

# Introduction to Assembler Programming Sessions 16317, 16318

Richard Cebula ([riccebu@uk.ibm.com](mailto:riccebu@uk.ibm.com)) IBM HLASM



## Who am I?

- Richard Cebula – HLASM, IBM Hursley, UK
- [riccebu@uk.ibm.com](mailto:riccebu@uk.ibm.com)
- Develop and support the following products:
  - HLASM
  - SuperC
  - XREF
  - IDF
  - DISASM
  - Structured Programming Macros

# Audience

- This course aims to provide the grounding knowledge for programming in assembler
- The audience should have a basic understanding of computers
- The audience should be new to the world of z Systems Assembler programming
- At the end of this course the attendee should be able to:
  - Understand the basics of assembler programming on z Systems
  - Understand a variety of simple machine instructions
  - Understand how to Assemble, Bind and run simple assembler programs

# Introduction to Assembler Programming

- Why program in assembler?
- Computer Organisation – Overview of z/Architecture
- Assemblers, Compilers, Binders – Building programs on z Systems
- Working with the High Level Assembler (HLASM)
  - Using HLASM to assemble your program
  - Syntax
  - Machine vs Assembler instructions
- Programming in Assembler
  - Moving data around
    - Loading, storing and moving data
  - Manipulating Data
    - Logical operations
    - Arithmetic
  - Making Decisions
    - Comparing
    - Branching

# Introduction to Assembler Programming

- Programming in Assembler
  - Forming High Level Language constructs
    - If...then...else
    - Looping
  - Addressing Data and some of its subtleties
  - Basics of Calling Conventions
  - Reading Principles of Operation
  
- Reading the HLASM listing
  
- The PSW and an introduction to Debugging Assembler Programs

# Why program in assembler?

## Why program in assembler?

- Assembler programming has been around since the very start of computer languages as an easy way to understand and work directly with machine code
- Assembler programming can produce the most efficient code possible
  - Memory is cheap
  - Chips are fast
  - So what?
- Assembler programming TRUSTS the programmer
  - Humans are smart (?)
  - Compilers are dumb (?)
- Assembler programming requires some skill
  - No more than learning the complex syntax of any high-level language, APIs (that change every few years), latest programming trends and fashions
  - Your favorite language will too become old, bloated and obsolete!

## Why program in assembler?

- Misconceptions of assembler programming
  - I need a beard right?
  - It's too hard...
  - Any modern compiler can produce code that's just as efficient now days...
  - I can do that quicker using...
  - But assembler isn't portable...



## Why program in assembler?

- Misconceptions of assembler programming
  - I need a beard right?
    - Assembler programmers tend to be older and more experienced and typically wiser
    - Experienced programmers that have used assembler know that they can rely on it for the most complex of programming tasks
  - It's too hard...
    - Learning assembler is just like learning any other language
    - Each instruction to learn is as easy as the next
    - Syntax is consistent
    - No difficult APIs to get to grips with
  - Any modern compiler can produce code that's just as efficient now days...
    - Compilers CAN produce efficient code but that is not to say that they WILL
    - Optimization in compilers is a double-edged sword – compilers make mistakes
  - I can do that quicker using...
    - Good for you, so can I...
  - But assembler isn't portable...
    - Neither is Java, nor C, nor C++... portability depends on your definition of it

## Why program in assembler?

- The assembler mindset
  - You are not writing code – you are programming the machine
  - You must be precise
  - Your assembler program is no better than your programming
- Assembler programming provides the programmer with TOTAL freedom
  - What you choose to do with that freedom is your choice and your responsibility
- WYWIWYG – What you write is what you get – even if you *think* you wrote something else...

## Why program in assembler? - I thought this was the 21<sup>st</sup> century...

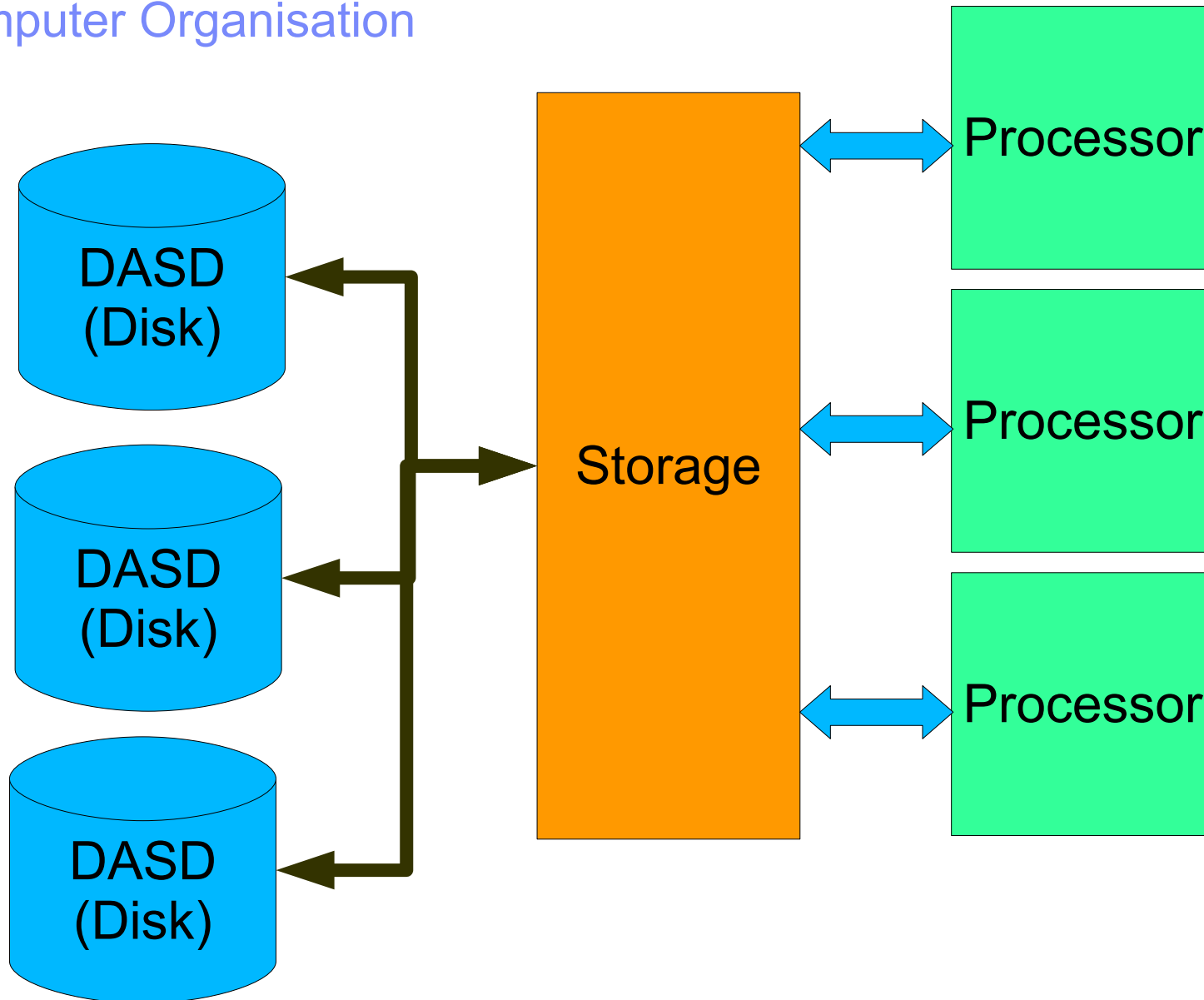
- Well there are a number of good reasons to program in assembler still:
  - Precise machine control – trust me...your compiler doesn't do what you think it does...
  - Precise definition of data – we have more data types than other languages
  - Self-writing code – the world's best macro facility
  - When someone's program goes wrong – finding the cause of the problem is much easier in assembler
  - Assembler is high-tech – no need to wait for compilers to catch up to what your chip can do

# Computer Organisation

# Computer Organisation

- The main parts of a computer when discussing programming are:
  - Processor – z Systems is a multiprocessor computer
  - Storage – (don't call it RAM on an EC12 and later it's RAIM!!)
  - Disks (more often referred to as DASD in z Systems)
- Programs are stored on disk since disks are non-volatile media, i.e. they do not lose their contents when the computer is not running
- A program is fetched from disk and placed into storage from where it can be executed by the processor
- The processor is the brain of the computer and is responsible for actually executing programs
- Operations inside a computer such as loading programs from disk are performed by a piece of software called an Operating System (OS)
- z Systems have 5 operating systems available z/OS, z/VM, z/VSE, z/TPF and z/Linux

# Computer Organisation



## Computer Organisation

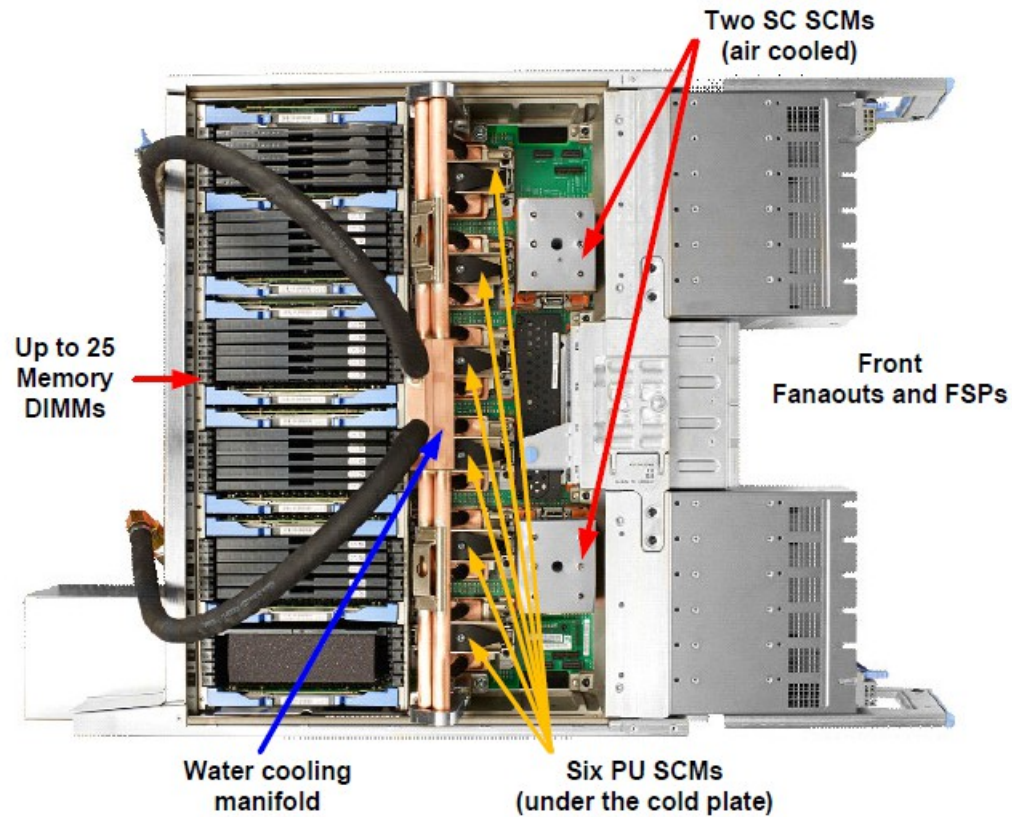
- The closer that data gets to the processor, the quicker it is to access the data both in terms of type of storage used and also due to the physical distance to the processor
- In order to improve the speed at which the program is accessed from storage, a special form of storage called *cache* is built into the processor
- Data and instructions are fetched from storage into the cache on the processor and then accessed from there
- Depending on how often the data requires to be accessed, how much data is to be accessed and whether the data needs to be shared between different processors or not, depends on which *cache level* the data is placed into on the processor with level 1 being the smallest and fastest and level 4 being the slowest but largest
- The implementation of cache is dependent on not only the processor architecture, e.g. Intel vs z/Architecture vs ARM etc. but also the model of processor itself, e.g. EC12 cache structure is different to z196

## Computer Organisation – z13

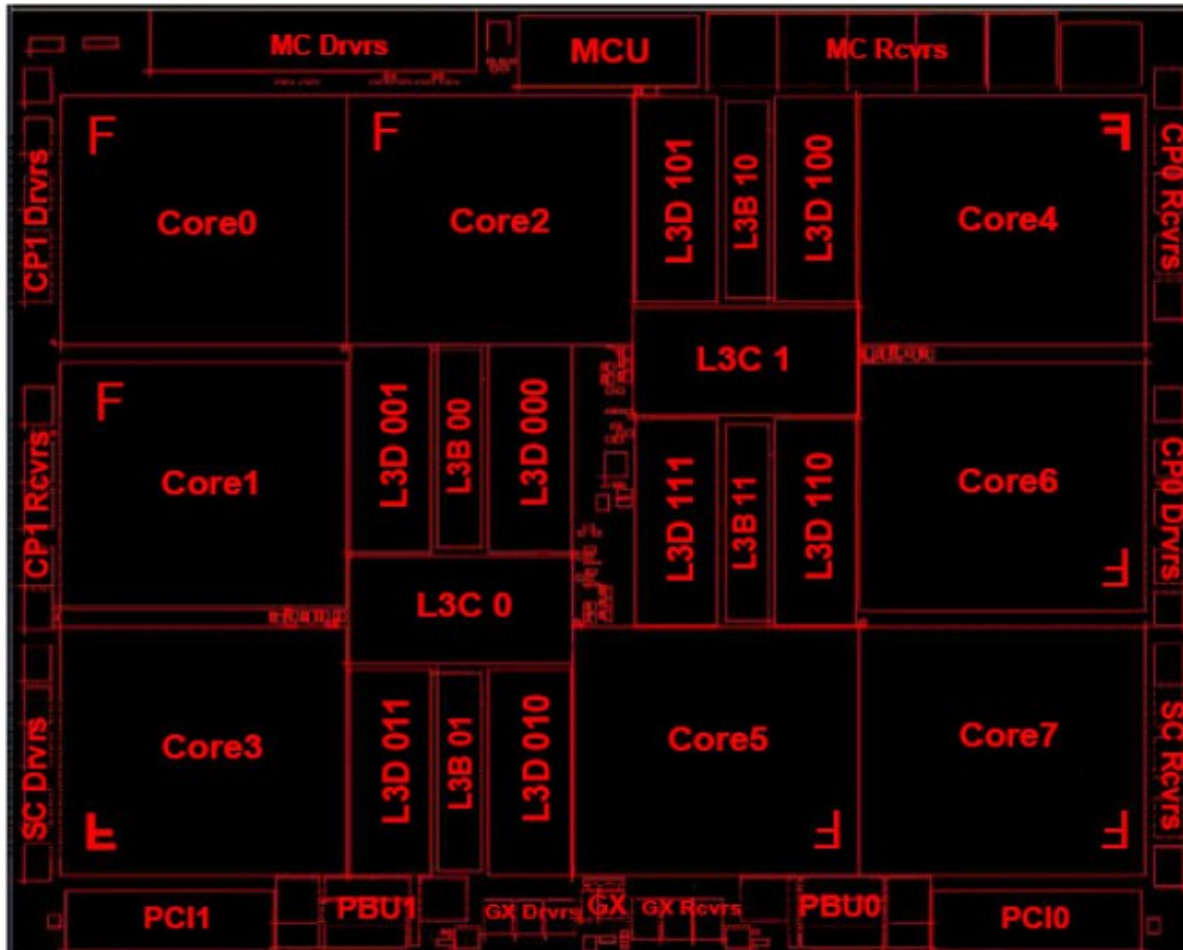
- Each processor in z Systems is split into a number of *processor cores* called PUs (processing units)
- Each PU (core) has its own L1 (96KB I, 128KB D) and L2 (2MB I, 2MB D) instruction and data caches
- Each processor chip contains up to 8 processing cores which share a 64MB L3 cache
- PU chips and storage control (SC) chips are packaged as single chip modules (SCMs). 3 PU chips and 1 SC chip form a CPC drawer node. Each SC has a 480MB L4 cache.
- 2 nodes are contained on a CPC drawer. Each system can have up to 4 drawers which contain all the storage for the system.
- All of the highly complex cache structure is transparent to the programmer
  - Experienced programmers will often change their code in order to improve cache performance
- Depending on how the machine is configured, each processor may run a different level of *microcode* which governs how it functions changing the PU into either a CP, zIIP or IFL
- Additionally, z Systems has specialist I/O processors called System Assist Processors (SAP) used to send signals to devices attached to the machine



# Computer Organisation – z/Architecture



# Computer Organisation – z/Architecture



## Computer Organisation – z/Architecture

- In order for data to be manipulated by the processor, it needs to be loaded into a *register*
- A register is the fastest storage available to the processor and is located inside each processor core
- In z/Architecture, each processor core has:
  - 16 64-bit General Purpose Registers (GPRs)
  - 16 32-bit Access Registers (ARs)
  - 16 64-bit Floating Point Registers (FPRs)\*
  - 16 64-bit Control Registers (CRs)
  - 32 128-bit Vector Registers (VRs)\*
  - 1 Program Status Word (PSW)
- Note that all registers are numbered 0-15 (or 0-31) in machine instructions – the instruction itself determines which type of register is being used
- z/Architecture – the processor architecture used for all z Systems Mainframes
- Processor specifications vary
  - Processor level – the physical (or virtual) chip used
  - Architecture level – the instruction specification of a chip
- \* Vector registers 0-15 bits 0-63 are mapped over the FPRs.

## Computer Organisation – z/Architecture

- z/Architecture is a big-endian, 64-bit, rich CISC processor architecture
  - Big-endian
    - Data is organised such that the most significant byte of a piece of data is stored in the lowest address of memory
  - 64-bit
    - The size of general purpose registers is 64-bits in length
  - CISC
    - Complex Instruction Set Computer
    - A single instruction comprises a number of micro-instructions which are executed by the processor
    - This scheme allows for a single machine instruction to perform a number of complex tasks
  
- For historical reasons, the size of data in z/Architecture is measured as:
  - 4 bits = 1 nibble
  - 8 bits = 2 nibbles = 1 byte
  - 16 bits = 2 bytes = a *halfword*
  - 32 bits = 4 bytes = 2 halfwords = a *word*
  - 64 bits = 8 bytes = 2 words = a *doubleword*
  - 128 bits = 16 bytes = 2 doublewords = a *quadword*

## Computer Organisation – Understanding Registers

- GPRs – used for arithmetic, logical operations, passing operands to instructions, calling subroutines etc
- ARs – used in “Access Register” mode – provides the ability to access another address space
- FPRs – used for floating point instructions, binary, decimal and hexadecimal floating-point arithmetic
  - Do not confuse decimal floating-point with packed decimal arithmetic – the latter is performed in storage not in registers
- VRs – used for SIMD (Single Instruction Multiple Data) operations including integer, string, floating-point and general operations
- CRs – used for controlling processor operations
- PSW – provides the status of the processor consisting of 2 parts:
  - PSW Flags – these show the state of the processor during instruction execution
  - Instruction address – this is the address of the next instruction to be executed
- GPRs and FPRs are sometimes operated on in pairs by certain instructions
  - GPRs form even-odd pairs, i.e. (0,1), (2,3),..., (14,15)
  - FPRs pair evenly / oddly, i.e. (0,2), (1,3),..., (13,15)

# Assemblers, Compilers and Binders

## Building programs on z Systems

## Assemblers, Compilers and Binders

- Typing in the code for a computer program will not mean that the computer can load and then execute the program
- The written code must be changed into binary form – this is done for the programmer by a compiler or assembler
  - Compilers are used for HLLs and often attempt to optimise the code written by the programmer
  - Assemblers are used for assembly language and do NOT optimise any code written by the programmer – assembler is a WYWIWYG language – What You Write Is What You Get
- Each language will require its own compiler / assembler to process it and change it into binary code
- *Cross compiling* is where a compiler runs on one machine architecture and produces machine code for another architecture

## Assemblers, Compilers and Binders

- The machine code produced by the assembler is called *object code*
- It is the job of a *binder* (sometimes called a *linker*) to create a complete program which can then be loaded by the operating system's *loader* into storage and then executed by the processor
- The binder works by ordering a set of objects and resolving any references between them
  - The bound object on z/OS is called a *program object* (for GOFF format) or a *load module* (for OBJ format)
  - Program objects and load modules vary in their capability – it is recommended that new programs use the GOFF format
- Some references are unable to be resolved by the binder at bind time and can only be resolved when the program is loaded into storage by the operating system's loader
  - In order to do this, the binder creates a list of these references and notes their location in the program object / load module
  - Examples of such unresolvable references are various operating system services and shared libraries
- The terms “load module” and “program object” are used interchangeably in this presentation



