Capping, Capping, and Capping:
A Comparison of Hard and Soft-capping Controls

Horst Sinram - STSM, z/OS Workload and Capacity Management

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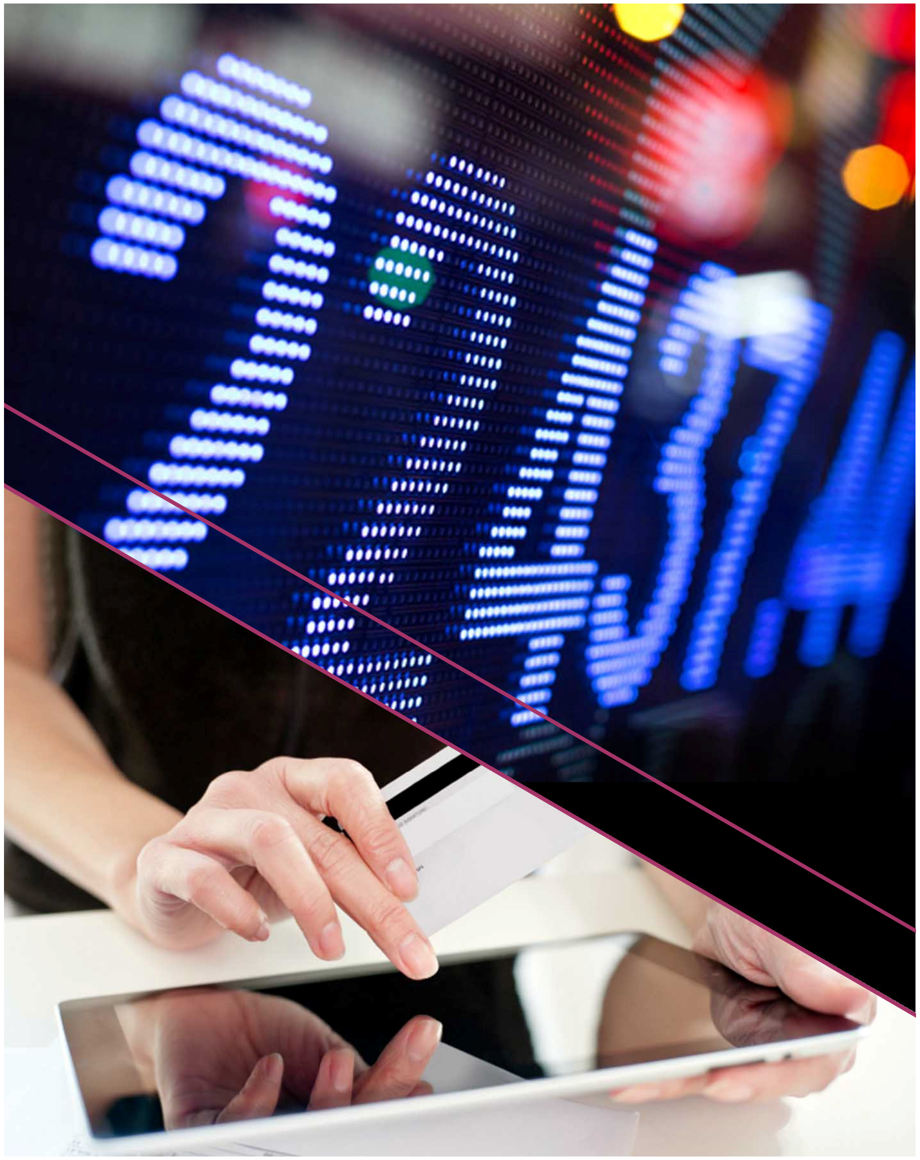
Agenda

- **Overview of capping types**
- Initial capping
- Absolute capping
- Defined capacity & group capacity
- Resource group capping
- 4HRA management

- Additional Material



Reasons you would consider capping techniques…

Technical motivation

- Protect LPARs against other LPARs, e.g. multi-tenancy
- Influence capacity-based workload routing
- Guarantee unused CPC processor capacity
- Protect workloads (sets of service classes) against other workloads

Financial motivation

- Limit software cost
  - Capacity limit for one or more LPARs
  - Four hour rolling average (4HRA) consumption

- Possible impact of capping needs to be monitored and accepted
- Cap limits should be adjusted as appropriate
  - Watch your SLAs

Comparison of LPAR capping types

Type of capping	Scope	Specification units	Proc types	Stability under configuration changes	Suitable to technically separate LPARs or groups of LPARs	Control point
Initial (hard capping)	LPAR	LPAR share of CPC capacity	Any	–	+	SE/HMC
Absolute capping	LPAR	Fractional #processors		O	+	
Defined capacity (DC, soft capping)	LPAR	MSU	CP	+	–	
LPAR group capacity (GC, soft capping)	Group of LPARs	MSU		+	–	
Resource group capping	Groups of service classes in Sysplex or per LPAR	Unweighted CPU SU/sec, fraction of LPAR share, or fractional #processors	CP*	+	N/A	WLM Policy
Logical configuration	LPAR	Integer #processors	Any	O	+ but coarse grain	HMC+ OS

PR/SM controlled

WLM controlled, PR/SM enforced

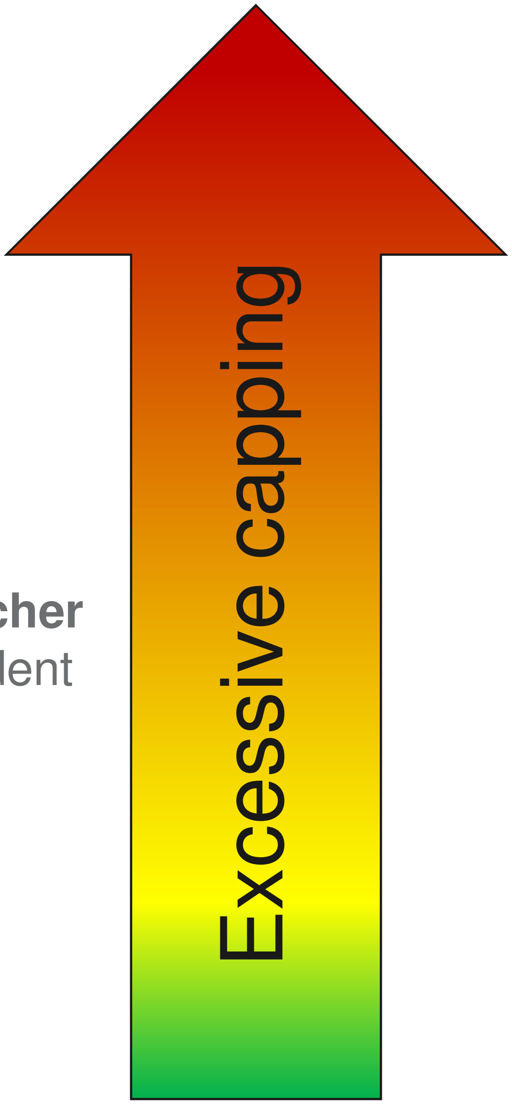
WLM controlled

Which capping techniques may be combined?

Type of capping	Initial (hard capping)	Absolute capping	Defined capacity (soft capping)	LPAR group capacity (soft capping)	Resource group capping
Initial (hard capping)		+	–	–	+
Absolute capping			+	+	+
Defined capacity (soft capping)				+	+
LPAR group capacity (soft capping)					+
Resource group capping					

Possible impacts of (excessive) capping

- **Sysplex / multi system outage**
  - E.g. for LOCKs or RESERVEs not being freed timely
- **System outage**
  - E.g. for resources not being freed timely
  - Storage shortages
  - Work (e.g. SRBs) backed up, common storage shortage
- **Important work displaced**
- **SLAs missed**
- **Contention and increased promotion by SRM or dispatcher**
  - Can be unproblematic if displaced work is truly independent from important work – no shared resources
- **Less important work displaced**
- **Goals missed**
- **Increased response times**
- **Increased CPU delays**



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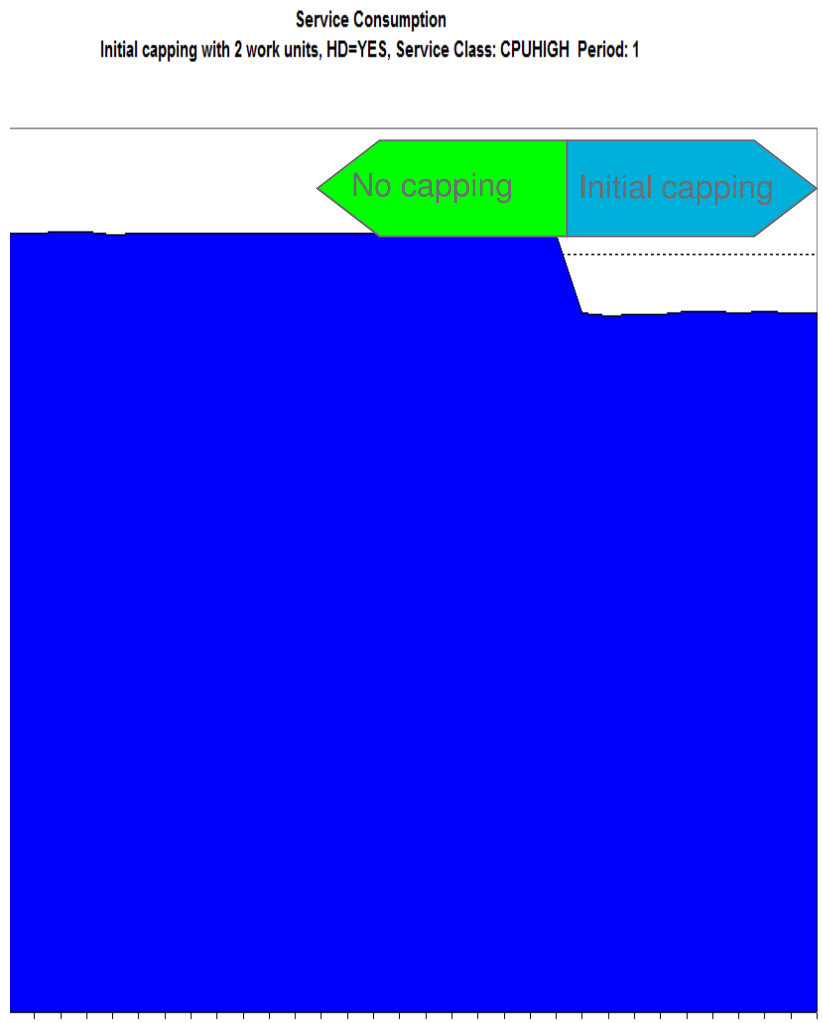
Initial capping (aka “hard capping”)

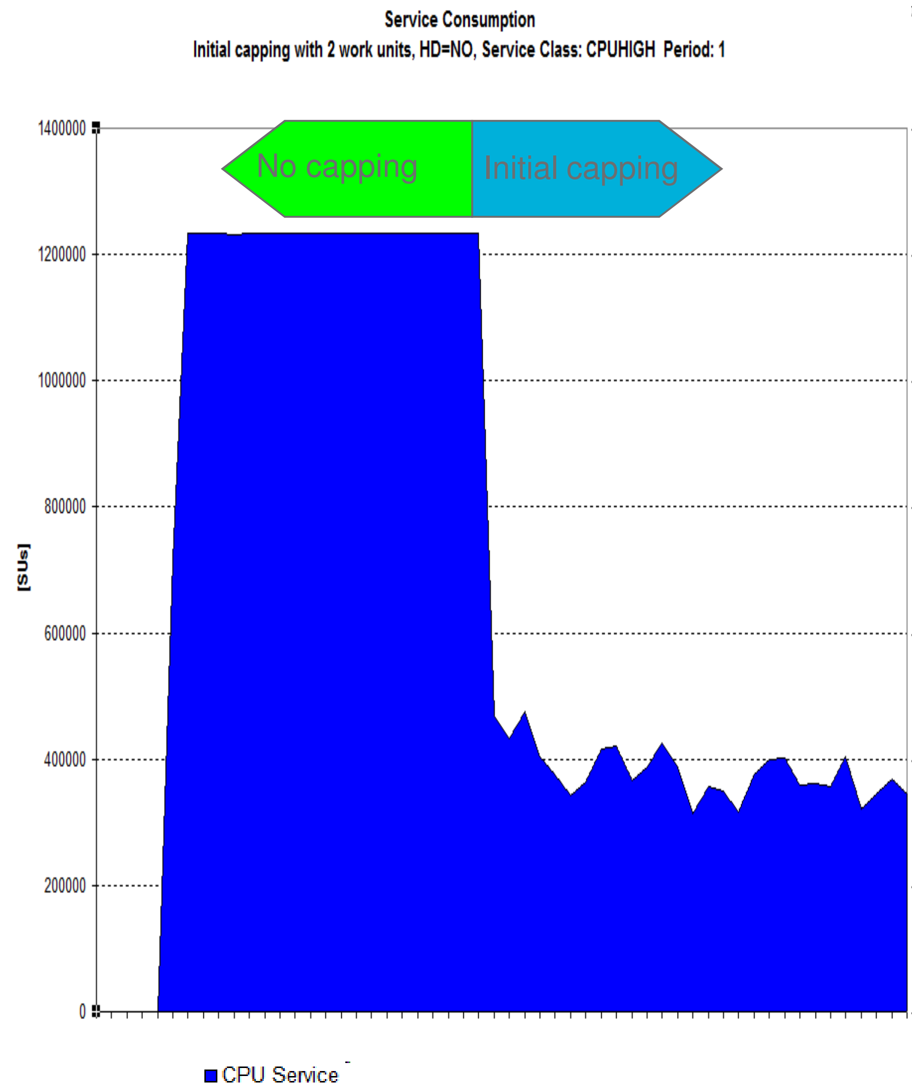
- Defined to PR/SM per processor type. Managed by PR/SM through limiting the processor time available to the LP’s logical processors
- The LPAR capacity is capped to LPAR share of CPC shared capacity

$$\text{LPAR}_i \text{ share} = \frac{\text{Weight}_i}{\sum_{\text{All activated LPARs}} \text{Weight}_j}$$

- LPAR weight is distributed across online CPs of the given type
- With HiperDispatch=NO an LP’s share is divided by the number of online logical CPs
  - Capping is done on a logical CP basis.
    May result in over capping if not all LCPs can be utilized
  - Consider following example:
    zEC12-732, 10 CPs online, Share=5.6%, low CPC utilization
    Workload: 2 TCBs

Initial Capping with HiperDispatch=Yes vs. No





Initial Capping with HiperDispatch=Yes vs. No

CPU Activity Reports

With HiperDispatch=Yes the high/medium processors receive a higher processor share.

