



Getting the most out of your OSA (Open Systems Adapter) with z/OS Comm Server

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Tuesday, August 11, 2015: 11:15 AM - 12:15 PM, Dolphin, Asia 2









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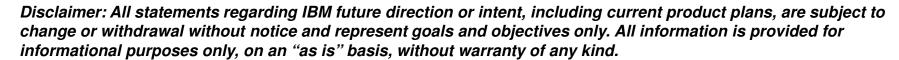
Agenda



- What is QDIO?
- OSA inbound "routing"
- Overview of selected recent OSA enhancements
- Appendix







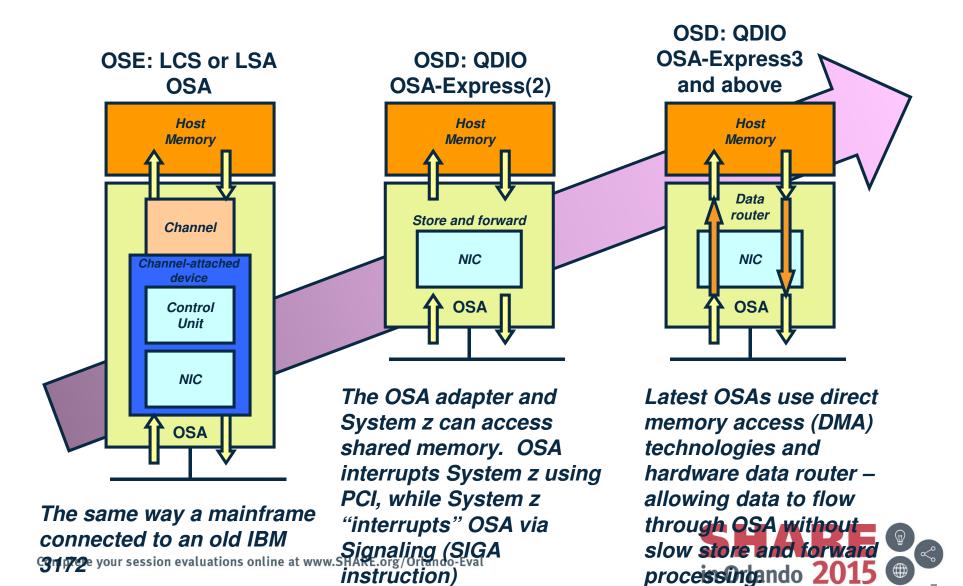




What is QDIO?



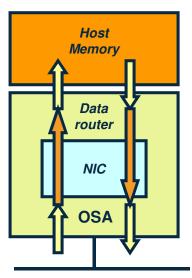
Queued Direct IO (QDIO) - How System z accesses a LANARE



instruction)

OSA and QDIO work together

- QDIO Layer 3 (z/OS)– for TCP/IP traffic only (IPv4 and IPv6)
 - OSA IP Processing Assist (IPA) offloads many TCP/IP functions to the OSA adapter, of which some are:
 - Performs all ARP processing (IPv4 only)
 - Provides Multicast support
 - Builds MAC and LLC headers
 - Performs checksum processing
 - Performs outbound TCP segmentation
 - For SNA/APPN/HPR traffic with QDIO layer 3: use TN3270 and Enterprise Extender



- OSA and QDIO are designed for high speed communication between System z and the Local Area Network
 - Reduced TCP/IP path length (reduced z/OS CPU)
 - **QDIO IP Processing Assist**
 - LPAR-to-LPAR communication with port sharing
 - Direct Memory Access (DMA) Protocol
 - Memory-to-memory communication
 - DMA allows data to move directly from the OSA-Express microprocessor to the host memory. This bypasses three layers of processing that are required when using ESCON and OSA-2 features, dramatically improving throughput.
 - **Dvnamic** customization
- 10 Gigabit Ethernet, Gigabit Ethernet, 1000BASE-T Ethernet (1000, 100, and 10 Mbit)



A little more on QDIO and devices

Control data between z/OS CS and the OSA port is exchanged over one read device and one write device (READ even address, WRITE odd address)

- The READ and WRITE device pair is shared among all the IP interface definitions in an

LPAR that use the same OSA port

Data is exchanged over data devices

One data device is used per:

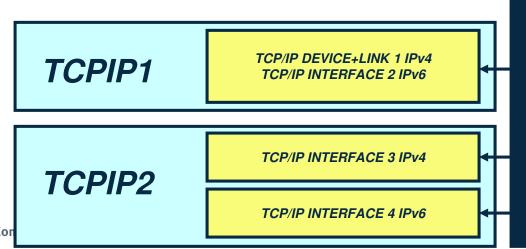
A set of exactly one IPv4 and one IPv6 interface

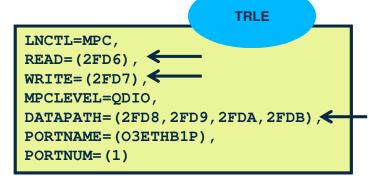
IPv4: DEVICE+LINK, IPv6: INTERFACE

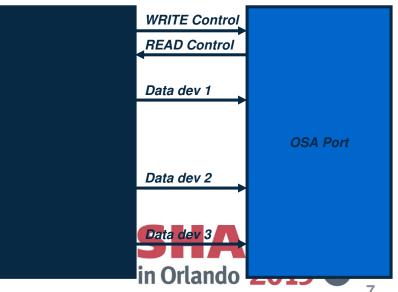
One IPv4 interface defined with INTERFACE

One IPv6 interface

One data device is used by OSAENTA trace, if that trace is activated









OSA inbound "routing"



Setting a few facts straight about OSA and "routing"



- OSA is not a full-function IP router.
 - OSA can analyze a destination IP address in a LAN frame to decide which LPAR an incoming IP packet belongs to.
 - OSA does not participate in dynamic routing updates.
 - OSA does not act as a general IP router
 - Depends on the TCP/IP stacks to do that and handle all IP routing issues.
- The OSA NIC has a single physical MAC address that is shared among all the LPARs that share the OSA port.
 - All IP packets to all LPARs can be destined for one and the same MAC address
 - In that case, OSA selects the correct stack based on destination IP address and the OSA Address Table – The OAT!
- If z/OS acts as an IP router to IP networks behind z/OS, the destination address may be any IP address on those networks behind z/OS.
 Primary
 - LPAR designated as default router will receive such packets
 - Can become extremely cumbersome to set up if LPARs that share an OSA port are connected to different back-end networks
- When a port is defined as an OSE/LCS port, the content of the OAT must be maintained manually through OSA/SF interaction.
- When a port is defined as a OSD/QDIO port, the content of the OAT is maintained automatically by the sharing TCP/IP stacks.

