



Preparing SMSVSAM for HSM and Catalog

VSAM/RLS Performance and Tuning

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Overview



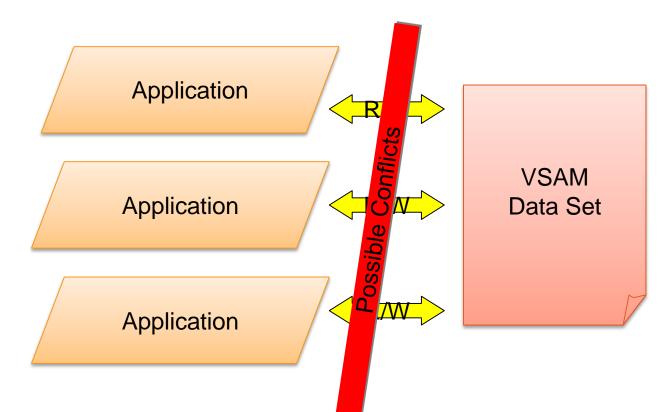
- Quick SMSVSAM Overview
- Why SMSVSAM for HSM or Catalog?
- Adding Support
 - Changes to SMSVSAM
 - Changes to HSM
 - Changes to Catalog
- Monitoring
 - SMSVSAM
 - HSM
 - Catalog



SMSVSAM Overview



Regular VSAM cannot have multiple updates simultaneously



Serialized at Data Set Level or application responsible for serialization.

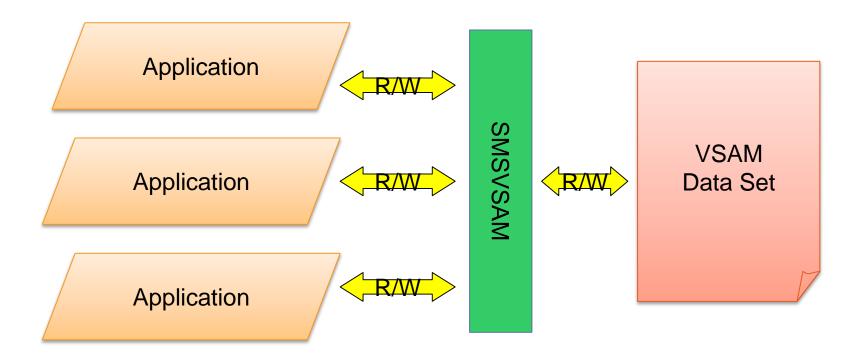
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SMSVSAM Overview



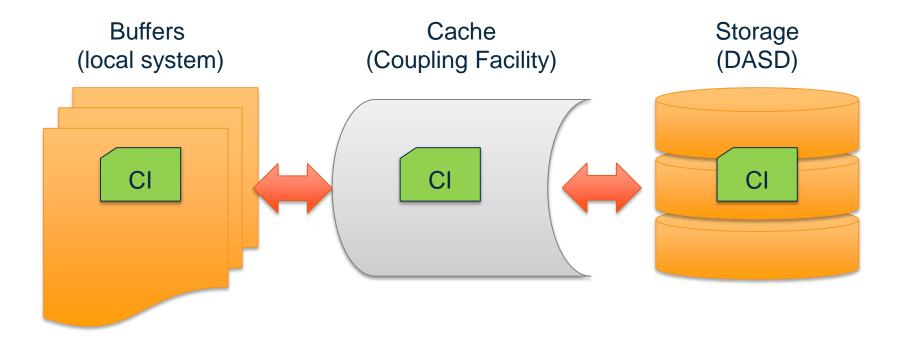
- SMSVSAM can handle all that serialization
- Serialization is at RECORD level instead of DATA SET





SMSVSAM Data Storage

- Data is staged across 3 layers
- Stored as full Control Intervals







SMSVSAM General Benefits

SHARE, Educate · Network · Influence

- Record-level serialization
 - Higher concurrent throughput
 - Multiple updates / erases / reads simultaneously
- Uses CF Locks
 - Maintains serialization, even past close in case of failure
 - Allows for transaction-level locking
- Cross-PLEX data sharing
 - Buffers and caching provide shortest path to data
 - Reduces I/O
- No single point of failure



Benefits for HSM CDS



- Reduces CDS contention during Primary Space Management and Automatic Backup
- Takes advantage of the data caching / reduces read time
- Avoids STRNO limits
- Performance Improvement*: (actual customer data, comparing NonRLS to RLS, 1yr elapsed)

Function	GB Moved Increase	Window Size Decrease
Auto Backup	33%	-25%
Migrate -> ML2	18%	-36%

- Actual Data #2 AUDIT Processing*
 - Before RLS: Could not complete in 24 hours
 - After RLS: Finished within 4 hours

*source: Glenn Wilcock, DFHSM Best Practices, Anaheim 2014, Session 15075

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Benefit for Catalog



- Reduced contention
 - Eliminates SYSIGGV2 'UCAT' ENQ contention
 - Plans to remove the SYSIGGV2 'sphere' ENQ
 - No need to split catalogs to lower contention
- Higher throughput
 - Significant improvement in elapsed time & CPU
 - Much shorter wait times
- Improved control
 - Suspend / resume ALL catalogs, PLEX-wide
 - Prevents un-serialized updates