- CPU 2827 CPC CAPACITY 3665 SEQUENCE CODE 00000000000
- MODEL 732 CHANGE REASON=NONE **HIPERDISPATCH=YES**
- H/W MODEL H43

---CPU--- NUM	TYPE	ONLINE	LPAR BUSY	MVS BUSY	PARKED	LOG PROC SHARE %	
0	CP	100.00	89.12	97.67	0.00	100.0	HIGH
1	CP	100.00	87.50	97.83	0.00	80.4	MED
2	CP	100.00	2.51	82.33	96.54	0.0	LOW
3	CP	100.00	1.87	63.68	96.54	0.0	LOW
4	CP	100.00	0.01	-----	100.00	0.0	LOW
5	CP	100.00	0.01	-----	100.00	0.0	LOW
6	CP	100.00	0.01	-----	100.00	0.0	LOW
7	CP	100.00	0.01	-----	100.00	0.0	LOW
A	CP	100.00	0.01	-----	100.00	0.0	LOW
B	CP	100.00	0.01	-----	100.00	0.0	LOW
TOTAL/AVERAGE			18.10	96.92		180.4	

- MODEL 732 CHANGE REASON=NONE **HIPERDISPATCH=NO**
- H/W MODEL H43

---CPU--- NUM	TYPE	ONLINE	LPAR BUSY	MVS BUSY	PARKED	LOG PROC SHARE %
0	CP	100.00	14.61	54.28	------	18.0
1	CP	100.00	13.00	46.80	------	18.0
2	CP	100.00	10.71	31.82	------	18.0
3	CP	100.00	6.77	18.55	------	18.0
4	CP	100.00	4.22	6.44	------	18.0
5	CP	100.00	4.87	13.16	------	18.0
6	CP	100.00	1.75	2.72	------	18.0
7	CP	100.00	4.54	13.05	------	18.0
A	CP	100.00	4.02	10.40	------	18.0
B	CP	100.00	3.08	6.88	------	18.0
TOTAL/AVERAGE			6.76	20.41		180.0

Stability of initial cap limits

- The **MSU equivalent** for an initial cap limit changes when…
  - The initial weight of the capped LPAR is changed
  - LPARs are de/activated or the total weight changes due to initial weight changes
  - Temporary capacity is de/activated
    - CBU, On/Off CoD…
- May require manual intervention when
  - A particular MSU/MIPS number is guaranteed for an LPAR
  - A particular MSU number must not be exceeded for licensing reasons

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- **Absolute capping**
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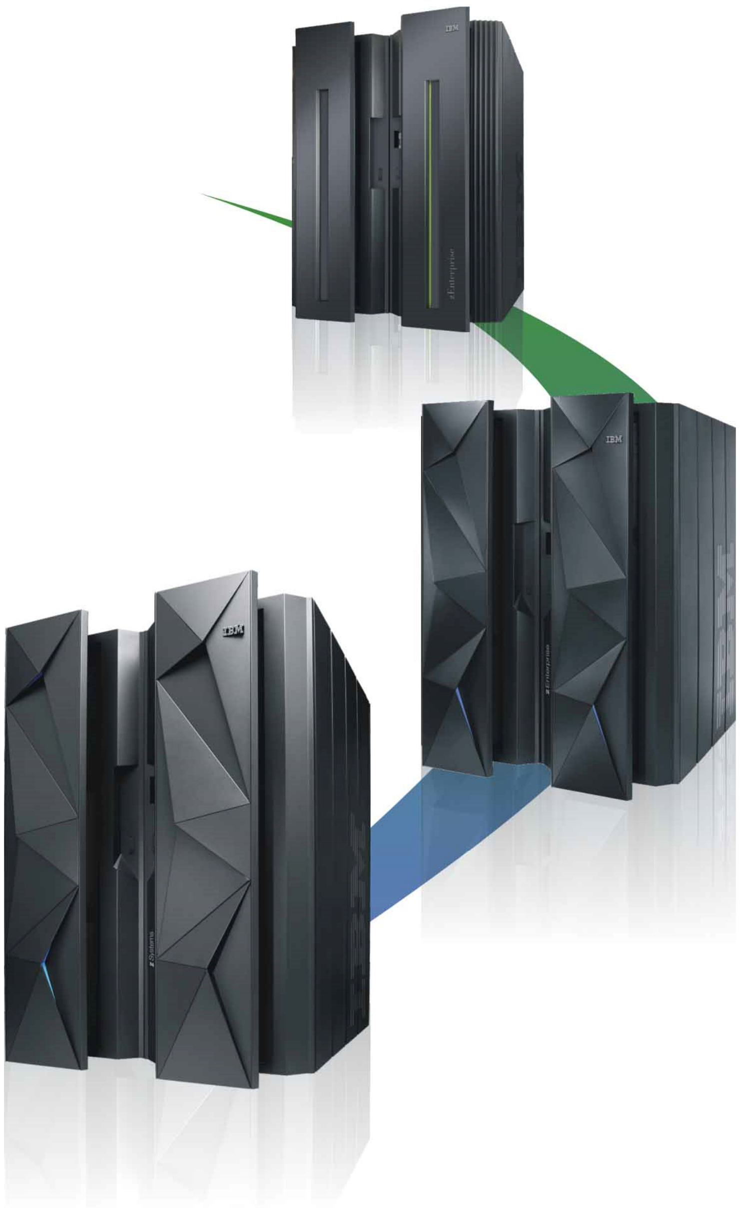

Absolute Capping Limit

- Defined to PR/SM per processor type. Managed by PR/SM through limiting the number of PR/SM time slices available to the LPAR's logical processors
- Specification in terms of (fractional) number of processors per processor type
  - E.g., 3.75 CPs
- Introduced with zEC12 GA2
- Primarily intended for non z/OS images
- Can be specified independently from the LPAR weight
  - But recommended to specify absolute cap above weight
  - WLM algorithms consider weight

Absolute Capping Limit

- Absolute capping may be used *concurrently* with defined capacity and group capacity management
  - The minimum of all limits becomes effective.
  - WLM/SRM is aware of the absolute cap, e.g. for routing decisions.
  - RCTIMGWU = MIN(absolute cap, defined capacity, group cap) when all capping types are in effect
    - RMF provides RCTIMGWU in SMF70WLA
    - In addition, SMF70HW_Cap_Limit value in hundredths of CPUs

Stability of absolute cap limits

- The effective limit for an absolute cap changes significantly when
  - the absolute cap value of the capped LPAR is changed, or
  - remporary capacity is de/activated AND the capacity level (processor speed) changes
    - I.e., general purpose processor CBU, On/Off CoD to/from subcapacity models
- The effective MSU rating for an absolute cap changes when the physical configuration changes

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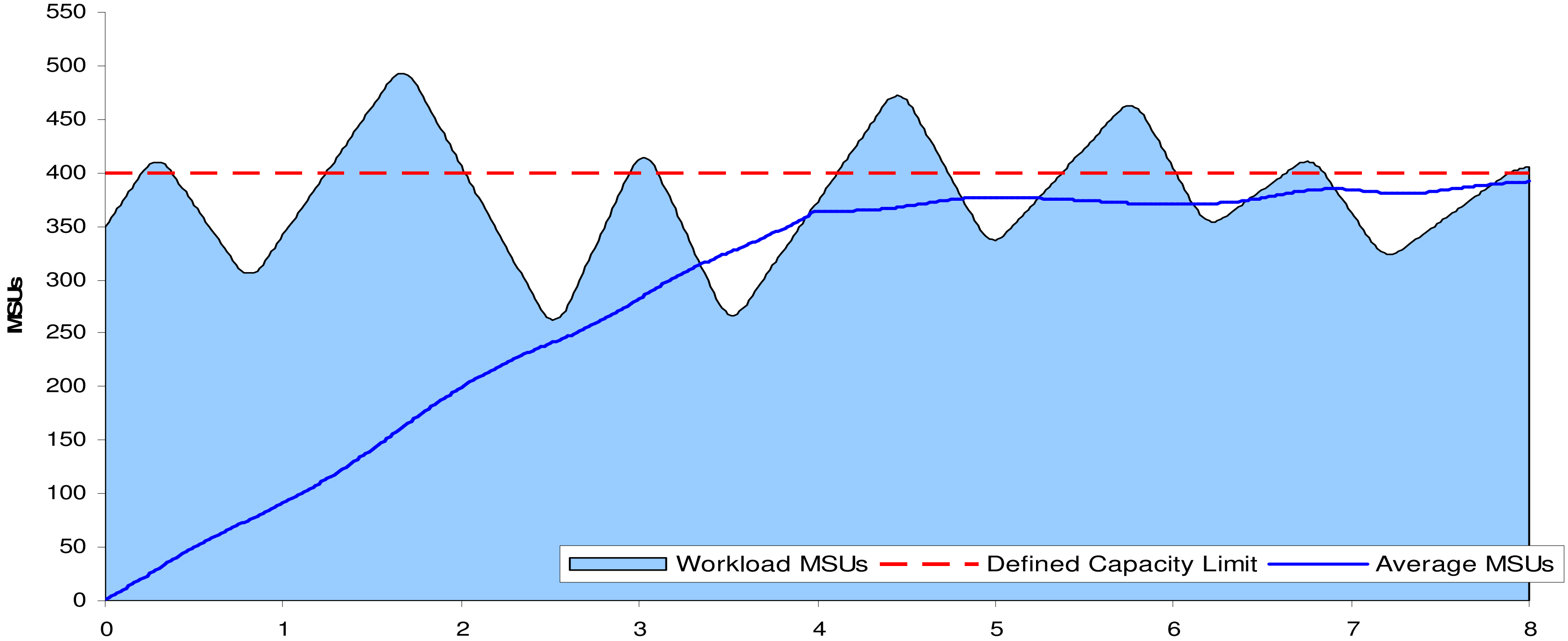
4 Hour Rolling Average

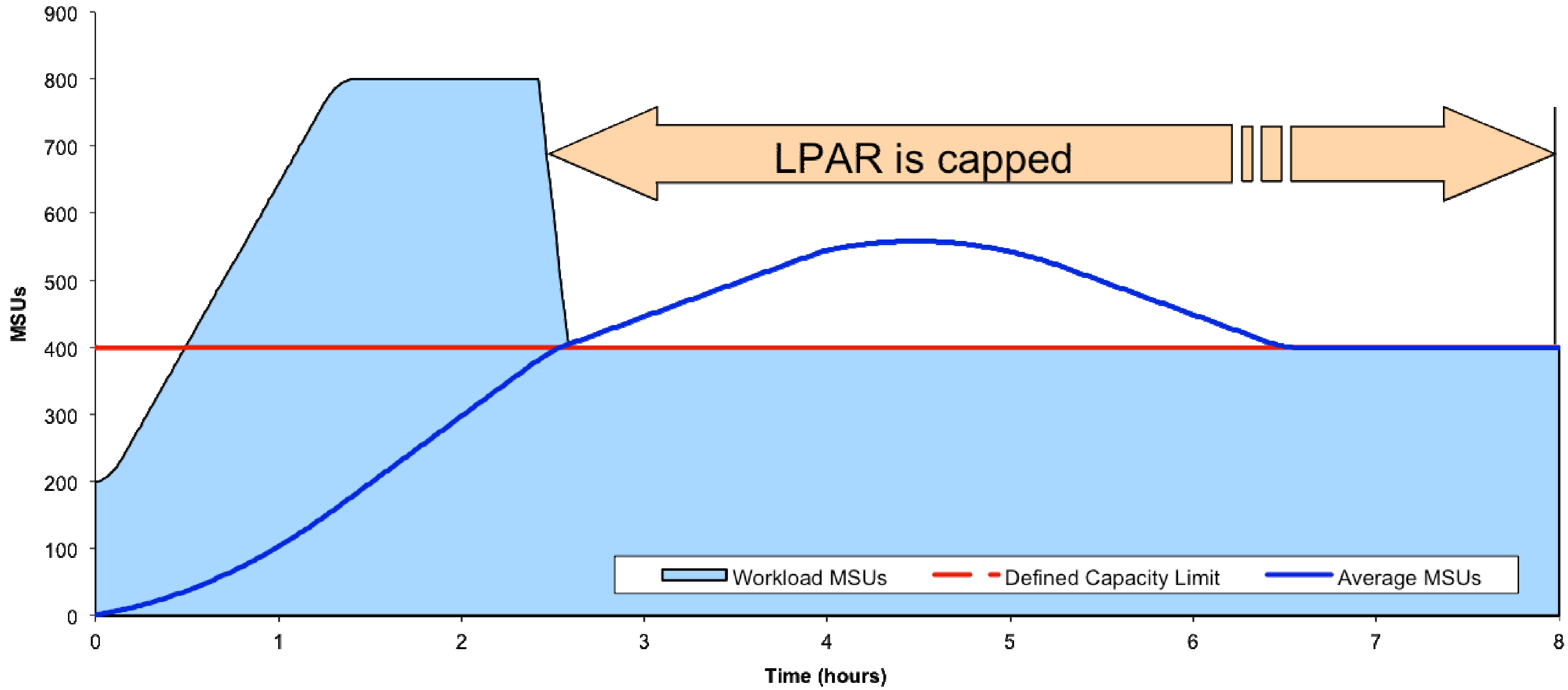


- Average consumption in LPAR in last 4h (rolling)
- MSU ≡ “Million Service Units per hour”

