DB2 11 for z/OS Application Functionality
(Check out these New Features)

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About This Presentation

- The main intent of this session is to highlight some of the new application related features of DB2 11. We want to introduce some of the new application functionality and get you thinking more about DB2 11!
Topics for Review

- Archive Transparency
- The Optimizer and RUNSTATS
- Autonomous Transaction Management
- Pseudo Deleted Index Entry Cleanup
- Further exploitation of RELEASE(DEALLOCATE)
Managing large amounts of data for a table

- Managing tables that contain a large amount of data is a common problem
  - Performance is a challenge
  - Many times, a large portion of the data is commonly accessed
- Many times application design / database design has been done to manage this challenge
  - Move “old” data to another table
  - Keep “current” data separate
Applications are designed to handle the access to the data
Still a challenge to manage and control
What is Archive Transparency in DB2 11?

- Data archive management from OLTP current table to “old” data table
- You can include or exclude archive table data access without a need to change the SQL statement
- Control the scope of the query with a global variable
- archive-enabled table and associated archive table
Transparent Archive

- Applications can query current + archive with no SQL changes
  - By default, data is retrieved from base table only, as usual
  - Set a new global variable when archive data is desired
  - DB2 automatically converts SQL to UNION ALL via dynamic plan switching technique (high performance)
- Archiving process is user-controlled
- Move_To_Archive global variable allows DELETEs to be automatically archived
- Leverages DB2 10 temporal constructs for archiving use cases
- Future potential for more IDAA synergy

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Optimizer input for RUNSTATS

• The better the RUNSTATS, the more effective the optimizer
• The standard or default statistics are commonly used
  – RUNSTATS TABLE(ALL) INDEX(ALL) KEYCARD
• Additional statistics will typically help queries perform better
• The challenge is in determining what RUNSTATS to collect
The flow of Optimizer input for RUNSTATS
DB2 11 and DB2 Optimizer Feedback

- During access path selection process, the optimizer identifies missing or conflicting statistics
  - Every BIND, REBIND or PREPARE
    - Writes recommendations to SYSIBM.SYSSTATFEEDBACK asynchronously
  - DB2 also provides statistics recommendations on EXPLAIN
    - Populates DSN_STAT_FEEDBACK synchronously
- Information from SYSSTATFEEDBACK or DSN_STAT_FEEDBACK can be used to generate input to RUNSTATS
  - Contents must be interpreted, not directly used by RUNSTATS
  - Requires DBA or tooling to convert to RUNSTATS input
Autonomous Transactions

- Useful for event logging, Audit data, special data management
- DB2 executes native SQL procedure in a unit of work that is independent of the calling program
  - May perform SQL, COMMITs, and ROLLBACK
  - No uncommitted changes from it’s caller are available
  - Autonomous SP and caller of SP do not share locks
- COMMIT is done when autonomous SP completes successfully, (SQLCODE >=0)
Definition of Autonomous Transactions

- Specify AUTONOMOUS keyword on ALTER or CREATE PROCEDURE:
  - DYNAMIC RESULT SETS 0 must be in effect.
  - Stored procedure parameters must not be defined as:
    - A LOB type
    - The XML data type
    - A distinct data type that is based on a LOB or XML value
    - An array type that is defined with array elements that are a LOB type
    - A value must not be assigned to a global variable when an autonomous procedure is executing.
Autonomous Transaction Basic Flow

Application Program

SQL INSERT CUST_TAB

CALL AUDIT_SP (autonomous SP)

“ROLLBACK”

Data inserted to CUST_TAB not visible to Autonomous SP

DB2DBM1 address space

AUDIT_SP;
SQL INSERT AUDIT_TAB END

Data is committed in AUDIT_TAB Regardless of what happens in calling program
Monitoring Autonomous Transactions

-DISPLAY THREAD(*) command

12.34.56 DB2P DISPLAY THREAD(*)
12.34.56 STC00090 DSNV401I DB2P DISPLAY THREAD REPORT FOLLOWS -
12.34.56 STC00090 DSNV402I DB2P ACTIVE THREADS -

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<th>PLAN</th>
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<td></td>
</tr>
</tbody>
</table>
Pseudo deleted index entries

• When table rows are deleted, a pseudo delete operation is done for the index RIDs
  – Index RIDS are “flagged” as deleted
  – unless the delete process has locked the entire table
• Impact of pseudo deleted index entries
  – Index size grows with increasing number of pseudo-deleted index entries
  – More getpages and lock requests required
  – Increased CPU cost and possibly longer elapsed times for access via index search
Pseudo deleted index entry clean up

• Prior to DB2 11, clean up of pseudo deleted index entries were part of main line processing and does impact application with increased costs
  – REORG INDEX takes care of pseudo-empty index pages and pseudo-deleted entries
• DB2 11 introduces additional action in the clean up of the pseudo deleted index entries
• Clean up pseudo-empty index pages and pseudo deleted index entries
  • Could reduce the size of indexes
  • Could improve SQL performance
  • Could reduce the need to run the REORG INDEX utility
DB2 11 Pseudo deleted index entry Process

• In Conversion Mode
• Automated cleanup of pseudo-deleted index entries in index leaf pages and of pseudo-empty index pages
• Mainline processing continues to do cleanup
• System tasks, running as enclave SRBs and are zIIP eligible are utilized
• Utilizes RTS information to identify indexes for cleanup
• Clean up work is only down for an index that is already open for “update”
• Page consistency is maintained
DB2 11 Pseudo deleted index management

• Cleanup threads are controlled
  – System parameter INDEX_CLEANUP_THREADS (0-128)
  – 0 disables index cleanup
  – Default is 10
  – Child clean up threads CORRELATION ID = 014.IDAEMKxx
• Specific index clean is done via SYSIBM.SYSINDEXCLEANUP table
  – indexes can be enabled or disabled for cleanup
  – Can specify:
    • Name of databases and indexes
    • Cleanup enabled or disabled
    • Day of week or day of month
    • Start time and end time
DB2 11 Pseudo deleted index monitoring

- IFCID 377 - tracks cleanup at the index page level
  - DBID, PSID, partition number, page number
  - pseudo empty page OR pseudo-deleted entries
  - Includes count of pseudo-deleted entries cleaned up

- **Not included in any trace class**

- RECORD TRACE used for reporting
Thread Management – Breaking in

DB2 10 provided a performance alternative with the use of RELEASE(DEALLOCATE)
- This creates a challenge in executing DDL, REORG, BIND

DB2 11 delivers a break-in mechanism for persistent RELEASE(DEALLOCATE) threads
- Persistent thread automatically detects operations that would like to break in
- If detected, then RELEASE(DEALLOCATE) will behave like RELEASE(COMMIT)
- zPARM PKGREL_COMMIT=YES must be set (parameter is online changeable)
  - Default is YES

The release of the resources after COMMIT/ROLLBACK occurs only if other DB2 operations (e.g. Bind, DDL, online REORG) are waiting for exclusive control to this package.

Packages resume normal RELEASE(DEALLOCATE) behavior after the break-in operation completes
Summary

- DB2 11 once again provides a wealth of new application functionality features
- Many features provide performance benefits like CPU reduction
- Many features provide options to handle application design

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Typical Utilization for Servers
Windows: 5-10%  Unix: 10-20%  System z: 85-100%

System z can help **reduce** your floor space up to **75%-85%** in the data center

Questions

**System z** can lower your total cost of ownership, requiring **as little as 30%** of the power of a distributed server farm running equivalent workloads

The cost of storage is typically **three times more** in distributed environments

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