASM Publications
 http://www.ibm.com/systems/z/os/zos/library/bkserv/v2r1pdf/#ASM
- HLASM Programmer's Guide (SC26-4941-06)
 http://publibz.boulder.ibm.com/epubs/pdf/asmp1021.pdf
- HLASM Language Reference (SC26-4940-06)
 http://publibz.boulder.ibm.com/epubs/pdf/asmr1021.pdf
- HLASM Toolkit Features User's Guide (GC26-8710-10) http://publibz.boulder.ibm.com/epubs/pdf/asmtug21.pdf
- z/Architecture Principles of Operation http://www.ibm.com/support/docview.wss?uid=isg2b9de5f05a9d57819852571c500428f9a