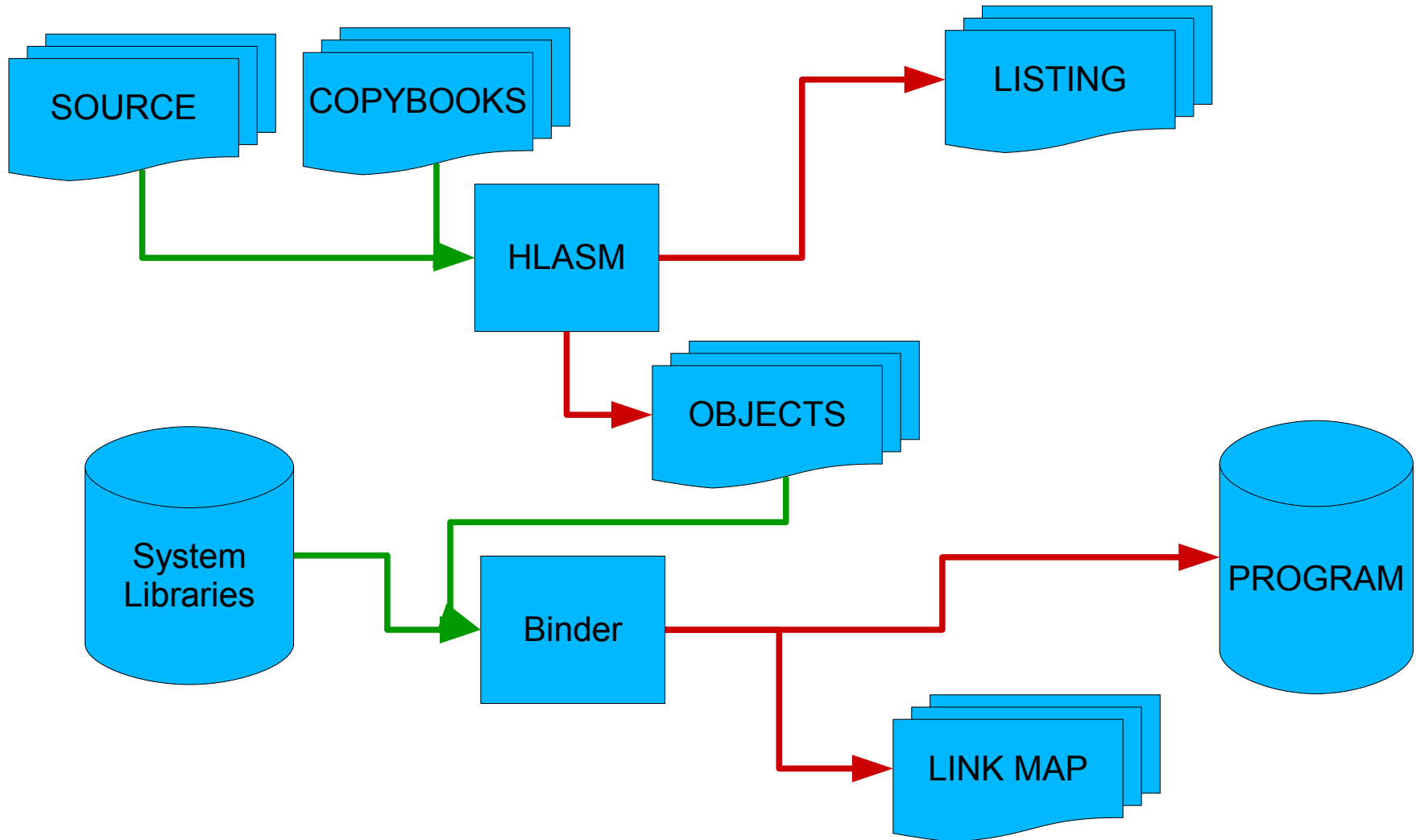
## Assemblers, Compilers and Binders – Working with HLASM

- HLASM – IBM's High Level Assembler
- Available on z/OS, z/VM, z/VSE, z/Linux and z/TPF
- *High Level Assembler*??? - YES!
  - Provides a wide range of assembler *directives*
    - An assembler *directive* is not a machine instruction
    - It is an instruction to the assembler during assembly of your program
  - A very powerful macro programming facility
  - Structured programming

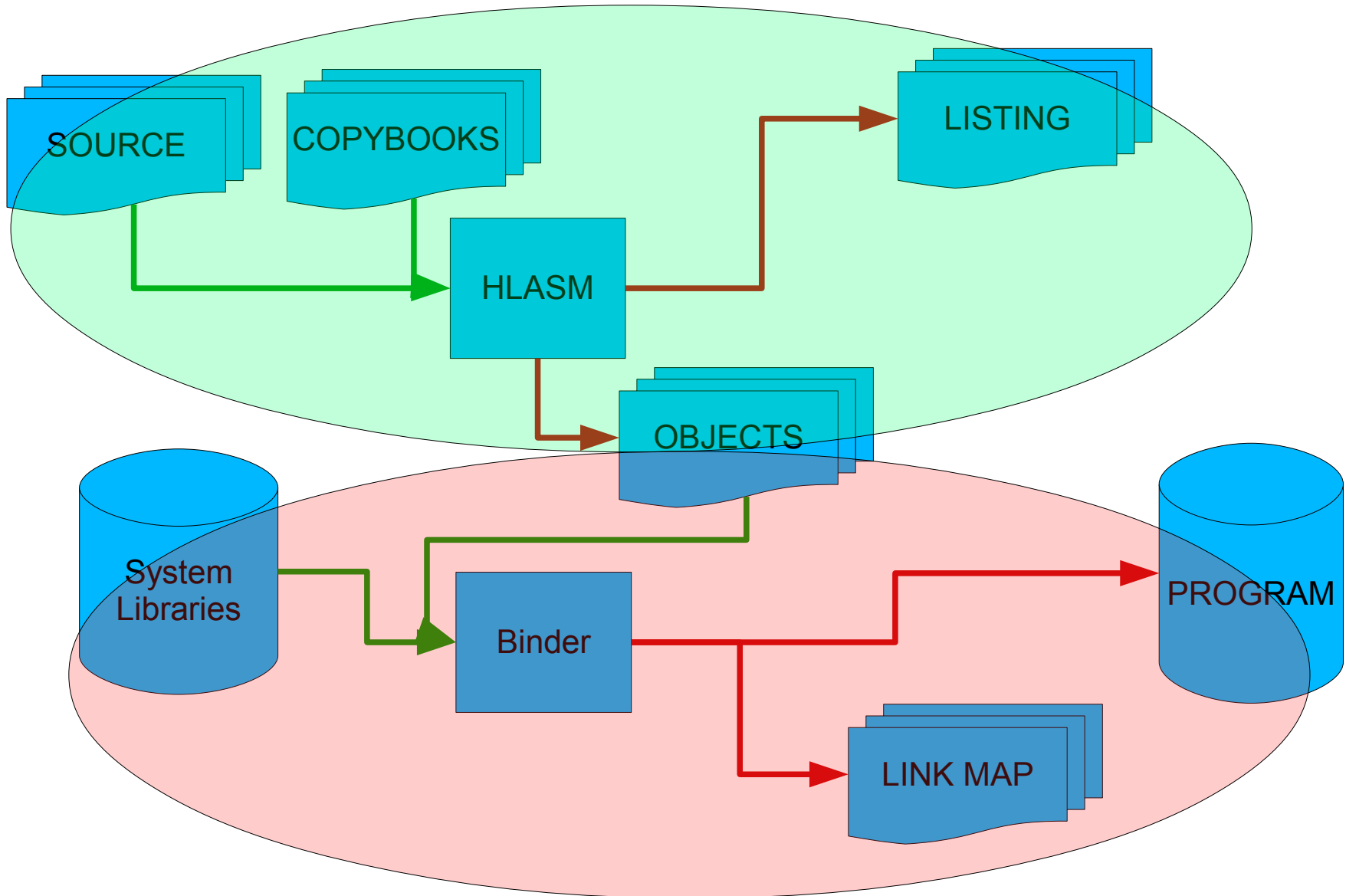
## Assemblers, Compilers and Binders – Working with HLASM

- HLASM produces 2 primary outputs
  - OBJECT DECKS – this is the object code that is used as input to binding
  - Listing – this shows any errors, all diagnostics and human readable output from the assemble phase
- The binder produces 2 primary outputs
  - LOAD MODULE – this is the bound object decks forming an executable program
  - A LOAD MAP – this is the Binder equivalent of an assembler listing
- A LOAD MODULE can be loaded into memory by the operating system and run

# Assemblers, Compilers and Binders – Working with HLASM



# Assemblers, Compilers and Binders – Working with HLASM



# Working with the High Level Assembler HLASM

## Working with the High Level Assembler HLASM

- These slides explain about using HLASM specifically for z/OS
  - The same principles apply on other platforms although HLASM is started differently, e.g. different command lines on z/VM and z/Linux, VSE JCL for z/VSE
  
- HLASM is started on z/OS via JCL and is shipped with some JCL PROCs to make using HLASM easier
  - ASMAC – Assembles a program
  - ASMACG – Assembles a program and invokes the loader to bind, load and execute the program (no load module is retained)
  - ASMACL – Assembles and invokes the binder to bind the program producing a load module
  - ASMACLG – Assembles and binds the program then runs the produced load module

## Working with the High Level Assembler HLASM

- The following JCL invokes ASMACL to assemble and bind the program into a load module:

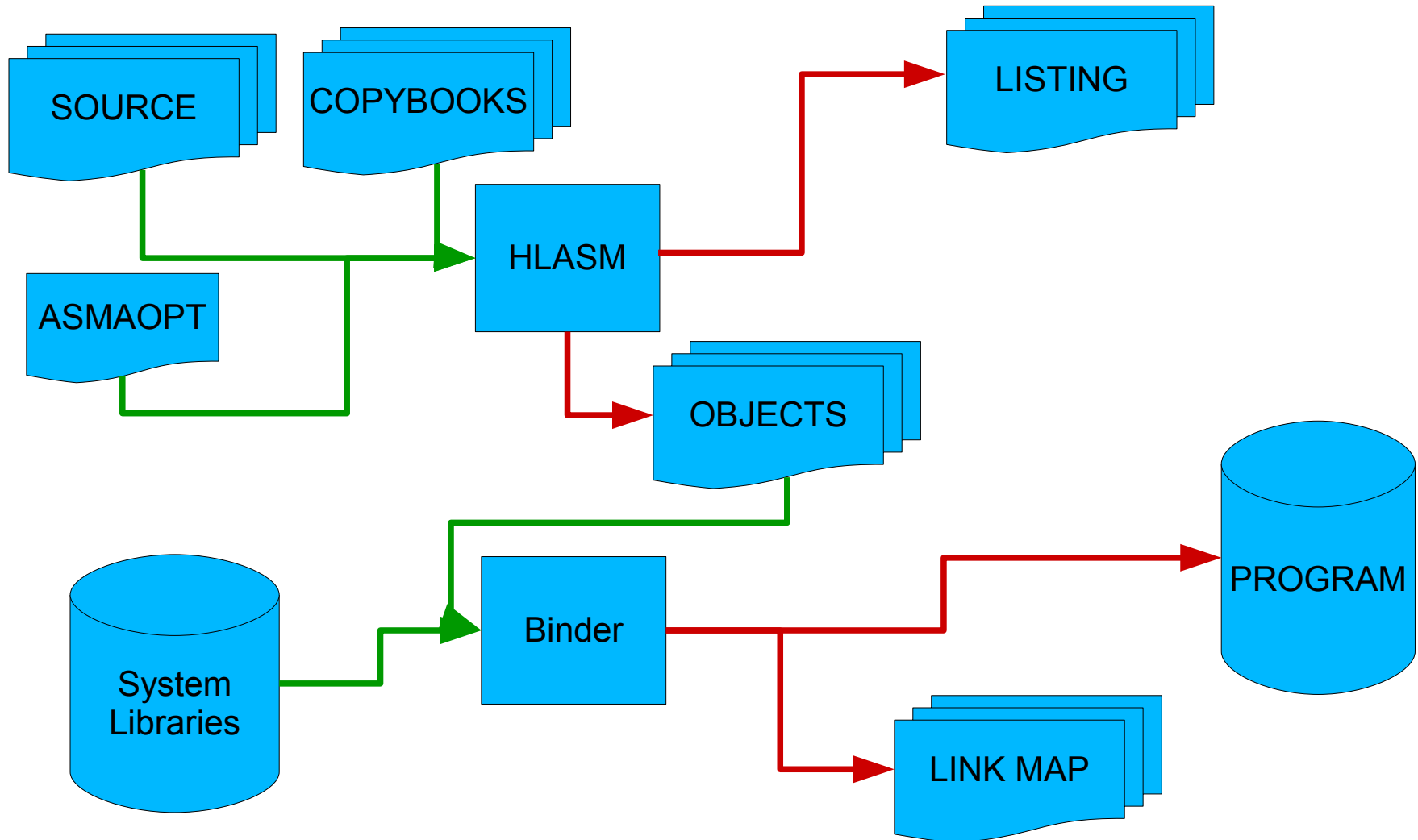
```
// SET OPTLIB=&SYSUID..AOPT
// SET SRCLIB=&SYSUID..SOURCE
// SET LSTING=&SYSUID..LISTINGS
// SET LODLIB=&SYSUID..LOAD
//*
// SET PRGNM=MYPROG
//*
//ASMMSAMP EXEC ASMACL, PARM.C=(OBJ, ADATA)
//C.ASMAOPT DD DSN=&OPTLIB, DISP=SHR
//C.SYSIN DD DSN=&SRCLIB (&PRGNM) , DISP=SHR
//C.SYSPRINT DD DSN=&LSTING (&PRGNM) , DISP=OLD
//L.SYSLMOD DD DSN=&LODLIB (&PRGNM) , DISP=SHR
```

## Working with the High Level Assembler HLASM

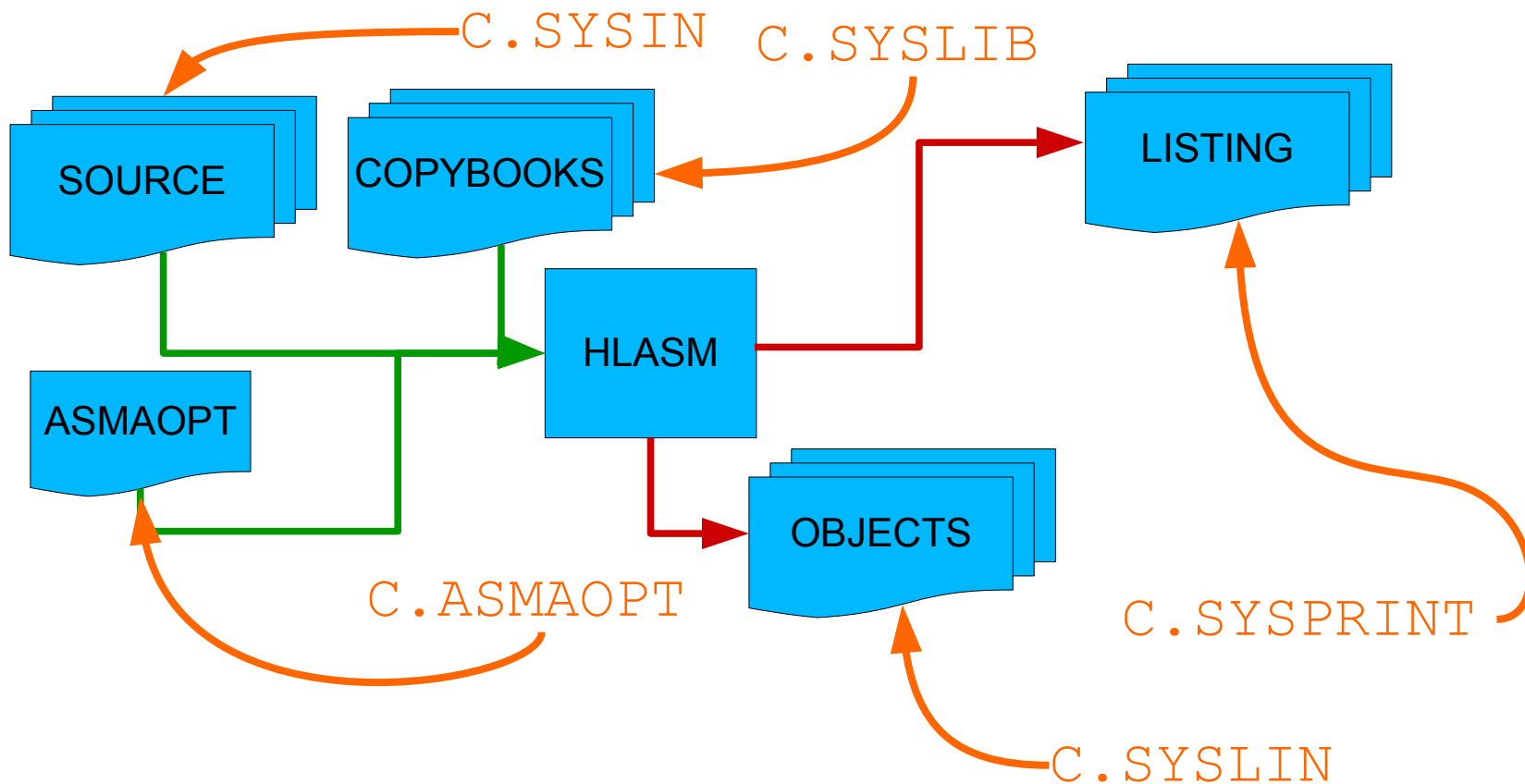
- The ASMACL JCL PROC has two steps:
  - C – Assemble the program
  - L – Bind the program
  - Using the notation *step.ddname* allows multiple DDNAMEs to be specified for the JCL steps
  
- DDNAMEs used by the Assemble step (C)
  - SYSIN → specifies the program's source
  - SYSLIN → specifies copybooks / macro libraries
  - SYSPRINT → specifies the destination of the program's *listing*
  - ASMAOPT → specifies a data set which contains assembly options\*
  - PARM → specifies some options for HLASM\*
  
- DDNAMEs used by the Bind step (L)
  - SYSLMOD → specifies the destination of the bound load module
  
- \* Options may be specified in up to 6 different locations when assembling a program – allowing for a hierarchy of options to be specified.



# Working with HLASM – Assembling and Binding a program



## Working with HLASM – Assembling and Binding a program



## Working with HLASM – Examining the listing

- The program listing produced by HLASM is an invaluable source of information for the programmer
- It lists in detail everything that HLASM has done when assembling a program including:
  - All the options that were used when the program was assembled
  - Any external references that the program produces or relies upon in the External Symbol Dictionary (ESD)
  - Each line of source code and the machine code that was produced
  - Any symbols in the code that require relocation in the Relocation Dictionary (RLD)
  - A summary of all lines in error in the program
  - A list of all data sets used
- The assembler uses the following return codes whenever it issues a message:
  - 0 – Success / Information
  - 2 – Notice – A condition which may need correcting – program appears to be correct
  - 4 – Warning – Program may not work as expected
  - 8 – Error – Error in program
  - 12 – Severe error – Unlikely that the program assembles as expected or runs
  - 16 – Critical – Unlikely that the program runs
  - 20 – Unrecoverable – Unable to continue

## Working with HLASM – A look at syntax

```

File Edit Edit_Settings Menu Utilities Compilers Test Help
VIEW
Command ==> Columns 00001 00080
Scroll ==> CSR
***** ***** Top of Data *****
000001 *****
000002 * SIMPLE DUMMY EXIT FOR HLASM
000003 *****
000004 *
000005 * MAIN PROGRAM STARTS HERE
000006 *
000007 LARLLOAD CSECT
000008 LARLLOAD AMODE 31
000009 LARLLOAD RMODE 24
000010 * USUAL PROGRAM SETUP
000011     STM    14,12,12(13)
000012     BALR  12,0      GET THE CURRENT ADDRESS
000013     USING *,12     USE 12 AS THE BASE REGISTER
000014     L     1,=F'12'
000015 LMRET   LM    14,12,12(13)
000016     XR    15,15
000017     BR    14
000018 * *****
000019 * END OF PROGRAM
000020 * *****
000021     END
***** ***** Bottom of Data *****

```

## Working with HLASM – A look at syntax

```

File Edit Edit_Settings Menu Utilities Compilers Test Help
VIEW Columns 00001 00080
Command ==> Scroll ==> CSR
***** Top of Data *****
000001 *****
000002 * SIMPLE DUMMY EXIT FOR HLASM
000003 *****
000004 *
000005 * MAIN PROGRAM STARTS HERE
000006 *
000007 LARLLOAD CSECT
000008 LARLLOAD AMODE 31
000009 LARLLOAD RMODE 24
000010 * USUAL PROGRAM SETUP
000011     STM    14,12,12(13)
000012     BALR  12,0      GET THE CURRENT ADDRESS
000013     USING *,12     USE 12 AS THE BASE REGISTER
000014     L     1,=F'12'
000015 LMRET   LM    14,12,12(13)
000016     XR    15,15
000017     BR    14
000018 * *****
000019 * END OF PROGRAM
000020 * *****
000021     END
***** Bottom of Data *****

```

Comments start with a \* in column 1 or appear after free-form instruction operands until column 72

## Working with HLASM – A look at syntax

```

File Edit Edit_Settings Menu Utilities Compilers Test Help
VIEW Columns 00001 00080
Command ==> Scroll ==> CSR
***** Top of Data *****
000001 *****
000002 * SIMPLE DUMMY EXIT FOR HLASM
000003 *****
000004 *
000005 * MAIN PROGRAM STARTS HERE
000006 *
000007 LARLLOAD CSECT
000008 LARLLOAD AMODE 31
000009 LARLLOAD RMODE 24
000010 * USUAL PROGRAM SETUP
000011     STM    14,12,12(13)
000012     BALR  12,0      GET THE CURRENT ADDRESS
000013     USING *,12     USE 12 AS THE BASE REGISTER
000014     L     1,=F'12'
000015 LMRET   LM    14,12,12(13)
000016     XR    15,15
000017     BR    14
000018 * *****
000019 * END OF PROGRAM
000020 * *****
000021     END
***** Bottom of Data *****

```

Labels start in column 1

## Working with HLASM – A look at syntax

```

File Edit Edit_Settings Menu Utilities Compilers Test Help
VIEW Columns 00001 00080
Command ==> Scroll ==> CSR
***** Top of Data *****
000001 *****
000002 * SIMPLE DUMMY EXIT FOR HLASM
000003 *****
000004 *
000005 * MAIN PROGRAM STARTS HERE
000006 *
000007 LARLLOAD CSECT
000008 LARLLOAD AMODE 31
000009 LARLLOAD RMODE 24
000010 * USUAL PROGRAM SETUP
000011 STM 14,12,12(13)
000012 BALR 12,0 GET THE CURRENT ADDRESS
000013 USING *,12 USE 12 AS THE BASE REGISTER
000014 L 1,=F'12'
000015 LMRET LM 14,12,12(13)
000016 XR 15,15
000017 BR 14
000018 * *****
000019 * END OF PROGRAM
000020 * *****
000021 END
***** Bottom of Data *****

```

Instructions start after column 1 or a label

# Working with HLASM – A look at syntax

```

File Edit Edit_Settings Menu Utilities Compilers Test Help
VIEW Columns 00001 00080
Command ==> Scroll ==> CSR
***** Top of Data *****
000001 *****
000002 * SIMPLE DUMMY EXIT FOR HLASM
000003 *****
000004 *
000005 * MAIN PROGRAM STARTS HERE
000006 *
000007 LARLLOAD CSECT
000008 LARLLOAD AMODE 31
000009 LARLLOAD RMODE 24
000010 * USUAL PROGRAM SETUP
000011 STM 14,12,12(13)
000012 BALR 12,0 GET THE CURRENT ADDRESS
000013 USING *,12 USE 12 AS THE BASE REGISTER
000014 L 1,=F'12'
000015 LMRET LM 14,12,12(13)
000016 XR 15,15
000017 BR 14
000018 * *****
000019 * END OF PROGRAM
000020 * *****
000021 END
***** Bottom of Data *****

```

Operands start after a space after instructions and are delimited by commas and brackets



## Working with HLASM – CSECTs and DSECTs

- CSECT → CONTROL SECTION (HLASM directive)
  - A CSECT contains machine instructions to be run on the machine
- DSECT → DUMMY SECTION (HLASM directive)
  - Used to define the structure of data

- Both CSECT and DSECT are terminated with the end statement

```
MYPROG    CSECT                START OF CODE
          ...awesome assembler program goes here...
MYSTRUCT  DSECT                START OF DATA STRUCTURE
          ...awesome data structure goes here...
          END                   END OF PROGRAM
```

## Working with HLASM – Defining Data

- Data is defined via the DC and DS HLASM directives
- DC – Define Constant
  - Defines data and initialises it to a given value
- DS – Define Storage
  - Defines storage for data but does not give it a value
- e.g.