  ≠ Service Units • 3600 / 1000000

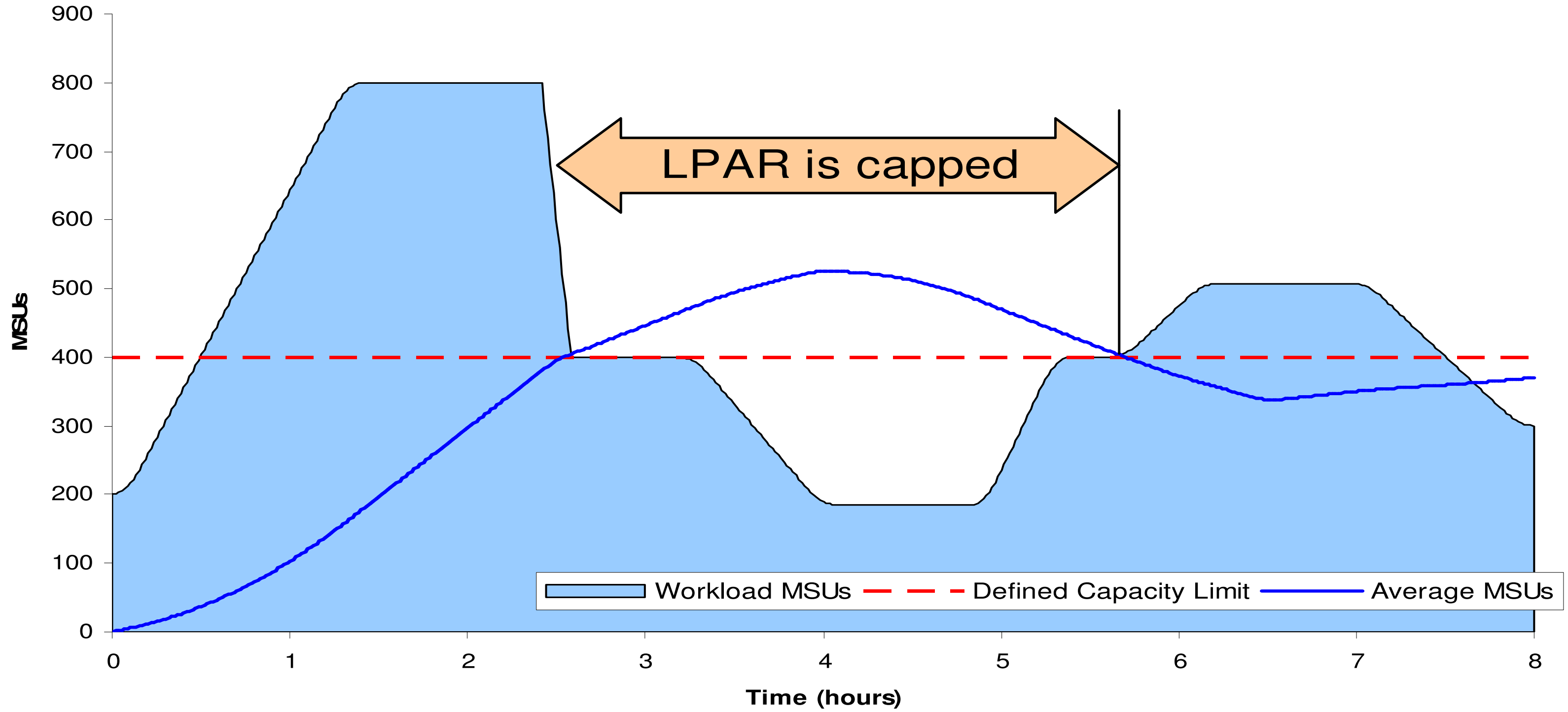
- Tracked as array of 48 intervals of 5 min = 4h

LPAR Capping



- An LPAR is -soft- capped when the 4HRA exceeds the defined capacity limit
- It remains capped until the 4HRA is below the defined limit
- When capped, the consumption is limited to the defined limit
- WLM advises PR/SM how to cap the LPAR

End of capping phase



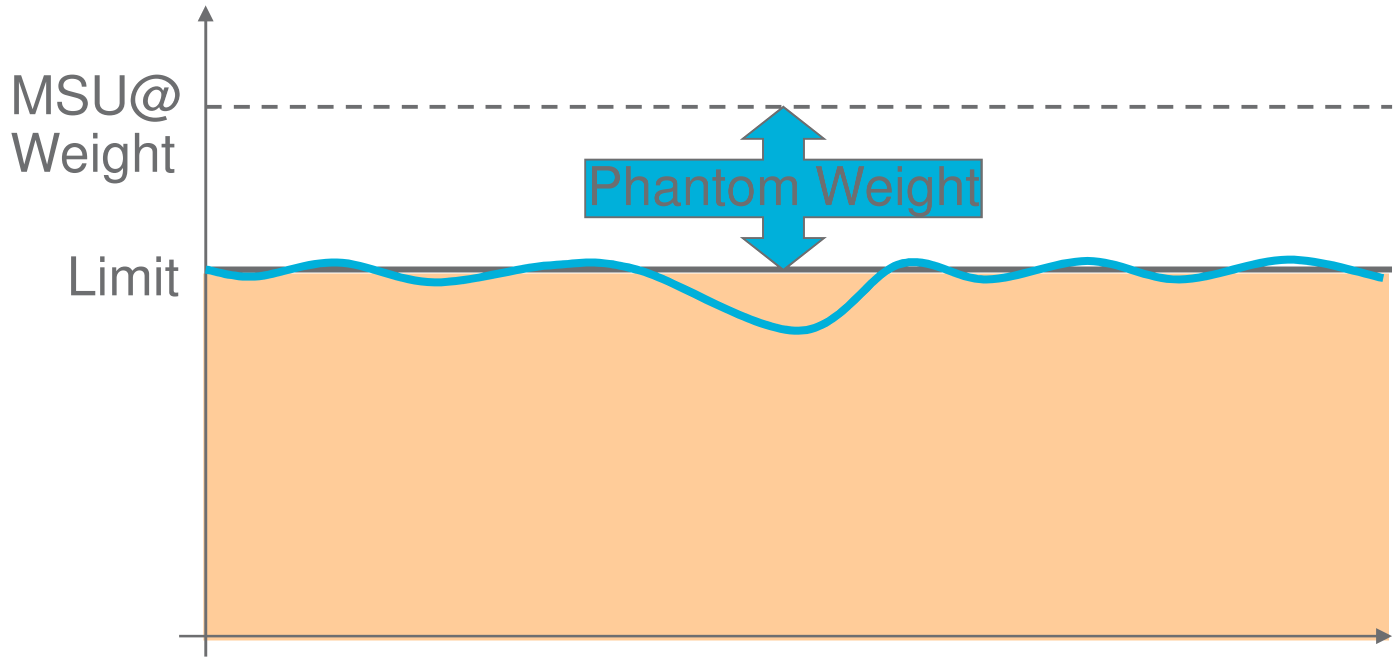
- Capping ends when the 4 hour average is below the softcap

Underlying soft capping techniques

- Historically, PR/SM algorithms were designed to cap a partition at its weight.
- Therefore, WLM and PR/SM use particular interfaces to cap a partition to an arbitrary MSU figure

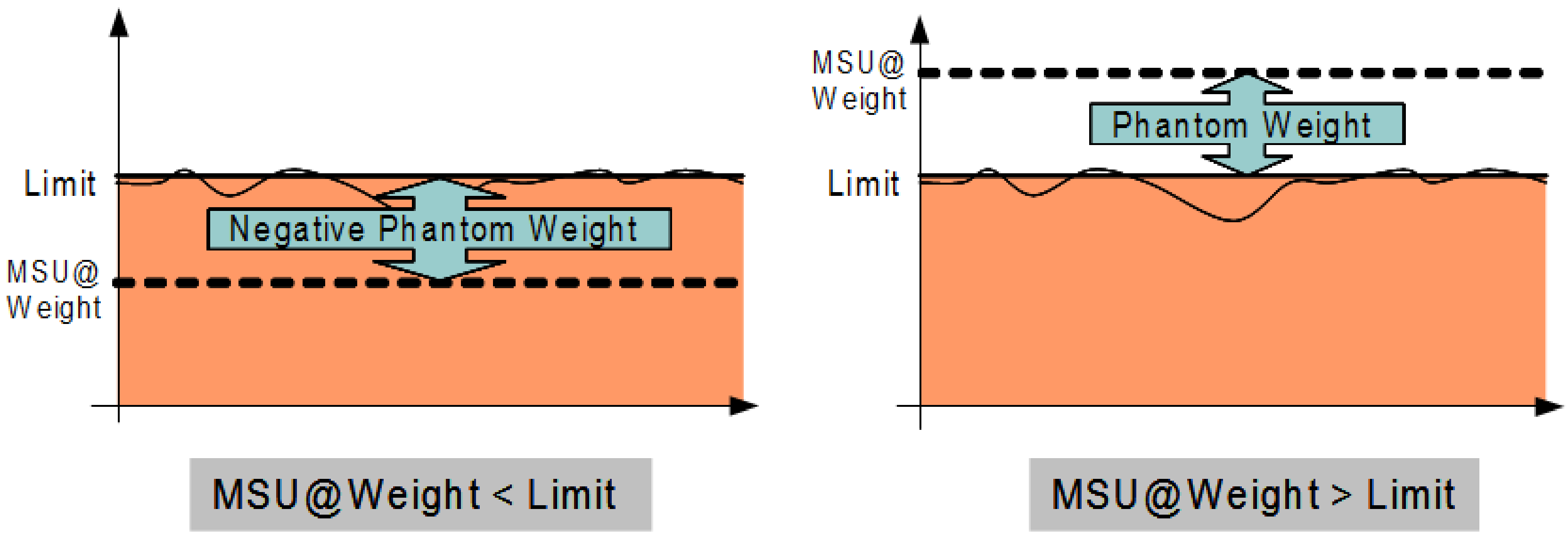
Weight vs. defined capacity limit	Hardware/Software level	Selected capping technique
MSU@weight > MSU imit	Any	**Phantom weight**
MSU@weight ≤ MSU limit	zEC12 GA2 and z/OS V2.1 or later	**Negative phantom weight**
	Other	Pattern capping

Phantom weight



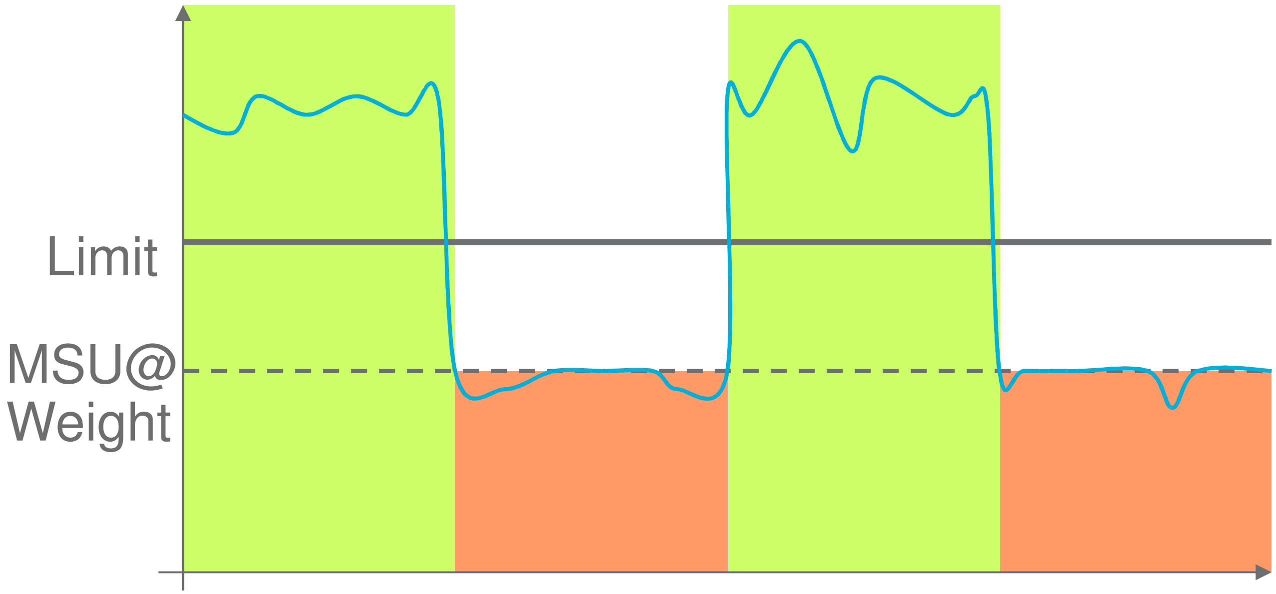
- Phantom weight is used to modify the PR/SM share of an LPAR
- WLM does not change a phantom weight as long as the limit and configuration do not change
  ➔smooth capping

Capping with phantom weight



- zEC12 with z/OS V2.1 and above support not only positive but also negative phantom weights.
  - Note: While a positive phantom weight changes the PR/SM priority of a partition, a negative phantom does not elevate the PR/SM dispatching priority.

Cap pattern

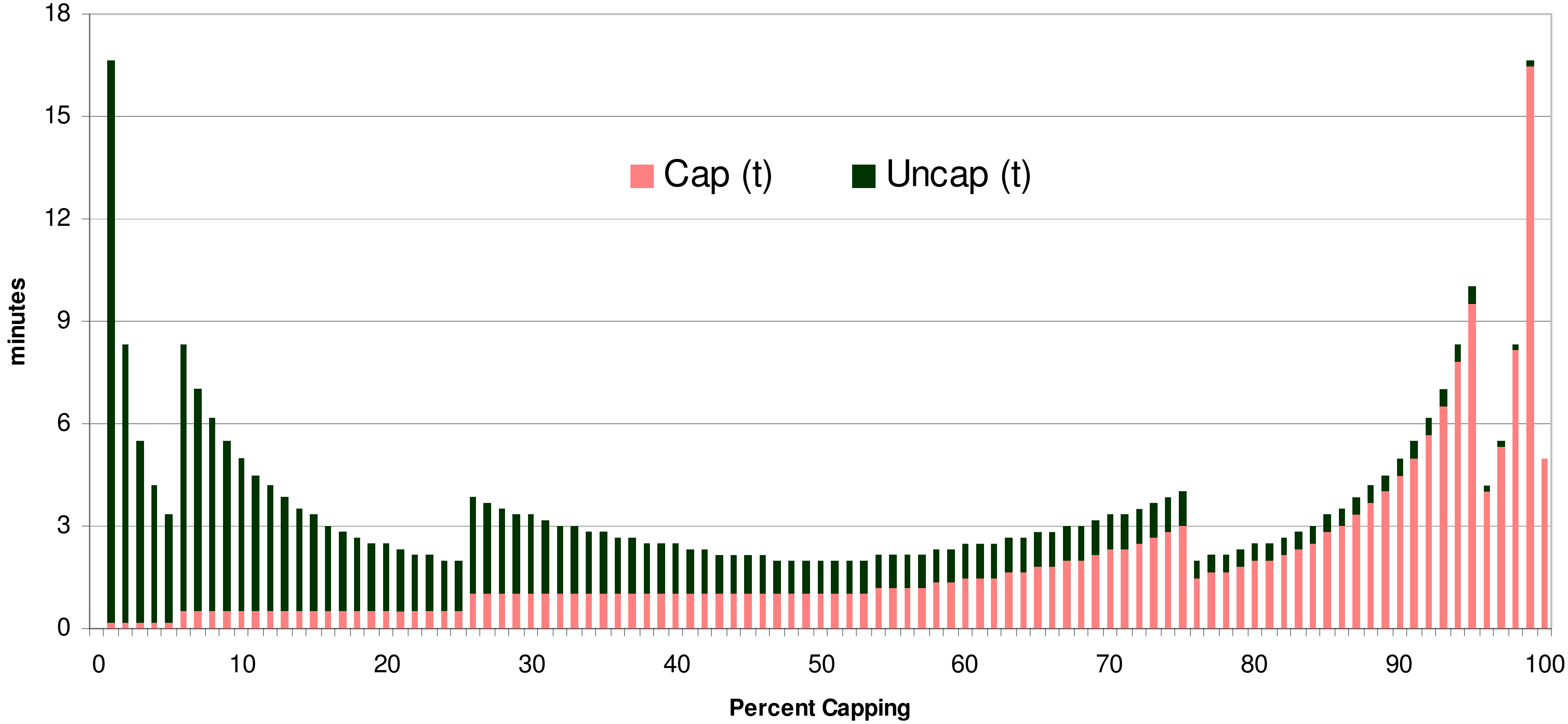


Prior to negative phantom weights WLM set up a cap pattern:
Alternating periods of

- LP capped to MSU@Weight, and
- LP uncapped

On average the MSU limit is enforced.