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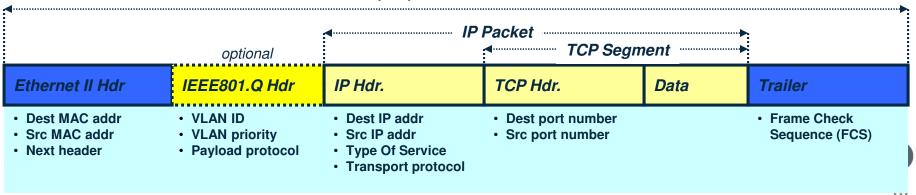
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Some basic LAN technology overview

- The LAN infrastructure transports "Frames" between Network Interface Cards (NICs) that are attached to the LAN media (Copper or fiber optic)
 - Ethernet II (DIX) MTU 1500 and jumbo frame MTU 9000 (most common)
- Each NIC has a physical hardware address
 - A Media Access Control (MAC) address
 - Burned in (world-wide unique by vendors) or alternatively locally administered
- Every frame comes from a MAC and goes to a MAC
 - There are special MAC values for broadcast and multicast frames
- Every frame belongs to the physical LAN or to one of multiple Virtual LANs (VLAN) on the physical LAN
 - A VLAN ID is in the IEEE801.Q header if VLAN technologies are in use
- A frame carries a payload of a specified protocol type, such as ARP, IPv4, IPv6, SNA LLC2, etc.

Ethernet II (DIX) LAN Frame

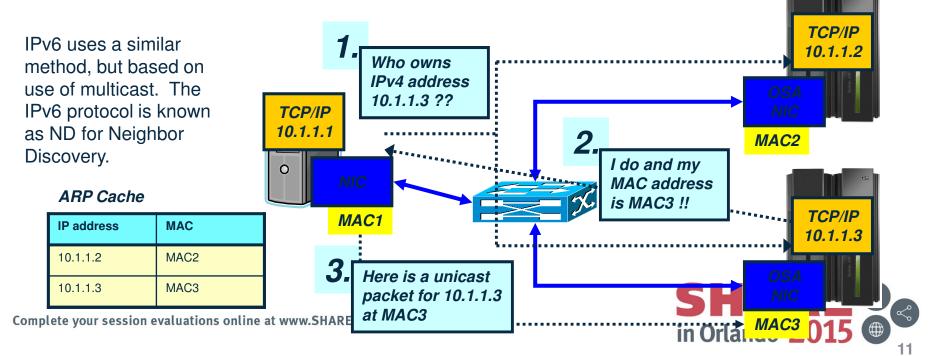




Correlation of IPv4 addresses and MAC addresses on a LAN SHI Address Resolution Protocol (ARP)

- An IPv4 node uses the ARP protocol to discover the MAC address of another IPv4 address that belongs to the same IPv4 subnet.
- ARP requests are broadcasted to all NICs on the LAN
- The one NIC that has a TCP/IP stack with the requested IPv4 address responds directly back to the IPv4 node that sent out the broadcast

 Each IPv4 node maintains a cache of IPv4 addresses and associated MAC addresses on their directly connected LANs





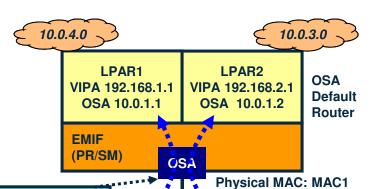
An example – formatted by Wireshark (http://www.wireshark.org)