NUMBER1	DC	F'12'	DEFINE A FULLWORD WITH VALUE 12
NUMBER2	DC	H'3'	DEFINE A HALFWORD WITH VALUE 3
TOTAL	DS	H	DEFINE A HALFWORD
MYSTR	DC	C'HELLO WORLD'	DEFINE A SERIES OF CHARACTERS
MYHEX	DC	X'FFFF'	DEFINE A SERIES OF HEX CHARACTERS

## Working with HLASM – Literals

- A literal is an inline definition of data used in an instruction but the space taken for the literal is in the nearest literal pool
- A literal pool collects all previous literals and reserves the space for them
- By default, HLASM produces an implicitly declared literal pool at the end of your CSECT
- To cause HLASM to produce a literal pool, use the LTORG directive

```
      L      1,=F'1'          LOAD REGISTER 1 WITH FULLWORD OF 1
      X      1,=H'2'          XOR REGISTER 1 WITH HALFWORD OF 2
...more awesome assembler code here...
      LTORG ,                THE LITERAL POOL IS CREATED
```

## Working with HLASM – Machine VS Assembler Instructions

- There are two distinct types of instruction accepted by HLASM – Machine instructions and Assembler instructions
- Machine instructions are changed into machine code by the assembler – these are the instructions that will be used when your program is executed
- Assembler instructions are instructions which affect the behavior of the assembler itself as it is assembling a program
- Some assembler instructions such as DC will cause the resultant program to contain certain pieces of data although the majority of assembler instructions do not
- The assembler also includes a macro language – macros are code which are used to generate other code programmatically and therefore can shorten development time

# Introduction to Assembler Programming

## Moving Data

## Moving Data

- Moving data is the most common operation that ANY computer program performs
- Before any computations can be performed, data must be moved to the correct places
  - Data is moved to the processor (from disk, memory, networks, devices etc)
  - Data is manipulated by the processor
  - The result is stored somewhere (back to disk, memory, networks, devices etc)
- Data is LOADED into the processor's registers via LOAD instructions
- Data is STORED to memory via STORE instructions

## Moving Data – Loading from Register to Register

- The LOAD REGISTER (LR) instruction is used to load the value stored in one register to another

```
LR 1,2      LOAD REGISTER 2 INTO REGISTER 1 (32-BITS)
```

- The instruction copies 32-bits from a register to another
- The instruction has a 64-bit variant LOAD GRANDE REGISTER (LGR)

```
LGR 1,2     LOAD REGISTER 2 INTO REGISTER 1 (64-BITS)
```

- The instruction has a 16-bit variant LOAD HALFWORD REGISTER

```
LHR 1,2     LOAD REGISTER 2 INTO REGISTER 1 (16-BITS)
```

## Moving Data – Loading from Memory to Register

- The LOAD (L) instruction is used to load the value stored in memory to a register

L 1,NUMBER      LOAD REGISTER 1 WITH THE VALUE NUMBER (32-BITS)

- The instruction copies 32-bits from memory to a register

- The instruction has a 64-bit variant LOAD GRANDE (LG)

LG 1,NUMBER      LOAD REGISTER 1 WITH THE VALUE NUMBER (64-BITS)

- The instruction has a 16-bit variant LOAD HALFWORD

LH 1,NUMBER      LOAD NUMBER INTO REGISTER 1 (16-BITS)



## Moving Data – Storing from a Register to Memory

- The STORE (ST) instruction is used to store the value in a register to memory

ST 1,NUMBER      STORE REGISTER 1 TO NUMBER (32-BITS)

- The instruction copies 32-bits from a register to memory

- The instruction has a 64-bit variant STORE GRANDE (STG)

STG 1,NUMBER      STORE REGISTER 1 TO NUMBER (64-BITS)

- The instruction has a 16-bit variant STORE HALFWORD

STH 1,NUMBER      STORE REGISTER 1 TO NUMBER (16-BITS)

## Moving Data – Moving data without registers

- The MOVE (MVC) instruction can be used to move data in memory without the need for a register

```
MVC OUTPUT, INPUT
```

```
MOVE INPUT TO OUTPUT
```

- The MVC instruction can move up to 256B from one area of memory to another
- The MVCL instruction can move up to 16M (but uses different parameters)
- The MVCLE instruction can move up to 2G (or up to 16EB in 64-bit addressing)
- In all cases, the move instruction moves 1 byte at a time (left to right in memory)

## Why no one writes assembler like this...

- Many instructions may appear very similar but may have very different (and unintended) consequences, e.g.:  

```
LR 1,2      Load register 1 with the value of register 2
```
- Unless you know what you're doing, don't do this:  

```
L  1,2      Load register 1 with the value at memory offset 2
```
- The very simple example above shows how confusing assembler programming can be as there is no distinction between writing the value of a number and the value of a register in an instruction and instead the distinction is made by which instruction was written by the programmer.
- Many programmers often name their registers r0-r15 to make it clearer to understand which values are registers and which values are numbers in an instruction. To do this, the best solution is to start a program with the statement `ASMDREG` which will include all the register names for your program to use.
  - From now on, this material will refer to GPRs registers as r0-r15
- Using r0-r15 also means that type checking can be performed by HLASM and they will also appear in the cross-reference section of the listing making it easier to find that rogue register...

# Introduction to Assembler Programming

## Logical Operations

## Logical Instructions – EXCLUSIVE OR (X, XG, XR, XGR, XC)

- The EXCLUSIVE OR instructions perform the EXCLUSIVE OR *bit-wise* operation

X	r1, NUMBER	XOR REGISTER 1 WITH NUMBER (32-BITS)
XG	r1, NUMBER	XOR REGISTER 1 WITH NUMBER (64-BITS)
XR	r1, r2	XOR REGISTER 1 WITH REGISTER 2 (32-BITS)
XGR	r1, r2	XOR REGISTER 1 WITH REGISTER 2 (64-BITS)
XC	NUM1, NUM2	XOR NUM1 WITH NUM2 (UP TO 256-BYTES)

- Notice a pattern with the instruction mnemonics?
  - Rules of thumb:
    - G → 64bits (DOUBLEWORD)
    - H → 16bits (HALFWORD)
    - R → register
    - C → character (byte / memory)
    - L → logical (i.e. unsigned)

## Logical Instructions – AND (Nx), OR (Ox)

- The AND instructions perform the AND *bit-wise* operation

N	r1, NUMBER	AND REGISTER 1 WITH NUMBER (32-BITS)
NG	r1, NUMBER	AND REGISTER 1 WITH NUMBER (64-BITS)
NR	r1, r2	AND REGISTER 1 WITH REGISTER 2 (32-BITS)
NGR	r1, r2	AND REGISTER 1 WITH REGISTER 2 (64-BITS)
NC	NUM1, NUM2	AND NUM1 WITH NUM2 (UP TO 256-BYTES)

- The OR instructions perform the OR *bit-wise* operation

O	r1, NUMBER	OR REGISTER 1 WITH NUMBER (32-BITS)
OG	r1, NUMBER	OR REGISTER 1 WITH NUMBER (64-BITS)
OR	r1, r2	OR REGISTER 1 WITH REGISTER 2 (32-BITS)
OGR	r1, r2	OR REGISTER 1 WITH REGISTER 2 (64-BITS)
OC	NUM1, NUM2	OR NUM1 WITH NUM2 (UP TO 256-BYTES)

## A word on instruction choice

- In 5 basic operations (loading, storing, AND, OR, XOR) we have already seen over 25 instructions!
- How do I decide which instruction to use?
  - The instruction should be chosen for:
    - Its purpose, e.g. don't use a STORE instruction to LOAD a register – it won't work!
    - Its data, e.g. 32-bits, 16-bits, 64-bits, bytes?

- Many instructions can perform *similar* operations, e.g.

XR	r15, r15	XOR REGISTER 15 WITH REGISTER 15
L	r15, =F'0'	LOAD REGISTER 15 WITH 0
LA	r15, 0	LOAD REGISTER 15 WITH ADDRESS 0

- Different instructions NEVER do the same thing even if you think they do
  - The result does not justify the means

# Introduction to Assembler Programming

## Arithmetic



# Arithmetic

- Arithmetic is performed in a wide variety ways on z/Architecture
  - Fixed point arithmetic (including logical) ← performed in GPRs
  - Packed Decimal arithmetic ← performed in memory
  - Binary and Hexadecimal Floating point arithmetic ← performed in FPRs
- Fixed point arithmetic
  - Normal arithmetic, e.g. adding the contents of 2 numbers together
  - Fixed point arithmetic is signed with numbers being stored in 2's complement form
  - Logical fixed point arithmetic is unsigned, i.e. both numbers are positive
- Pack Decimal arithmetic
  - Performed in memory
  - Numbers are in packed decimal format

## Arithmetic – Fixed point arithmetic operations

- **ADD instructions**

AR	r1, r2	ADD REGISTER 2 TO REGISTER 1 (32-BIT SIGNED)
ALR	r1, r2	ADD REGISTER 2 TO REGISTER 1 (32-BIT LOGICAL)
A	r1, NUMBER	ADD NUMBER TO REGISTER 1 (32-BIT SIGNED)
AL	r1, NUMBER	ADD NUMBER TO REGISTER 1 (32-BIT LOGICAL)
AFI	r1, 37	ADD 37 TO REGISTER 1 (IMMEDIATE 32-BIT SIGNED)

- Note that for immediate instructions, the operand is included in the instruction rather than needing to be obtained from memory
- At the end of the addition, the CC is updated (as specified in POPs)
  - CC → 0 → Result is 0; no overflow
  - CC → 1 → Result less than 0; no overflow
  - CC → 2 → Result greater than 0; no overflow
  - CC → 3 → Overflow occurred

## Arithmetic – Fixed point arithmetic operations

- **SUBTRACT instructions**

```
SR    r1, r2        SUBTRACT REGISTER 2 FROM REGISTER 1 (SIGNED)
SLR   r1, r2        SUBTRACT REGISTER 2 FROM REGISTER 1 (LOGICAL)
S     r1, NUMBER    SUBTRACT NUMBER FROM REGISTER 1 (SIGNED)
SL    r1, NUMBER    SUBTRACT NUMBER FROM REGISTER 1 (LOGICAL)
AFI   r1, -37       ADD -37 TO REGISTER 1 (IMMEDIATE 32-BIT SIGNED)
```

- **At the end of the subtraction, the CC is updated (as specified in POPs)**

- CC → 0 → Result is 0; no overflow
- CC → 1 → Result less than 0; no overflow
- CC → 2 → Result greater than 0; no overflow
- CC → 3 → Overflow occurred

## Arithmetic – Fixed point arithmetic operations

- **MULTIPLY instructions**

```
MR    r2, r7          MULTIPLY REGISTER 2 BY REGISTER 7
```

```
M     r2, NUMBER      MULTIPLY REGISTER 2 BY NUMBER
```

- The first operand is an even-odd pair – the result of the MULTIPLY is stored in:
  - The even register (of the pair) – top 32-bits of result
  - The odd register (of the pair) – bottom 32-bits of the result
- At the end of the multiplication, the CC is UNCHANGED

## Arithmetic – Fixed point arithmetic operations

- **DIVIDE instructions**

```
DR    r2, r7          DIVIDE REGISTER 2 BY REGISTER 7
D     r2, NUMBER      DIVIDE REGISTER 2 BY NUMBER
```

- The first operand is an even-odd pair
  - The even register (of the pair) – top 32-bits of dividend
  - The odd register (of the pair) – bottom 32-bits of the dividend
- The result is stored in the first operand:
  - The quotient is stored in the odd register of the pair
  - The remainder in the even register of the pair
- At the end of the division, the CC is UNCHANGED

# Introduction to Assembler Programming

## Branching

## Branching

- Branching allows control flow in the program to move nonsequentially
- Branches are performed via the BRANCH instructions
- Most branch instructions are *conditional* – i.e. they will pass control to the *branch target* if a condition is met otherwise control will continue sequentially
- The condition on which the branch will take place is called the CONDITION CODE (CC)
  - The CC is 2-bits stored in the PSW; thus the value is 0-3
  - Each instruction may (or may not) set the CC
- A branch instruction provides a *branch mask*
  - The *branch mask* instructs the processor that the branch will be taken if any of the bits in the CC match those in the branch mask
- Fortunately, HLASM provides extended-mnemonics which provide branch masks for most branch instructions

## Branching – Using HLASM's extended-mnemonics

- B – Branch (unconditionally)
- BE – Branch on condition Equal
- BL – Branch on condition Lower than
- BH – Branch on condition Higher than
- BNL – Branch Not Low
- BNH – Branch Not High
- BZ – Branch on Zero
- BNZ – Branch Not Zero
- There are also other extended-mnemonics which HLASM provides



## Branching – How does a branch mask work

- B – Branch (unconditionally)
  - This is translated to the BRANCH ON CONDITION (BC) instruction with a mask of 15

Condition Code	0	1	2	3
Mask value	8	4	2	1

- So, 15  $\rightarrow$  b'11111'  $\rightarrow$  8+4+2+1
- Thus the branch is taken if CC 0, 1, 2 or 3 is met, i.e. ALWAYS

## Branching – How does a branch mask work

- BE – Branch on Equal
  - This is translated to the BRANCH ON CONDITION (BC) instruction with a mask of 8

Condition Code	0	1	2	3
Mask value	8	4	2	1

- So, 8 → b'1000' → 8
- Thus the branch is taken if CC 0 is met

## Branching – Using a branch to form an *if* statement

```
LT    r1,NUMBER    LOAD NUMBER INTO REGISTER 1 AND SET CC
BNZ   NONZERO      BRANCH TO 'NONZERO' IF REGISTER 1 IS NOT ZERO
      ...code where register 1 is zero goes here...
B     COMMONCODE    REJOIN COMMON CODE
```

```
NONZERO DS 0H
```

```
      ...code where register 1 is non-zero goes here...
```

```
COMMONCODE DS 0H
```

## Branching – Using a branch to form an *if* statement

```
//Example C-like equivalent
if(register_1==0){
    //Code for register_1 being 0 goes here
}
else{
    //Code for register_1 being non-zero goes here
}

//Common code goes here
```

# Introduction to Assembler Programming

## Looping

## Looping

- A simple loop is formed by using a counter, a comparison and a branch, e.g.

```

        LA      r2,0           INITIALISE COUNTER REGISTER TO 0
MYLOOP  AHI     r2,1           INCREMENT COUNTER
        WTO    'HELLO'        SAY HELLO
        CHI    r2,10          IS THE COUNTER 10?
        BL     MYLOOP         IF IT'S LESS THAN 10, GO TO MYLOOP

```

- That's simple – but there's a better way – use **BRANCH ON COUNT (BCT)**

```

        LA      r2,10          INITIALISE COUNTER REGISTER TO 10
MYLOOP  WTO    'HELLO'
        BCT    r2,MYLOOP      SUBTRACTS, COMPARES & BRANCHES

```

- There are other instructions similar to BCT that subtract/add values and then branch depending on the result, e.g. BCTR, BXH etc...

# Introduction to Assembler Programming

## Addressing Data

## Addressing Data

- There are 2 ways to access data for manipulation
  - Base-Displacement (and index) addressing
  - Relative addressing
- Relative addressing is a new form of addressing which calculates the data's relative position from the current PSW (in half-word increments)

```
LRL    r1,NUMBER    LOAD RELATIVE REGISTER 1 WITH NUMBER
...more awesome assembler code here...
NUMBER DC    F'23'
```

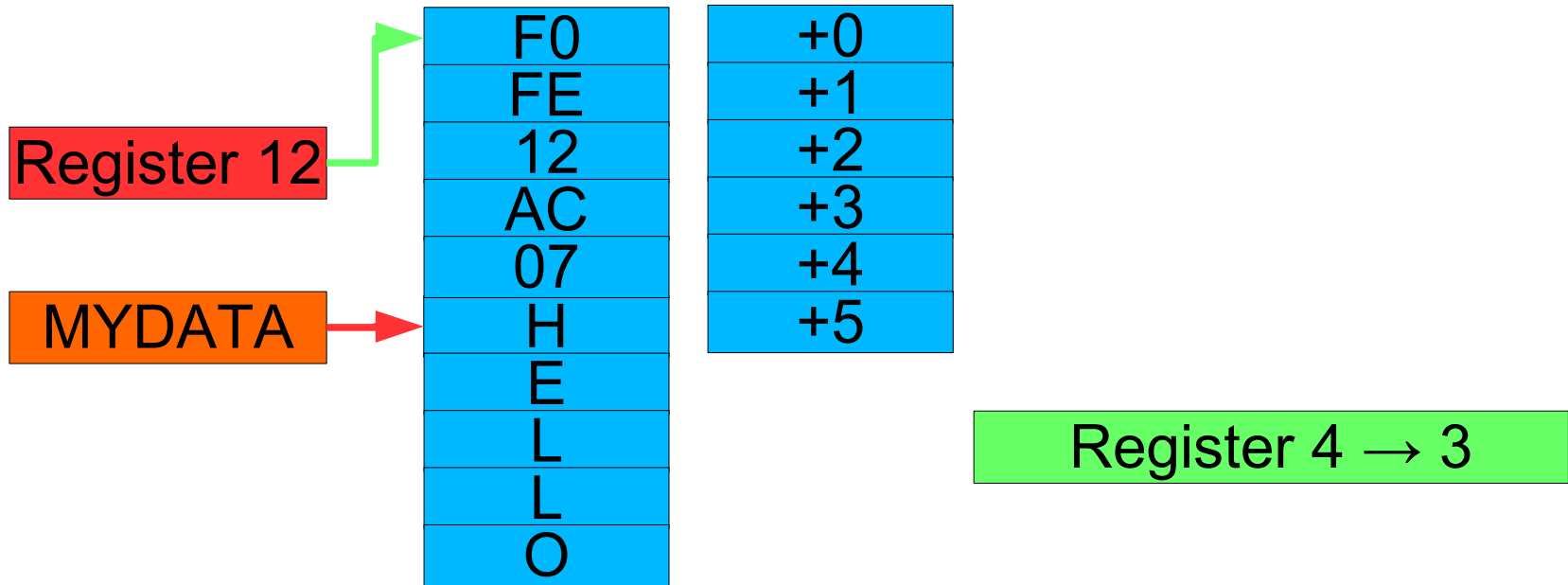


## Addressing Data - Base-Displacement-Index

- Base-Displacement(-index) addressing involves using a register as a pointer to memory – this is called the BASE register
  - Base (and index) registers, are GPRs and the term *base* and *index* are used purely for referring to the different usage in a particular set of instructions
- A displacement is usually between 0 and 4095 bytes allowing a single base register to address 4K of memory
- An index register is an additional register whose value is added to the base and displacement to address more memory
- Incrementing an index register allows the assembler programmer to cycle through an array whilst maintaining the same base-displacement
- Note that register 0 cannot be used as a base or index register
  - Register 0 used in this way means that the *value* 0 will be used as a base / index and NOT the contents of register 0
- Base, displacement and indexes are optionally specified on an instruction
  - Implicit default value for each is 0