Cap pattern length

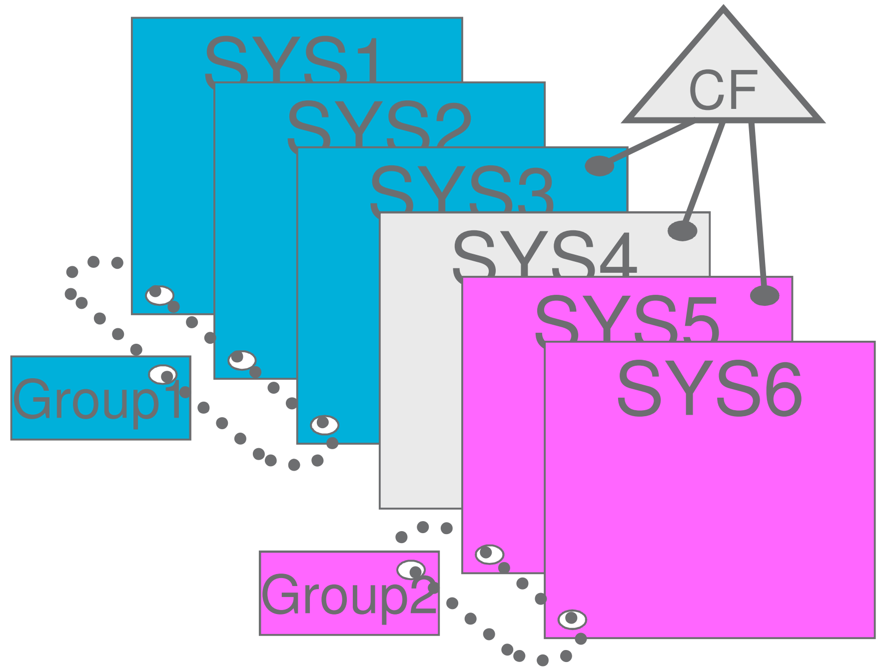


The LPAR cap pattern changes usually at an order of a few minutes.
The extreme cases are

01% WLM capping = 10 sec capped / 16.5min uncapped
99% WLM capping = 10 sec uncapped/ 16.5min capped

Group Capping

An LPAR capacity group can be used to enforce a MSU limit for a set of one or more LPARs.



- A capacity group is limited to a **single CPC** but independent from the Sysplex
- A system can be joined to one group at most
- A system will not join or will leave the capacity group when requirements not met
  - Namely, initial capping must not be active

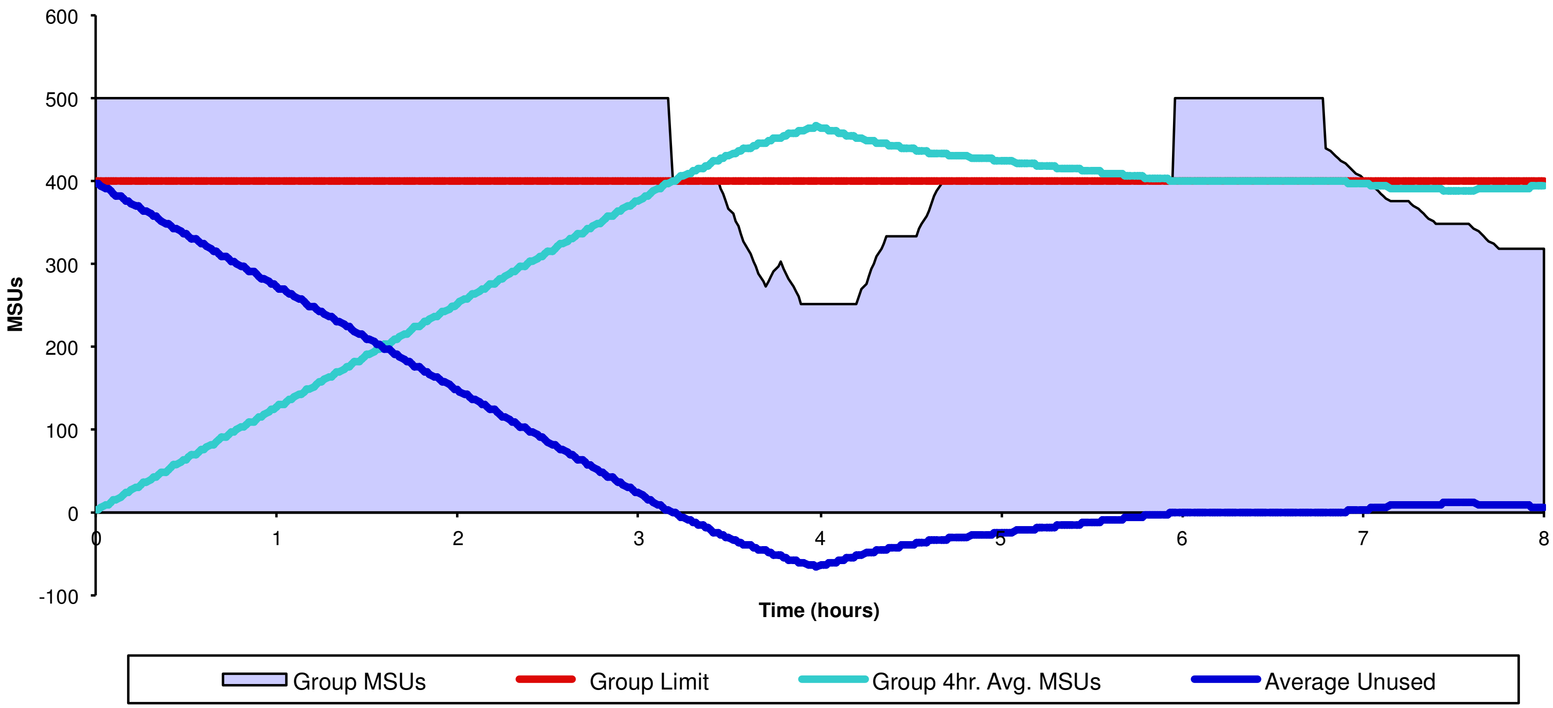
Group capping example

System	Weight	DC (MSU)	GC (MSU)	Initial GC Share (MSU)	Donation at full demand (MSU)	GC Entitlement (MSU)
SYS1	600	-	400	200	–	240
SYS2	300	-		100	–	120
SYS3	300	40		100	60	40

- The share of a group member is based on its *weight*
  - With IRD with zEC12 GA2 & z/OS V2.1: initial weight
  - With IRD in prior environments: current weight

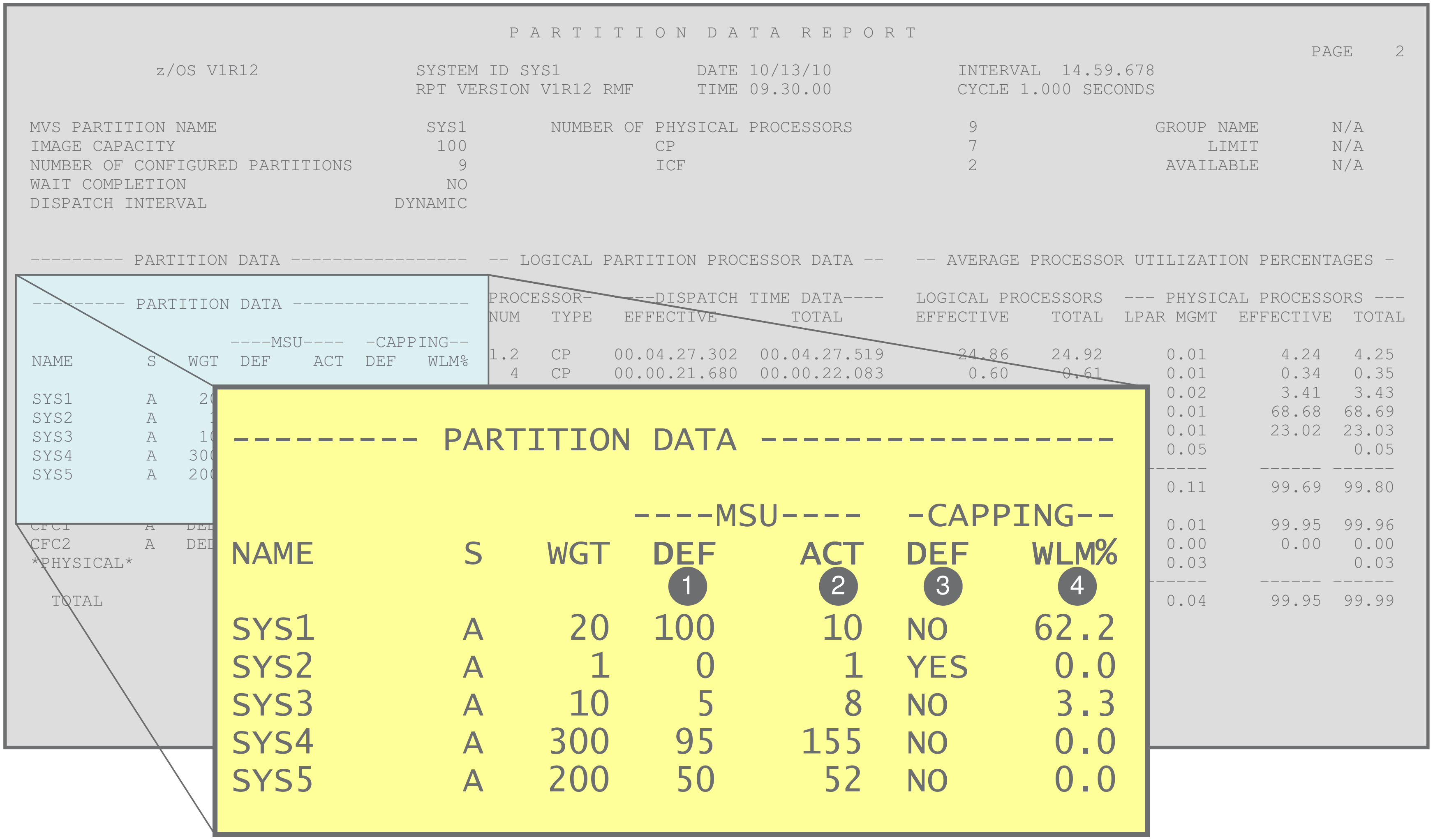
- Unused capacity is donated to other group members
  - …and re-distributed based on weight
- The minimum of DC and GC entitlement is used for capping an LPAR

Unused vector (group capping)



- Group capacity is tracked via an **un**used group capacity array of 48 intervals of 5 min
- Group capping is active when average unused group capacity negative
- Each system tracks unused capacity while joined to a capacity group
  - Not synchronized upon group changes: systems may have a different view for up to 4h

RMF: Partition Data Report



RMF: Partition Data Report

1. **MSU DEF** DC limit for this partition in MSU as specified on HMC
2. **MSU ACT** Actual avg. MSU consumption of this LPAR
3. **CAPPING DEF** Indicates whether this partition uses initial capping
4. **CAPPING WLM%** Portion of time the LPAR was capped during the RMF interval
   - Does not necessarily imply that the cap constrained the LPAR's consumption.

RMF: Group Capacity report

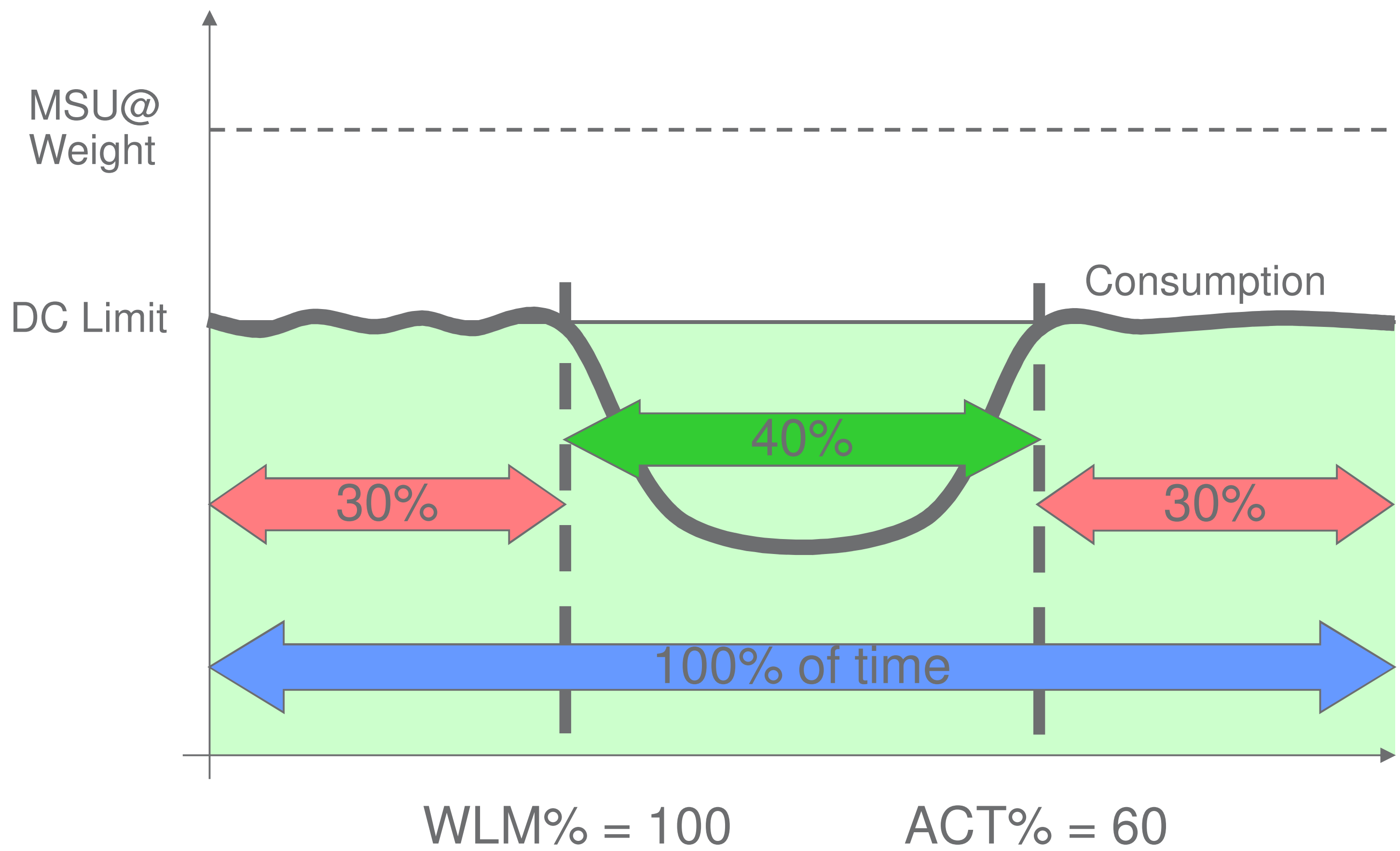
```
                          G R O U P   C A P A C I T Y   R E P O R T

       z/OS V1R12          SYSTEM ID SYS1              DATE 10/13/2010             INTERVAL 14.59.968
                           RPT VERSION V1R12 RMF       TIME 15.15.00               CYCLE 1.000 SECONDS

 ----GROUP-CAPACITY----   PARTITION   SYSTEM          -- MSU --    WGT    ---- CAPPING ----    - ENTITLEMENT -
 NAME       LIMIT    AVAIL                            DEF    ACT          DEF   WLM%    ACT%    MINIMUM MAXIMUM
  (1)        (2)      (3)                             (4)    (5)          (6)    (7)     (8)        (9)    (10)
 GROUP1     1500      -22  SYS1        SYS1            80      3    600   NO      25      23         80      80
                           SYS2        SYS2            80      3    500   NO     100      46         80      80
 -----------------------------------------           -----------------------------------------------------------
                                       TOTAL                   6   1100
```