RLS Catalog Performance Benchmark Test



	Elapsed Time	e (min)	CPU* (sec)		Deltas	
Test	Non-RLS	RLS	Non-RLS	RLS	Elapsed	CPU*
DELETE	80.42	8.42	1269.3	298.7	89.5%	77.0%
DEFINE	48.84	21.42	685.6	130.8	56.1%	80.9%
SEQ READ	7.40	5.03	65.1	75.2	32.0%	-15.5%
DIR READ (first sys)	26.77	20.33	94.0	109.6	24.1%	-16.6%
DIR READ (second sys)	26.86	20.29	95	109.9	24.5%	-15.7%

*CPU in GRS, CATALOG may see a small increase – best to compare per request

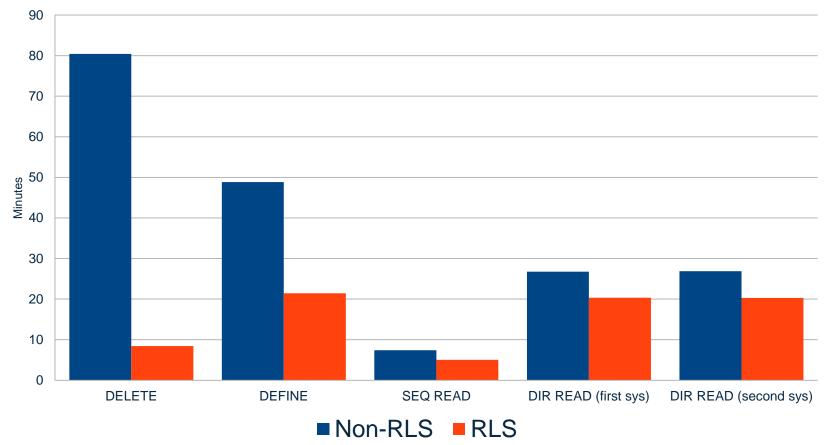
Test environment: Z10 2097 E12, 3 LPARs, 7 CPUs, 1 CF, z/OS 2.1 Catalog parms: TASKMAX=180, CISIZE(32768) and CISIZE(4096), STRNO(255) RLSABOVETHEBAR(NO) RLSCFCACHE(ALL) RLSMAXPOOLSIZE(100M) CF Cache size 1G Catalog RLS vs Catalog VLF at z/OS 2.1 Tests: 300,000 data sets, 100 jobs using 1000 data sets on each LPAR Source: "Unclog your Systems with z/OS 2.1 – Something New and Exciting for Catalog" by Terri Menendez, IBM SHARE San Francisco, Spring 2013 Session #12977, 12978



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Making the Changes

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Gather information

- SMSVSAM
 - Current buffer goal / size
 - Current cache names
- HSM
 - Size of CDS
 - Expected growth
- Catalog
 - Number of catalogs to use RLS
 - Size of catalogs
 - Expected growth



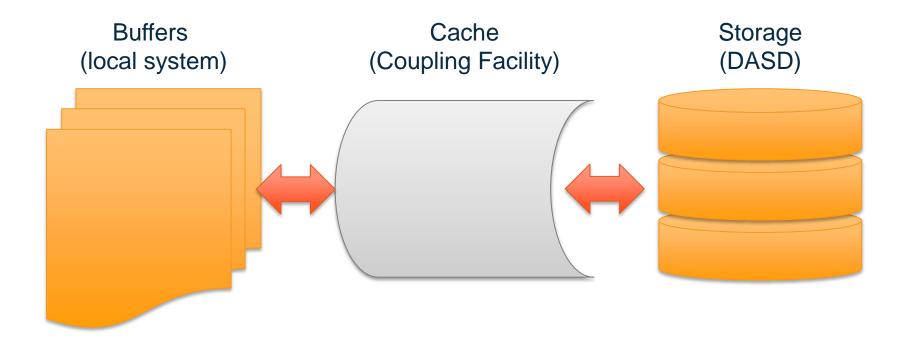




SMSVSAM Changes



- We'll probably need to expand Buffers and Cache
- May also want to expand the lock structure









Buffering

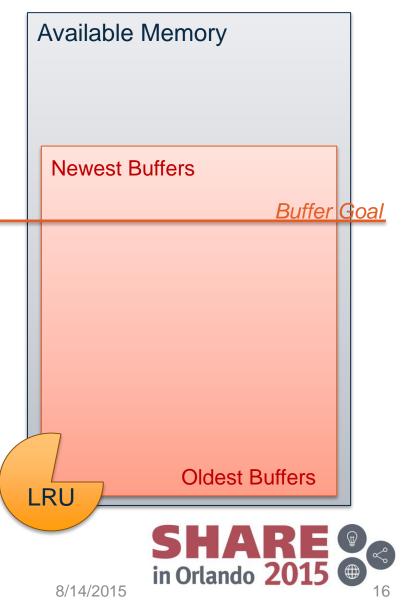
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Buffer Overview



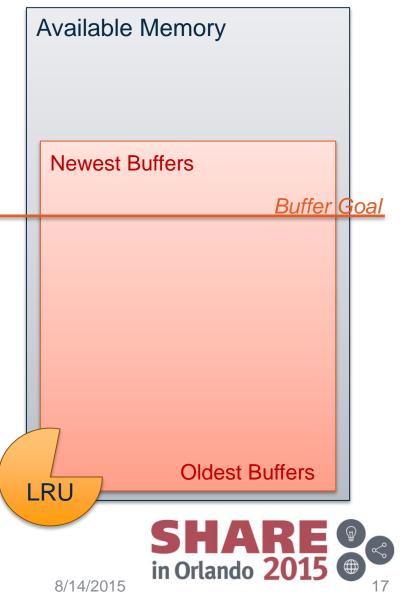
- Once buffers obtained, they are put on the top of the stack
- New buffers pile on top
- Re-referencing a buffer will pull it back to the top
- RLS Least Recently Used (LRU)
 routine purges old buffers
- As usage approaches or crosses the goal, LRU speeds up
- Below-the-bar: Panic @ 200% Above-the-bar: Panic @ 100%



