

# Look What I Found Under The Bar!



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Session 17696

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# Agenda

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- Overview of Language Environment storage areas
  - Control blocks, stack, heap
  - Which can you control?
- How to control Language Environment Storage
- Tuning Storage
- More advanced tuning
- Sources of Additional Information



# Language Environment Storage Areas

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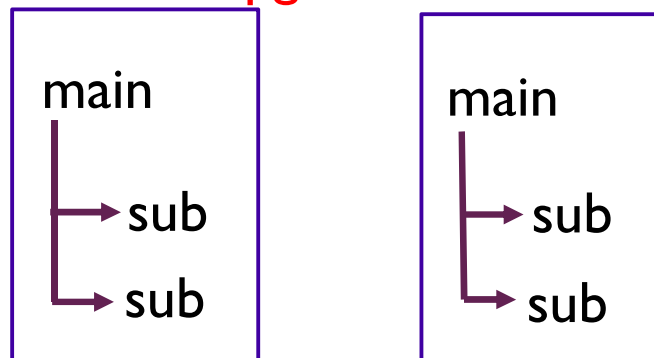
- Language Environment Control Blocks
  - Region level
    - Normally 1 region per address space
  - Process level
    - Normally 1 process per address space
  - Enclave level
    - Potentially many per address space
  - Thread level
    - Potentially very many per address space