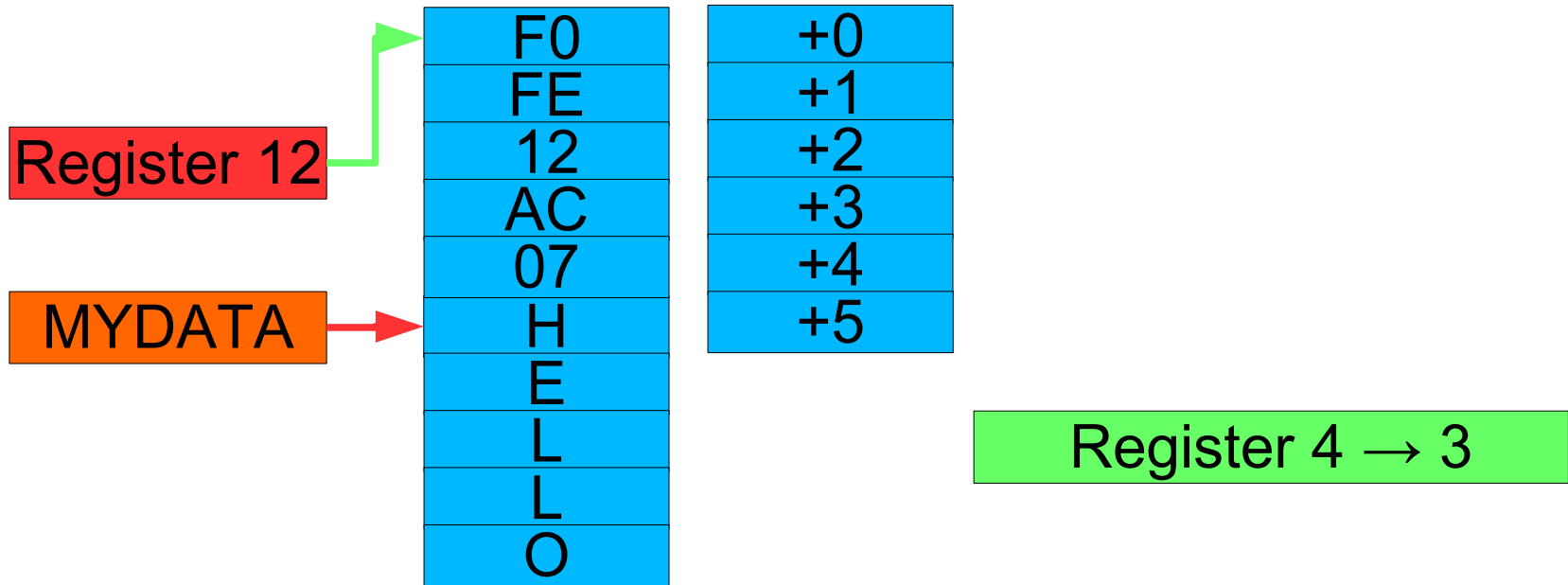
## Addressing Data - Base-Displacement-Index

- $\text{Address} = \text{BASE}(\text{register}) + \text{INDEX}(\text{register}) + \text{DISPLACEMENT}$



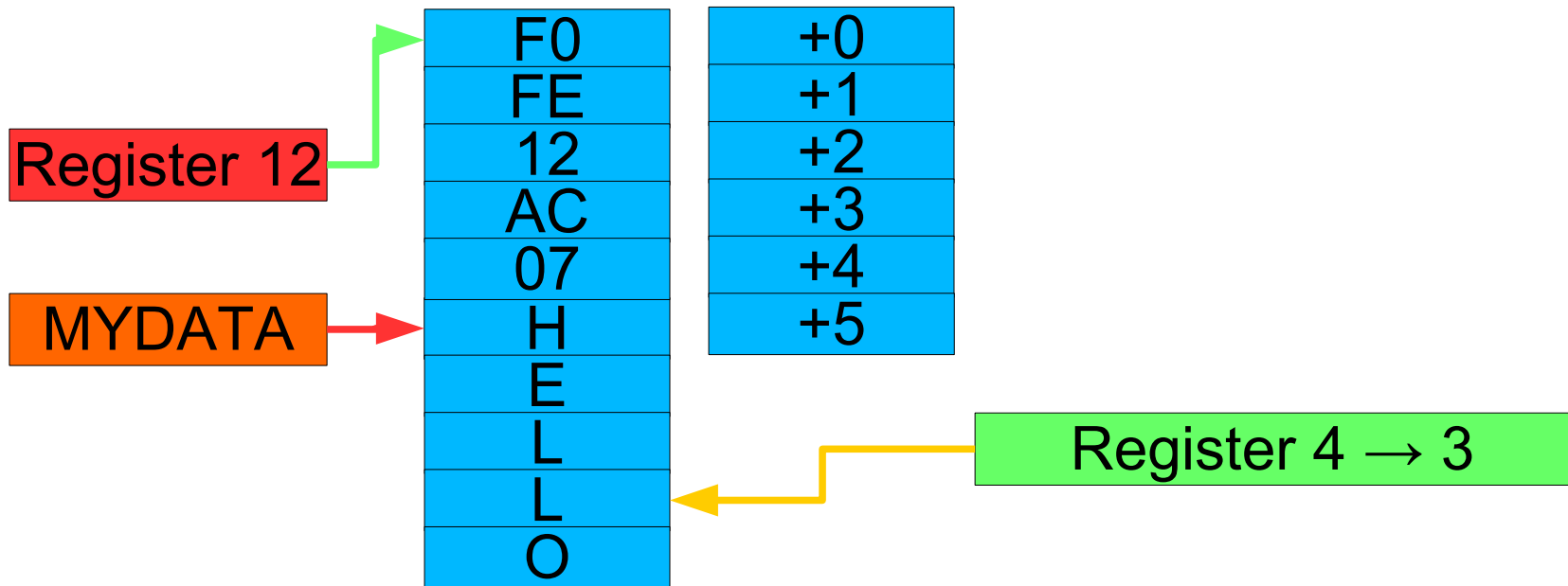
## Addressing Data - Base-Displacement-Index

- Address of MYDATA =  $5(r0,r12) \rightarrow$  displacement 5 + index (register) 0 + base (register) 12



## Addressing Data - Base-Displacement-Index

- Address of 'L' in 'HELLO' =  $5(r4,r12)$  → *displacement 5 + index (register) 4 + base (register) 12*



## Addressing Data – Loading addresses

- To load an address into a register, use the LOAD ADDRESS (LA) instruction

```
LA    r1,DATA          LOAD ADDRESS OF DATA INTO REGISTER 1
```

- The LA instruction can be used to set a register to between 0 and 4095 by specifying a base and index register of 0 – these are automatically implicitly specified, e.g.

```
LA    r1,12           base=0, index=0, displacement=12
```

- To store a 'L' in 'HELLO' in the previous example:

...some setup for REGISTER 12...

```
LA    r4,3           LOAD ADDRESS 3* INTO REGISTER 4
IC    r3,=C'L'       LOAD CHARACTER 'L' INTO REGISTER 3
STC   r3,MYDATA(r4) base=12, index=4, displacement=5
```

## Addressing Data – Base VS Index registers

- Both base and index registers can be dropped from an instruction and can *almost* be used interchangeably. Consider the following:

LA	r4,MYDATA	LOAD ADDRESS OF MYDATA INTO REGISTER 4
IC	r3,=C'H'	LOAD CHARACTER 'H' INTO REGISTER 3
STC	r3,0(,r4)	STORE H AT 0 DISP + (0 INDEX) + BASE 4

- Contrast with:

LA	r4,MYDATA	LOAD ADDRESS OF MYDATA INTO REGISTER 4
IC	r3,=C'H'	LOAD CHARACTER 'H' INTO REGISTER 3
STC	r3,0(r4)	STORE H AT 0 DISP + 4 INDEX + (BASE 0)

- Both store the character H at the same address, i.e. at a displacement 0 from the start of MYDATA.
- However, if the program is running in Access Register Mode, then using a *base* register will cause its corresponding access register to be involved in the address calculation too!

## Addressing Data – When a base register is not a base register...

- Certain instructions, e.g. various shift instructions, specify a base-displacement operand but do not address storage
- Instead, the calculated address forms a value for the operand which is used by the instruction
- The advantage of instructions such as this is that by varying the contents of a register, the same instruction in a program can be used repeatedly but it will modify varying amounts of data

## Addressing Data – the HLASM USING instruction

- HLASM can be used to generate base-displacement values for your program automatically.
- In order to do this, the USING instruction is used to specify a base register for program / data addressability.
- HLASM will then calculate any displacements necessary in order to address parts of the program or data, e.g.:

```
BALR    r12,0          LOAD ADDRESS OF NEXT INSTRUCTION INTO R12
USING   *,r12         USE R12 AS A BASE REGISTER FROM HERE
```

- If a register that is being used as a base register by HLASM needs to be used for another purpose, it is good practice to use the DROP instruction to cancel the previous USING instruction for that register, e.g.:

```
DROP    r12           DROP THE USE OF R12 AS A BASE
```

- Note that HLASM has different forms for the USING instruction – the HLASM Language Reference should be consulted for more information.



## Addressing Data – using multiple base registers

- For most instructions, the displacement field is only 12-bits in length and therefore a maximum displacement value of 4095 can be achieved.
- If a section of code or data is bigger than 4096-bytes in size, then the programmer must use another base register to address the data further into the code. This can be done by issuing another USING instruction and assigning another register as a base register.
  - It is important that the other base register's contents point at the correct place in the code which it will be addressing
- A common use of multiple base registers is to assign one set of base registers for a piece of code and another for the data in the program.
- The HLASM PUSH instruction can be used to save the current state of the USINGs for a program and the POP instruction can be used to restore the previous USING state.

## Addressing Data – using multiple base registers

```
* Call the subroutine MYSUB
      LA      r15,MYSUB      Prepare to call subroutine
      BALR   r14,r15        Call it
      LTR    r15,r15        Check return code
```

```
* *****
* START OF MY SUBROUTINE MYSUB
* This subroutine starts with the contents of register 15 pointing to the
* start of the code. The mainline code of the program uses register 12 as
* its base register.
* *****
```

```
MYSUB      DC      0H
* Subroutine code goes here...
      L       r15,RETURN_CODE    Set return code
      BR      r14                Branch back to caller
RETURN_CODE DC      F'1'
```

## Addressing Data – using multiple base registers

```

000000 05C0                                11          BALR   r12,0
                                R:C 00002 12          USING *,r12
000002 41F0 C008                          0000A      13          la     r15,mysub
000006 05EF                                14          balr   r14,r15
000008 12FF                                15          ltr    r15,r15
                                16 * *****
                                17 * MYSUB
                                18 * *****
00000A                                19 mysub    dc     0h
00000A 58F0 C00E                          00010      20          l      r15,return_code
00000E 07FE                                21          br     r14
000010 00000001                          22 return_code dc f'1'

```

- Register 12 is being used as a base register for both the main program and the subroutine.

## Addressing Data – using multiple base registers

```

* Call the subroutine MYSUB
      LA    r15,MYSUB           Prepare to call subroutine
      BALR r14,r15             Call it
      LTR   r15,r15            Check return code

* *****
* START OF MY SUBROUTINE MYSUB
* This subroutine starts with the contents of register 15 pointing to the
* start of the code.  The mainline code of the program uses register 12 as
* its base register.
* *****
MYSUB      DC    0H
* Subroutine code goes here...
      PUSH USING
      USING *,r15              Set new base register*
      L     r15,RETURN_CODE    Set return code
      BR   r14                 Branch back to caller
      POP  USING
RETURN_CODE DC    F'1'

```

## Addressing Data – using multiple base registers

```

000000 05C0                                11          BALR   r12,0
                                12          USING *,r12
000002 41F0 C008 0000A                        13          la     r15,mysub
000006 05EF                                14          balr  r14,r15
000008 12FF                                15          ltr   r15,r15
                                16          DROP  r12
                                17 * *****
                                18 * MYSUB
                                19 * *****
00000A                                20 mysub   dc   0h
                                21          push  using
                                22          using *,r15
00000A 58F0 F006 00010                        23          l     r15,return_code
00000E 07FE                                24          br    r14
                                25          pop   using
000010 00000001                            26 return_code dc f'1'

```

- Register 12 is being used as a base register for the main program
- The current state of the USINGs is saved via the PUSH USING statement
- Register 15 is established as the base register for the subroutine

## Addressing Data – Addressing beyond 4096 bytes

- There are a number of methods for addressing code and data beyond 4096 bytes – we have already seen how we could use more than one base register to do this.
- The LONG DISPLACEMENT family of instructions can address up to 512KB (both positive and negative from the base register) of data using 20-bits. This is only available if the long-displacement facility is installed.
- Using relative instructions, an immediate field is encoded into the instruction which specifies the number of half-words that the target of the instruction is from the instruction itself.
  - Relative instructions can be used to eliminate the need for base registers altogether – so long as your code has been designed for it

# Introduction to Assembler Programming

## Calling conventions

## Calling Conventions

- A calling convention is a convention used between programs and subroutines to call each other
- The calling convention is not enforced, but if it is disregarded undesirable and unpredictable results may occur
- In general, when programming in assembler, the *caller* will provide a *save area* and the *called* program or routine will save all GPRs into that save area.
- The subroutine will then execute its code
- To return control to the caller, the subroutine will typically:
  - Set a return code in a register
  - Prepare the register on which it should branch back on
  - Restore all other registers
  - Branch back



## Calling Conventions – Typical register usage on z/OS

- Although free to do as they please, most assembler programs on z/OS use the following register convention during initialisation
  - Register 1 → parameter list pointer
  - Register 13 → pointer to register save area provided by caller
  - Register 14 → return address
  - Register 15 → address of subroutine
- Once the registers are saved, the called subroutine will:
  - Update register 13 to point to a new save area (so that it can call other programs / routines)
  - Establish register 12 as a base register for the program
- Upon termination, the called subroutine will:
  - Set a return code in register 15
  - Restore register 13 to the value it was previously
  - Restore registers 14,0,1,...,12 from the save area pointed to by register 13
  - Branch back on register 14

## Calling a subroutine in code – Going in...

- The caller calls the subroutine

```
LA      r1,PARAMS          POINT TO PARAMETERS
LA      r15,SUB1           LOAD ADDRESS OF SUBROUTINE
BALR    r14,r15            BRANCH AND LINK
LTR     r15,r15            CHECKS RETURN CODE 0?
...caller code continues here...
```

- The subroutine saves the caller's registers and establishes a base register

```
STM     r14,r12,12(r13)    STORE REGISTERS
LR      r12,r15            GET ENTRY ADDRESS
...subroutine code continues here...
```

## Calling a subroutine in code – Getting out...

- The subroutine restores the caller's registers, sets the return code and branches back

```
LM      r14, r12, 12(r13)    RESTORE REGISTERS
XR      r15, r15             SET RETURN CODE 0
BR      r14                 BRANCH BACK TO CALLER
```

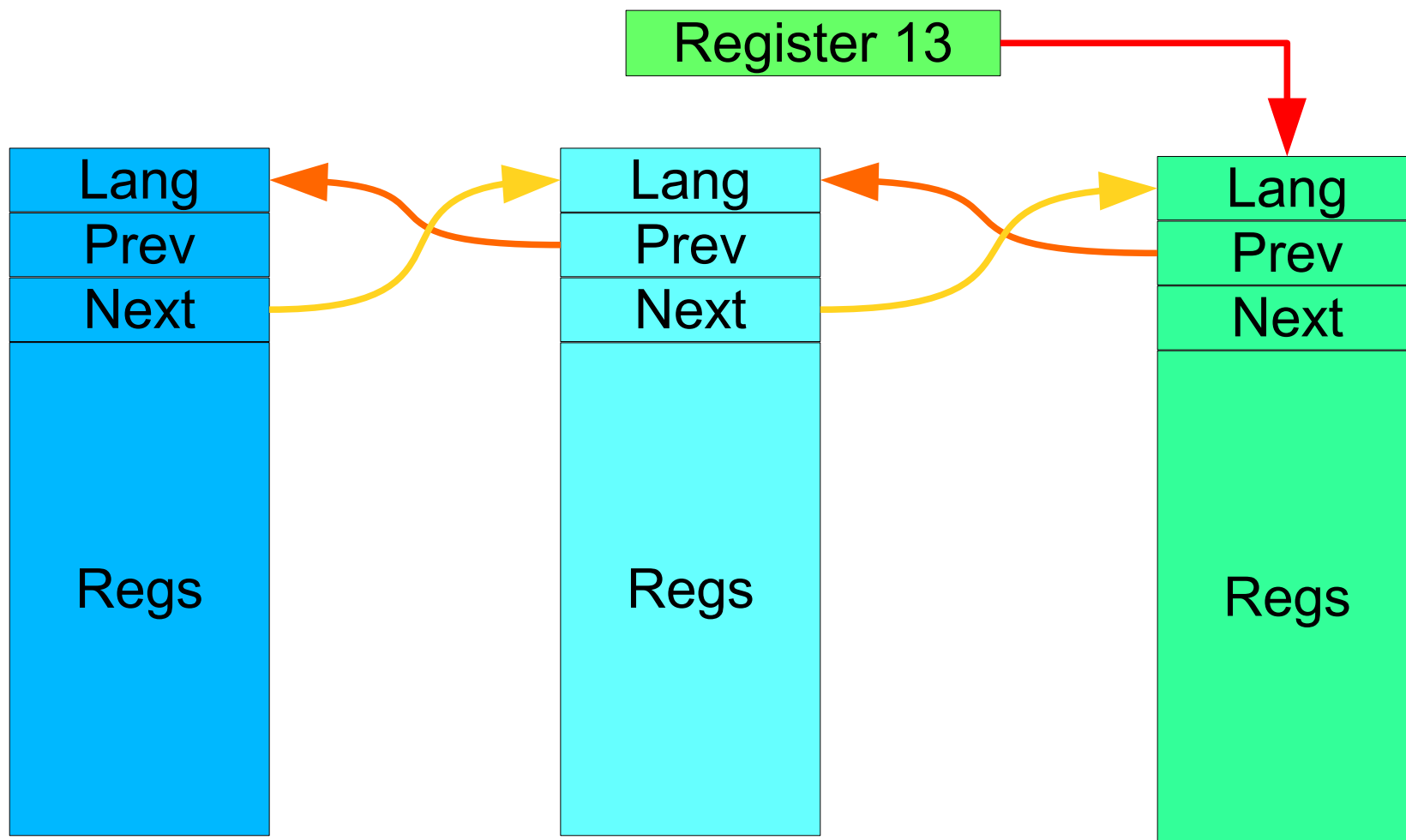
- Due to this calling convention, during epilog and prologue of a program or subroutine or when calling or having control returned from a program or subroutine, avoid using registers 0, 1, 12, 13, 14, 15
- z/OS services, typically will use registers 0, 1, 14, 15
- Not sure which registers are used by a service?
  - The manuals explain in detail

## What exactly is a “save area”?

- As with most things in software, the answer is, “it depends”
  - On z Systems, there are different types of save area which are used differently depending on the calling convention in use.
  - Your program must use the correct type of save area
- The standard z/OS linkage save area has the following format:

Offset in save area	Purpose
0	Used by language products
4	Address of previous save area
8	Address of next save area
12	Register 14
16	Register 15
20	Registers 0-12

## Chaining Save areas



## Chaining save areas – Going in...

- The caller calls the subroutine

```

LA      r1,PARAMS          POINT TO PARAMETERS
LA      r15,SUB1           LOAD ADDRESS OF SUBROUTINE
BALR    r14,r15            BRANCH AND LINK
LTR     r15,r15            CHECKS RETURN CODE 0?
...caller code continues here...

```

- The subroutine saves the caller's registers and establishes a base register

```

STM     r14,r12,12(r13)    STORE REGISTERS
GETMAIN RU,LV=72          Get storage for save area
ST      r13,4(,r1)         Chain previous save area to new
ST      r1,8(,r13)         Chain new to previous
LR      r13,r1             Set r13 to new save area
LR      r12,r15            GET ENTRY ADDRESS
...subroutine code continues here...

```

## Chaining save areas – Getting out...

- The caller calls the subroutine

```

LA      r1,PARAMS           POINT TO PARAMETERS
LA      r15,SUB1            LOAD ADDRESS OF SUBROUTINE
BALR    r14,r15             BRANCH AND LINK
LTR     r15,r15             CHECKS RETURN CODE 0?
...caller code continues here...