Buffer Overview

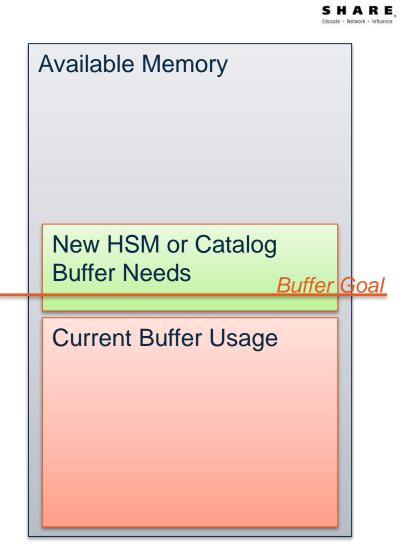
- What if I don't change my buffer sizes?
- More data fighting for the same space
- Buffers will be flushed more quickly
- RLS LRU CPU may rise
- Extra caching requests
- Extra I/O





Size Your Buffers

- Calculate Needed Size
 - No easy formula
 - How big are your CDS / Catalogs?
 - Prefer big enough to hold:
 - DS Index
 - Several Data CI
 - Note: buffers are shared
- Add New Size to Old Usage
 - If you are below goal, you may not need to change
 - If you are above goal already, you might want to increase even more





Example #1



Current usage: Goal: 750MB Usage: 500MB New Usage: 5 Catalogs: 100MB

- HSM CDS: 100MB
- No Need to Change Buffer Goal

Available Memory	
Buffer Goal: 75	OMB
500MB used	



Example #2

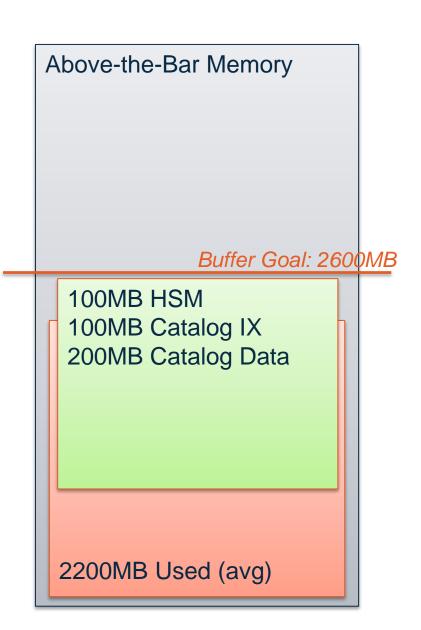
- Current usage:
 - Goal: 2000MB
 - Usage: 2200MB
- New Usage:
 - 5 Catalogs: 500MB
 - HSM CDS: 1000MB
- Solution 1:
 - Raise Goal to 3500MB
- That's a big change.. Consider:
 - How much HSM data is needed concurrently?
 - How much Catalog data is needed concurrently?
 - Does HSM run at the same time as your other workload?

Above-the-Bar Memory	
1500MB needed	
Buffer Goal: 200	0MB
2200MB Used (avg)	

Example #2

• Solution 2:

- HSM PSM runs overnight when other RLS data is offline / batch processing
- 100MB of concurrent data needed during day for HSM
- Catalog index total: 100MB
- Highly active catalogs, so make space for 50% of data
- Final Changes:
 - Raise goal to 2600MB
 - Monitor performance



Buffer Sizing



• **RLS_MAX_POOL_SIZE({nnnn|100})** in IGDSMSxx

- Below the bar GOAL
- One size for all LPARs
- Set by first system to start SMSVSAM
- 10 to 999, Recommendation: 850MB

• RLSABOVETHEBARMAXPOOLSIZE({ALL, size}) in IGDSMSxx

- Above the bar GOAL
- Can set per-system using (SYS1,size1),(SYS2,size2)
- 500MB to 2TB.. Recommended to keep < 32GB
- Relies on real storage
- **RLSFIXEDPOOLSIZE({ALL,size})** in IGDSMSxx
 - Amount of the total real storage, both above and below the bar, that will be permanently fixed (pinned)
 - Does not allow paging of the buffers
- RLS Above the 2-GB Bar (YES | NO) in Data Class definition
 - Must be YES to use above the bar buffering
 - Defaults to NO







Caching

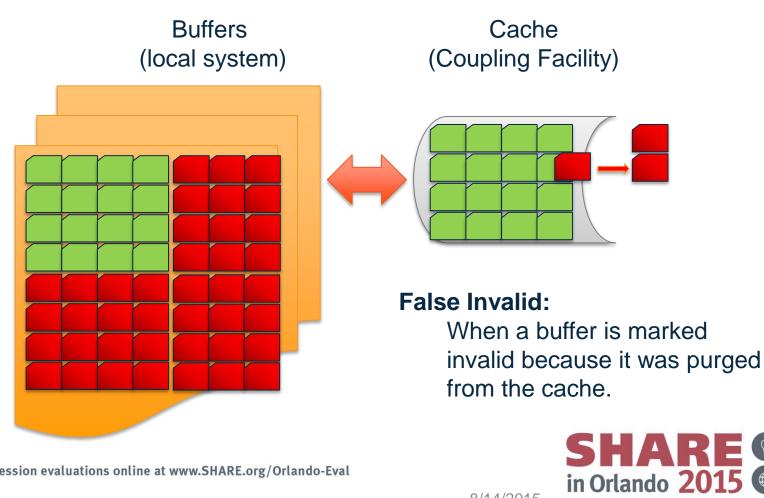
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Cache Needs