# Language Environment Storage Areas

region - address space

process - application

enclave - pgm - enclave



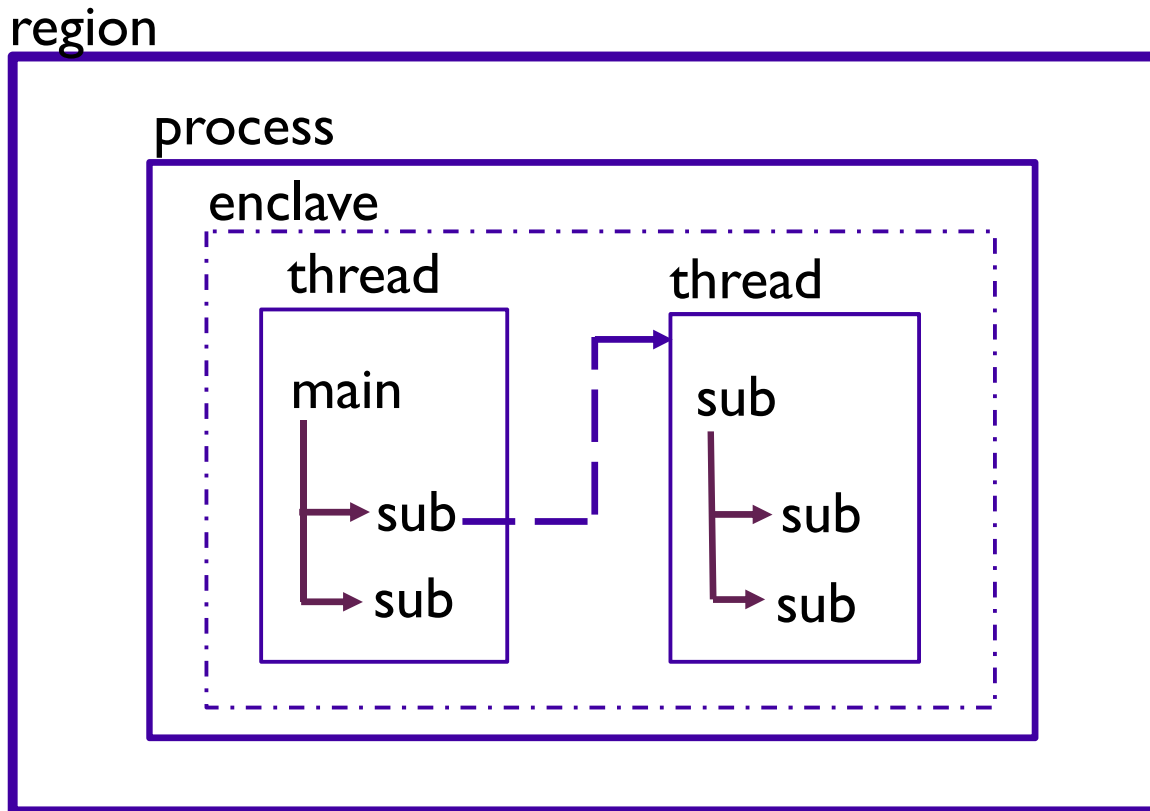


# Language Environment Storage Areas

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- Language Environment Enclave
  - Every “main” program is a new enclave
  - Every “link” is a new enclave
  - Contains
    - CEECAA, CEEEDB, CEEOCB, stacks, heaps, environment variables, and much more!
  - “Extra” enclaves are expensive both to initialize and in storage usage.

# Language Environment Storage Areas





# Language Environment Storage Areas

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- Language Environment Thread
  - Created by
    - pthread\_create() (C/C++ pthread support)
    - ATTACH statement (PL/I multithreading support)
  - Contains
    - CEECAA, stacks, and a little more
  - Threads are not nearly as expensive as enclaves.





# Language Environment Storage Areas

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- Language Environment Stacks
  - Stacks
    - Last In, First Out structures
    - Allow programs to be reentrant
    - Thread level structures
      - “Main” programs have separate stacks
      - “Linked” programs have separate stacks
      - pthreads have separate stacks



# Language Environment Storage Areas

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- Language Environment supports 2 independent stacks
  - User stack – (poorly named)
    - Used by user programs and Language Environment
  - Library stack
    - Used “rarely” by Language Environment
    - Always below the 16M line



# Language Environment Storage Areas

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- DATA in stacks
  - “Chunks” are called stack segments
    - Made up of 1 or more DSAs
  - DSA – Dynamic Save Area
    - Also called a “stack frame”
  - DSAs contain
    - Register Save Area (RSA)
    - NAB – Next Available Byte
    - Automatic (local) variables
      - C – int i;
      - PL/I – declare i fixed;
      - NOT COBOL WORKING-STORAGE
        - COBOL LOCAL-STORAGE in stack



# Language Environment Storage Areas

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- Language Environment Heaps
  - Heaps
    - Completely random access
    - Allows storage to be dynamically allocated at runtime
  - Enclave level control structures
    - Each 'main' has a separate heap
    - Each 'link' causes a separate heap
    - pthreads share a single heap for all threads



# Language Environment Storage Areas

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- Language Environment Heaps
  - Four independently maintained sets of heap segments all with similar layouts:
    - User Heap
      - COBOL WORKING-STORAGE
      - C/C++ (malloc or operator new)
      - PL/I dynamic storage (allocate)
    - LE Anywhere Heap
      - COBOL and LE above the line CBs
    - LE Below Heap
      - COBOL and LE below the line CBs
    - Additional Heap
      - Defined by the user



# Controlling Storage

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- Run-time options dealing with stacks
  - STACK(init,inc,ANY|BELOW,KEEP|FREE,dsInit,dsInc)
    - Init - Initial size of storage "chunk" allocated and managed by LE for user stack
    - Inc - When init is full, size of next storage "chunk" (increment)
    - ANY|BELOW - Location of storage
      - ANY Anywhere in 2G virtual storage
      - Below Always below 16M line
        - Required when all31(OFF)
    - KEEP|FREE - What to do when done with inc
      - KEEP Do not free the storage "chunks"
      - FREE Free the storage "chunks"
    - DsInit - Initial size of storage "chunk" (XPLINK)
    - DsInc - When initial full, size of next "chunk" (XPLINK)



# Controlling Storage

---

- Run-time options dealing with stacks
  - LIBSTACK(init,inc,KEEP|FREE)
    - Init - Initial size of storage "chunk" allocated and managed by LE for library stack
    - Inc - When init is full, size of next storage "chunk" (increment)
    - KEEP|FREE - What to do when done with inc
      - KEEP Do not free the storage "chunks"
      - FREE Free the storage "chunks"