```

- The subroutine restores frees its save area, restores the caller's registers, sets a return code and branches back

```

LR      r1,r13              Address of save area to free
LA      r0,72               Length of save area
L       r13,4(,r13)         Point at previous save area
FREEMAIN R,LV=(0),A=(1)    Free save area
LM      r14,r12,12(13)     Restore registers
XR      r15,r15             Set return code
BR      r14                 Branch back to caller

```

# Introduction to Assembler Programming

## How to read Principles of Operation



## Reading POPs

- Principles of Operation (better known as POPs) is the z/Architecture manual
- It explains everything from system organisation and memory, to instructions and number formats
- It provides a useful set of appendices some of which provide good detailed examples of instruction use, including programming techniques
- The vast majority of POPs is instruction descriptions
  - Hint – Appendix A contains examples of instructions

## Reading POPs – Understanding Instruction Descriptions

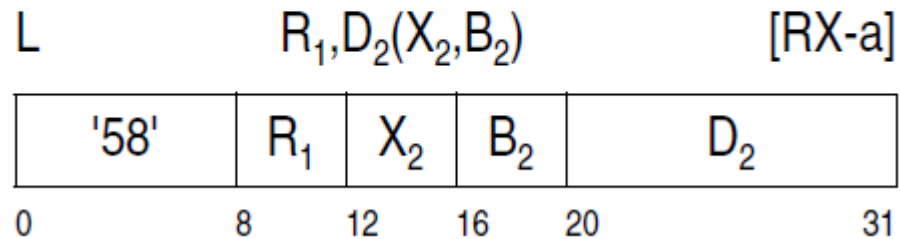
- Each instruction is described in exact detail including:
  - The instruction's syntax
  - Machine code
  - Operation
  - Condition code settings
  - Programming Exceptions
  
- There are 2 forms of syntax provided for each instruction
  - The syntax for the assembler, i.e. what is written in your assembler program
  - The machine code for the instruction, i.e. the binary code run on the processor
  
- The instruction's machine code is grouped together with other instructions which share a similar machine code layout called an *instruction format*

## Reading POPs – Instruction Formats

- The instruction format used, is generally related to
  - The assembler syntax used to code the instruction
  - The operation that the instruction performs
  
- Instructions that we've used have had the following formats:
  - RR – Register-Register – this form usually manipulates registers, e.g. LR, MR, DR
  - RX – Register, Index, base displacement – usually moving data between memory and registers, e.g. L, LA, ST, A, X, S, D, M
  - SS – Storage-Storage – acts on data in memory, e.g. MVC



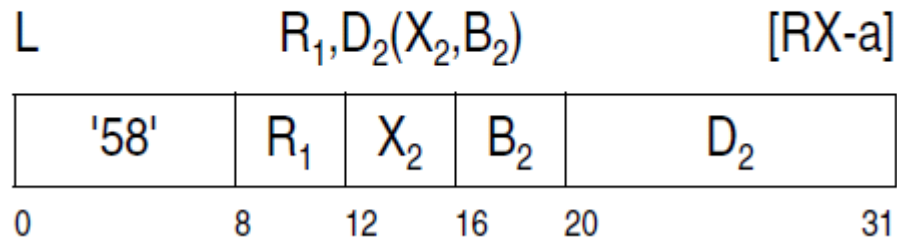
## Reading POPs – Instruction Formats – RX – L instruction

***Register-and-storage formats:***

## Looking at Instruction Formats

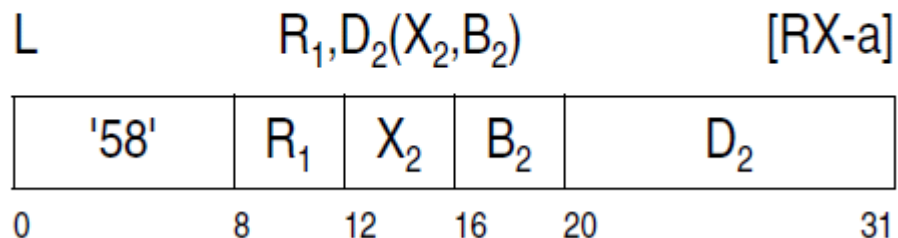
- Why do I need to know an instruction's format?
  - You don't...but it might come in useful when debugging...
- Consider having a dump and the failing instruction was 5810 0004
- Examining the “Principles of Operation” z/Architecture manual tells me:
  - 58 = LOAD Instruction
  - Format = RX-a

### *Register-and-storage formats:*

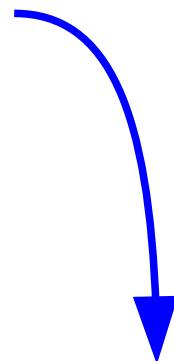


# Looking at Instruction Formats

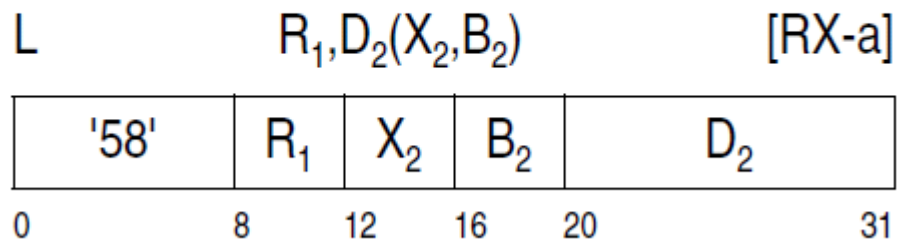
## *Register-and-storage formats:*



5810 0004



## *Register-and-storage formats:*

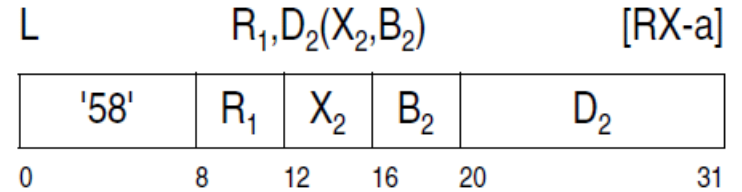




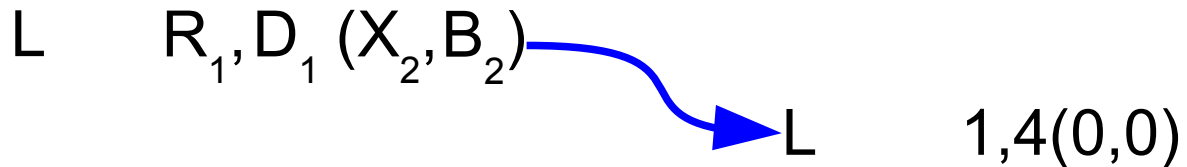
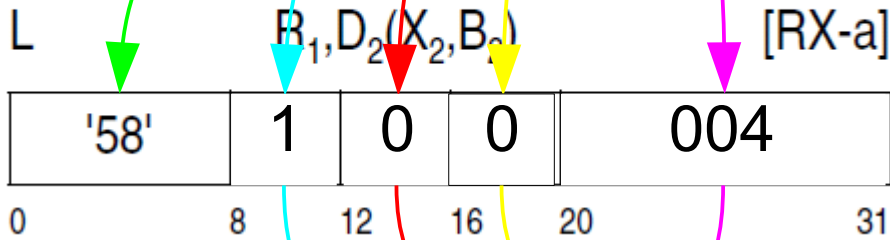


# Looking at Instruction Formats

58 100004



*Register and storage formats:*



# Reading HLASM Listings

## Reading HLASM Listings

- One of the outputs produced by HLASM is the “listing” - it explains in detail all that transformations that have happened to change your source code into the produced object code.
- In certain programs, it is not possible to use a debugger and it is in these circumstances where relying on a dump and the assembler listing proves invaluable.
  - Examining a dump and a listing can also be much quicker to solve the problem than trying to work through the program with a debugger
- Unlike a lot of compilers, the HLASM listing contains much more information than just error messages that were produced when attempting to assemble your source code:

## Reading HLASM Listings

- The HLASM listing is divided into:
  - High Level Assembler Option Summary
  - External Symbol Dictionary
  - Source Statements
  - Relocation Dictionary
  - Ordinary Symbol and Literal Cross Reference
  - Unreferenced Symbols Defined in CSECTs
  - Macro and Copy Code Source Summary
  - Macro and Copy Code Cross Reference
  - DSECT Cross Reference
  - Using Map
  - General Purpose Register Cross Reference
  - Diagnostic Cross Reference and Assembler Summary

## Specifying options

- HLASM options can be specified in:
  - \*PROCESS OVERRIDE
  - ASMAOPT
  - Invocation Params
  - \*PROCESS
  - Installation Defaults
- The location in which an option is specified since the location of each option specifies its precedence over other options.
  - Also some options cannot be specified in certain places, e.g. \*PROCESS OVERRIDE VS \*PROCESS



## Specifying options

Our example specifies:

In source:

```
*process override (adata, mxref (full))
*process align
*process nodbcs
*process mxref (full), nolibmac
*process flag(0)
*process nofold, language (ue)
*process nora2
*process nodbcs
*process xref (full)
```

In JCL procedure - parms:

```
OPTS1='NOOBJECT, language (en), size (4meg) ',
OPTS2='xref (short, unrefs) ',
OPTS3='nomxref, norxref, adata, noadata '
//C EXEC PGM=ASMA90,
// PARM='&OPTS1, &OPTS2, &OPTS3'
```

In ASAMOPT DD:

```
//ASMAOPT DD *
* My ASMAOPTS Overrides
sysparm(thisisatestsysparm)      pass this to assembler
goff                               create GOFF object code
/*
```

# HLASM Option Summary

## High Level Assembler Option Summary

(PTF R160 ) Page 1  
HLASM R6.0 2013/11/01 13.45

### Overriding ASMAOPT Parameters -

>\* My ASMAOPTS Overrides

>sysparm(thisisatestsysparm)

>goff

pass this to assembler  
create GOFF object code

### Overriding Parameters-

### Process Statements-

NOBJECT,language(en),size(4meg),xref(short,unrefs),nomxref,norxref,adata,noadata  
override(adata,mxref(full))  
align  
nodbcs  
mxref(full),nolibmac  
flag(0)  
nofold,language(ue)  
nora2  
nodbcs  
xref(full)

\*\* ASMA434N GOFF/XOBJECT option specified, option LIST(133) will be used

\*\* ASMA400W Error in invocation parameter - size(4meg)

\*\* ASMA423N Option adata, in a \*PROCESS OVERRIDE statement conflicts with invocation or default option.  
Option is not permitted in a \*PROCESS statement and has been ignored.

\*\* ASMA422N Option language(ue) is not valid in a \*PROCESS statement.

\*\* ASMA437N Attempt to override invocation parameter in a \*PROCESS statement. Suboption full of xref  
option ignored.

## HLASM Option Summary

```

Options for this Assembly
3 Overriding Parns NOADATA
  5 *Process ALIGN
    NOASA
    NOBATCH
    CODEPAGE(047C)
  5 *Process NOCOMPAT
    NODBCS
    NODECK
    DXREF
    ESD
    NOEXIT
  5 *Process FLAG(0,ALIGN,NOCONT,EXLITW,NOIMPLEN,NOPAGE0,PUSH,RECORD,NOSUBSTR,USING0)
  5 *Process NOFOLD
  2 ASMAOPT GOFF(NOADATA)
    NOINFO
3 Overriding Parns LANGUAGE(EN)
  5 *Process NOLIBMAC
    LINECOUNT(60)
    LIST(133)
    MACHINE(,NOLIST)
1 *Process Override MXREF(FULL)
3 Overriding Parns NOOBJECT
    OPTABLE(UNI,NOLIST)
    NOPCONTROL
    NOPESTOP
    NOPROFILE
  5 *Process NORA2
    NORENT
    RLD
3 Overriding Parns NORXREF
    SECTALGN(8)
    SIZE(MAX)
    TYPECHECK(MAGNITUDE,REGISTER)
    USING(NOLIMIT,MAP,NOWARN)
    NOWORKFILE
3 Overriding Parns XREF(SHORT,UNREFS)

```



## External Symbol Dictionary (ESD)

### External Symbol Dictionary

Symbol	Type	Id	Address	Length	Owner Id	Flags	Alias-of
LISTINGB	SD	00000001					
B_IDRL	ED	00000002			00000001		
B_PRV	ED	00000003			00000001		
B_TEXT	ED	00000004	00000000	00000084	00000001	08	
LISTINGB	LD	00000005	00000000		00000004	08	
EXTERNAL_FUNCTION							
	ER	00000006			00000001		
FUNCY	ER	00000007			00000001		
listme	ER	00000008			00000001		LISTINGZ
COMMON_DATA							
	SD	00000009					
B_IDRL	ED	0000000A			00000009		
B_PRV	ED	0000000B			00000009		
B_TEXT	ED	0000000C	00000000	00000018	00000009	00	
COMMON_DATA							
	CM	0000000D	00000000		0000000C	00	

This section of the listing contains the External Symbol Dictionary information passed to the Binder

## External Symbol Dictionary

- Each entry in the ESD has a particular type:
  - SD – Section Definition
    - The symbol appeared in the name field of a START, CSECT or RSECT instruction
  - LD – Label Definition
    - The symbol appeared as the operand of an ENTRY statement. When you specify the GOFF assembler option on z/OS or CMS, the assembler generates an entry type of LD for each CSECT and RSECT name.
  - ER – External Reference
    - The symbol appeared as the operand of an EXTRN statement or appeared as an operand of a V-type address constant.
  - CM – Common control section definition
    - The symbol appeared in the name field of a COM statement

## Source Statement

- This section of the listing documents source statements of the module and the resulting object code.
- The TITLE, CEJECT and EJECT assembler instructions can be used to control when the page title is printed

```

00000034                                     360 *
00000034 9836 2014                             361 continue_code_again dc 0h'0'   define a label
00000038 58B0 A038                             362      lm r3,r6,l_regs         reload registers
0000003C 50B0 3014                             363      l  r11,=f'1504'        load constant
00000040 58C0 A028                             364      st r11,comm_country    ... and save in common
00000044 07FE                                 365      l  r12,pDateFormat      load address constant
                                           366      br  r14                and exit program
                                           367 *

```

Active	Usings:	linkage_data(X'54'),R2	common_data,R3	static_data(X'3C'),R10					Page	4
R-Loc	Object Code	Addr1	Addr2	Stmt	Source Statement	HLASM	R6.0	2013/11/01	15.23	
00000046	0000									
00000048				369	static_data dc 0d'0'					start of static
00000048	839697A899898788			370	copyright dc c'copyright IBM(UK) Ltd 2013'					
00000062	0000									
00000064	00000000			371	listingx dc v(external_function)					declare external
00000068	E8E8E8E8D4D4C4C4			372	dateFormat dc c1(date_len)'YYYYMMDD'					define format

## Source Statement – Location Counter and Statement number

00000034				360 *				
00000034	9836	2014		361	continue_code_again	dc	0h'0'	define a label
00000038	58B0	A038	00000014	362	lm	r3,r6,l_regs		reload registers
0000003C	50B0	3014	00000080	363	l	r11,=f'1504'		load constant
00000040	58C0	A028	00000014	364	st	r11,comm_country		... and save in common
00000044	07FE		00000070	365	l	r12,pDateFormat		load address constant
				366	br	r14		and exit program
				367 *				

Active Usings:	linkage_data(X'54'),R2	common_data,R3	static_data(X'3C'),R10					Page	4
R-Loc	Object Code	Addr1	Addr2	stmt	Source Statement			HLASM R6.0	2013/11/01 15.23
00000046	0000								
00000048				369	static_data	dc	0d'0'		start of static
00000048	839697A899898788			370	copyright	dc	c'copyright IBM(UK) Ltd 2013'		
00000062	0000								
00000064	00000000			371	listingx	dc	v(external_function)		declare external
00000068	E8E8E8E8D4D4C4C4			372	dateFormat	dc	c1(date_len)'YYYYMMDD'		define format

## Source Statement – Location Counter and Statement number

```

Active Usings: linkage_data(X'54'),R2 common_data,R3 static_data(X'3C'),R10
R-Loc  Object Code      Addr1  Addr2  Stmt  Source Statement      HLASM R6.0  2013/11/01 15.23
00000074 00000000          374 pZERO      dc a(0)
00000078 00000000          375 ListData   dc v(FUNCY)
          376          extrn listingz
          377 listingz  alias c'listme'
          378 *
00000080          379          ltorg
00000080 000005E0          380          =f'1504'
          381 *
          00000084 382 static_data_end equ *
          383 *
          384          drop r2
          385          drop r3
          386 *
00000000          00000000 00000018 387 common_data com
00000000 A8A8A8A894948484 388 comm_date   dc cl(date_len)'yyy....'
00000008 A8A8A8A894948484 389 comm_user    dc cl(10)'yyyymmdd'
00000012 0000
00000014 00000000          390 comm_country dc f'0000'
          391 *
00000000          00000000 00000054 392 linkage_data dsect ,
00000000 C140D58194854040 393 l_name       dc cl(10)'A Name'
0000000A A8A8A8A894948484 394 l_date       dc cl(date_len)'yyyymmdd'

```

## Source Statement – Location Counter and Statement number

		19	print on,gen	print statements
		20	sysstate archlvl=1	set arch level
		21+*	THE VALUE OF SYSSTATE IS NOW SET TO ASCENV=P AMODE64=NO	
		+	L=1 OSREL=00000000	
		22	ieabrcx DEFINE	use relative branching
	00000048	341	larl r10,static_data	
00000048		342	using (static_data,static_data_end),r10	
		343 *		

The column following the statement number contains one of these values:

A space ( ) indicates open source

A plus sign (+) indicates that the statement was generated as the result of macro call processing.

An unnumbered statement with a plus sign (+) is the result of open code substitution.

A minus sign (-) indicates that the statement was read by a preceding AREAD instruction.

An equals sign (=) indicates that the statement was included by a COPY instruction.

A greater-than sign (>) indicates that the statement was generated as the result of a preceding AINSERT instruction. If the statement is read by an AREAD instruction, this takes precedence and a minus sign is printed.

## Source Statement - Addr1 and Addr2 Fields and USING Statements

R-Loc	Object Code	Addr1	Addr2	Stmt	Source Statement
				21+*	THE VALUE OF SYSSTATE IS NOW SET TO ASCENV=P..
				22	ieabrcx DEFINE use relative branching
00000000	C0A0 0000 0024		00000048	341	larl r10,static_data
		R:A 00000048		342	using (static_data,static_data_end),r10
			343 *		
00000006	17FF			344	xr r15,r15 initialise register
		R:2 00000000		345	using (linkage_data,l_end),r2 set using scope
		R:3 00000000		346	using common_data,r3 set using scope
00000008	C040 0000 000A		0000001C	347	larl r4,address_constant address of....
0000000E	D207 3000 200A	00000000	0000000A	348	mvc comm_date,l_date copy
00000014	9036 2014		00000014	349	stm r3,r6,l_regs save a copy of the reg..
00000018	A7F4 0008		00000028	350	j continue_code jump over constant
				351 *	
0000001C	C1C4C4D9C5E2E240			352	address_constant dc c112'ADDRESS' constant value

## Relocation Dictionary (RLD)

Relocation Dictionary				
Pos.Id	Rel.Id	Address	Type	Action
00000004	00000004	<u>00000070</u>	A 4	+
00000004	00000006	00000064	V 4	ST
00000004	00000007	00000078	V 4	ST

---

00000048	839697A899898788	370	copyright	dc c'copyright IBM(UK) Lt
00000062	0000			
00000064	00000000	371	listingx	dc v(external_function)
00000068	E8E8E8E8D4D4C4C4	372	dateFormat	dc cl(date_len)'YYYYMMDD'
<u>00000070</u>	<u>00000068</u>	<u>373</u>	<u>pDateFormat</u>	<u>dc a(dateFormat)</u> po
00000074	00000000	374	pZERO	dc a(0) po
00000078	00000000	375	ListData	dc v(FUNCY) de
		376	extrn listingz	de

*Location counter 70, and ADCON symbol pDateFormat has a value of 00000068 – which is the location counter for symbol dateFormat*

- This section of the listing describes the relocation dictionary information passed to the Binder.
- The entries describe the address constants in the assembled program that are affected by relocation.
- This section helps you find relocatable constants in your program.



## Ordinary Symbol and Literal Cross Reference

Ordinary Symbol and Literal Cross Reference									
Symbol	Length	Value	Id	R	Type	Asm	Program	Defn	References
address_constant	12	0000001C	00000004		C	C		352	347
comm_country	4	00000014	0000000C		F	F		390	364M
comm_date	8	00000000	0000000C		C	C		388	348M
common_data	1	00000000	0000000C		J			387	346U
continue_code	2	00000028	00000004		H	H		354	350B
continue_code_again	2	00000034	00000004		H	H		361	357B
date_len	1	00000008	00000004	A	U			4	372 388 394
r10	1	0000000A	00000004	A	U	GR32		12	341M 342U
r11	1	0000000B	00000004	A	U	GR32		13	363M 364
r12	1	0000000C	00000004	A	U	GR32		14	365M
r14	1	0000000E	00000004	A	U	GR32		16	366B
r15	1	0000000F	00000004	A	U	GR32		17	344M 344

Symbol name , length, value, Assembler type, Definition

References – no suffix, **B**Branch, **M**Modified, **U**Using, **D**Drop and **eX**ecution

This section of the listing concerns symbols and literals that are defined and used in the program.