A sample Address Resolution Request frame

Broadcast

```
■ Ethernet II, Src: 02:00:00:13:89:01 (02:00:00:13:89:01), Dst: Broadcast (ff:ff:ff:ff:ff)
  ■ Destination: Broadcast (ff:ff:ff:ff:ff:ff) 
      Address: Broadcast (ff:ff:ff:ff:ff)
      .... 1 .... = IG bit: Group address (multicast/broadcast)
      .... ..1. .... .... = LG bit: Locally administered address (this is NOT the factory default)
  ■ Source: 02:00:00:13:89:01 (02:00:00:13:89:01)
      Address: 02:00:00:13:89:01 (02:00:00:13:89:01)
      .... ... 0 .... .... = IG bit: Individual address (unicast)
      .... ..1. .... .... = LG bit: Locally administered address (this is NOT the factory default)
    Type: 802.10 Virtual LAN (0x8100)
■ 802.1Q Virtual LAN
    000. .... = Priority: 0
    ...0 .... = CFI: 0
    .... 0001 0101 1111 = ID: 351
    Type: ARP (0x0806) 🔔
Address Resolution Protocol (request)
                                                                              ARP type protocol
    Hardware type: Ethernet (0x0001)
    Protocol type: IP (0x0800)
    Hardware size: 6
    Protocol size: 4
    opcode: request (0x0001)
    Sender MAC address: 02:00:00:13:89:01 (02:00:00:13:89:01)
    Sender IP address: 10.3.51.16 (10.3.51.16)
    Target MAC address: 00:00:00_00:00:00 (00:00:00:00:00:00)
    Target IP address: 10.3.51.17 (10.3.51.17)
                                                         Addr in question
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So how does this work in a System z environment with shared A R E (virtualized) OSA adapters?



10.0.1.0

USA Adaress
Table (OAT)

IP address	LPAR / Device
192.168.1.1	LPAR1/Dx
192.168.2.1	LPAR2/Dy
10.0.1.1	LPAR1/Dx
10.0.1.2	LPAR2/Dy
Primary	LPAR2/Dy
Secondary	LPAR1/Dx

If the destination IP address in a frame that arrives at the OSA adapter doesn't belong to any of the LPARs that share the adapter, that IP packet is given to the LPAR that acts as the Default OSA router. There are separate IPv4 and IPv6 default OSA router specifications.

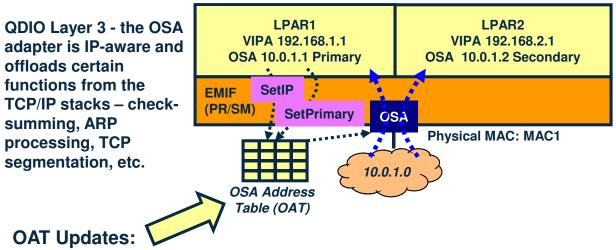
LAN frames

Dest_MAC	Dest_IPv4	Who gets it?	Why?
MAC1	10.0.1.1	LPAR1	Registered
MAC1	192.168.1.1	LPAR1	Registered
MAC1	10.0.1.2	LPAR2	Registered
MAC1	192.168.2.1	LPAR2	Registered
MAC1	10.0.3.5	LPAR2	Default router
MAC1	10.0.4.6	LPAR2	Default router

- An OSA NIC has a physical MAC address, just like all NICs
- An OSA NIC is often used by multiple LPARs and TCP/IP stacks
 - The NIC is virtualized so it functions as the NIC for multiple TCP/IP stacks, each with their own IP address
- If someone ARPs for 10.0.1.1 in LPAR1, OSA will return MAC1
- If someone ARPs for 10.0.1.2 in LPAR2, OSA will also return MAC1
- So what does OSA then do when a unicast frame arrives with a destination MAC address of MAC1?
 - It peeks into the IP header inside the frame, and consults a table known as the OSA Address Table (OAT) to see which LPAR the IP address inside the frame belongs to



QDIO layer-3: less administration, more dynamics – but OAT remains an important element of OSA inbound routing



OAT is maintained dynamically by TCP/IP

IP address	LPAR / Device	
192.168.1.1	LPAR1/Dx	
192.168.2.1	LPAR2/Dy	
10.0.1.1	LPAR1/Dx	
10.0.1.2	LPAR2/Dy	
Primary	LPAR1/Dx	
Secondary	LPAR2/Dy	

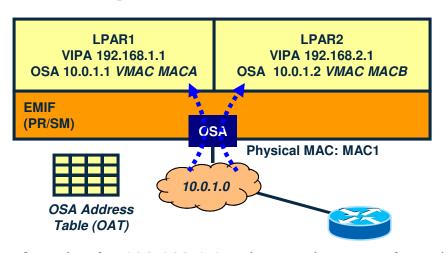
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- LCS: operator intervention via OSA/SF commands
- QDIO: dynamic by software via QDIO control channel
- QDIO Interface definitions in the TCP/IP stacks are used to dynamically establish the stack as the OAT default router, secondary router, or non-router.
- Whenever a QDIO device is activated or the TCP/IP home list is modified (through OBEYFILE command processing or through dynamic changes, such as dynamic VIPA takeover), the TCP/IP stack updates the OAT configuration dynamically with the HOME list IP addresses of the stack.
- The OAT includes all (non-LOOPBACK) HOME IP addresses of all the stacks that share the OSA adapter.
- The fact that the OSA microcode is IP address-aware (as it is in this scenario) is the reason for referring to this as QDIO layer 3 processing (layer 3 is generally the networking layer in an OSI model, the IP networking layer when using TCP/IP)

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Wouldn't it be so much easier with a MAC address per z/OS TCP/IP network interface?







Router's ARP cache

IP Address	MAC Address
10.0.1.1	MACA
10.0.1.2	MACB

The whole mess with PRIROUTER and SECROUTER is gone !!

It also removes issues with external load balancers that use MAC-level forwarding. It makes a system z LPAR look like a "normal" TCP/IP node on α Ι ΔΝ

- A packet for 192.168.1.1 arrives at the router from the downstream network
 - Router determines the destination is not directly attached to the router
 - Router looks into its routing table and determines next-hop IP address for this destination is 10.0.1.1, which is on a network that is directly attached to the router
 - Router looks into its ARP cache and determines the associated MAC address is MACA