NOTE: No ANY|BELOW, LIBSTACK always below the 16M line



# Controlling Storage

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- Run-time options dealing with stacks
  - THREADSTACK(ON|OFF,init,inc,ANY|BELOW,KEEP|FREE,dsInit,dsInc)
    - ON|OFF – Whether or not to use THREADSTACK for pthreads
    - Init - Initial size of storage “chunk” (like STACK)
    - Inc - Increment size of storage “chunk” (like STACK)
    - ANY|BELOW - Location of storage
      - ANY Anywhere in 2G virtual storage
      - Below Always below 16M line
        - Required when all31(OFF)
    - KEEP|FREE - What to do when done with inc
      - KEEP Do not free the storage “chunks”
      - FREE Free the storage “chunks”
    - DsInit, Dsinc – XPLINK “chunk” sizes





# Controlling Storage

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- Runtime options dealing with the heaps
  - HEAP(init,inc,ANY|BELOW,KEEP|FREE,int24,inc24)
    - User heap - mostly application use
    - init - Initial size of the "chunk" of storage obtained to be managed by LE for user heap
    - Inc - When initial "chunk" is full, size of next "chunk" (minimum)
    - ANY|BELOW - Location of "chunk"
      - Not sensitive to ALL31 setting
    - KEEP | FREE - What to do when done with the increment when empty
      - KEEP - Do not free the storage "chunks"
      - FREE - Free the storage "chunks"
    - int24 - Initial size of the "chunk" of storage obtained
      - (if ANY specified but BELOW requested (minimum))
    - inc24 - Size of next "chunk"
      - (if ANY specified but BELOW requested (minimum))



# Controlling Storage

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- Runtime options dealing with the heaps...
  - ANYHEAP(init,inc,ANY|BELOW,KEEP|FREE)
    - Thread stack storage lives in anyheap!!! Tune if multi-threaded
    - LE use - normally above the line
    - init - Same as HEAP.
    - inc - Same as HEAP. (minimum)
    - ANY | BELOW - Location of storage
    - KEEP | FREE - Same as HEAP
  - BELOWHEAP(init,inc,KEEP|FREE)
    - LE use - always below the line
    - init - Same as HEAP.
    - inc - Same as HEAP. (minimum)
    - KEEP | FREE - Same as HEAP



# Initializing Storage

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- STORAGE(getheap, freeheap, dsa alloc)
  - Getheap – Initialize heap storage
    - NONE – no overhead
    - One byte hex value to initialize storage with when heap element obtained
      - 00 similar to WSCLEAR option
      - Relatively low overhead
  - Freeheap – Overwrite heap storage
    - NONE – no overhead
    - One byte hex value to initialize storage with when heap element freed
      - Useful for debug purposes or security
      - Relatively low overhead



# Initializing Storage

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- STORAGE(getheap, freeheap,dsa alloc)
  - DSA alloc – Initialize stack storage
    - NONE – No initialization – no overhead
    - CLEAR – Entire unused initial stack segment is cleared just before the main program is given control – low overhead
    - A one byte hex value to initialize storage with when stack frame (DSA) is obtained
      - EXTREMELY HIGH OVERHEAD
      - EXTREMELY HIGH OVERHEAD
      - EXTREMELY HIGH OVERHEAD

# Initializing Storage

- Simple program that makes lots of calls

- STORAGE(,,none)

```
-----  
-                               REGION          --- ST  
- STEPNAME PROCSTEP PGMNAME      CC      USED      CPU TIME  
- GO                STORRTO      00      60K      0:00:00.56
```

- STORAGE(,,00)

```
-----  
-                               REGION          --- ST  
- STEPNAME PROCSTEP PGMNAME      CC      USED      CPU TIME  
- GO                STORRTO      00      60K      0:00:02.15
```

- STORAGE(,,CLEAR)

```
-----  
-                               REGION          --- ST  
- STEPNAME PROCSTEP PGMNAME      CC      USED      CPU TIME  
- GO                STORRTO      00      60K      0:00:00.57
```

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# Initializing Storage

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- Best ways to ensure the proper initial value for your variables
  - Use compiler initialization
  - Set them prior to use in your program



# Tuning storage

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- Objectives
  - Use as little storage as possible
  - Have program run as efficiently as possible
- The above objectives are often at odds with each other. (But not always)
- One way to make a program run faster is to “throw” more storage at it.
  - Care must be taken to use storage wisely
  - Much of what we will talk about can be done without recompiling or reworking the program.