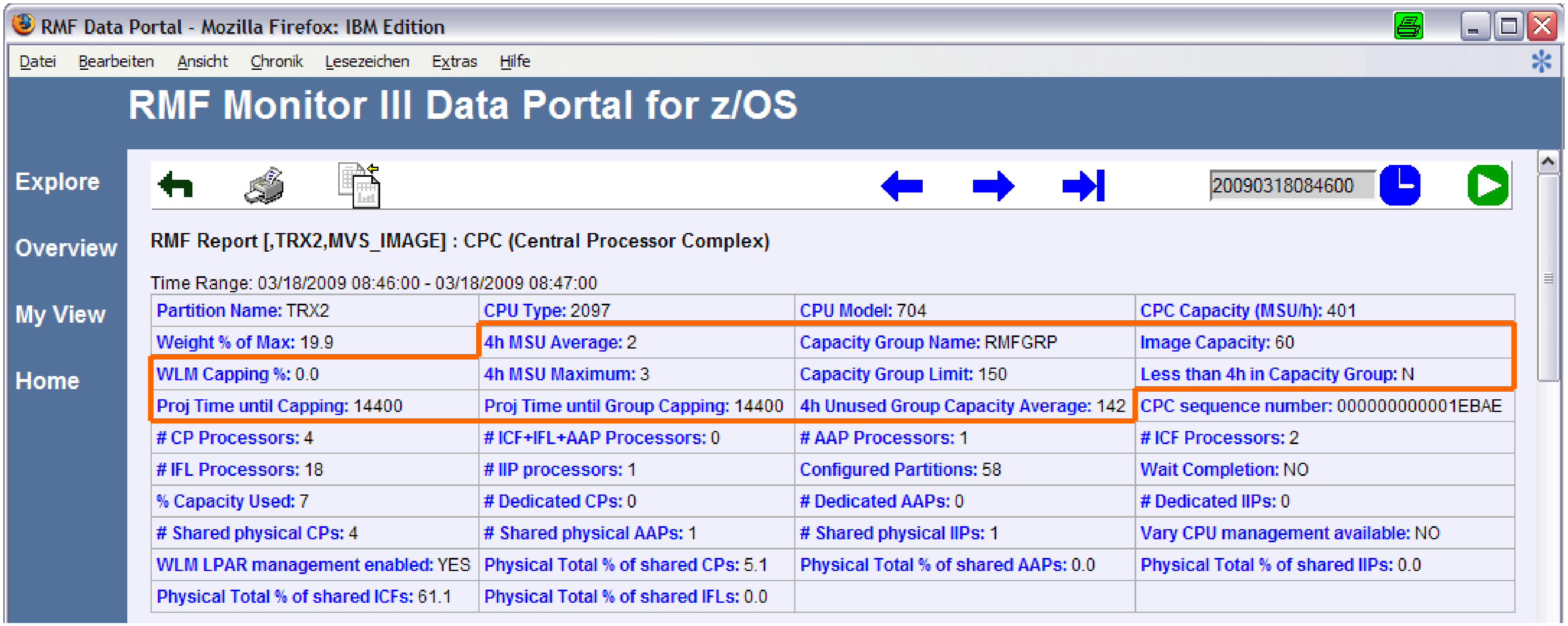
1. **NAME** Name of the WLM capacity group
2. **LIMIT** Group limit
3. **AVAIL** Average unused capacity in MSUs (avg. unused vector)
4. **MSU DEF** Defined capacity limit
5. **MSU ACT** Average used capacity
6. **CAPPING DEF** YES indicates that initial capping is active
7. **CAPPING WLM%** Percentage of time that WLM had set up a cap for the partition
8. **CAPPING ACT%** Percentage of time found capping actually limited the usage of processor resources for the partition
9. **MINIMUM ENT.** Minimum of the GC member share and the DC limit
10. **MAXIMUM ENT.** Minimum of the GC limit and the DC limit

Phantom weight: WLM% vs. ACT% in RMF



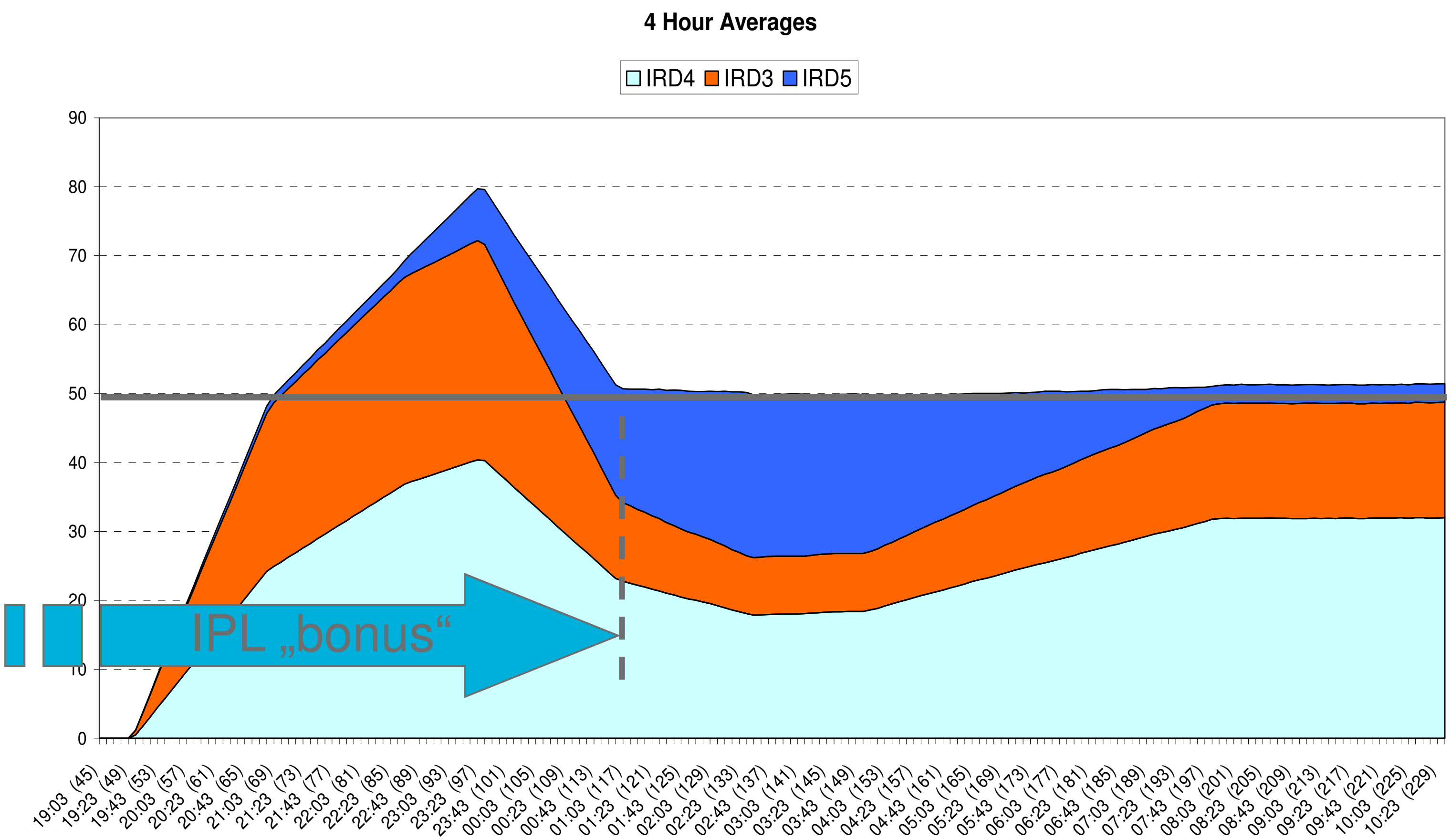
- RMF: WLM% is always 100 in case of phantom weight

RMF Data Portal



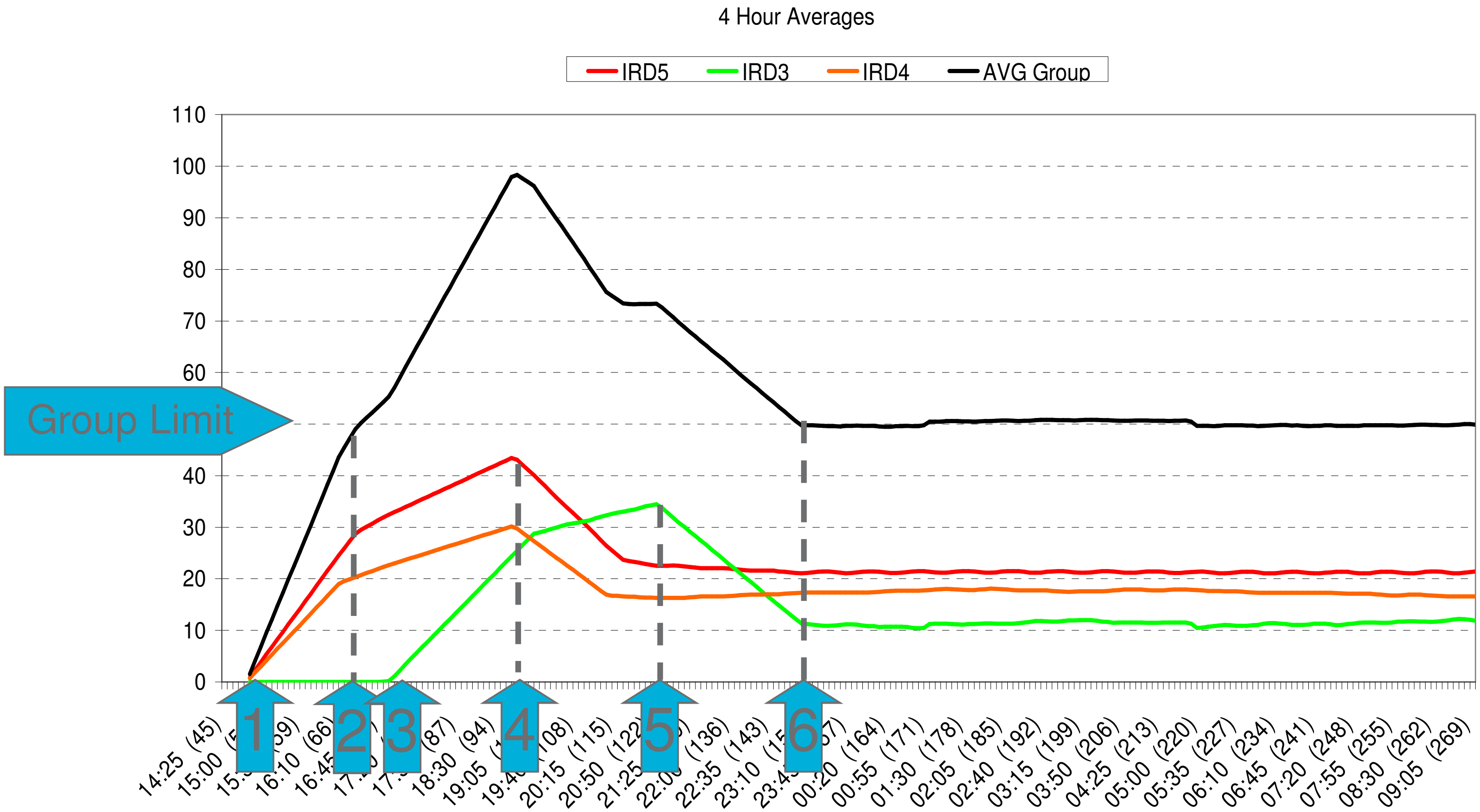
Many capping related fields are available in RMF Monitor III Data Portal

4 hour rolling average at IPL



Average is always for 4 hours even when the IPL was less than 4 hours ago

A member joins the capacity group



1. Workloads begin on IRD4 & 5
2. Group limit reached
3. System IRD3 joins group
4. IRD4 & 5: Four hours since (1.)
5. IRD3: Four hours since (3.). All systems have same GC view.
6. Group Avg. = Group limit

Capping and HiperDispatch

- z/OS V2.2 and V2.1 with APAR OA43622 provide some HiperDispatch enhancements that become effective when running capped, or when capped LPARs are present on the CPC.