 - Router forwards the IP packet in a LAN frame to MACA
- Frame arrives in OSA
 - OSA determines which LPAR it belongs to based on the virtual MAC address
 - The OAT may optionally still play a role in inbound routing decisions by the OSA adapter:
 - If the destination address (192.168.1.1) were not in the home list of this LPAR (and hence not in the OAT for this LPAR), should OSA still send it up to the LPAR or not?
 - You decide via a configuration option
 - ROUTEALL: send all packets with my VMAC to me
 - ROUTELCL: send only packets to me if they are in my HOME list
- The OAT remains in use for other functions, such as ARP ownership, etc.





OSA-Express virtual MAC while operating in QDIO layer-3 mode

- OSA MAC sharing problems do not exist if each stack has its own MAC
 - A "virtual" MAC
 - To the network, each stack appears to have a dedicated OSA port (NIC)
- MAC address selection
 - Coded in the TCP/IP profile
 - Generated and assigned by the OSA adapter
- All IP addresses for a stack are advertised with the virtual MAC
 - by OSA using ARP for IPv4
 - by the stack using ND for IPv6
- All external routers now forward frames to the virtual MAC
 - OSA will "route" to an LPAR/Stack by virtual MAC instead of IP address
 - All stacks can be "routing" stacks instead of 1 PRIROUTER stack
- Simplifies configuration greatly
 - No PRIROUTER/SECROUTER!



No reasons not to use VMACs. Use them!!!

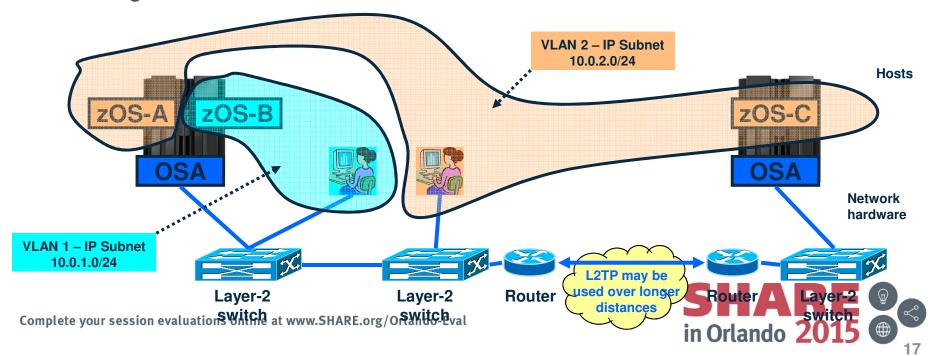


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What is a Virtual LAN (a VLAN)?

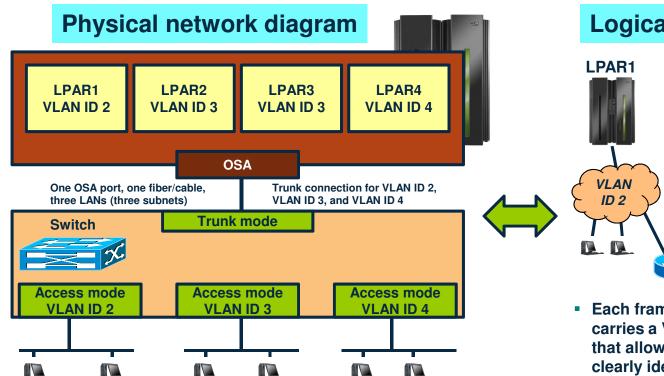
Wikipedia:

- A virtual LAN, commonly known as a VLAN, is a group of hosts with a common set of requirements that communicate as if they were attached to the same broadcast domain, regardless of their physical location.
- A VLAN has the same attributes as a physical LAN, but it allows for end stations
 to be grouped together even if they are not located on the same network switch.
- Network reconfiguration can be done through software instead of physically relocating devices.



z/OS and VLANs

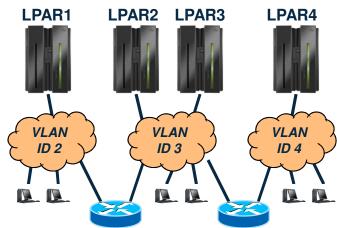




- Depending on switch configuration, the switch may interconnect the VLANs using a layer-3 IP router function.
- The subnets may belong to different routing domains or OSPF areas:
 - Test, production, demo
- The subnets may belong to different security zones:

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Logical network diagram



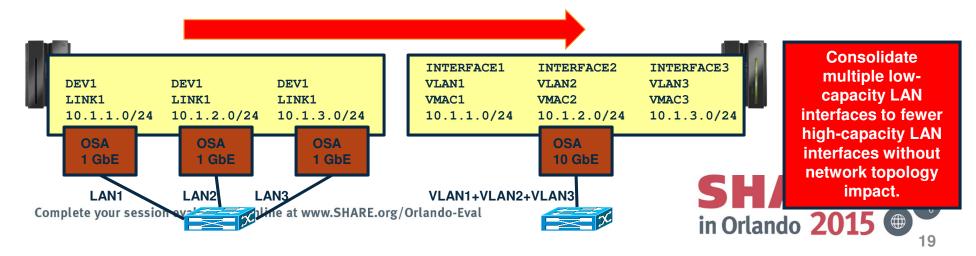
- Each frame on the trunk mode connection carries a VLAN ID in the IEEE802.3 header that allows the network equipment to clearly identify which virtual LAN each frame belongs to.
- On an access mode connection, the switch will transport frames belonging to the configured VLAN ID for that access mode connection only.

VLAN is a LAN media virtualization technology that allows multiple independent IP networks (IP subnets) to share one physical media, such as a cable, an adapter, or a layer-2 switch. Connectivity between VLANs is under control of IP routers.

Multiple network interfaces (VLANs) per OSA port per stack per IP protocol version



- As installations consolidate multiple OSA Gigabit Ethernet ports to a smaller number of 10 Gigabit Ethernet ports, low limits on VLANs has become too restrictive:
 - Not possible to retain existing network interface and IP subnet topology
 - Consolidating multiple LANs to one LAN requires IP renumbering
- z/OS V1R10 added support for eight VLANs per IP protocol per OSA port:
 - Each VLAN on the same OSA port must use unique, non-overlapping IP subnets or prefixes
 - Will be enforced by the TCP/IP stack
 - Each VLAN must be defined using the INTERFACE configuration statement
 - IPv4 INTERFACE statement only supports QDIO interfaces
 - Each VLAN must use layer-3 unique virtual MAC addresses with ROUTEALL
- z/OS V1R13 raised the number of VLANs to 32 per IP protocol per OSA port



A few more words on VLANs with z/OS CS



- z/OS Communications Server supports one VLAN ID per defined network interface
 - This is known as a global VLAN ID
 - Remember: z/OS CS now supports up to 32 interfaces per OSA port each with its own VLAN ID
 - OSA performs inbound routing based on VLAN IDs in the incoming frames and only sends frames to z/OS with the global VLAN ID z/OS has registered
- Linux on System z supports multiple VLAN IDs per interface
 - OSA sends multiple VLAN IDs up to Linux and Linux then de-multiplexes the different virtual LANs
- Some interfaces that share an OSA port may select to not specify a VLAN ID, while others do specify a VLAN ID (a mixed VLAN environment)
 - Warning: be very careful mixed VLAN environments are not recommended due to complexities with OSA default router selection and other functional issues. Use VLAN IDs on all TCP/IP interfaces that share an OSA port - or not at all.