# Tuning storage

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- Simple example
  - In a test environment (not production) use the RPTSTG run-time option.
  - A report will be generated describing the storage used by the program.
  - This information can be used to assist with better settings of Language Environment run-time options





# Tuning storage

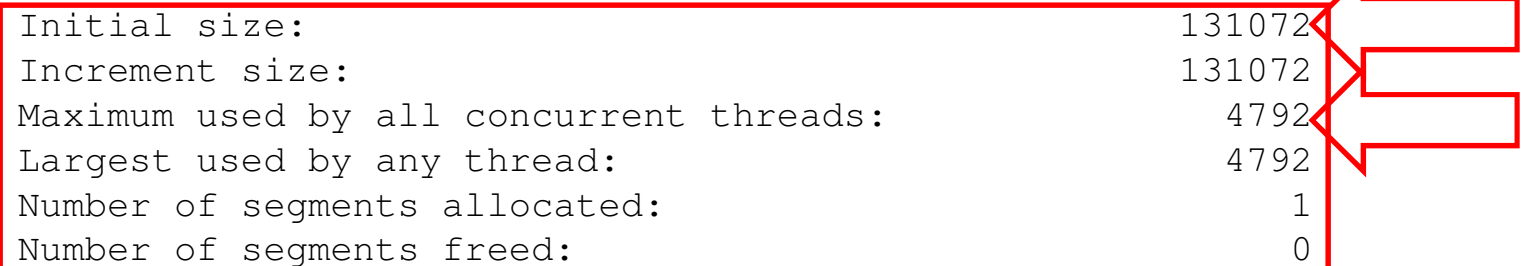
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## ■ Simple example

Storage Report for Enclave main 02/07/11 5:12:26 PM  
Language Environment V01 R12.00

### STACK statistics:

Initial size:	131072
Increment size:	131072
Maximum used by all concurrent threads:	4792
Largest used by any thread:	4792
Number of segments allocated:	1
Number of segments freed:	0



### THREADSTACK statistics:

Initial size:	0
Increment size:	0
Maximum used by all concurrent threads:	0
Largest used by any thread:	0
Number of segments allocated:	0
Number of segments freed:	0

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# Tuning storage

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- Simple example...

LIBSTACK statistics:

Initial size:	4096
Increment size:	4096
Maximum used by all concurrent threads:	0
Largest used by any thread:	0
Number of segments allocated:	0
Number of segments freed:	0

THREADHEAP statistics:

Initial size:	4096
Increment size:	4096
Maximum used by all concurrent threads:	0
Largest used by any thread:	0
Successful Get Heap requests:	0
Successful Free Heap requests:	0
Number of segments allocated:	0
Number of segments freed:	0



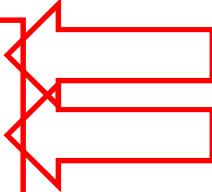
# Tuning storage

---

- Simple example...

HEAP statistics:

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	3328
Successful Get Heap requests:	4
Successful Free Heap requests:	2
Number of segments allocated:	1
Number of segments freed:	0



HEAP24 statistics:

Initial size:	8192
Increment size:	4096
Total heap storage used (sugg. initial size):	0
Successful Get Heap requests:	0
Successful Free Heap requests:	0
Number of segments allocated:	0
Number of segments freed:	0



# Tuning storage

---

## ■ Simple example...

ANYHEAP statistics:

Initial size:	16384
Increment size:	8192
Total heap storage used (sugg. initial size):	616
Successful Get Heap requests:	6
Successful Free Heap requests:	2
Number of segments allocated:	1
Number of segments freed:	0

BELOWHEAP statistics:

Initial size:	8192
Increment size:	4096
Total heap storage used (sugg. initial size):	0
Successful Get Heap requests:	0
Successful Free Heap requests:	0
Number of segments allocated:	0
Number of segments freed:	0



# Tuning storage

- Now using stack(8k,4k) heap(4k,4k)

STACK statistics:

Initial size:	8192
Increment size:	4096
Maximum used by all concurrent threads:	4792
Largest used by any thread:	4792
Number of segments allocated:	1
Number of segments freed:	0

...snip...