- Cache needs multiply with more systems
- Undersized cache reduces buffer effectiveness

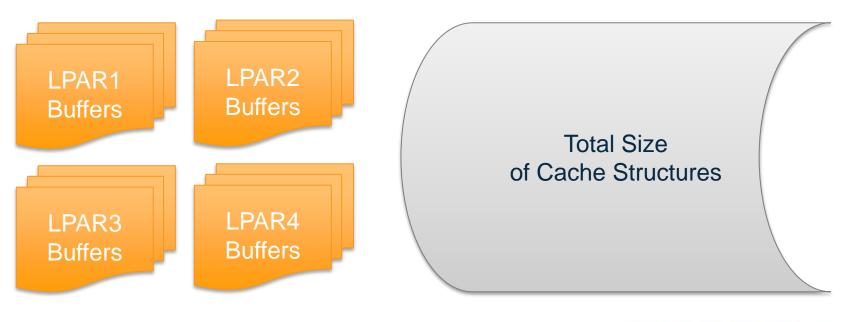


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Cache Needs



- "Best" Sizing:
 - Buffer Size x Number of Systems
- Optimized Sizing:
 - Sized of Shared Data + Size of Unique Data from each LPAR



8/14/2015

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RLS Caching Modes



ALL Data CIs and Index CIs stored READ or WRITE will add CI to cache

NONE | Cache on index Cis READ or WRITE will add CI to cache

UPDATES Data Cis and Index CIs stored Only WRITES will update cache

DIRONLY | No CI data is stored READ or WRITE will update interest



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Inside a Cache



Directory Entry

- Holds control information
- Holds interest information
- One per Cl
- Relatively Small



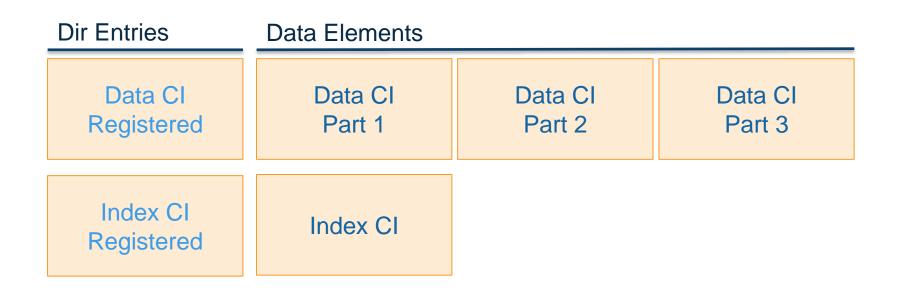
- Holds the stored data
- Up to 2k in size
- Many per CI (depending on CISIZE)
- Ex: 6k CI would require 3 DE







Mode: ALL or UPDATESONLY

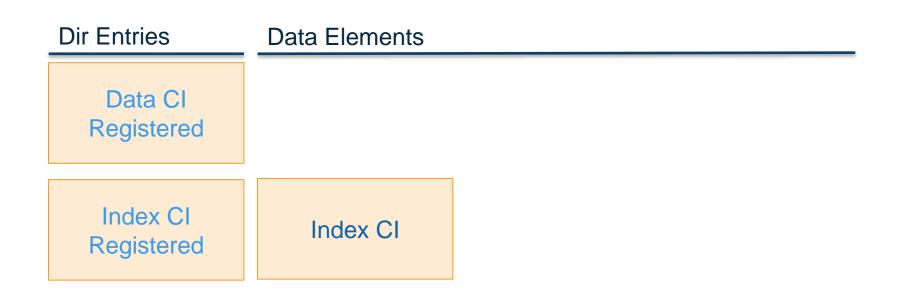








Mode: NONE



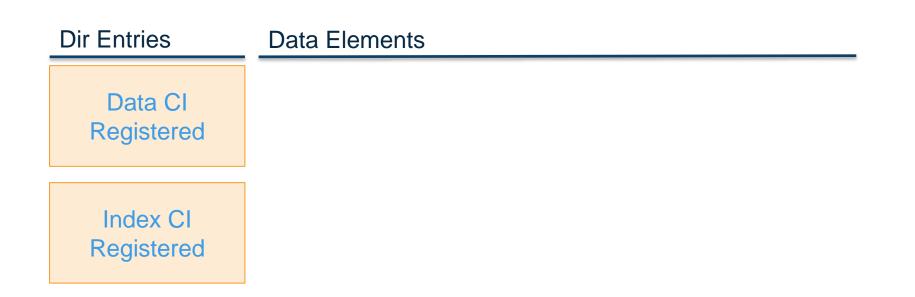








Mode: **DIRONLY**







Caching Modes



- · Can have a big impact on how cache is used
- Can reduce cache needs
- RLS CF Cache Value in Data Class
 - -A = AII
 - N = None
 - U = Updates
 - D = Directory Only
- Requires RIs_MaxCFFeatureLevel(A)
- Toleration: OA36443, OA36415