- Mixed VLAN environment without VMAC LPAR1 LPAR2 LPAR3 LPAR4 No VLAN VLAN ID 2 VLAN ID 2 VLAN ID 3 No VMAC No VMAC No VMAC No VMAC **PRIROUTER PRIROUTER SECROUTER PRIROUTER** OSA
- LPAR1 acts as PriRouter for:
 - Untagged frames
 - Frames tagged with a null VLAN ID
 - Frames tagged with an unregistered (anything but VLAN 2 or 3) VLAN ID
 - Frames tagged with a registered VLAN ID, an unknown destination IP address but no VLAN Pri/SecRrouter
- LPAR2 acts as PriRouter for frames tagged with VLAN 2 while LPAR3 acts as SecRouter for that same VLAN.
- · LPAR4 acts as PriRouter for frames tagged with VLAN 3



OSA inbound routing

Start -

This is correct if all interfaces that share the OSA port use VLAN IDs or none of them do. If

some do and others do not, this becomes utterly complex, rather

Dest MAC

- VMACs are unique across all interfaces and VLANs that share an OSA port
- z/OS CS registers VMAC with the OSA port

VLAN ID

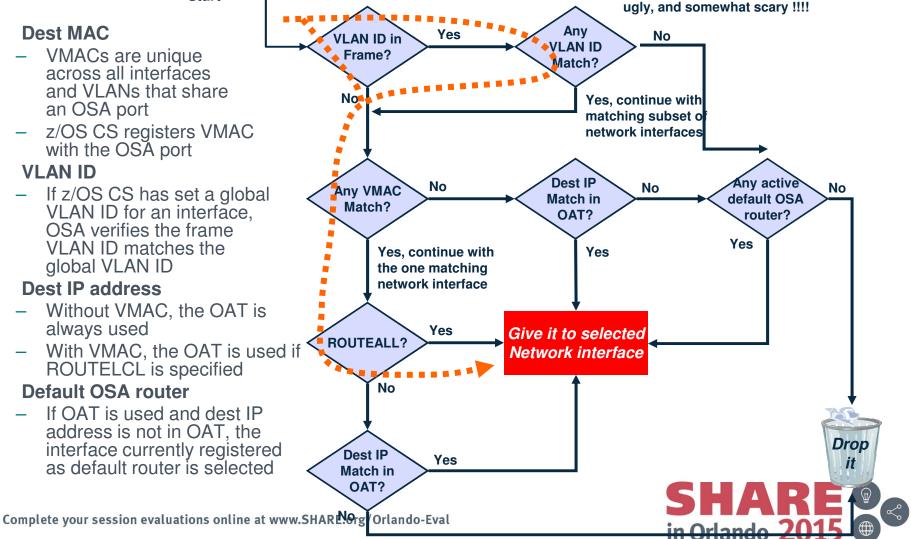
If z/OS CS has set a global VLAN ID for an interface, OSA verifies the frame VLAN ID matches the global VLAN ID

Dest IP address

- Without VMAC, the OAT is always used
- With VMAC, the OAT is used if ROUTELCL is specified

Default OSA router

If OAT is used and dest IP address is not in OAT, the interface currently registered as default router is selected



This elusive Maximum Transmission Unit – what is the



size really?

Interface MTU

- Configured on the INTERFACE statement or learned from the device
 - For IPv4, OSA reports 1492/8992 even when you use Ethernet II frames
 - For IPv6, OSA reports 1500/9000 because IPv6 always uses Ethernet II frames
- Reported as ActMtu: in a DEVLINKS netstat report

Configured route MTU

- For static routes: the value coded in the BEGINROUTES section for the destination
- For dynamic routes: the MTU value from the OMPROUTE interface over which the route was learned
- If no MTU size defined in OMPROUTE, a default of 576 will be used for IPv4

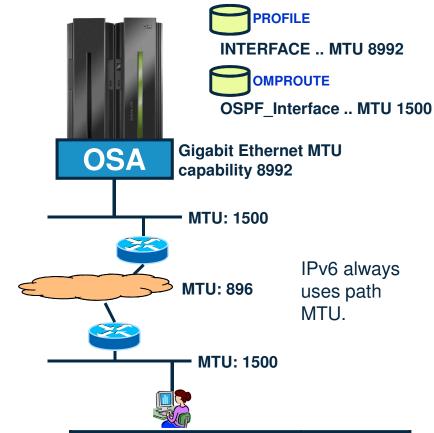
Actual route MTU

- Minimum of the interface MTU and the configured route MTU
- If path MTU is not enabled, this is the MTU that will be used for that route

Path MTU

- With path MTU enabled, TCP/IP starts out trying to use the actual route MTU and may then lower the MTU to what is discovered
 - With path MTU enabled, all frames are sent with the Do Not Fragment bit

complete your sessiath Multiple after a certain amount of time to accommodate topology changes that allows a higher MTU



MTU type	MTU value
Interface MTU	8992
Configured route MTU	1500
Actual route MTU	1500
Path MTU to this destination	896

Some examples on MTU specifications



MTU 1500 configured on an IPv4 INTERFACE stmt:

- CfgMtu: 1500 ActMtu: 1500

- OSA reports 8992 in this case, but our configured MTU overrides that and determines the actual MTU
- No MTU configured on an IPv4 INTERFACE stmt:

- CfgMtu: None ActMtu: 8992

- We have no configured MTU to override what OSA reports, so the actual MTU ends up being 8992 as reported by OSA for an IPv4 interface
- No MTU configured on an IPv6 INTERFACE stmt:

- CfgMtu: None ActMtu: 9000

- Since this is an IPv6 interface, OSA reports the full jumbo frame size of 9000 bytes as the MTU it supports. Since we did not override that with a configured MTU, the 9000 ends up being the actual MTU size.
- In the sample setup, an MTU of 1500 is coded in the OMPROUTE configuration file, so all actual route MTUs end up being 1500

- Default

9.42.105.65

UGO

0000000002 QDIO4A

- Metric: 00000001 MTU: 1500

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Overview of selected recent OSA enhancements



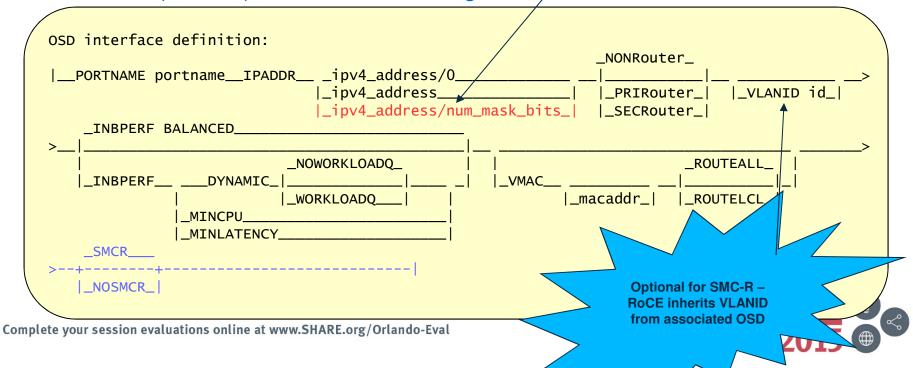
SMC-R TCP/IP Configuration – OSA Interface

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- IPAQENET INTERFACE statements
 - SMCR only valid for CHPIDTYPE OSD



- SMCR cannot be used with IPv4 OSD interfaces defined using DEVICE and LINK statements
- z/OS Communications Server: Shared Memory Communications over RDMA (SMC-R) Protocol – Wed, Aug 12 10:00 Asia 2



Inbound OSA/QDIO performance: Dynamic LAN idle timer



- Dynamic LAN idle timer is designed to reduce latency and improve network performance by dynamically adjusting the inbound blocking algorithm.
 - When enabled, the z/OS TCP/IP Stack is designed to adjust the inbound blocking algorithm to best match the application requirements.
- For latency sensitive applications, the blocking algorithm is modified to be "latency sensitive." For streaming (CPU sensitive) applications, the blocking algorithm is adjusted to minimize CPU. In all cases, the z/OS TCP/IP stack dynamically detects the application requirements, making the necessary adjustments to the blocking algorithm.
 - The monitoring of the application and the blocking algorithm adjustments are made in real-time, dynamically adjusting the application's LAN performance.
- System administrators can authorize the z/OS TCP/IP stack to enable the DYNAMIC setting, which was previously a static setting.