HEAP statistics:

Initial size:	4096
Increment size:	4096
Total heap storage used (sugg. initial size):	3328
Successful Get Heap requests:	4
Successful Free Heap requests:	2
Number of segments allocated:	1
Number of segments freed:	0



# Tuning storage

---

- A bit more meaty!
  - Testcase requests 50000 random pieces of storage of sizes from 1 to 32K in size
  - Then the storage is freed.
  - We'll run the program without tuning
  - We'll then re-run the program (same random values) after tuning
  - Have we saved storage? Performance?

# Tuning storage

## ■ Untuned example

STACK statistics:

Initial size:	131072
Increment size:	131072
Maximum used by all concurrent threads:	204184
Largest used by any thread:	204184
Number of segments allocated:	3
Number of segments freed:	0

...Snip...

HEAP statistics:

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	819229056
Successful Get Heap requests:	50002
Successful Free Heap requests:	50000
Number of segments allocated:	27949
Number of segments freed:	0

- Note: 27949 segments of 32k each – 915,832,832 bytes

# Tuning storage

- Tuned HEAP(100M,100M) STACK(256K,256K)

STACK statistics:

Initial size:	262144
Increment size:	262144
Maximum used by all concurrent threads:	200944
Largest used by any thread:	200944
Number of segments allocated:	1
Number of segments freed:	0

...Snip...

HEAP statistics:

Initial size:	104857600
Increment size:	104857600
Total heap storage used (sugg. initial size):	818334944
Successful Get Heap requests:	50002
Successful Free Heap requests:	50000
Number of segments allocated:	8
Number of segments freed:	0

- Note: 8 segments of 100M each – 838,860,800 bytes!!!

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# Tuning storage

- Look what else happened!

- Untuned

```

=====
-                                     REGION          --- STEP TIMINGS ---
- STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME    ELAPSED TIME    EXCP
- CLPG      COMPILE   CBCDRVR      00      72K      0:00:00.06   0:00:02.84     1590
- CLPG      PLKED     EDCPRLK     04      60K      0:00:00.01   0:00:00.99     534
- CLPG      LKED       HEWL        00      92K      0:00:00.01   0:00:00.63     174
- CLPG      GO         PGM=* .DD   00      60K      0:00:15.44   0:00:18.02     505
  
```

- Tuned

```

=====
-                                     REGION          --- STEP TIMINGS ---
- STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME    ELAPSED TIME    EXCP
- CLPG      COMPILE   CBCDRVR      00      72K      0:00:00.06   0:00:03.67     1555
- CLPG      PLKED     EDCPRLK     04      60K      0:00:00.01   0:00:01.25     535
- CLPG      LKED       HEWL        00      92K      0:00:00.01   0:00:00.50     170
- CLPG      GO         PGM=* .DD   00      60K      0:00:00.12   0:00:01.45     501
  
```

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# Tuning storage

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- What about KEEP vs FREE
  - Testcase requests 50000 random pieces of storage of sizes from 1 to 32K in size
  - Free 20000 pieces, then get 20000 more
  - Free everything
  - We'll run the program without tuning and FREE
  - We'll run the program without tuning and KEEP
  - What have we done to storage and performance?

# Tuning storage

## ■ Untuned

HEAP statistics: **(Using FREE)**

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	819982896
Successful Get Heap requests:	70002
Successful Free Heap requests:	69999
Number of segments allocated:	39122
Number of segments freed:	39120

HEAP statistics: **(Using KEEP)**

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	819983152
Successful Get Heap requests:	70002
Successful Free Heap requests:	70000
Number of segments allocated:	27952
Number of segments freed:	0

- Note: You can't determine storage used to back segments now

# Tuning storage

- Performance – not a huge difference but KEEP is faster!