## Unreferenced Symbols Defined in CSECTs

### Unreferenced Symbols Defined in CSECTs

```
Defn Symbol
370 copyright
375 ListData
371 listingx
374 pZERO
  6 r0
  7 r1
 15 r13
```

This section of the listing shows symbols that have been defined in CSECTs but not referenced. This may help you to remove unnecessary data definitions, and reduce the size of your program.

## Macro and Copy Code Source Summary

Con	Source	Macro and Copy Code Source Summary	Page
	PRIMARY INPUT	Volume Members	8
		B BAL BAS BC HLASM R6.0 BCT 2013/11/01 15.23 BE BH	
		BL BM BNE BNH BNL BNM BNO	
		BNP BNZ BO BP BXH BXLE BZ	
L2	SYS1.MACLIB	37SY01 IEABRC IEABRCX SYSSTATE	

*In section 'Diagnostic Cross Reference and Assembler Summary'*

Data Sets Allocated for this Assembly

Con	DDname	Data Set Name	Volume	Member
A1	ASMAOPT	SMORSA.LISTINGC.JOB63822.D0000101.?		
P1	SYSIN	SMORSA.ASM.ASM	37P001	LISTINGC
L1	SYSLIB	SMORSA.ASM.ASM	37P001	
L2		SYS1.MACLIB	37SY01	
	SYSLIN	SYS13305.T152320.RA000.LISTINGC.OBJ.H01		
	SYSPRINT	SMORSA.LISTINGC.JOB63822.D0000102.?		

This section of the listing shows the names of the macro libraries from which the assembler read macros or copy code members, and the names of the macros and copy code members that were read from each library.

## Macro and Copy Code Cross Reference

Macro and Copy Code Cross Reference				
Macro	Con	Called By	Defn	References
B		PRIMARY INPUT	117	355
BAL			273	
BAS			297	
BC		PRIMARY INPUT	177	356
BCT			201	
BE			121	
BH			129	
BL			125	
BM			133	
BNE			137	
.....				
BXLE			225	
BZ			173	
IEABRC	L2	IEABRCX	-	22C
IEABRCX	L2	PRIMARY INPUT	-	22
SYSSTATE	L2	PRIMARY INPUT	-	20

---

```

19      print on,gen          print statements
20      sysstate archlvl=1    set arch level
21+*    THE VALUE OF SYSSTATE IS NOW SET TO ASCENV=P AMODE64=N
      +      L=1 OSREL=00000000
22      ieabrcx DEFINE        use relative branching
341     larl r10,static_data
342     using (static_data,static_data_end),r10

```

---

This section of the listing shows the names of macros and copy code members and the statements where the macro or copy code member was called.

## DSECT Cross Reference

				Dsect Cross Reference	
Dsect	Length	Id	Defn		
linkage_data	00000054	FFFFFFFF	392		
<hr/>					
00000000		00000000	00000054	391 *	
00000000	C140D58194854040			392 linkage_data dsect ,	data passed in to me
0000000A	A8A8A8A894948484			393 l_name dc cl(10)'A Name'	users name
00000012	0000			394 l_date dc cl(date_len)'yyyymmdd'	users date
00000014	0000000000000000			395 l_regs dc 4f'0,0,0,0'	return registers values
		00000054		396 l_end equ *	end
				397 *	
<hr/>					

This section of the listing shows the names of all internal or external dummy sections defined in the program, and the number of the statement where the definition of the dummy section began.

## Using Map

Stmt	-----Location-----		Action	Using Map			
	Count	Id		Type	Value	Range	Id
342	00000006	00000004	USING	ORDINARY	00000048	0000003C	00000004
345	00000008	00000004	USING	ORDINARY	00000000	00000054	FFFFFFFF
346	00000008	00000004	USING	ORDINARY	00000000	00001000	0000000C
384	00000084	00000004	DROP				
385	00000084	00000004	DROP				

Stmt	-----Location-----	Action	-----	.....	Reg	Max	Last Label and Using Text
	Count	Id	Type	.....	Disp		Stmt
342	00000006	00000004	USING	ORDINARY	.....	10 00038	365 (static_data,static_data_end),r10
345	00000008	00000004	USING	ORDINARY	.....	2 00014	362 (linkage_data,l_end),r2
346	00000008	00000004	USING	ORDINARY	.....	3 00014	364 common_data,r3
384	00000084	00000004	DROP		.....	2	r2
385	00000084	00000004	DROP		.....	3	r3

	Addr1	Addr2	Stmt	Source Statement
R:A	00000048	00000048	341	larl r10,static_data
			342	using (static_data,static_data_end),r10

Active Usings: linkage\_data(X'54'),R2 common\_data,R3 static\_data(X'3C'),R10

R-LOC	Object Code	Addr1	Addr2	Stmt	Source Statement
00000046	0000				
00000048				369	static_data dc 0d'0' start of
00000048	839697A899898788			370	copyright dc c'copyright IBM(UK) Ltd 2
00000062	0000				
00000064	00000000			371	listingx dc v(external_function)

This section of the listing shows a summary of the USING, DROP, PUSH USING, and POP USING instructions used in your program.

## General Purpose Register Cross Reference

### General Purpose Register Cross Reference

Register References (M=modified, B=branch, U=USING, D=DROP, N=index)

```

0(0) (no references identified)
1(1) (no references identified)
2(2) 345U 384D
3(3) 346U 349 362M 385D
4(4) 347M 349 362M
5(5) 349 362M
6(6) 349 362M
7(7) (no references identified)
8(8) (no references identified)
9(9) (no references identified)
10(A) 341M 342U
11(B) 363M 364
12(C) 365M
13(D) (no references identified)
14(E) 366B
15(F) 344M 344

```

#### Register 10 – no DROP statement?

Loc	Object	Code	Addr1	Addr2	Stmt	Source	Statement
00000000	C0A0	0000	0024	00000048	341		larl r10,static_data
			R:A 00000048		342		using (static_data,static_data_end),r10

#### Register 5?

00000014	9036	2014		00000014	349		stm r3,r6,l_regs
00000034	9836	2014		00000014	362		lm r3,r6,l_regs

This section of the listing shows all references in the program to each of the general registers. Additional flags indicate the type of reference.

# Diagnostic Cross Reference and Assembler Summary

## Diagnostic Cross Reference and Assembler Summary

Statements Flagged  
24(P1,24)

1 Statement Flagged in this Assembly 8 was Highest Severity Code  
HIGH LEVEL ASSEMBLER, 5696-234, RELEASE 6.0, PTF R160  
SYSTEM: z/OS 01.13.00 JOBNAME: LISTINGB STEPNAME: B PROCSTEP: C

---

R-Loc	Object Code	Addr1	Addr2	Stmt	Source	Statement
0000000C	50B0 3014		00000014	22		st r11,comm_country
00000010	58C0 F040		00000040	23		l r12,pDateFormat
00000014	0000 0000		00000000	24		l r7,someData
** ASMA044E Undefined symbol - r7						
** ASMA029E Incorrect register specification - r7						
** ASMA044E Undefined symbol - someData						
** ASMA435I Record 24 in SMORSA.ASM.ASM(LISTINGB) on volume: 37P001						
00000018	07FE			25		br r14
				26	*	

---

Data Sets Allocated for this Assembly

Con	DDname	Data Set Name	Volume	Member
A1	ASMAOPT	SMORSA.LISTINGB.JOB63825.D0000101.?		
P1	SYSIN	SMORSA.ASM.ASM	37P001	LISTINGB
L1	SYSLIB	SMORSA.ASM.ASM	37P001	
L2		SYS1.MACLIB	37SY01	
	SYSLIN	SYS13305.T173626.RA000.LISTINGB.OBJ.H01		
	SYSPRINT	SMORSA.LISTINGB.JOB63825.D0000102.?		

---

This section of the listing summarises the error diagnostic messages issued during the assembly, and provides statistics about the assembly.



## Diagnostic Cross Reference and Assembler Summary

### Diagnostic Cross Reference and Assembler Summary

No Statements Flagged in this Assembly

HIGH LEVEL ASSEMBLER, 5696-234, RELEASE 6.0, PTF R160

SYSTEM: z/OS 01.13.00

JOBNAME: LISTINGC

STEPNAME: B

PROCSTEP: C

Data Sets Allocated for this Assembly

Con	DDname	Data Set Name	Volume	Member
A1	ASMAOPT	SMORSA.LISTINGC.JOB63822.D0000101.?		
P1	SYSIN	SMORSA.ASM.ASM	37P001	LISTINGC
L1	SYSLIB	SMORSA.ASM.ASM	37P001	
L2		SYS1.MACLIB	37SY01	
	SYSLIN	SYS13305.T152320.RA000.LISTINGC.OBJ.H01		
	SYSPRINT	SMORSA.LISTINGC.JOB63822.D0000102.?		

4096k allocated to Buffer Pool

Storage required 200k

76 Primary Input Records Read

1093 Library Records Read

0 Work File Reads

2 ASMAOPT Records Read

312 Primary Print Records Written

0 Work File Writes

23 Object Records Written

0 ADATA Records Written

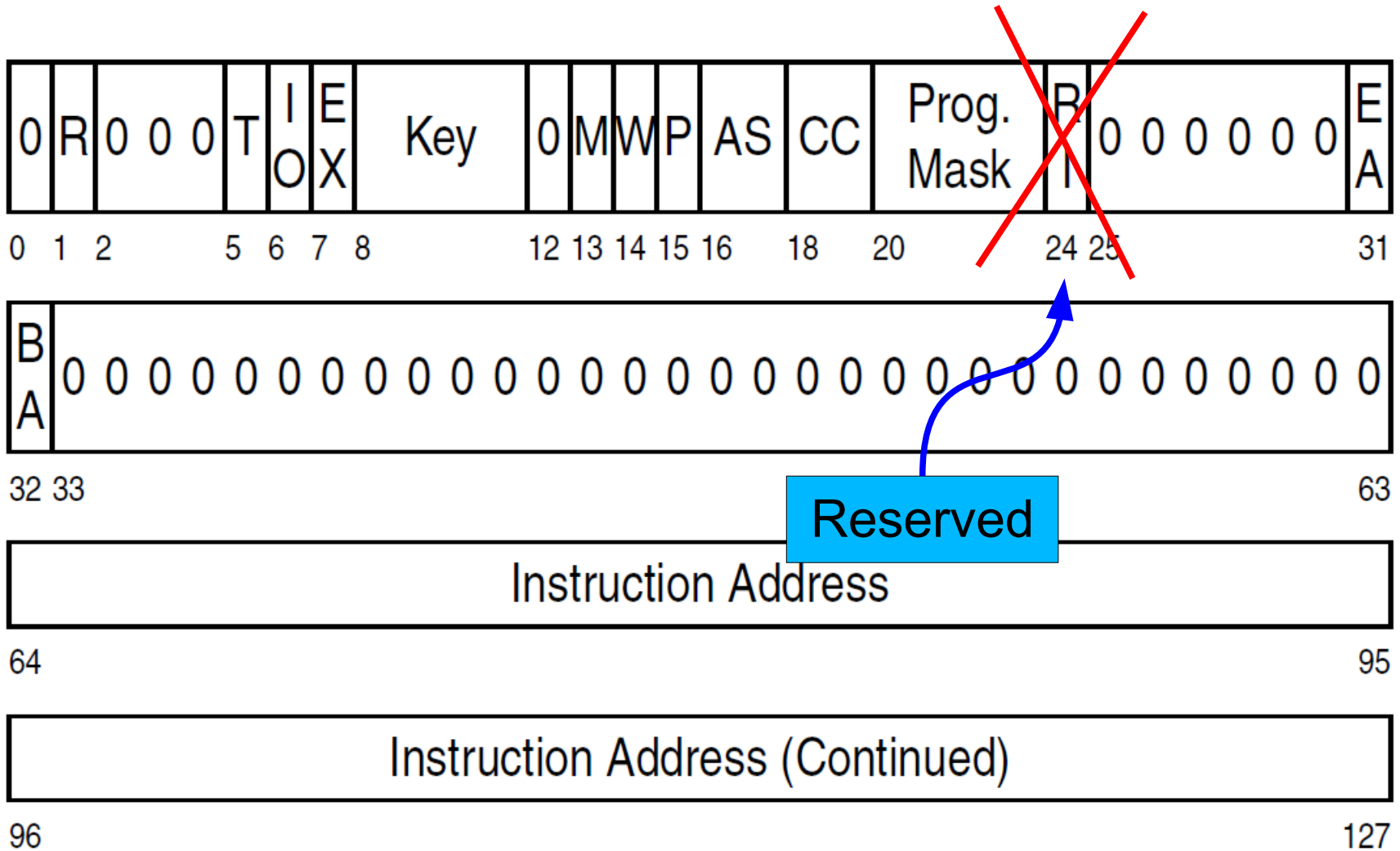
Assembly Start Time: 15.23.20 Stop Time: 15.23.20 Processor Time: 00.00.00.0044

Return Code 000

# The PSW and an introduction to debugging assembler programs



# Program Status Word (PSW)



## Program Status Word (PSW)

- The Program Status Word (PSW) is a register in the processor which includes control information to determine the state of the CPU.
- The z/Architecture PSW is 128-bits in length
  - Bits 0-32 contain flag bits indicating control information for the CPU
  - Bits 33-63 are 0
  - Bits 64-127 contain the instruction address
- EPSW – Extract PSW
  - Obtain bits 0-63 of the PSW and place them into operands of the instruction
- LPSW(E) – Load PSW (Extended)
  - Replace the entire PSW with the contents of storage
  - This means that the instruction branches – well might do...



## Program Status Word (PSW) – Addresses

- Since the z/Architecture can run in a number of addressing modes, the instruction address is determined by a variable number of bits in the PSW. The current addressing mode is determined by bits 31-32 of the PSW with the following combinations:
  - 00 → 24-bit mode
  - 01 → 31-bit mode
  - 10 → invalid
  - 11 → 64-bit mode
  
- Bits 64-127 are used to determine the address of the next instruction to be executed
  - However, some instructions may be interrupted and therefore the PSW may point at the same instruction which was being executed so that it is redriven





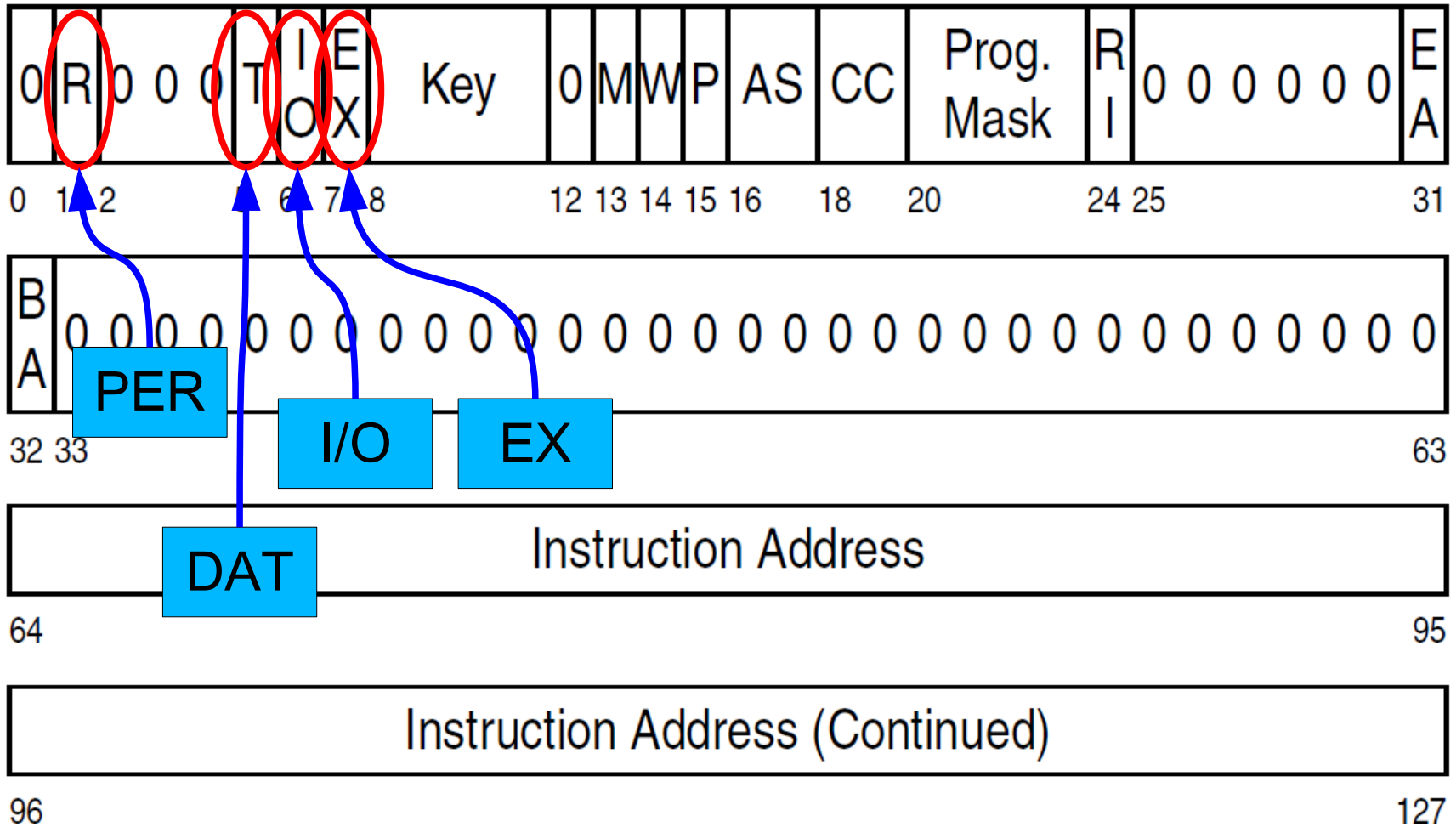


## Program Status Word (PSW) – Flag bits

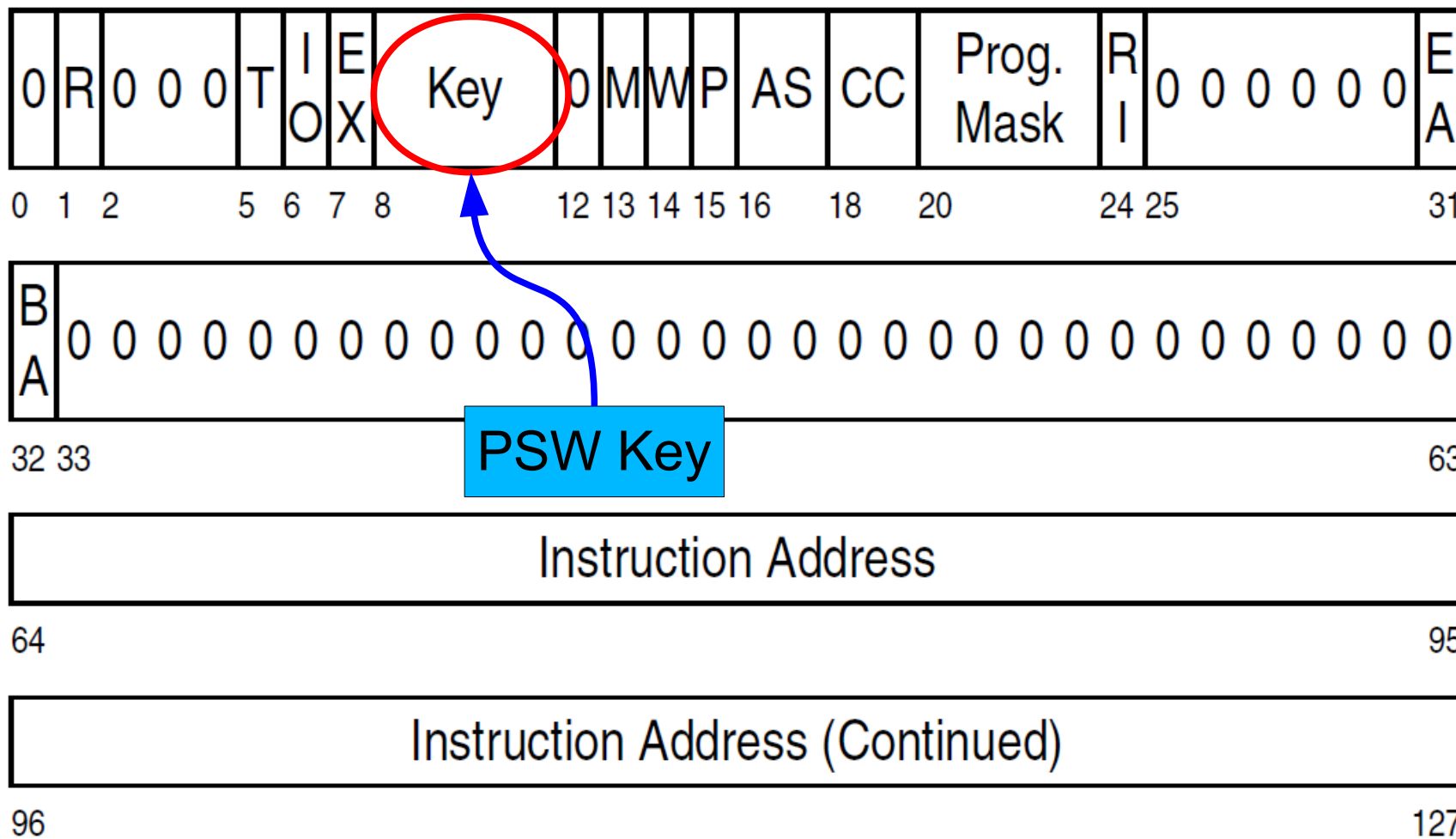
(A bit value of 1 indicates that the CPU is enabled for a function unless stated otherwise)