 - The z/OS TCP/IP stack is designed to dynamically determine the best setting for the current running application based on system configuration, inbound workload volume, CPU utilization, and traffic patterns.
- For an OSA Express port in OSD/QDIO mode the z/OS TCP/IP profile INBPERF parameter can be specified with one of the following options:
 - MINCPU a static interrupt-timing value, selected to minimize host interrupts without regard to throughput
 - MINLATENCY a static interrupt-timing value, selected to minimize latency
 - BALANCED (default) a static interrupt-timing value, selected to achieve reasonably high throughout and reasonably low CPU
 - throughput and reasonably low CPU

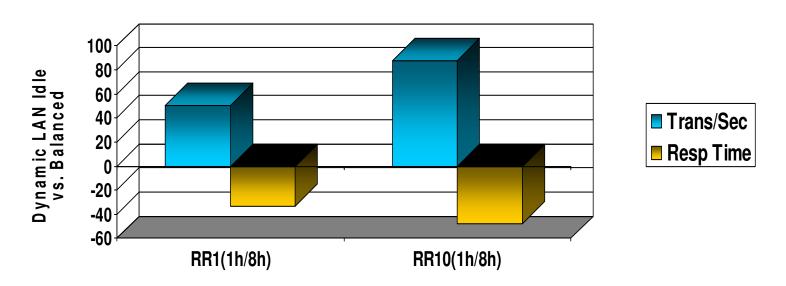
 DYNAMIC implements the new dynamic LAN idle algorithm ← recommended! → ■



Dynamic LAN Idle Timer: Performance Data

Dynamic LAN Idle improved RR1 TPS 50% and RR10 TPS by 88%. Response Time for these workloads is improved 33% and 47%, respectively.

RR1 and RR10 Dynamic LAN Idle

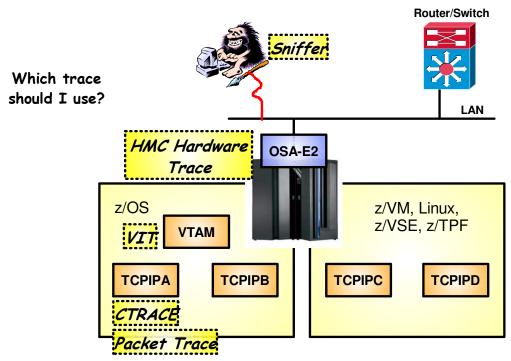


1h/8h indicates 100 bytes in and 800 bytes out



OSA-Express Network Traffic Analyzer (OSAENTA) function

- Diagnosing OSA-Express QDIO problems can be somewhat difficult
 - TCP/IP stack (CTRACE and/or packet trace)
 - VTAM (VIT)
 - OSA (hardware trace)
 - Network (sniffer trace)
- Often not clear where the problem is and which trace(s) to collect
 - Offloaded functions and shared OSAs can complicate the diagnosis
 - ARP frames
 - Segmentation offload



- Allows z/OS Communications Server to collect Ethernet data frames from the OSA adapter
 - Not a sniffer trace (but similar in some aspects)
 - Minimizes the need to collect and coordinate multiple traces for diagnosis
 - Minimizes the need for traces from the OSA Hardware Management Console (HMC)
 - Formatted with the standard z/OS CS packet trace formatter
 - Including the ability to convert to Sniffer/Wireshark format





QDIO priority queue selection based upon WLM importance level - Outbound

Use workload **WLM** service objective of Basic principle is that if class importance **Application** application address QoS policies are active. level (SCIL) they will determine spaces to select which priority queue to QDIO outbound use. priority queue **Optional TCPIP** TOS Profile QoS **Policy** WLMPRIORITYO SetSubnetPrioTosMask SubnetTosMask 11100000 PriorityTosMapping 1 11100000 PriorityTosMapping 1 11000000 PriorityTosMapping 1 10100000 PriorityTosMapping 1 10000000 PriorityTosMapping 2 01100000 PriorityTosMapping 2 01000000 PriorityTosMapping 3 00100000 Q4 OSA PriorityTosMapping 4 00000000

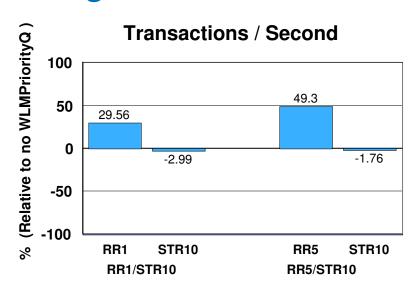
Complete

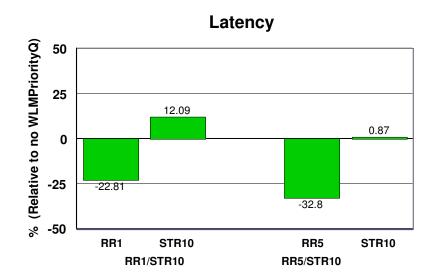
An easy way to take advantage of QDIO's four outbound priority queues!





OSA Express (QDIO) WLM Outbound Priority Queuing



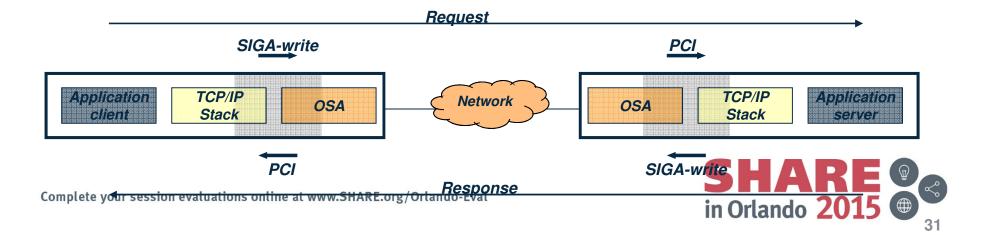


- Request-Response and Streaming mixed workload
- RR1/STR10: 1 RR session, 100 / 800 and 10 STR sessions, 1 / 20 MB
- RR5/STR10: 5 RR sessions, 100 / 800 and 10 STR sessions, 1 / 20 MB
- WLMPRIORITYQ assigned importance level 2 to interactive workloads and level 3 to streaming workloads
- The z/OS Workload Manager (WLM) system administrator assigns each job a WLM service class
- z/OS with WLM I/O Priority provides 30 to 50% higher throughput for interactive workloads compared to without WLM I/O Priority
- z/OS with WLM I/O Priority provides 23 to 33% lower latency compared to without WLM I/O Priority

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OSA-Express optimized latency mode (OLM)

- OSA-Express4/5 has significantly better latency characteristics than OSA-Express3
- The z/OS software and OSA microcode can further reduce latency:
 - If z/OS Communications Server knows that latency is the most critical factor
 - If z/OS Communications Server knows that the traffic pattern is not streaming bulk data
- Inbound
 - OSA-Express signals host if data is "on its way" ("Early Interrupt")
 - Host looks more frequently for data from OSA-Express
- Outbound
 - OSA-Express does not wait for SIGA to look for outbound data ("SIGA reduction")



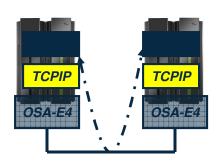
Performance indications of OLM for interactive workloads

were collected using a dedicated system environment. The results obtained in other configurations or operating system environments may vary.

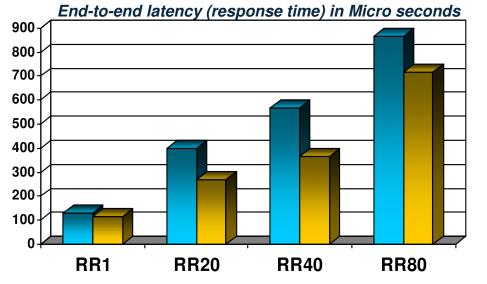


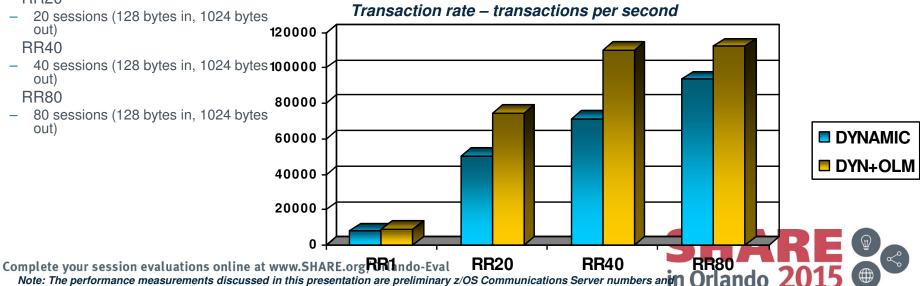
DYNAMIC

DYN+OLM



- Client and Server
 - Has close to no application logic
- - 1 session (1 byte in 1 byte out)
- **RR20**
 - 20 sessions (128 bytes in, 1024 bytes out)
- **RR40**
 - 40 sessions (128 bytes in, 1024 bytes 100000 out)
- **RR80**
 - 80 sessions (128 bytes in, 1024 bytes out)

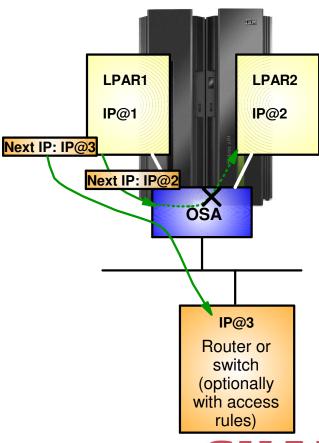




OSA interface isolation



- Allow customers to disable shared OSA local routing functions
 - ISOLATE/NOISOLATE option on QDIO network INTERFACE definition
- OSA local routing can in some scenarios be seen as a security exposure