Create New Structures



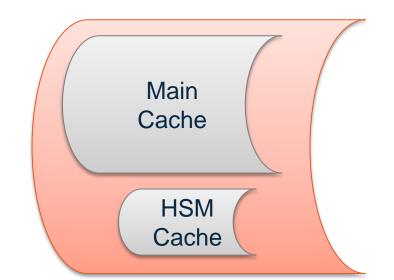
- Recommendation: Use new cache structures
 - Allows for separation between workload and CDS
 - Keeps load off of current structures
 - Makes sizing easier
- Define in CFRM Policy
 - Note: RLS does not grow structures, so use INITSIZE=MAXSIZE
- Update cache sets / storage classes in SMS via ISMF
 - May want to create new cache set
 - May want to create new Storage Class
- Recommendation: Separate Catalog / HSM / Work SC



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Caching Example

- Two Systems
- RLS CF Cache = ALL
- Shared Data: 2048MB
- Unique Data:
 - System1: 400MB
 - System2: 2448MB
- New HSM Data = 100MB



- Total Cache Size:
 2048 (shared) + 400 + 2448 (unique) = 4896MB (~4.8GB)
- CICS/Online Cache Size = 4896MB
- HSM Cache Size = 100MB









Lock Structures

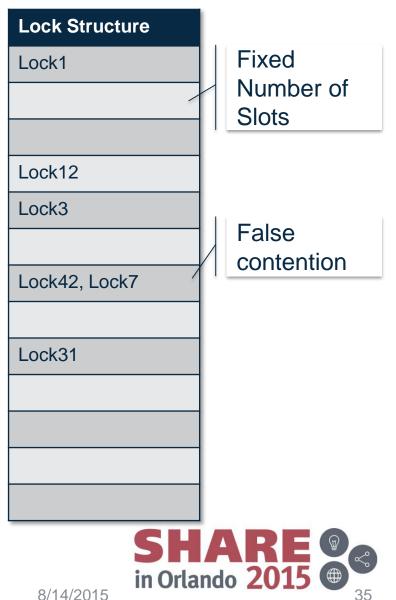
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Lock Structure Considerations



- Lock Structure
 - Divided into slots
 - Uses hash to place lock
 - Too few slots results in hash collisions
 - "False contention"
- Recommendation: Use secondary lock structures
 - TEST, WORK, HSM, Catalog
 - Prevents overload in one from infecting others
- Monitor False Contention rates



Multiple Lock Structures



- Allows separation of workloads
- SMS allows up to 256 different lock sets
- Can use any name. IGWLOCK01 is easy
- Update SMS:
 - Storage Class \rightarrow SMS Lock Set \rightarrow Lock Structure
- Holds all record locks for that data set (does not hold Component or Special locks)





SHCDS considerations

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SHCDS Considerations



- Share Control Data Sets (SHCDS)
- Maintains data set level information in case of failure
- Used to build Lost Locks information
- Hold:
 - Data set / cache set / lock structure pairing info
 - Data set / region pairings
- You may want to grow the SHCDS **only** if you are adding a lot of data sets (lots of catalogs)

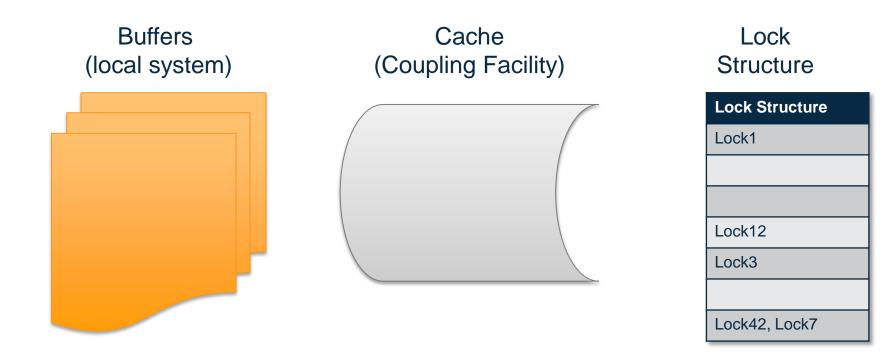


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SMSVSAM Changes Summary

- Examined each in detail
- Made changes based on new information









HSM Changes

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HSM Changes



- Update HSM Data Sets:
 - IDCAMS ALTER cdsname LOG(NONE)
 - Consider SHAREOPTIONS(2 3) if you want non-RLS read access while open
 - Make sure Storage Class points to a Cache Set
- Startup Keyword **CDSSHR** = {YES | **RLS** | NO}
- Make sure SYSVSAM ENQs are SYSTEMS (not RNL excluded)





Catalog Changes

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Catalog Changes



- To ENABLE for a single catalog:
 - IDCAMS ALTER ucat LOG(NONE)
 - Ensure that Storage Class points to Cache Set
 - F CATALOG, RLSENABLE(ucat)
 - IEC352I MODIFY CATALOG cat.name TO STATE RLSENABLE SUCCESSFUL
 - F CATALOG, RLSQUIESCE(ucat)
- Notes:
 - Requires that SMSVSAM be up and active
 - Only available on z/OS 2.1 and up



IDCAMS Tools Support



- REPRO, PRINT, IMPORT, EXPORT supported
- To use, specify
 - **RLSSOURCE**(YES | NO | QUIESCE)
 - RLSTARGET(YES | NO | QUIESCE)
- Options:
 - YES use RLS mode to access data set
 - NO use Non-RLS (regular VSAM) to access data set
 - QUIESCE Use Non-RLS mode, but QUIESCE first.