Function \ z/OS release		V2.2	V2.1	V1.13
Hiper-Dispatch z13 & zEC12	***Unpark while capped*** ***Unused capacity refinement*** *Prime cycle elimination*	**+**	***OA43622***	

HiperDispatch "Unpark while capped"

- Previously, HiperDispatch
  - Parked all Vertical Low (VL) processors when a system capped via positive phantom weight
    - VLs are used for discretionary capacity and not required to absorb the LPAR weight
    - However, it was seen that, for some workloads, the reduced number of logical processors made it difficult to fully utilize the cap target capacity.
  - Unparked all VL processors when a system was capped by negative phantom weight, or some cases of PR/SM absolute capping

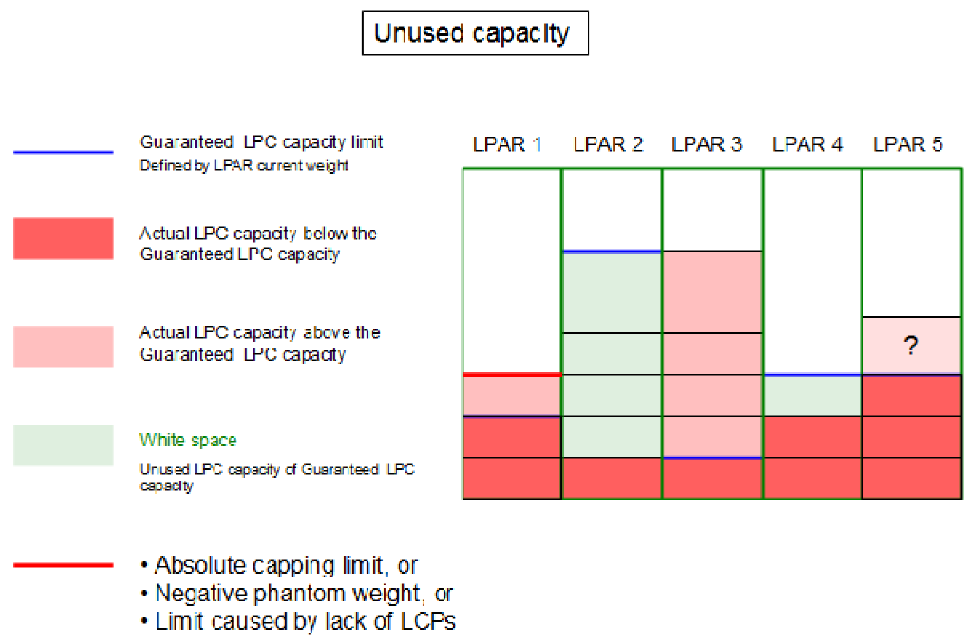
- Now, HiperDispatch can unpark VL processors <u>if</u> the processors can be used efficiently.

HiperDispatch refinement of “unused capacity” use

- HiperDispatch decisions are based on the CPC-wide unused capacity situation
- The ‘unused capacity share’ calculation was enhanced to also consider the LPAR configuration values
  - absolute capping value
  - negative phantom weight
  - number of logical processors
  - effective defined capacity and group capacity limit

  of possible ‘unused capacity’ receivers

CPC with 5 LPARs. LPAR1 has an absolute capping limit, which is indicated with the red line. LPAR2, and LPAR4 are unused capacity donors, while LPAR1 / 3 / 5 are unused capacity receivers.



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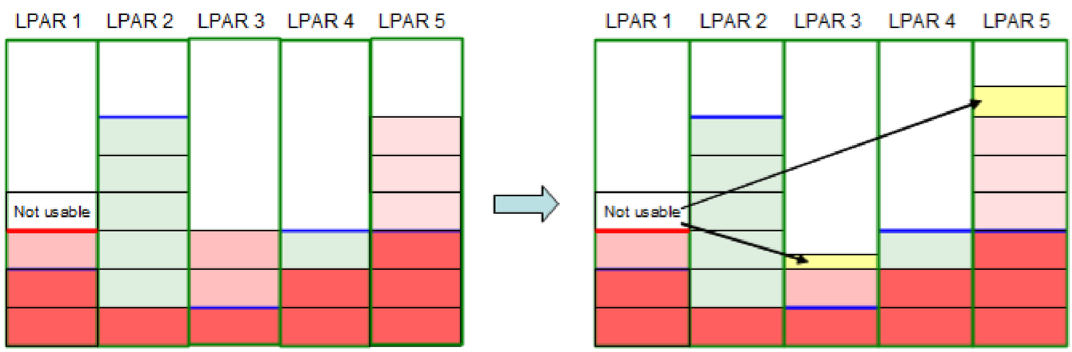
HiperDispatch refinement of “unused capacity” use

Enhanced unused capacity calculation



Today

Enhanced unused capacity calculation

- Figure on the left shows today’s unused capacity calculation, which does not consider LPAR capping limits.
- Unused capacity calculation is only based on the receiver’s weight share.
- Figure on the right shows an example of enhanced unused capacity calculation. It considers the capping limits of the receivers.
- Because LPAR1 is not able to use its total unused capacity share its ‘not usable’ unused capacity share portion increases the unused capacity share of LPAR5.

* Statements regarding IBM future direction and intent are subject to change or withdrawal, and represent goals and objectives only.

Agenda

- Overview of capping types
- Initial capping
- Absolute capping
- Defined capacity & group capacity
- **Resource group capping**
- 4HRA management

- Additional Material

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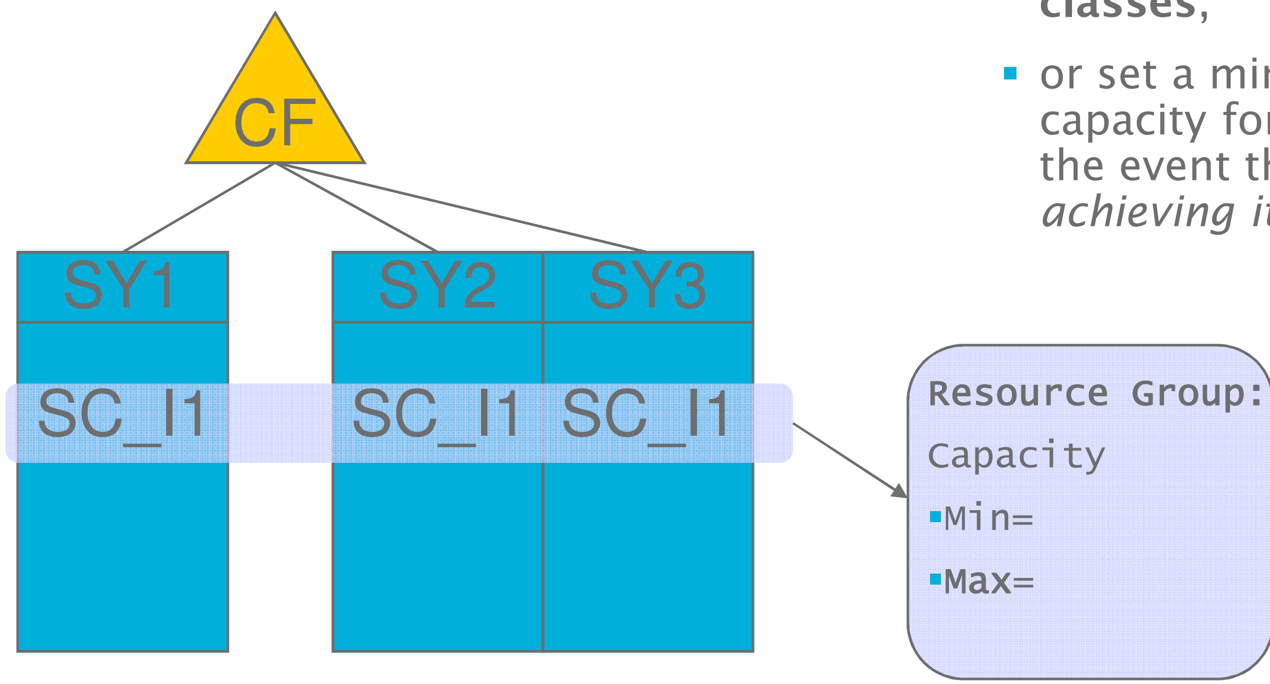
What is a Resource Group?

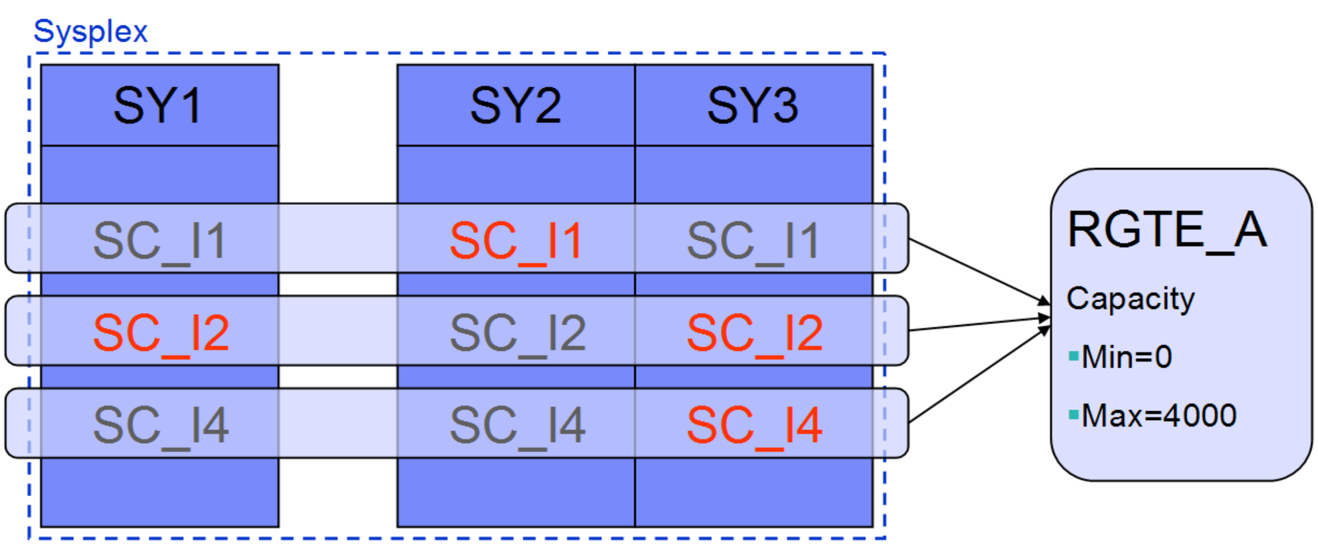
- Resource groups are a means to limit or protect work *when proper classification, goals and importance are not sufficient.*



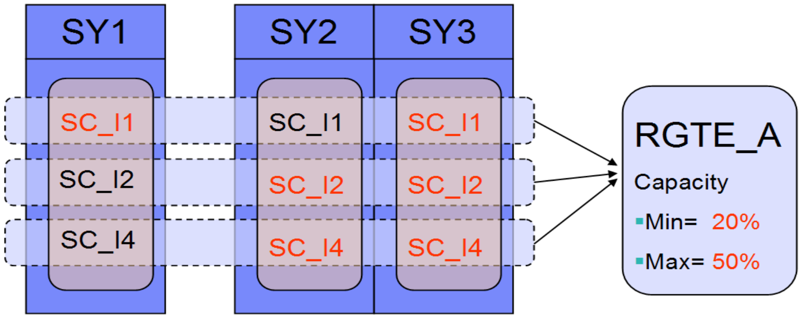
- A Resource Group is associated to one or more Service Classes
- Defines the service that the related Service Class(es) are managed to. Either
  - **limit the amount of processing capacity available to the service classes,**
  - or set a minimum processing capacity for the service classes in the event that the work is *not achieving its goals*

Type 1 Resource Groups

- Sysplex-wide defined in unweighted service units per second
  - "Unweighted" or "raw" meaning that the CPU and SRB service definition coefficients are not applied
- Sysplex-wide managed
- General Considerations
  - Multiple service classes may be assigned to a resource group
    - Different utilizations on the different systems and mix of importance levels make it difficult to predict actual consumption
  - Systems may have different capacities



Type 2 and 3 Resource Groups



- Sysplex-wide defined, but definition applies to each system
- Managed by each system
- General Considerations
  - Multiple service classes can be assigned to a resource group but this has no sysplex-wide effect
  - Definition is based on one of two possible units:
    - **Type 2: Percentage of LPAR capacity**
    - **Type 3: In number of processors (100 = 1 CP)**

Locating LPAR SU/sec Numbers

The service units that

- The Service Unit information can be located in the “z/OS MVS Planning: Workload Management” manual CPU Capacity Table
- Or on IBM Resource Link https://ibm.biz/BdFHFv :

IBM zEnterprise EC12

Processor	STIDP Type	STSI Model Name	CPs	SU/SEC	SRMsec/RealSec
2827-701	2827	701	1	78048.7805	1811.5932
2827-702	2827	702	2	73394.4954	1811.5932
2827-703	2827	703	3	71428.5714	1811.5932
2827-704	2827	704	4	69868.9956	1811.5932
2827-705	2827	705	5	68085.1064	1811.5932
2827-706	2827	706	6	66945.6067	1811.5932
2827-707	2827	707	7	65843.6214	1811.5932

A 4-way LPAR on a zEC12 model 7xx server can deliver approx.
4 * 69869
~ **279476 SU/sec**

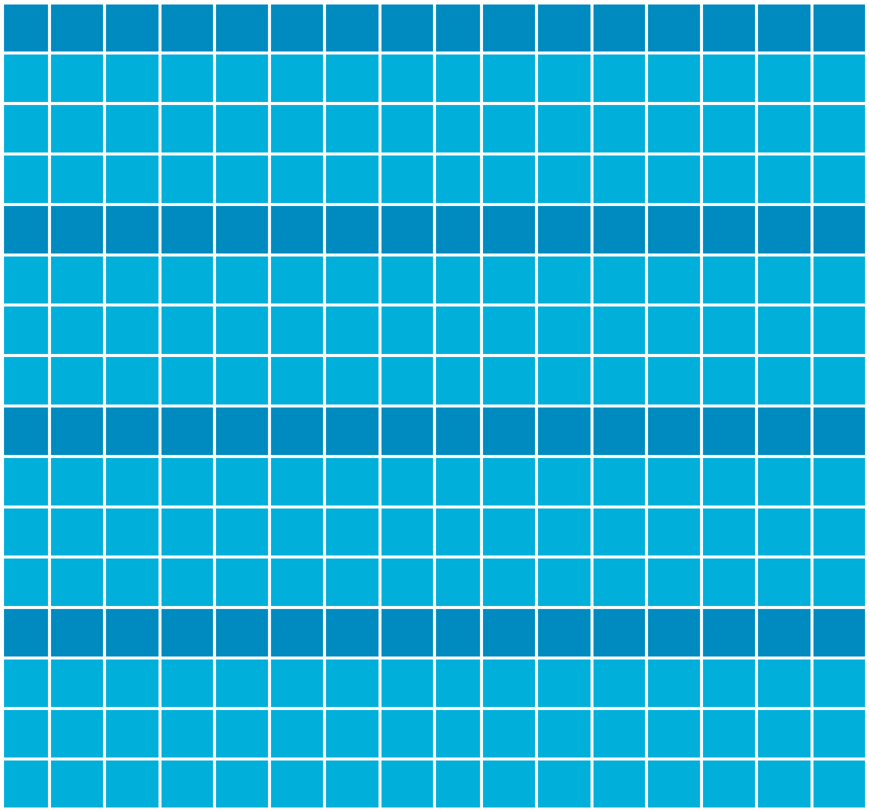
Resource Group Management

- To implement capping, the elapsed time is divided into 256 or 64 (pre-z/OS V2.1) slices. Each cap slice then represents $1/256^{th}$ or $1/64^{th}$ of the total elapsed time.