- Bit 1 – Program Event Recording (PER) Mask
  - Controls whether the CPU is enabled for interrupts associated with PER
- Bit 5 – DAT Mode
  - Controls whether the CPU has Dynamic Address Translation turned on
- Bit 6 – I/O Mask
  - Controls whether the CPU is enabled for interrupts
- Bit 7 – External Mask
  - Controls whether the CPU is enabled for external interrupts

## Program Status Word (PSW) – Flag bits



## Program Status Word (PSW) – PSW Key



## Program Status Word (PSW) – PSW Key

- Bits 8-11 are used to represent the PSW Key (value of 0-15)
- PSW Keys are used to provide a security mechanism over various regions of memory with key 0 being the most secure
- Whenever an instruction attempts to access a storage location that is protected against that type of reference (read/write of storage) and the storage key does not match the access key, a protection exception is recognised.
- Programs running in PSW key 0 have read write access to storage in every storage key
- Programs in keys 1 – 15 have read access to:
  - Storage which matches their PSW key
  - Storage (in any key) that's not fetch protected
  - Storage in key 9 if the hardware feature “subsystem storage protection override” is installed
- Programs in keys 1-15 have write access to:
  - Storage whose key matches their PSW key
  - Storage in key 9 if subsystem storage protection override is installed

## Program Status Word (PSW) – PSW Key – Manipulation

- IPK – Insert PSW Key
  - Used to *insert* the PSW Key *into* register 2
  - Used to store a copy of the current PSW Key typically before a switch to another key.
  - Bits 56-59 of register 2 are updated to contain the PSW Key, bits 60-63 are set to 0 and all other bits remain unmodified
- IPK cannot be used when bit 36 of CR0 is set to 0 and in problem state
- SPKA – Set PSW Key from Address
  - Used to set the PSW Key from an *address value*
  - Bits 56-59 of the 2<sup>nd</sup> operand are inserted into the PSW Key
- SPKA can only be used to set a key to which the current task is allowed to set a key determined by the PSW Key mask in CR3
- Both IPK and SPKA are privileged instructions

## Program Status Word (PSW) – PSW Key – Manipulation

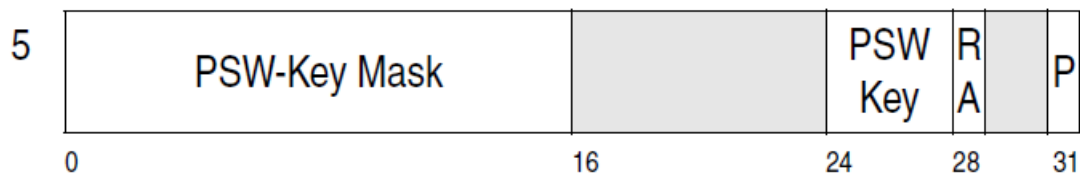
- MVCK – Move with Key
  - Moves an operand with an access key specified as part of the instruction
  - If the program is not enabled to use that access key, then a privileged operation exception is raised
  - Can be a slow instruction
- BSA – Branch and Set Authority
  - Used to branch to another place in code and set the PSW key at the same time
  - Works as a flip-flop branching from “base authority” state to “reduced authority” state

## Program Status Word (PSW) – PSW Key – BSA Operation

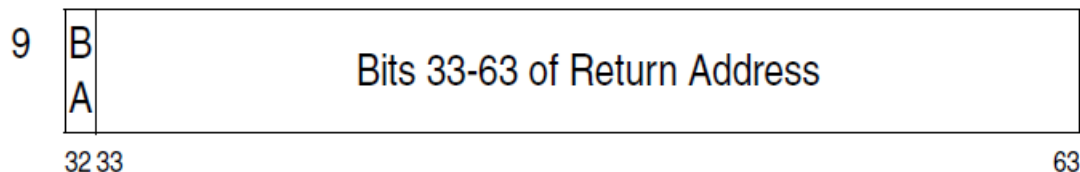
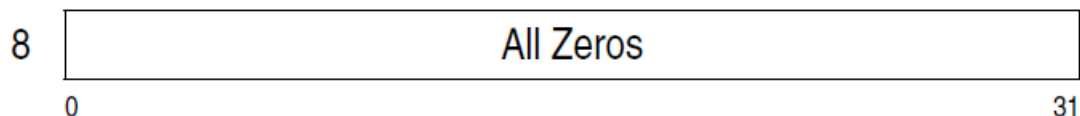
- BSA – Branch and Set Authority – example scenario
  - A service routine, e.g. a middleware service begins in the *base authority* state
  - The routine issues a BSA to switch to running a user routine
  - The user routine runs in *reduced authority* state
  - When the user routine wants to invoke the middleware service, it issues a BSA which branches back to a fixed location in the middleware and the state is returned to running in base authority state
- The control of the states is determined by the Dispatchable Unit Control Table (DUCT)
  - The BSA instruction uses words 5, 8 and 9 of the DUCT.



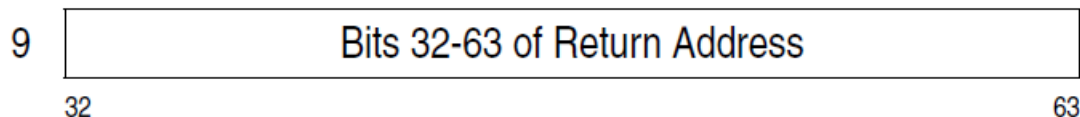
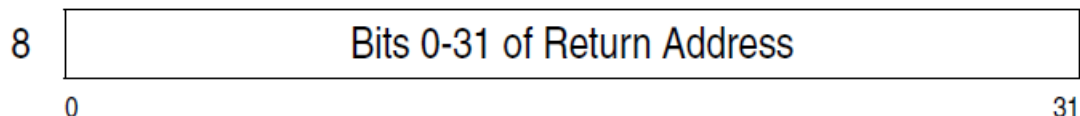
# Program Status Word (PSW) – PSW Key – BSA Operation - DUCT



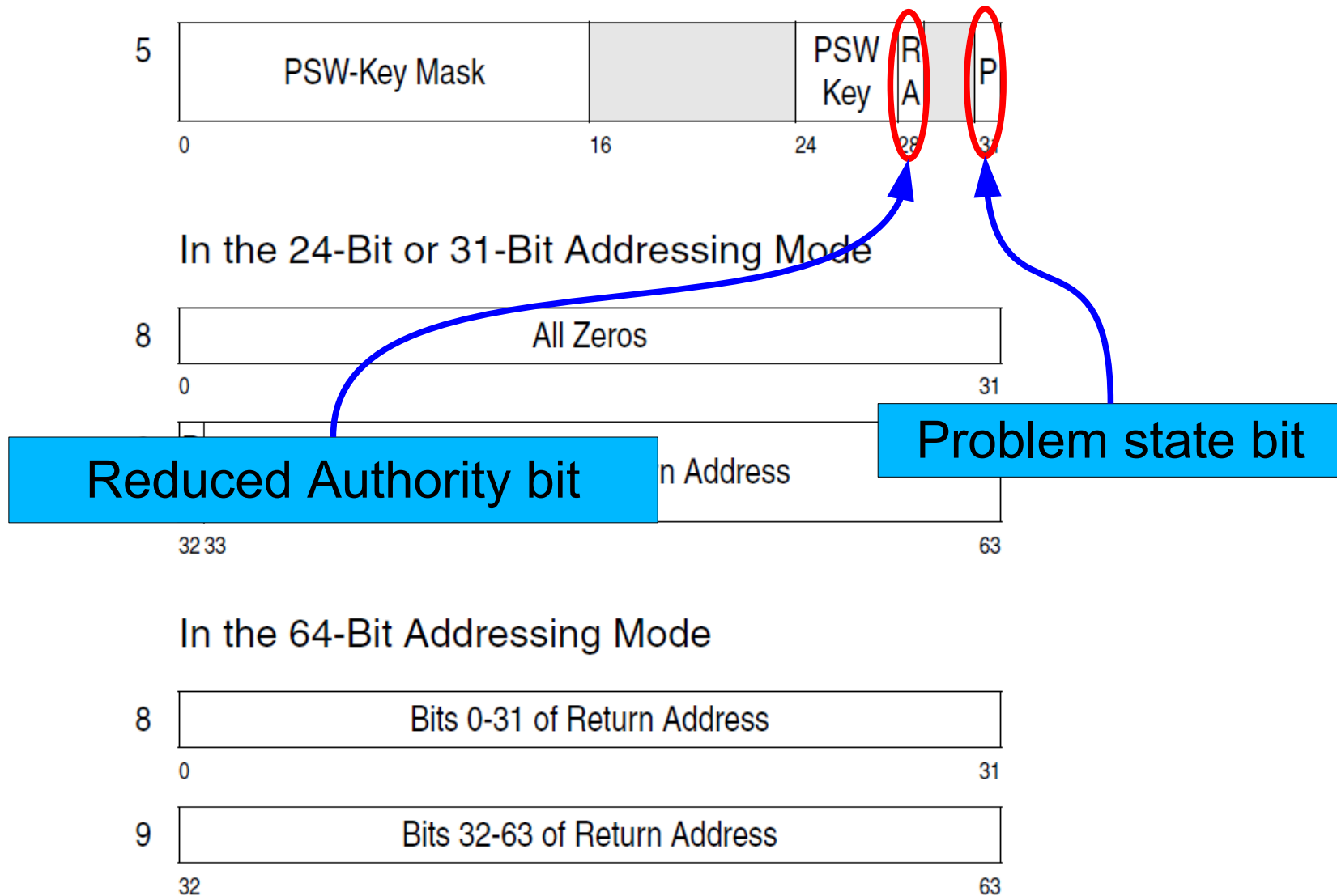
## In the 24-Bit or 31-Bit Addressing Mode



## In the 64-Bit Addressing Mode



# Program Status Word (PSW) – PSW Key – BSA Operation



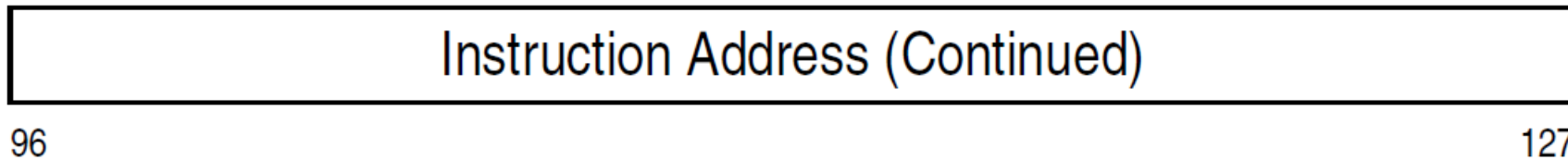
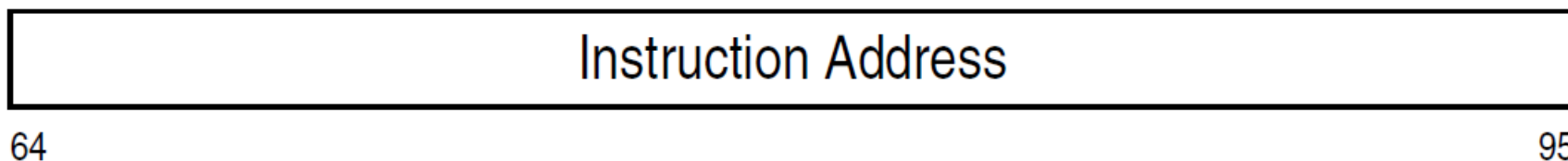
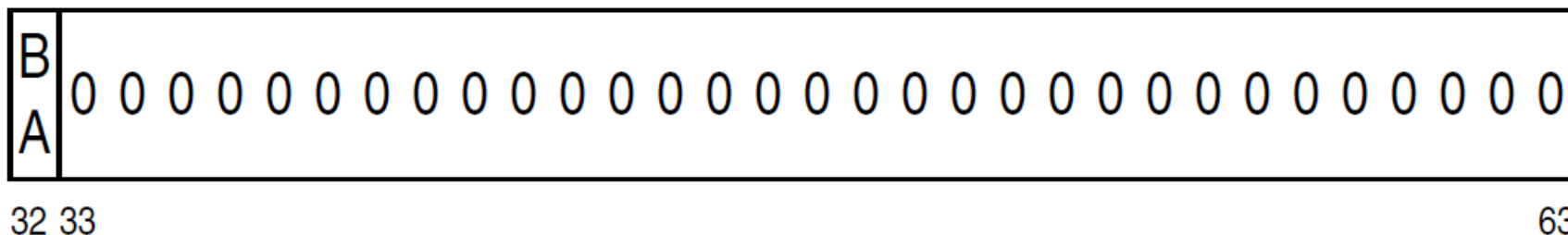
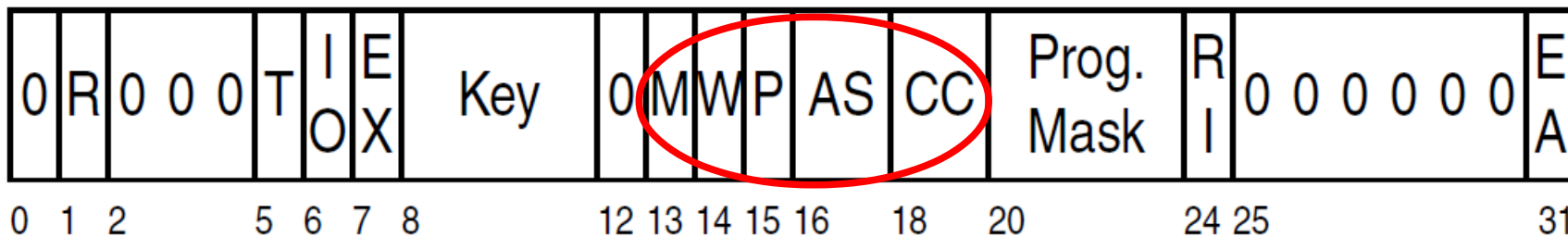
## Program Status Word (PSW) – PSW Key – BSA Operation – BA

- When the BSA instruction is used in base authority, the following is stored in the DUCT:
  - The PSW-key Mask (from CR3)
  - The current PSW Key
  - Problem state bit
  - The return address
  
- BSA then sets the Reduced Authority bit (RA) to 1 and loads:
  - The PSW-key Mask into CR3 from operand 1
  - The PSW Key from operand 1
  - The branch address into the PSW

## Program Status Word (PSW) – PSW Key – BSA Operation – RA

- When the BSA instruction is used in reduced authority, the following is restored from the DUCT:
  - The PSW-key Mask (to CR3)
  - The current PSW Key (to the PSW)
  - Problem state bit (to the PSW)
  - The return address (to the PSW – therefore the machine branches...)
  
- BSA then sets the Reduced Authority bit (RA) to 0

## Program Status Word (PSW) – More flag bits



## Program Status Word (PSW) – More flag bits

- Bit 13 – Machine Check Mask
  - Controls whether the CPU is enabled for interrupts by machine check conditions
- Bit 14 – Wait State
  - If on, the machine is waiting and no instructions are processed but interrupts may take place.
- Bit 15 – Problem State
  - The machine operates in two states – problem state (used for user code) and supervisor state (used for privileged code)
  - If an attempt is made to execute a privileged instruction in problem state, then a privileged operation exception occurs.  
Some instructions are *semi-privileged* and may or may not be permitted to execute in problem state depending on the outcome of other flags
  - All instructions are valid in supervisor state

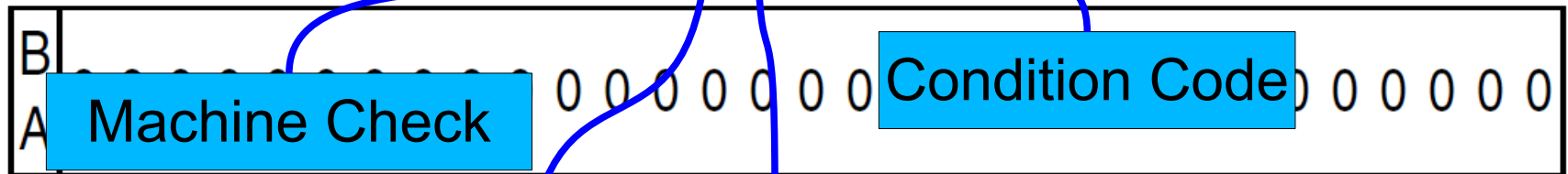
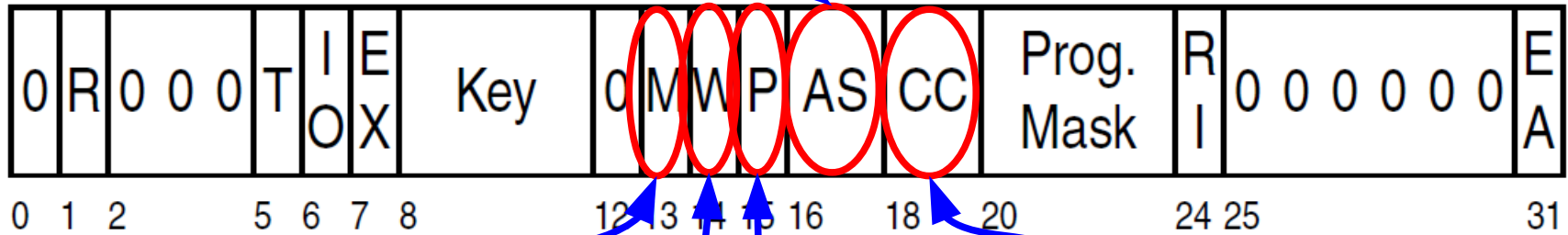
## Program Status Word (PSW) – More flag bits

- Bit 16-17 – Address-Space Control
  - Determines how addresses are handled in conjunction with bit 5 (DAT) via the following table:

5	16	17	DAT	Mode	Instruction Addresses	Logical Addresses
0	0	0	Off	Real Mode	Real	Real
0	0	1	Off	Real Mode	Real	Real
0	1	0	Off	Real Mode	Real	Real
0	1	1	Off	Real Mode	Real	Real
1	0	0	On	Primary-space Mode	Primary virtual	Primary virtual
1	0	1	On	Access-register Mode	Primary virtual	AR specified vrt
1	1	0	On	Secondary-space Mode	Primary virtual	Secondary vrt
1	1	1	On	Home-space Mode	Home virtual	Home virtual

# Program Status Word (PSW) – More flag bits

**Address Space Control**



32 33 **Problem state** 63

**Wait** Instruction Address 95

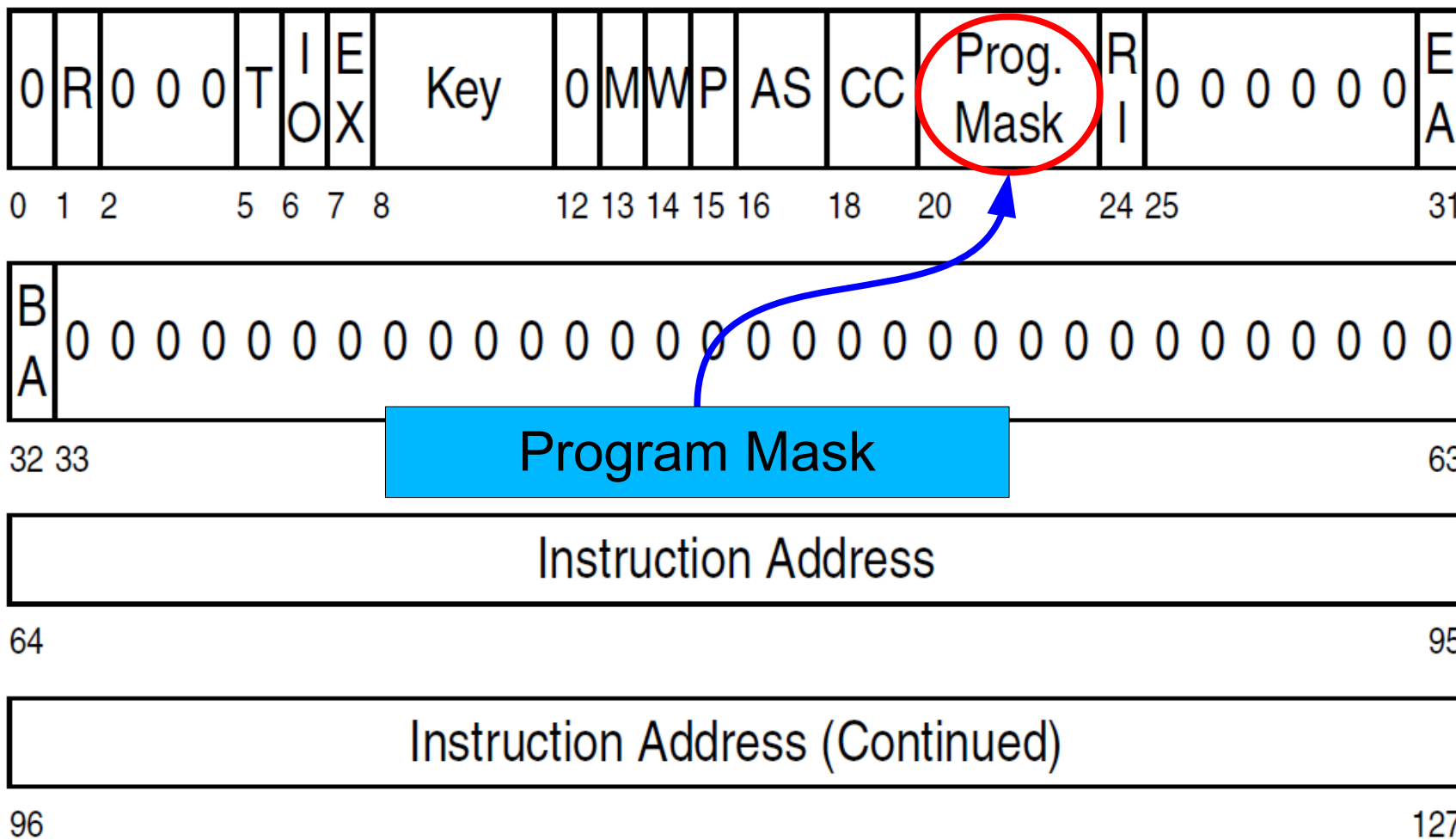
64

Instruction Address (Continued)