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Monitoring and Tuning

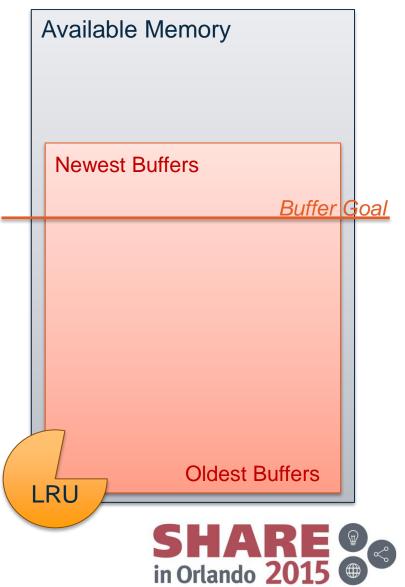
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Monitoring Buffers

- Measurements to Watch:
 - LRU Mode
 - Buffer usage vs buffer goal
 - BMF Hit Rate
 - Reclaim Rate
- SMF 42 subtype 15, 16, 19
- RMF Monitor III
 - RLSSC panel
 - RLSLRU panel
- OMEGAMON XE





SMF42 Buffer-related Fields



- Subtype 19 LRU Summary
 - Available at SYSPLEX-level or system-level
 - Split into below-the-bar and above-the-bar sections
 - Buffer goal: SMF42JQI
 - Buffer used: SMF42JQM
 - Accelerated/panic mode rate: SMF42JPH / SMF42JPG
- Subtype 15, 16 Storage Class & Data Set
 - Also split into PLEX / system
 - Also split into below / above the bar
 - Additionally split into SEQ / DIR
 - BMF Hit rate = SMF42FII / SMF42FIG



RMF MON III - RLSLRU





OMEGAMON XE for Storage

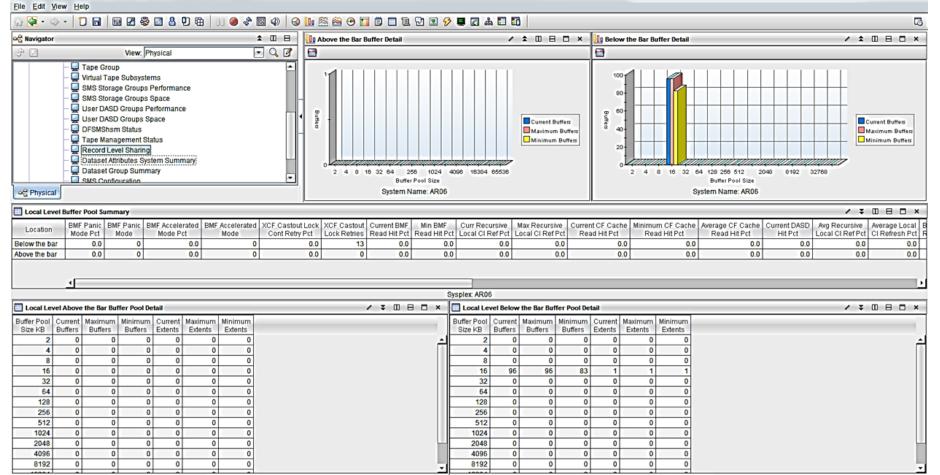


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IGWLOCK00 DEMOPLX	0	0.00 0.0	0.00 0.0	0.00 0.0		0.00		0.0 0.00	0.0 0	0.00 0.0	0.00	
Buffer LRU Summary Location System BMF	В	uffer LRU Summa				_						
Above the bar DEMOPLX Below the bar DEMOPLX		Location	System Sysplex Name	BMF Panic Mode Pct		F Panic lode	BMF Accelerate Mode Pct	d BMF Accele Mode		F Castout L Cont Retry P		CF Ca
	Ø	Above the bar	DEMOPLX	0.0		0	0.	0	0		0.0	
4	Ø	Below the bar	DEMOPLX	0.0		0	0.	0	0	1	5.6	
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4	() Hut) Time: Fri, 08/01/2014 07:27 PM		Server Available			RLS Sum	nary - dem21lnx.democentral.it	om.com - Vickie Dault			Þ
Complete your sessi			www.SHARE.org				8/14/20/	in Or		2015		

OMEGAMON XE for Storage RLS Workspaces - Buffer LRU System Detail Workspace





Sysplex: AR06

Sysplex: AR06

Buffer Pool Statistics for a single z/OS image

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- 0 -X



Monitoring the Cache

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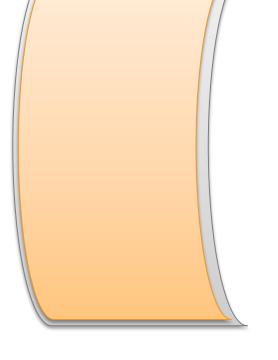


Monitoring the Cache

- Caches:
 - Entries claimed until structure is full
 - Once full, entries are RECLAIMed when needed
 - Even after close, cache entries are "used" until they are reclaimed
 - Reclaims based on age like buffers
- As a results:
 - Percent Full is not relevant
 - Reclaim Rate is not relevant
- Measurement to Watch: BMF FALSE INVALIDS