- Dispatchable units from address spaces or enclaves belonging to a resource group are made nondispatchable during some slices in order to reduce access to the CPU to enforce the resource group maximum.

- The time where address spaces or enclaves in a resource group are set non-dispatchable is called a CAP SLICE.
- The time where address spaces or enclaves in a resource group are set dispatchable is called an AWAKE SLICE.

Resource Group Maximum continued…



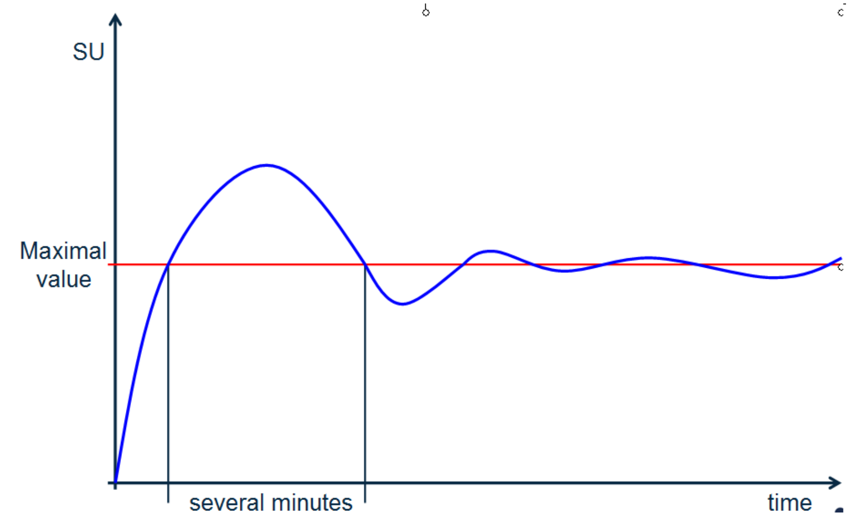
- This table is an example of a cap pattern with 64 awake slices and 192 cap slices.
- The active slices are distributed equally over the pattern

Resource Group Maximum continued…

- Every 10 seconds the policy adjustment code re-evaluates the resource groups and adjusts the cap pattern accordingly

- The forecast for the next 10 seconds is based on the average data from the last minute

- Because of the 1 minute average data, during a ramp up period, the max may be exceeded. Also, during periods of workload oscillation WLM may tend to under cap on the up swing but over cap when the workload is dropping off.



Resource Group Maximum continued…

Under certain conditions work may continue consuming service even while being capped

- Any locked work will continue to be dispatched as long as the lock is held
  - Check promoted times in RMF workload activity report
- The region control task is exempt from this nondispatchability.
- The address space will not be marked nondispatchable until the next dispatch.

Resource Group Considerations with zAAP/zIIPs

- Resource Groups are managed based on their general purpose processor consumption (TCB+SRB)

- Difficult to predict result of assigning RGs to service classes that execute on specialty processors
  - Especially when IFAHONORPRIORITY=YES or IIPHONORPRIORITY=YES is in effect.

1	9	17	25	33	41	49	57
2	10	18	26	34	42	50	58
3	11	19	27	35	43	51	59
4	12	20	28	36	44	52	60
5	13	21	29	37	45	53	61
6	14	22	30	38	46	54	62
7	15	23	31	39	47	55	63
8	16	24	32	40	48	56	64

Other considerations for Resource Groups

- **Not valid for transaction oriented work, such as CICS or IMS transactions.**
  - In order to assign a minimum or maximum capacity to CICS or IMS transactions, the region service classes can be assigned to a resource group.
    - Such interactive work can respond harshly to CPU bottlenecks: Evaluate what cap level can be tolerated

- **Given the combination of the goals, the importance level, and the resource capacity, some goals may not be achievable when capacity is restricted.**

- Unless there is a specific need for limiting or protecting capacity for a group of work, it is best to not define resource groups and to just let workload management manage the processor resources to meet performance goals.

Identifying Resource Group Capping

- In the RMF Workload Activity report, RG capping is identified in the Execution Delays section as CAP delays

- CAP delays may also be incurred by service classes that have not been associated with resource groups
  ➔ Discretionary Goal Management (DGM)

```
GOAL: EXECUTION VELOCITY 20.0%     VELOCITY MIGRATION:    I/O MGMT  93.9%     INIT MGMT 90.1%

          RESPONSE TIME EX    PERF  AVG    --EXEC USING%--   ------------- EXEC DELAYS % -----------   -USING%-
SYSTEM                  VEL%  INDX ADRSP  CPU AAP IIP I/O   TOT CPU CAP I/O                              CRY CNT

SYS1          --N/A--    93.9   0.2   0.0   46 N/A N/A  43   5.8 4.3 1.2 0.3                              0.0 0.0
```

Discretionary Goal Management (DGM)

- Allows an *eligible over-achieving* service class to donate CPU to a discretionary period
  - Objective is to improve service that discretionary periods receive when no non-discretionary periods need help and goals are vastly overachieved

- The donation is implemented through resource group capping.

- To be considered as a donor a period must meet several requirements, including
  - Not a member of a Resource Group (RG)
  - Non-aggressive goal:
    - If it has a velocity goal, the goal must be ≤ 30
    - If it has a response time goal, the goal must be > 60 sec
  - The performance index PI must be < 0.7

- If a period should never donate due to DGM, define appropriately:
  - Velocity goal > 30 or response time goal ≤ 60 sec, or
  - Define resource group with MIN=MAX=0 and associate service classes to be protected with that RG

Agenda

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- Additional Material

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4HRA business aspects

- Peak value of MIN(4HRA, defined capacity limit) over billing period determines software charges
  - 4HRA peaks may exceed the defined limit

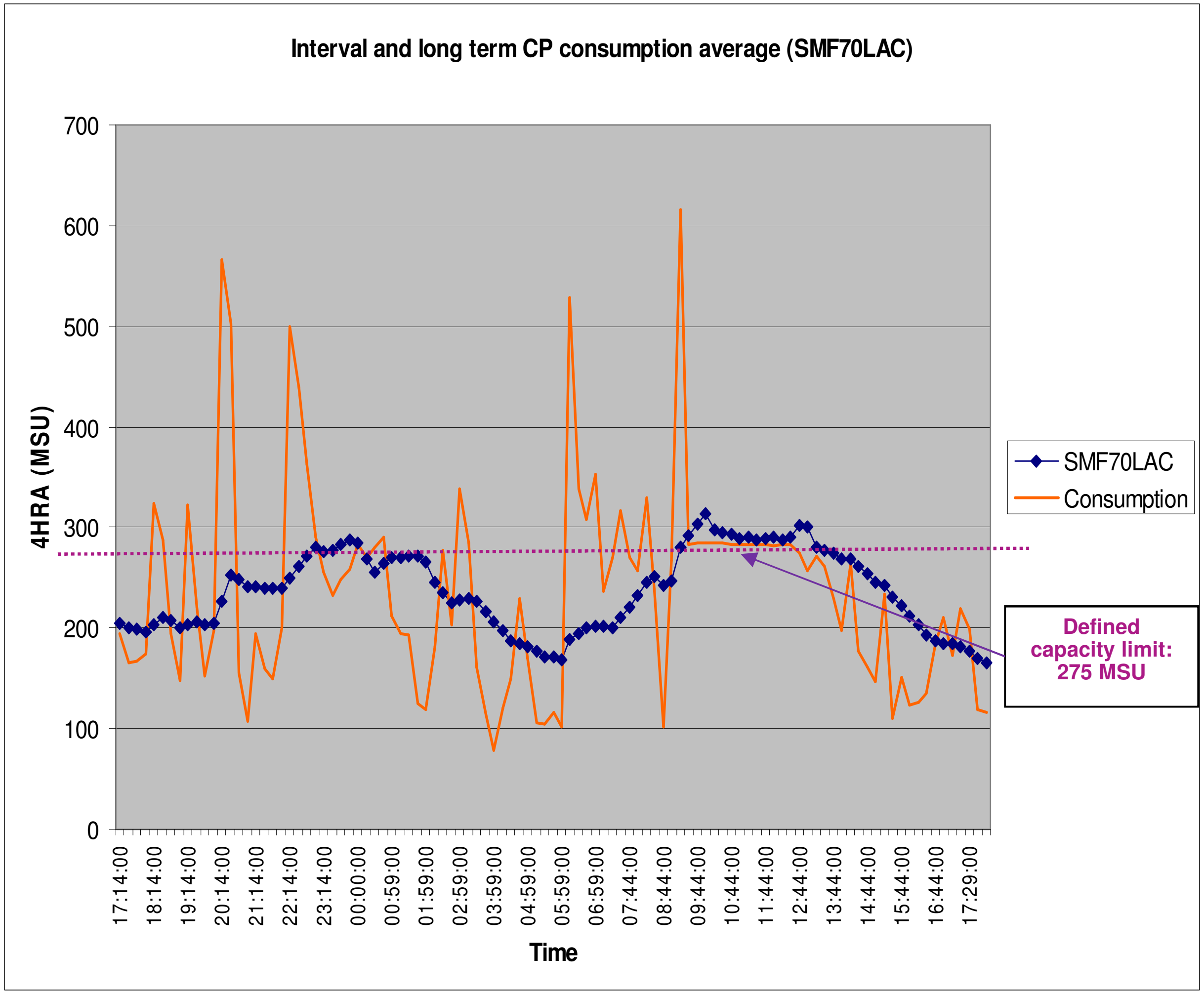
- Periods of low utilization can be used to “save” capacity for subsequent peak times
  - No capping when 4HRA < limit

- Utilization peaks drive up the 4HRA

- From a cost perspective it is usually desirable to **limit the peak consumption**

- Seek for technical means to
  - Limit consumption (→peak consumption)
    - Primarily of less important work
    - Also during -previously uncapped- periods
  - Maintain service levels, responsiveness and system integrity
    - Especially for important work

Interval consumption and the 4 hour rolling average: A sample day



Techniques for managing the 4HRA

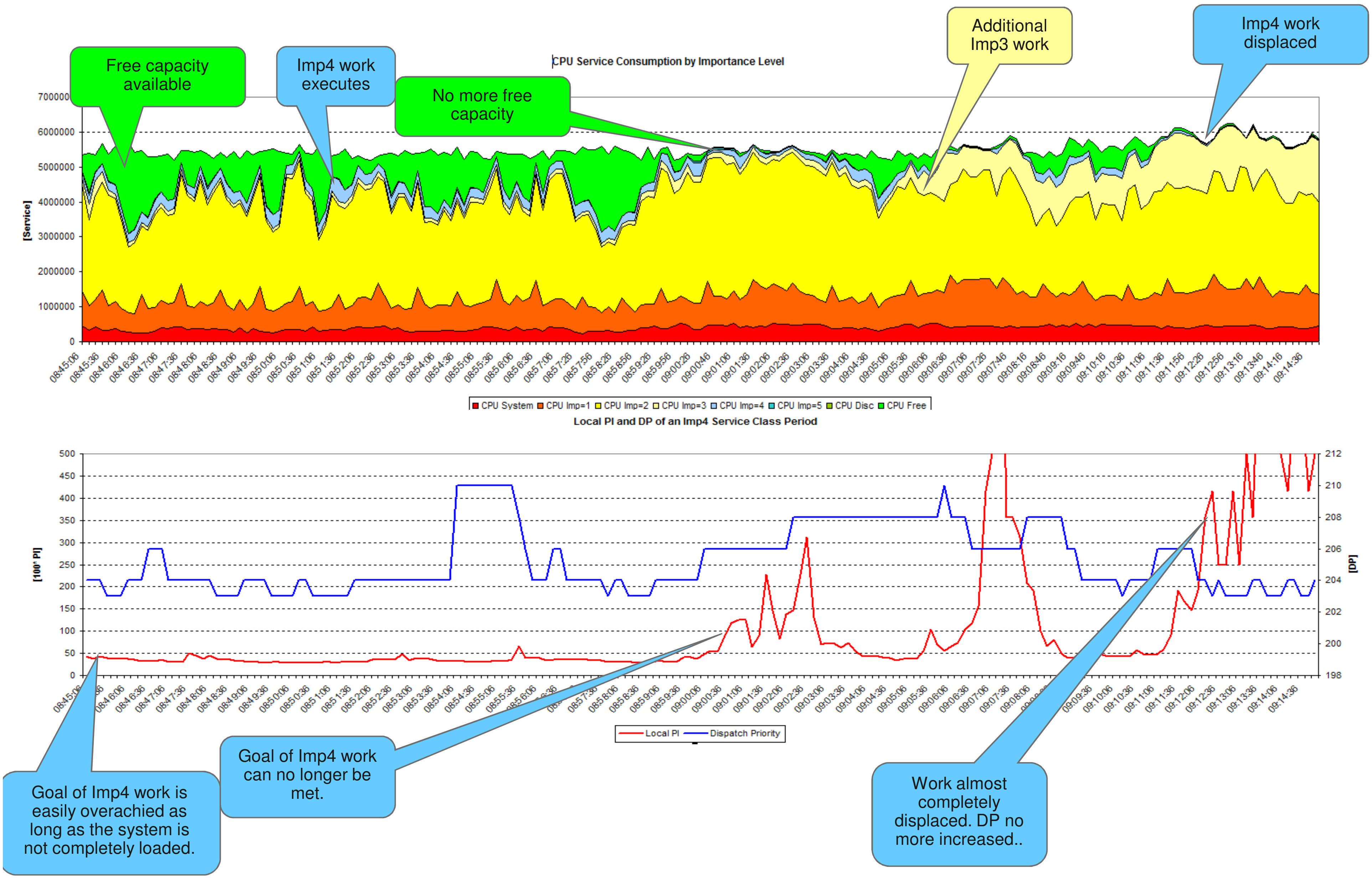
- Schedule work into off-peak hours

- Limit consumption at an LPAR level
  - Defined or group capacity
  - WLM importance level determines what work gets sacrificed first
    - 4HRA-wise irrelevant, but technically - beware of reduced preemption, promotion

- Selectively limit work within a system
  - Limit demand or parallelism
    - E.g. number of initiators
  - Resource groups
    - But not suitable for every work.