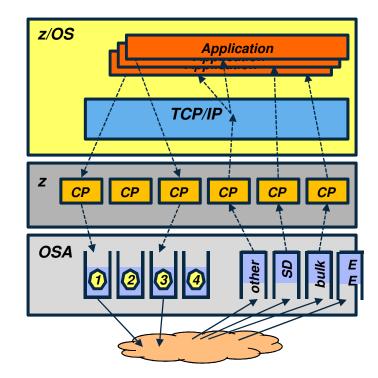
If you enable ISOLATE, packets with a nexthop IP address of another stack that share the OSA port, will be discarded.





OSA multiple inbound queue support: improved bulk transfer, EE and Sysplex Distributor connection routing performance

- Allow inbound QDIO traffic separation by supporting multiple read queues
 - "Register" with OSA which traffic goes to which queue
 - OSA-Express Data Router function routes to the correct queue
- Each input queue can be serviced by a separate process
 - Primary input queue for general traffic
 - One or more ancillary input queues (AIQs) for specific traffic types
- Supported traffic types
 - Bulk data traffic queue
 - Serviced from a single process reduces out of order delivery issues
 - Sysplex distributor traffic queue
 - SD traffic efficiently accelerated or presented to target application
 - Enterprise Extender (EE) queue
 - Avoids memory copy of data
 - All other traffic not backed up behind bulk data, EE or SD traffic (Primary Queue)
- Dynamic LAN idle timer updated per queue
- Can use OLM on Primary Queue



TCP/IP defines and assigns traffic to queues dynamically based on local IP address and port





Inbound Workload Queuing: Performance Data

Aix 5.3

p570

IWQ: Mixed Workload Results vs DYNAMIC:



Streaming Throughput also improved in this test: +5%

OSA-Express's 1GBe In Dynamic or 10GBe or IWQ mode

network

z/OS

(3 CP

LPARs)

For z/OS outbound streaming to another

platform, the degree of performance boost (due to IWQ) is relative to receiving platform's

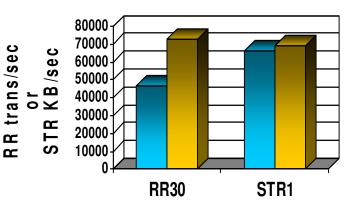
sensitivity to out-of-order packet delivery. For

streaming INTO z/OS, IWQ will be especially

beneficial for multi-CP configurations.

z/OS

Mixed Workload (IWQ vs Dynamic)





RR (z/OS to AIX)

STR (z/OS to z/OS)



Complete your session evaluations online at www.SHARE.org/Orlando-Eval

Operator command to query and display OSA information



- OSA/SF (Support Facility) has been used for years to configure OSA and display the configuration.
 OSA/SF has played a more central role for OSE devices (pre-QDIO) than for today's OSD devices (QDIO).
- OSD devices exclusively use IPA signals exchanged with the host to enable and configure features and register IP addresses to OSA.
- However, there was no mechanism to display the information directly from OSA without OSA/SF.
- D TCPIP, OSAINFO command for use with OSA-3 and above:
 - Base OSA information
 - OSA address table information
 - Information related to the new multiple inbound queues

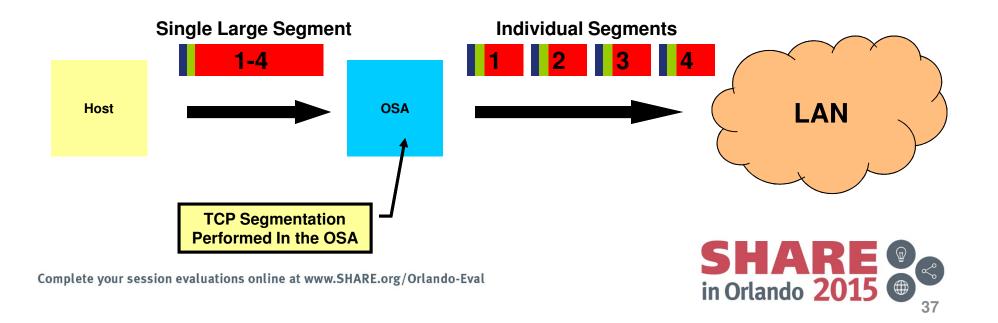
```
D TCPIP, OSAINFO, INTFN=V6O3ETHG0, MAX=100
EZZ0053I COMMAND DISPLAY TCPIP, OSAINFO COMPLETED SUCCESSFULLY
EZD0031I TCP/IP CS V2R2 TCPIP Name: TCPSVT
                                             15:39:52
Display OSAINFO results for Interface: V603ETHG0
PortName: O3ETHGOP PortNum: 00 DevAddr: 2D64 RealAddr: 0004
                                CHPID Type: OSD OSA code level: 5D76
PCHID: 0270
                   CHPID: D6
                   Active speed/mode: 10 gigabit full duplex
Gen: OSA-E3
Media: Singlemode Fiber
                              Jumbo frames: Yes Isolate: No
PhysicalMACAddr: 001A643B887C LocallyCfgMACAddr: 00000000000
Oueues defined Out: 4 In: 3
                             Ancillary queues in use: 2
Connection Mode: Layer 3
                              IPv4: No
                                        IPv6: Yes
SAPSup: 00010293
```

Complete yo

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TCP Segmentation Offload

- Segmentation consumes (high cost) host CPU cycles in the TCP stack
- OSA-Express (QDIO mode) feature called Segmentation Offload (also referred to as "Large Send")
 - Offload most IPv4/IPv6 TCP segmentation processing to OSA
 - Decrease host CPU utilization
 - Increase data transfer efficiency for packets
 - Configured in IPCONFIG statement



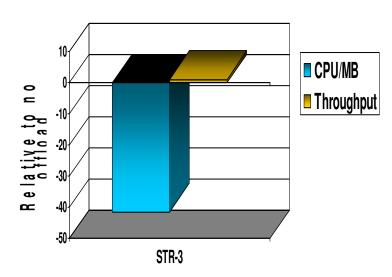


z/OS V2R2 segmentation offload performance measurements

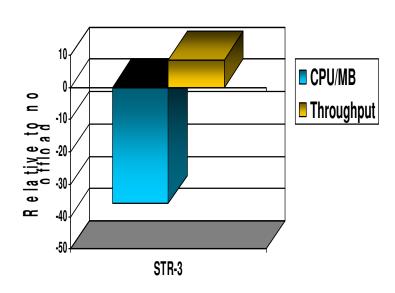


Note: The performance measurements discussed in this presentation were collected using a dedicated system environment. The results obtained in other configurations or operating system environments may vary.

OSA-Express3 10Gb



OSA-Express4/5 10Gb



Send buffer size: 180K for streaming workloads

Segmentation offload may significantly reduce CPU cycles when sending bulk data from z/OS!

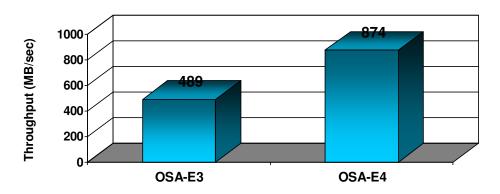




OSA Express 4/5 Enhancements

- Improved on-card processor speed and memory bus provides better utilization of 10GB network
- Enterprise Extender inbound workload queue provides internal optimizations
 - EE traffic processed quicker
 - Avoids memory copy of data
- Checksum Offload support for IPv6 traffic
- Segmentation Offload support for IPv6 traffic

OSA 10GBe - Inbound Bulk traffic







Please fill out your session evaluation

- Getting the most out of your OSA Adapter with z/OS Comm Server
- QR Code:











```
CHPID PATH=(CSS(0,1),FD),SHARED,
PCHID=581,TYPE=OSD

CNTLUNIT CUNUMBR=2FD0,PATH=((CSS(0),FD),(CSS(1),FD)),UNIT=OSA

IODEVICE ADDRESS=(2FD0,14),CUNUMBR=2FD0,UNIT=OSA,
UNITADD=00

IODEVICE ADDRESS=(2FDE,1),CUNUMBR=2FD0,UNIT=OSAD,
UNITADD=FE
```





```
O3ETHBOT TRLE
                LNCTL=MPC,
                                                                             X
                READ=(2FD0),
                                                                             X
                WRITE=(2FD1),
                                                                             X
                MPCLEVEL=QDIO,
                                                                             X
                DATAPATH= (2FD2, 2FD3, 2FD4, 2FD5),
                                                                             X
                PORTNAME = (O3ETHBOP),
                                                                             X
                PORTNUM=(0)
O3ETHB1T TRLE
                LNCTL=MPC,
                                                                             X
                READ = (2FD6),
                                                                              X
                WRITE=(2FD7),
                                                                             X
                MPCLEVEL=QDIO,
                                                                             X
                DATAPATH= (2FD8, 2FD9, 2FDA, 2FDB),
                                                                              X
                PORTNAME = (O3ETHB1P),
                                                                             X
                PORTNUM=(1)
```





OSA Express TCP/IP Interface definitions

```
INTERFACE O3ETHB0 DEFINE IPAQENET
PORTNAME O3ETHB0P
IPADDR 16.11.16.105/20
SOURCEVIPAINT LGEVIPA1
MTU 1500
VLANID 3
READSTORAGE GLOBAL
INBPERF DYNAMIC WORKLOADQ
IPBCAST
MONSYSPLEX
DYNVLANREG
OLM
VMAC 0204100B1069 ROUTEALL
INTERFACE O3ETHB1 DEFINE IPAQENET
PORTNAME O3ETHB1P
IPADDR 16.11.17.105/20
SOURCEVIPAINT LGEVIPA1
MTU 1500
VLANID 3
READSTORAGE GLOBAL
INBPERF BALANCED
IPBCAST
MONSYSPLEX
DYNVLANREG
OLM
VMAC 0204100B1169 ROUTEALL
```

Complete your session