 If > 5%, time to raise cache sizes









RMF MON III – RLSSC & RLSDS



	RMF V	2R1 VSA	AM RLS /	Activi	ty - C	APTKEN1	Line	1 of 5
Samples: 100	Systems	:1 Da	ate: 03	/08/14	Time:	15.23.20	0 Range: 10	0 Se
LRU Status : Contention % : False Cont % :	0.0 /	ood 0.0						
Stor Class Acc							BMF False Inv%	
RLSSTOR								
Below 2GB DI	IR 0.00	0 0.00	0.0	0.0	0.0	0.0	0.00	0.00
SI	EQ 0.00	0 0.00	0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB DI	IR 0.00	1 2254	90.4	9.0	0.6	100	0.10	1137
SI	EQ 0.00	0 0.00	0.0	0.0	0.0	0.0	0.00	0.00



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Monitoring the Lock Structure

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Monitoring Lock Structures

False Contention

- Causes performance degradation
- Two locks incorrectly stored in the same spot
- Indicates that Lock Structure is too small
- Target rate < 5%</p>
- To review false contention
 - Health Check
 - D SMS,CFLS command
 - RMF CF activity report
 - SMF 42 subtype 17









Health Check



- VSAMRLS_CFLS_FALSE_CONTENTION
 - Average of the last hour
 - Samples taken every second
 - Has configurable threshold

HZS0002E CHECK(IBMVSAMRLS,VSAMRLS_CFLS_FALSE_CONTENTION): IGWRH0131E FALSE LOCK CONTENTION HAS BEEN DETECTED. THE CURRENT FALSE CONTENTION RATE IS 7.120%.



D SMS,CFLS



		splay SMS,C	•	•		
		GWLOCK00 VE S :10358 USE		5425A4759D195	5 SIZ Total Av	ailable Locks
				C6425AD06E2	L9A SIZE:4096	K
RECORD T	ABLE ENTRIE	S:10358 USE	D:2			
LOCK STR	UCTURE MODE	: DUPLEXED	STATUS: E	NABLE		
System	Interval	LockRate	ContRate	FContRate	WaitQLen	
SYSTEM1	1 Minute	585.6	0.000	0.031	0.00	
SYSTEM1	1 Hour	95.1	0.000	0.000	0.00	
SYSTEM1	8 Hour	11.9	0.000	0.000	0.00	
SYSTEM1	1 Day	8.9	0.000	0.000	0.00	
*** No o	ther system	ns provided	data			

```
LockRate = number of lock requests per second
CONTRATE = % of lock requests globally managed
FCONTRATE = % of lock requests falsely globally managed
WaitQLen = Average number of requests waiting for locks
```



Monitoring False Contention



- RMF Mon III Structure details on IGWLOCK00
 S.7 "Coupling Facility Activity"
- RMF Mon III RLSSC and RLSDS
- SMF 42 subtypes 15,16,17
 - SMF42FOA Number of Lock Requests
 - SMF42FOC Number of Lock Requests causing False Contention



RMF MON III – RLSSC & RLSDS



	RMF V2R1	L VSA	MRLS	Activi	ty - C	APTKEN1	Line	1 of 5
Samples: 100 S	ystems: 1	L Da	te: 03,	/08/14	Time:	15.23.20	0 Range: 10	0 Se
< 2	GB / > 20	βB						
LRU Status : Goo	d / Good	ł						
Contention % : 0	.0 / 0.	0						
False Cont % : 0	.0 / 0.	0						
Stor Class Access	Resp		Read	d			BMF	Write
	Time	Rate	BMF%	CF%	DASD%	Valid%	False Inv%	Rate
RLSSTOR								
Below 2GB DIR	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB DIR	0.001	2254	90.4	9.0	0.6	100	0.10	1137
SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00

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Monitoring Lock Structure



- If structure fills, you'll see these messages:
 - IEC161I 248-ccc indicates failure during OPEN
 - IGW326W *** Warning *** DFSMS SMSVSAM RECORD TABLE IN lockstructurename IS percent % FULL.
- Once full, you'll need to increase the size and rebuild
- Track usage via D SMS,CFLS command

D SMS,CFLS – Monitoring Lock Structure



		splay SMS,C	•	*		
				425A4759D195	5 SIZ Total Ava	ailable Locks
RECORD TA	BLE ENTRIE	S:10358 USE	D:2			
SECONDARY	STRUCTURE	:IGWLOCK00	VERSION:CC	C6425AD06E21	L9A SIZE:4096k	r •
RECORD TA	BLE ENTRIE	S:10358 USE	D:2			
LOCK STRU	ICTURE MODE	: DUPLEXED	STATUS: E	NABLE		
System	Interval	LockRate	ContRate	FContRate	WaitQLen	
SYSTEM1	1 Minute	585.6	0.000	0.031	0.00	
SYSTEM1	1 Hour	95.1	0.000	0.000	0.00	
SYSTEM1	8 Hour	11.9	0.000	0.000	0.00	
SYSTEM1	1 Day	8.9	0.000	0.000	0.00	
*** No ot	her system	s provided	data			

LockRate = number of lock requests per second CONTRATE = % of lock requests globally managed FCONTRATE = % of lock requests falsely globally managed WaitQLen = Average number of requests waiting for locks



Using SHCDS



IDCAMS LISTSUBSYS(ALL)