- **FREE**

```
-----  
-  
- REGION                      --- STEP TIMINGS ---  
- STEPNAME  PROCSTEP  PGMNAME    CC    USED    CPU TIME    ELAPSED TIME    EXCP  
- CLPG      COMPILE   CBCDRVR     00    72K    0:00:00.06   0:00:04.95     1496  
- CLPG      PLKED     EDCPRLK     04    60K    0:00:00.01   0:00:02.46     504  
- CLPG      LKED       HEWL        00    92K    0:00:00.01   0:00:01.13     171  
- CLPG      GO         PGM=* .DD   00    60K    0:00:25.79   0:01:02.34     474
```

- **KEEP**

```
-----  
-  
- REGION                      --- STEP TIMINGS ---  
- STEPNAME  PROCSTEP  PGMNAME    CC    USED    CPU TIME    ELAPSED TIME    EXCP  
- CLPG      COMPILE   CBCDRVR     00    72K    0:00:00.06   0:00:03.15     1493  
- CLPG      PLKED     EDCPRLK     04    60K    0:00:00.01   0:00:00.87     505  
- CLPG      LKED       HEWL        00    92K    0:00:00.01   0:00:00.46     171  
- CLPG      GO         PGM=* .DD   00    60K    0:00:22.34   0:00:24.85     469
```

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# Tuning storage

- Look what happens when we tune.

HEAP statistics:

```
Initial size: 104857600
Increment size: 104857600
Total heap storage used (sugg. initial size): 819088944
Successful Get Heap requests: 70002
Successful Free Heap requests: 69999
Number of segments allocated: 8
Number of segments freed: 0
```

```
-----
-                                     REGION          --- STEP TIMINGS ---
- STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME  ELAPSED TIME  EXCP
- CLPG      COMPILE   CBCDRVR      00      72K      0:00:00.06  0:00:02.67   1499
- CLPG      PLKED     EDCPRLK     04      60K      0:00:00.01  0:00:00.81    547
- CLPG      LKED      HEWL        00      92K      0:00:00.01  0:00:00.32    171
- CLPG      GO        PGM=*.DD    00      60K      0:00:00.15  0:00:01.04   496
```

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# More advanced tuning

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- What about those pesky Language Environment control blocks?
  - No externals to help
  - Effort can be made to reduce the number of enclaves
  - Use dynamic calls rather than linking to next program
  - Hard to see the results without using system tools... but let's try



# More advanced tuning

---

- Simple program does a LINK to another program
  - A new enclave is created
  - This 2<sup>nd</sup> program continues to get storage until it runs out
  - It is able to obtain 21568K of storage



# More advanced tuning

---

- Add to program to call down through 5 nested enclaves
  - Last enclave is able to obtain 20576K of storage
- Add to program to call down through 10 nested enclaves
  - Last enclave is able to obtain 19808K of storage
- Storage being consumed is to:
  - Load programs
  - Create enclave control blocks
    - This includes stacks and heaps
  - 1760K of storage usage (21568K-19808K)





# More advanced tuning

---

- Change programs to use dynamic call rather than LINK
  - One call case – 21728K could be obtained
  - Five call case – 21664K could be obtained
  - Ten call case – 21600K could be obtained
- Note how much less storage is consumed.
  - Basically just the amount to load the programs
  - 128K for 10 calls deep (21728K-21600K)



# A Real Life Example

- Software vendor's multi-threaded transaction server
  - Experienced significant delay due to storage thrashing
  - Followed similar methodology
    - Stacks tuned
    - Heaps tuned
    - Heap Pools turned on, tuned

	Real CPU Consumed (sec)	# API Calls	Avg Elapsed Time per call (sec)	Avg CPU per call (ms)	Min elapsed time (sec)	Max elapsed time (sec)
Baseline	405.68	53349	0.5302	6.91	0.0003	265.9225
After Tuning	106.51	61943	0.0033	1.56	0.0002	0.6142

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# Summary

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- Storage run-time option has high overhead for initializing the stack
- Use RPTSTG to tune your stack, heap and other storage sizes
- KEEP is faster than FREE
- Use dynamic call versus LINK
  - Requires program update or recompile



# Sources of Additional Info

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- All Language Environment documentation available on the Language Environment Web site



- Language Environment Debugging Guide
- Language Environment Programming Reference
- Language Environment Programming Guide
- Language Environment Web site
  - [http://www.ibm.com/systems/z/os/zos/features/lang\\_environment](http://www.ibm.com/systems/z/os/zos/features/lang_environment)