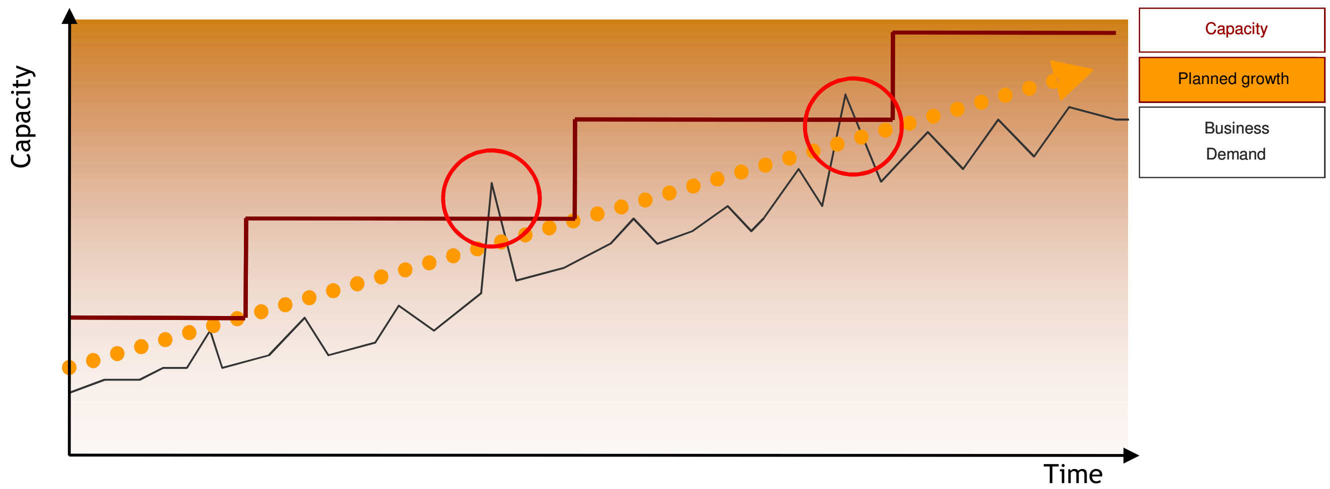
- Any combination of the above
  - Can also help to mitigate impacts of capping

Importance Distribution and Displacement of Work



IBM z/OS Capacity Provisioning Basics

- **Contained in z/OS base component free of charge**
  - Requires a monitoring component, such as z/OS RMF, or equivalent
  - Base element since z/OS V1.9
- **Exploits on System z On/Off Capacity on Demand Feature**
  - IBM zEnterprise System z10 or later
  - If On/Off CoD is not used CPM "analysis" mode may be used for monitoring and alerts
- **Exploits Defined Capacity and Group Capacity**
  - Defined Capacity with IBM System z10 or later
  - Group Capacity with IBM zEnterprise z196 or later



Capacity Provisioning Capabilities Overview

– The Capacity Provisioning Manager (CPM) can control additional capacity on IBM zEC12, z196, or z10 (plus BC10 and later)
  - Number of temporary zAAPs or zIIPs
  - Temporary general purpose capacity
- Considers different capacity levels (i.e. effective processor speeds) for subcapacity processors (general purpose capacity)
  - Can advise on logical processors
  - **Defined capacity and group capacity limits**
  - Can control one or more IBM zEnterprise or System z10 servers
    - Including multiple Sysplexes
  - Provides commands to control z196 and later static power save mode
  - Provides commands to control temporary IFLs

CPM allows for different types of provisioning requests:

- Manually at the z/OS console through Capacity Provisioning Manager commands
- Via user defined policy at specified schedules
- Via user defined policy by observing workload performance on z/OS

Policy Approach

The Capacity Provisioning policy defines the circumstances under which additional capacity may be provisioned:

- Three "dimensions" of criteria considered:
  - **When** is provisioning allowed
  - **Which** work qualifies for provisioning
  - **How much** additional capacity may be activated
- These criteria are specified as "rules" in the policy:

```
If
{   in the specified time interval
    the specified work "suffers"
}
Then up to
{       - the defined additional capacity
        may be activated

}
```

- The specified rules and conditions are named and may be activated or deactivated selectively by operator commands

Key benefit of CPM is the real time in-depth analysis of bottlenecks

```
MODIFY CPOSERV,APPL=REPORT WORKLOAD TYPE=DETAILED

Workload is analyzed for 1 system(s)
Workload for system PROD1 of sysplex PRODPLEX on CPC CPC1
   CICSHIGH.1 PL/PD/DL/DD/S 1.8 5 1.2 12 System
       PI from 11/16/2012 07:43 is 2.76
           Last limit crossing was 12/16/2012 07:27
       Demand for additional physical zIIPs not recognized
           System zIIP-utilization too low
       Demand for additional physical zAAPs not recognized
           System zAAP-utilization too low
       Demand for additional defined capacity recognized
       Demand for additional physical CPs not recognized
       Demand for capacity level increase not recognized
       Demand for additional logical CPs not recognized
           CPC-wide CP-utilization too low
```

- Key benefit of CPM is the real time in-depth analysis of workload constraints and demands
  - Based on WLM-provided metrics
- Can identify what type of capacity (if any) will help
- Timely reaction, even before capping begins

Capacity Provisioning Policy Strategies… for cost optimization

- Baseline defined or group capacity (DC/GC) limit relatively low
  - but still realistic for periods of low to average utilization

- Use Capacity Provisioning Manager rules to increase DC/GC limit
  - only when required by a qualifying workload during a qualifying time period
    - Time & workload conditions:
      Allow for higher DC/GC limits as required by workload

  - unconditionally during a qualifying time period
    - Time conditions without workload conditions:
      Unconditionally provision full rule scope

- When needed, can differentiate between different systems, service definitions, or override policies

Capacity Provisioning Policy sample scenario for cost optimization with LPAR defined capacity

- Sample scenario defines two qualifying workloads
  - Important online work
    - Monday through Friday, 07:45 – 18:00
    - Comprised of two service classes
      - DB2HIGH
      - ONLSTC
    - Up to +300 MSU may be provided in addition
  - Early evening batch
    - Monday through Friday, 20:00 – 22:00
    - Comprised of one service classe
      - BATCRIT
    - Up to +70 MSU may be provided in addition

Capacity Provisioning Policy Sample…
… with LPAR defined capacity (1)

- Two workloads that may warrant higher DC limits during different times of day:

Maximum Processor Scope | Logical Processor Scope | Maximum Defined Capacity Scope | Maximum Group Capacity Scope | Rules

Actions

Name	Description	Default Status
WeekNight	Weekdays DC pre midnight batch	Enabled
WeekdayDC	Weekdays DC for online work	Enabled

- WeekdayDC rule scope allows for up to +300 (additional) MSU:

Processor Scope | Defined Capacity Scope | Group Capacity Scope | Conditions

Actions

System	Sysplex	Max. Increase (MSU)
SYS1	PLEX1	300

Capacity Provisioning Policy Sample…
… with LPAR defined capacity (2)

- Rule is enabled for all weekdays prime time

Nonrecurring Time Conditions | Recurring Time Conditions | Workload Conditions

Actions

	Name	Start Date	End Date	Mon	Tue	Wed	Thu	Fri	Sat	Sun	Start Time ▲	Deadline	End Time
	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter
☐	AllWeekD	Jan 2, 2014	Dec 31, 2014	☑	☑	☑	☑	☑	☐	☐	7:45 AM	6:00 PM	6:30 PM

- Workload is defined by specific service classes

Importance Filters | Included Service Classes | Excluded Service Classes

Actions

	Service Definition	Service Policy	Service Class	Period	Provisioning PI	Provisioning Duration (Minutes)	Deprovisioning PI	Deprovisioning Duration (Minutes)
	Filter	Filter	Filter	Filter	Filter	Filter	Filter	Filter
☐	Any service definition	Any service policy	DB2HI	1	1.4	2	1.1	10
☑	Any service definition	Any service policy	ONLSTC	1	1.5	5	1.1	10

Capacity Provisioning Policy Sample…
… with LPAR defined capacity (3)

- Similarly, another rule is defined to cover a batch workload
  - Up to +70 MSU for a single batch service class

Nonrecurring Time Conditions | **Recurring Time Conditions** | Workload Conditions

Actions

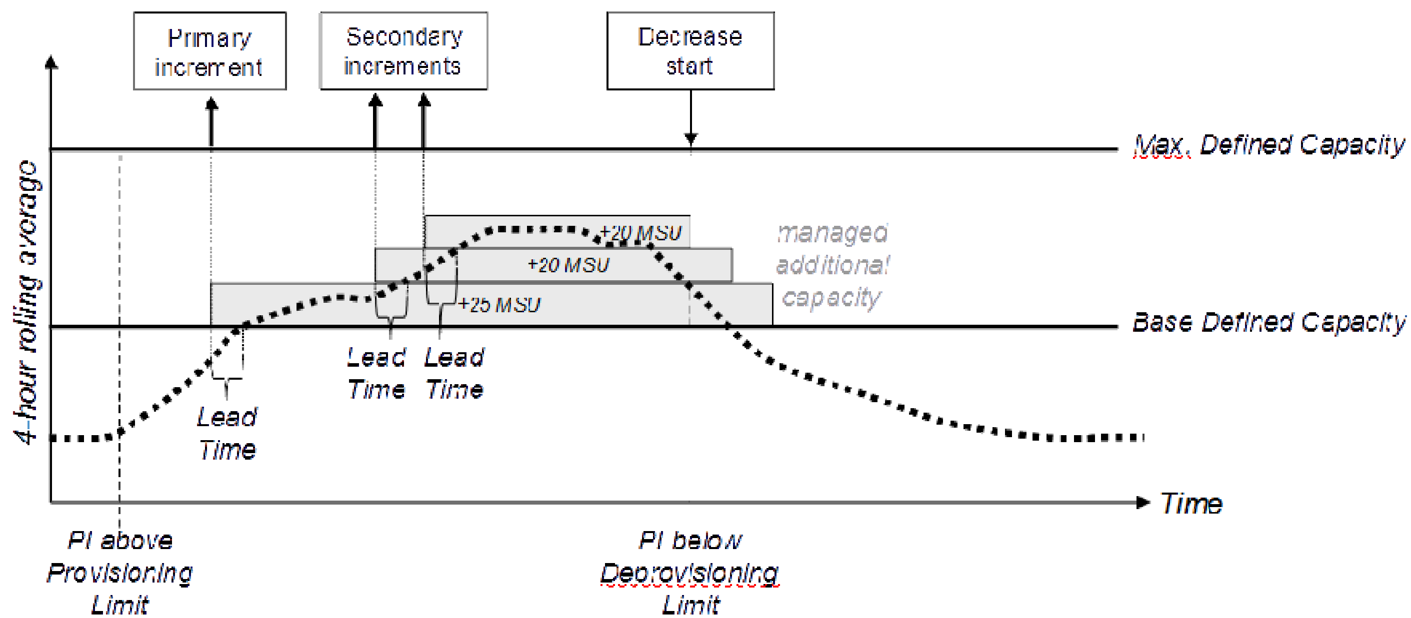
	Name Filter	Start Date Filter	End Date Filter	Mon Filter	Tue Filter	Wed Filter	Thu Filter	Fri Filter	Sat Filter	Sun Filter	Start Time ▲ Filter	Deadline Filter	End Time Filter
	AllWeekN	Jan 2, 2014	Dec 31, 2014	✓	✓	✓	✓	✓			8:00 PM	10:00 PM	10:00 PM

Importance Filters | **Included Service Classes** | Excluded Service Classes

Actions

	Service Definition Filter	Service Policy Filter	Service Class Filter	Period Filter	Provisioning PI Filter	Provisioning Duration (Minutes) Filter	Deprovisioning PI Filter	Deprovisioning Duration (Minutes) Filter	PI Scope Filter
	Any service definition	Any service policy	BATCRIT	1	1.8	5	1.3	10	System

Capacity Provisioning Defined Capacity Management



- When required by the defined workload the CPM will increase the defined capacity limit while the workload criteria are met
- The additional defined capacity will be managed down as the workload permits
  – Or deferred, based on user specification
- Additional user-initiated DC/GC activations are recognized and tolerated.

z/OS Capacity Provisioning Documentation

- For more information contact: IBMCPM@de.ibm.com
- z/OS Capacity Provisioning: Introduction and Update for z/OS V2.1, SHARE in Anaheim, Session 14210, 8/2013
- Website http://www.ibm.com/systems/z/os/zos/features/cpm
- z/OS MVS Capacity Provisioning User's Guide, SC34-2661, at http://publibz.boulder.ibm.com/epubs/pdf/iea3u100.pdf
- ITSO Redbook:
  System z10 Enterprise Class Capacity on Demand, SG24-7504
  http://www.redbooks.ibm.com/abstracts/sg247504.html?Open

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धन्यवाद
Hindi

多謝
Traditional Chinese

ขอบคุณ
Thai

Спасибо
Russian

Gracias
Spanish

Thank You
English

شكرا
Arabic

Obrigado
Brazilian Portuguese

多谢
Simplified Chinese

Danke
German

Bedankt
Dutch

Grazie
Italian

Merci
French

நன்றி
Tamil

ありがとうございました
Japanese

감사합니다
Korean

z/OS Workload Management - More Information -

- z/OS Capacity Provisioning homepage:
  http://www.ibm.com/systems/z/os/zos/features/cpm/

- z/OS WLM homepage:
  http://www.ibm.com/systems/z/os/zos/features/wlm/
  WLM Capping Technologies: https://ibm.biz/BdF4Lr

- z/OS MVS documentation
  – z/OS MVS Capacity Provisioning User's Guide
    http://publibz.boulder.ibm.com/epubs/pdf/iea3u110.pdf
  – z/OS MVS Planning: Workload Management:
    http://publibz.boulder.ibm.com/epubs/pdf/iea3w101.pdf
  – z/OS MVS Programming: Workload Management Services:
    http://publibz.boulder.ibm.com/epubs/pdf/iea3w201.pdf

- *IBM Redbooks publications*:
  – System Programmer's Guide to: Workload Manager:
    http://publib-b.boulder.ibm.com/abstracts/sg246472.html?Open
  – ABCs of z/OS System Programming Volume 12
    http://publib-b.boulder.ibm.com/abstracts/sg247621.html?Open