Show who holds locks

SHCDS LISTSUBSYS(ALL) ----- LISTING FROM SHCDS ----- IDCSH03

SUBSYSTEM NAME	STATUS				LOCKS WAITING	
IYNXV	ONLINEACTIVE	NO		200	0	0
DATA SETS IN	LOST LOCKS		0			
DATA SETS IN	NON-RLS UPDATE S	STATE	0			
TRANSACTION	COUNT		1			
IYNX5	ONLINEACTIVE	NO		0	3	0
DATA SETS IN	LOST LOCKS		0			
DATA SETS IN	NON-RLS UPDATE S	STATE	0			
TRANSACTION	COUNT		1			
SMSVSAM	BATCHACTIVE	NO		0	0	0
DATA SETS IN	LOST LOCKS		0			
DATA SETS IN	NON-RLS UPDATE S	STATE	0			
TRANSACTION	COUNT		0			
IDC0001I FUNCTIO	N COMPLETED, HIGH	IEST CONDITION	CODE W	AS 0		



Using SHCDS



IDCAMS LISTSUBSYSDS(ALL)

SHCDS LISTSUBSYSD	•	•	ð4			
SUBSYSTEM NAME	RETLKØ5A	SUB	SYSTEM STAT	TUSONLIN	EFAILED	
DATA SET NAME / CACHE STRUCTURE / LOCK STRUCTURE	RETAINED LOCKS	LOST LOCKS	Locks Not Bound	RECOVERY REQUIRED	NON-RLS UPDATE PERMITTED	PERMIT FIRST TIME SWITCH
SYSPLEX.KSDS.RETLK CACHE01 TESTLOCK01 IDC0001I FUNCTION	YES	NO HIGHEST	NO CONDITION	NO CODE WAS Ø	NO	NO

Show region / data set pairings



Monitoring HSM



- DFSMSrmm Report Generator
 - Generate reports via ISPF panels
 - Customizable reports
 - Reports based on FSR, WWFSR, Inventory via DCOLLECT
 - See sessions 17592, 17717 (Monday) for more information

• FSRSTAT

- REXX program that formats FSR records
- Shipped with HSM
- Easy to use and customize



Monitoring Catalog



F CATALOG, ALLOCATED

IEC348I ALLOCATED CATALOGS 118 *CAS***********************************	***	k *
* FLAGS -VOLSER-USER-CATALOG NAME	%	
* YSU-R- XP0301 0001 BOHLING.RLS.UCAT	1	*
* Y-I USRPAK 0001 SYS1.MVSRES9.MASTCAT	1	*
***************************************	***	**
* Y/N-ALLOCATED TO CAS, S-SMS, V-VLF, I-ISC, C-CLOSED, D-DELETED,		*
* R-SHARED, A-ATL, E-ECS SHARED, K-LOCKED, U-RLS, W-SUSPENDED		*
*CAS***********************************	***	**



Monitoring Catalog

F CATALOG, REPORTS, PERFORMANCE

*	Statistics since 21:45:1	2.84 on 03	/07/2014	*
*	CATALOG EVENT	COUNT	AVERAGE	*
*	Entries to Catalog	840	27.624 MSEC	*
*	BCS ENQ Shr	636	0.036 MSEC	*
*	BCS ENQ Shr Sys	380	9.266 MSEC	*
*	BCS ENQ Excl	5	0.013 MSEC	*
*	BCS ENQ Excl Sys	5	9.830 MSEC	*
*	BCS DEQ	1,255	3.879 MSEC	*
*	BCS Allocate	13	58.862 MSEC	*
*	SMF Write	35	0.020 MSEC	*
*	CAS MLA Lock	1	0.288 MSEC	*
*	VVDS Format	2	10.147 MSEC	*
*	MVS Allocate	6	127.281 MSEC	*
*	SMS Active Config	3	0.073 MSEC	*
*	SYSVSAM S ENQ Excl	13	36.419 MSEC	*
*	SYSVSAM S DEQ	13	8.585 MSEC	*
*	SYSVSAM D ENQ Shr	13	4.689 MSEC	*
k	SYSVSAM D DEQ	13	10.067 MSEC	*





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Monitoring Catalog



F CATALOG, REPORT, CATSTATS

	ADDS	UPDATES	GETS	GETUPD	DELETES	BUFNI	BUFND	STRNO	*
									*
BC	HLING.R	LS.UCAT							*
	4	1	36	3	0	1	2	180	*
SY	S1.MVSR	ES9.MASTCA	λT						*
	2	3	2,352	3	0	4	4	2	*







Benefits of RLS

Gathered Information

Calculated Changes

Made Changes

Monitored Performance

8/14/2015

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Preparing SMSVSAM for HSM and Catalog

VSAM/RLS Performance and Tuning

Neal Bohling

Session #17832

IBM Development bohling@us.ibm.com

CTTEBRATIVO 60 *YEARS* OF SHARE



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Appendix - Lock Structures - Sizing



- Size requirement depends on how many locks are active
- Lock_Structure_Size = 10M * Number_Of_Systems * Lock_Entry_Size
 - Where Lock_Entry_Size, depends on the MAXSYSTEM
 - MAXSYSTEM <= 7 Lock_Entry_Size = 2
 - MAXSYSTEM >= 8 & <24 Lock_Entry_Size = 4
 - MAXSYSTEM >=24 & <=32 Lock_Entry_Size = 8
- Example: MAXSYSTEM = 16, with 8 systems in sysplex:
- Lock_Structure_Size = 10M * 4 * 8 = 320M
 - Note: Minimum size of 13M is recommended
 - You can also use CFSIZER to help determine appropriate size: http://www-947.ibm.com/systems/support/z/cfsizer/vsamrls/



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