96 127



## Program Status Word (PSW) – PSW Key – Program Mask



## Program Status Word (PSW) – Program Mask

- Bits 20-23 – Program Mask
  - Controls a set of program exceptions
  - When the corresponding bit is on, the exception results in an interrupt

Program Mask PSW bit	Program Exception
20	Fixed-point overflow
21	Decimal overflow
22	HFP exponent underflow
23	HFP significance

- The Program Mask can be manipulated by using the instruction SET PROGRAM MASK (SPM)
- The contents of the Program Mask can be examined using the instruction INSERT PROGRAM MASK (IPM)

## Program Status Word (PSW) – Using the PSW for debugging

- The PSW stores invaluable information about the general state of the machine during a program's execution
- The most interesting time to examine a PSW is when something goes wrong. Even a summary dump will provide the programmer with:
  - The contents of the PSW
  - The contents of the general purpose registers
  - The next instruction to be executed

## Program Status Word (PSW) – Something has gone wrong

```
SYSTEM COMPLETION CODE=0C4 REASON CODE=00000010
TIME=10.58.19 SEQ=01263 CPU=0000 ASID=002C
PSW AT TIME OF ERROR 078D0000 80007FF6 ILC 4 INTC 10
ACTIVE LOAD MODULE ADDRESS=00007FF0 OFFSET=00000006
NAME=GO
DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C00607FE
GR 0: FD000008 1: 00006FF8
2: 00000040 3: 008DAD6C
4: 008DAD48 5: 008FF130
6: 008CAFC8 7: FD000000
8: 008FCE28 9: 008D88F0
A: 00000000 B: 008FF130
C: FD000008 D: 00006F60
E: 80FDA688 F: 80007FF0
END OF SYMPTOM DUMP
```

- Running a job has resulted in an 0C4 ABEND occurring. The summary dump in the job may be enough information to work out what has gone wrong.

## Program Status Word (PSW) – Something has gone wrong

```

SYSTEM COMPLETION CODE=0C4 REASON CODE=00000010
TIME=10.58.19 SEQ=01263 CPU=0000 ASID=002C
PSW AT TIME OF ERROR 078D0000 80007FF6 ILC 4 INTC 10
ACTIVE LOAD MODULE ADDRESS=00007FF0 OFFSET=00000006
NAME=GO
DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C00607FE
GR 0: FD000008 1: 00006FF8
2: 00000040 3: 00BDAD6C
4: 00BDAD48 5: 008FE130
6: 008CAFCE8 7: FD000000
8: 008FCE28 9: 008D88F0
A: 00000000 B: 008FF130
C: FD000008 D: 00006F60
E: 80FDA688 F: 80007FF0
END OF SYMPTOM DUMP

```

- First, look at the active load module
  - In this example, the load module name is GO since the LKEDG JCL procedure was used. From this we already know that the error occurred in our load module and not in either the assembler, linkage editor nor other part of z/OS

## Program Status Word (PSW) – Something has gone wrong

```

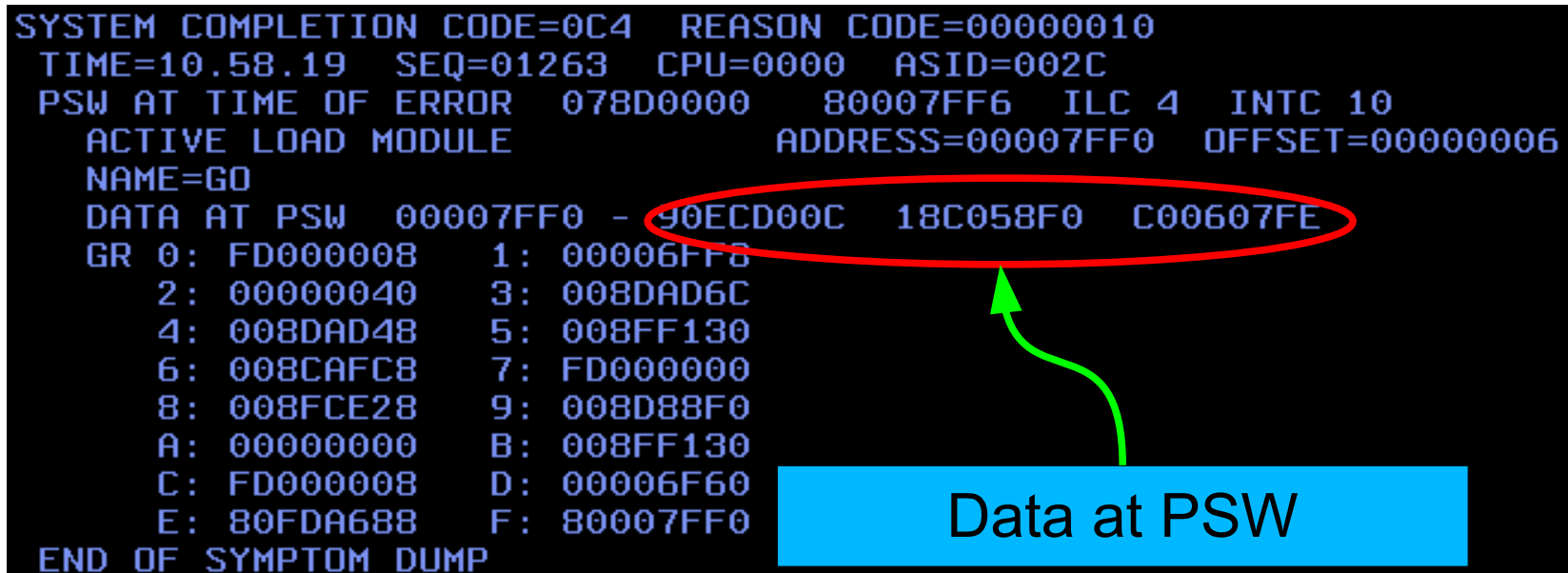
SYSTEM COMPLETION CODE=0C4 REASON CODE=00000010
TIME=10.58.19 SEQ=01263 CPU=0000 TSTD=002C
PSW AT TIME OF ERROR 078D0000 80007FF6 ILC 4 INTC 10
ACTIVE LOAD MODULE ADDRESS=00007FF0 OFFSET=00000006
NAME=GO
DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C0607FE
GR 0: FD000008 1: 00006FF8
B: 008CFE28 7: FD000000
B: 008FCE28 9: 008D
A: 00000000 B: 008F
C: FD000008 D: 00006FF8
E: 80FDA688 F: 80007FF0
END OF SYMPTOM DUMP

```

- Next, double check the information in the PSW against the other information in the summary dump
  - The PSW shows that the next instruction address to be executed is X'7FF6'
  - This agrees with the data in the dump showing the address of the load module (X'7FF0') and the offset into the load module (X'0006')

## Program Status Word (PSW) – Something has gone wrong

```
SYSTEM COMPLETION CODE=0C4 REASON CODE=00000010
TIME=10.58.19 SEQ=01263 CPU=0000 ASID=002C
PSW AT TIME OF ERROR 078D0000 80007FF6 ILC 4 INTC 10
ACTIVE LOAD MODULE ADDRESS=00007FF0 OFFSET=00000006
NAME=GO
DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C00607FE
GR 0: FD000008 1: 00006FF8
2: 00000040 3: 008DAD6C
4: 008DAD48 5: 008FF130
6: 008CAFC8 7: FD000000
8: 008FCE28 9: 008D88F0
A: 00000000 B: 008FF130
C: FD000008 D: 00006F60
E: 80FDA688 F: 80007FF0
END OF SYMPTOM DUMP
```



Data at PSW

- The data at the PSW shows the instructions which were, are being, and will be executed

## Program Status Word (PSW) – Something has gone wrong

**DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C00607FE**

```

000000          00000 00010      15 *
                                16 * MAIN PROGRAM STARTS HERE
                                17 *
000000          90EC D00C          0000C      18 EX1      CSECT
000004          18C0              R:C 00006      19 EX1      AMODE 31
                                20 EX1      RMODE 24
                                21 * USUAL PROGRAM SETUP
                                22          STM   R14,R12,12(R13)
                                23          LR    R12,0
                                24          USING *,12
000006          58F0 C006          0000C      25 * SET RETURN CODE IN REGISTER 15
                                26          L     R15,RET_CODE
00000A          07FE              27 * RETURN
                                28          BR   R14
                                29 * *****
                                30 * END OF PROGRAM
                                31 * *****
00000C          00000000          32 RET_CODE DC   F'0'
000010          00000000          33 LTORG ,

```

- Examining the program listing at offset 6 shows where the error occurred. Using the data at the PSW and looking at the machine code generated by HLASM in the listing confirms this and that so far our diagnosis of the problem is correct



## Program Status Word (PSW) – Something has gone wrong

DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C00607FE

Offset 6

```

000000 000000 00010 18 EX1 CSECT
000004 90EC D00C 0000E 19 EX1 AMODE 31
R:C 00006 20 EX1 RMODE 24
000006 58F0 C006 0000C 21 * USUAL PROGRAM SETUP
00000A 07FE 22 STM R14,R12,12(R13)
23 LR R12,0
24 USING *,12
25 * SET RETURN CODE IN REGISTER 15
26 L R15,RET_CODE
27 * RETURN
28 BR R14
29 * *****
30 * END OF PROGRAM
31 * *****
00000C 00000000 32 RET_CODE DC F'0'
000010 33 LTORG ,
    
```

- Examining the program listing at offset 6 shows where the error occurred. Using the data at the PSW and looking at the machine code generated by HLASM in the listing confirms this and that so far our diagnosis of the problem is correct
- We now know the instruction which caused the error was:  
 58F0 C006 → L R15,RET\_CODE

## Program Status Word (PSW) – Something has gone wrong

- At this stage of our debugging we know:
  - The load module name that caused the error
  - The offset into the load module at which the error occurred
- We have also double-checked that what was printed in the summary dump is confirmed by the data at the PSW
- Examining the instruction at fault, we determine the following:
  - 58F0 C006 → L R15,RET\_CODE
  - 58 – OPCODE = LOAD
  - F – Register 15, the register to be loaded
  - 0 – Index register (unused since it has a value of 0)
  - C – Base register is register 12
  - 006 – Displacement from the base register from which the data will be loaded
- So, the instruction is attempting to load register 15 with the contents of memory at an offset of 6 from register 12.

## Program Status Word (PSW) – Something has gone wrong

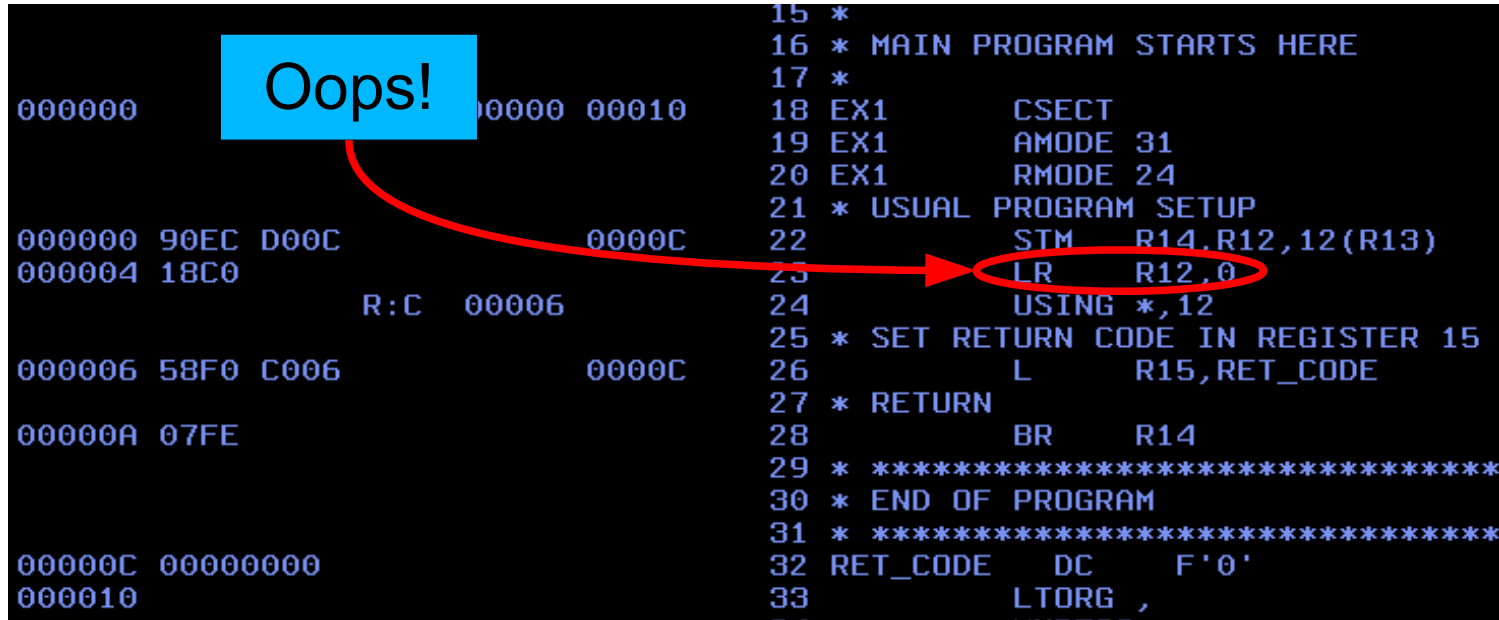
```

SYSTEM COMPLETION CODE=0C4 REASON CODE=00000010
TIME=10.58.19 SEQ=01263 CPU=0000 ASID=002C
PSW AT TIME OF ERROR 078D0000 80007FF6 ILC 4 INTC 10
ACTIVE LOAD MODULE ADDRESS=00007FF0 OFFSET=00000006
NAME=GO
DATA AT PSW 00007FF0 - 90ECD00C 18C058F0 C00607FE
GR 0: FD000008 1: 00006FF8
 2: 00000040 3: 008DAD6C
 4: 008DAD48 5: 008FF130
 6: 008CAFC8 7: FD000000
 8: 008FCE28 9: 008D88F0
 A: 00000000 B: 008FF130
 C: FD000008 D: 00006F60
 E: 80FDA688 F: 80007FF0
END OF SYMPTOM DUMP

```

- The summary dump also shows the contents of the general purpose registers
- The value in register 12 is X'FD000008'
- The instruction at fault is attempting to load a value from address X'FD00000E' – which is unaddressable by our program and therefore the cause of the error
- Note that the value of register 12 is the same as the value of register 0...

## Program Status Word (PSW) – Something has gone wrong



```

0000000 000000 000010 15 *
0000000 90EC D00C 0000C 16 * MAIN PROGRAM STARTS HERE
0000004 18C0 R:C 00006 17 *
0000006 58F0 C006 0000C 18 EX1 CSECT
000000A 07FE 23 LR R12,0 19 EX1 AMODE 31
000000C 00000000 24 USING *,12 20 EX1 RMODE 24
0000010 25 * USUAL PROGRAM SETUP
26 * SET RETURN CODE IN REGISTER 15
27 * RETURN
28 BR R14
29 * *****
30 * END OF PROGRAM
31 * *****
32 RET_CODE DC F'0'
33 LTORG ,

```

- Looking back through the program, we can see that register 12 was loaded with the value of register 0 during program startup
- It looks as if the programmer made a typo and instead of using LR 12,0 should have used BALR 12,0 in order to load the address of the next instruction into register 12. This would make sense since they are using register 12 to establish addressability to the program's data
- Correcting this mistake fixes the program

## Summary

- Introductory topics
  - Computer organisation and z/Architecture
  - Building programs on z Systems
  - Working with HLASM
- Programming in Assembler
  - Loading, storing and moving data
  - Manipulating data – logic and arithmetic
  - Making decisions
  - Branching and looping
  - Reading Principles of Operation
- Reading the HLASM Listing
- The PSW and an introduction to debugging assembler programs

## Where can I get help?

- z/OS V2R1 Elements and Features  
<http://www.ibm.com/systems/z/os/zos/bkserv/v2r1pdf/#IEA>
- HLASM Programmer's Guide (SC26-4941-06)  
<http://publibz.boulder.ibm.com/epubs/pdf/asmp1021.pdf>
- HLASM Language Reference (SC26-4940-06)  
<http://publibz.boulder.ibm.com/epubs/pdf/asmr1021.pdf>
- z/Architecture Principles of Operation  
<http://www.ibm.com/support/docview.wss?uid=isg2b9de5f05a9d57819852571c500428f9a>

# Putting it all together

## Putting it all together

- The following slides show a small demo program which determines whether or not an employee is eligible for a pay increase
- The slides are ordered as:
  - JCL to assemble, bind and run the program called SALARY
  - Assembler source code for the SALARY program
  - Job output from the program



## Putting it all together

# JCL to assemble, bind and run SALARY program

```
//XXXXXXX JOB NOTIFY=&SYSUID
//S1 EXEC PGM=ASMA90
// SET PRGM=SALARY
//* *****
// SET SRCLIB=&SYSUID..STAGE1.ASM0.ANSWERS
//SYSPRINT DD SYSOUT=*
//SYSLIN DD DSN=&&TEMP,DISP=(,PASS),SPACE=(CYL,1)
//SYSIN DD DSN=&SRCLIB(&PRGM),DISP=SHR
//SYSLIB DD DSN=SYS1.MACLIB,DISP=SHR
// DD DSN=PP.HLASM.ZOS201.SASMMAC1,DISP=SHR
// DD DSN=PP.HLASM.ZOS201.SASMMAC2,DISP=SHR
//S2 EXEC LKEDG,COND=(8,LE),
// PARM.LKED='XREF,LIST,NCAL,MAP'
//SYSLIN DD DISP=OLD,DSN=*.S1.SYSLIN
//SYSPRINT DD SYSOUT=*
//GO.SYSUDUMP DD SYSOUT=*
//* *****
//* *****
//* *****
```

## Putting it all together

### Assembler source code for SALARY program (1)

```

*****
* PUTTING IT ALL TOGETHER PROGRAM
*****
*
* The purpose of this small demo program is to demonstrate some small
* parts of assembler programming.
* The demo pretends that it has been passed an employee record via
* register 1.
* It will copy this record to some working storage and then proceed
* to determine whether or not the employee is eligible for a pay
* increase by comparing the employee's annual salary to the target
* salary.
*
* The employee's annual salary is calculated as:
*   12 x (MONTHLY_PAY-BENEFITS) + BONUS
*
*
*          ASMDREG ,
SALARY    CSECT
SALARY    AMODE 31
SALARY    RMODE 24
* USUAL PROGRAM SETUP
*          STM    14,12,12(13)
*          BALR   12,0
*          USING *,12
* Point register 1 at the first employee to process...
*          la     r1,employee_id_1

```

## Putting it all together

### Assembler source code for SALARY program (2)

```

*
* The data for the employee record is passed in from register 1.
* This small program will determine whether or not the employee
* is eligible for a pay rise or not.
*
        xc      employee(employee_rec_len),employee          Clear WS
        mvc     employee(employee_rec_len),0(r1)             Copy record to WS
*
* Output the name of the employee that is being processed
*
        xc      wto_text,wto_text                            Clear text buffer
        mvc     wto_text(1'process_text),process_text
        mvc     wto_text+1'process_text(1'employee_name),employee_name
* Calculate the amount of text to output
        lhi    r5,1'process_text+1'employee_name
        sth    r5,wto_buf                                    Store length in buffer
        la     r5,wto_buf                                    Load address of buffer
        WTO    TEXT=(5)                                     Output text
*
* Calculate employee's yearly pay as 12*(MONTHLY_PAY-BENEFITS)+BONUS
* We will use register 3 as a work register
*
        l      r3,employee_monthly_pay
        s      r3,employee_benefits                        MONTHLY-BENEFITS
        m      r2,=f'12'                                    Multiply by 12
        ah    r3,employee_bonus                            Add yearly bonus
        c      r3,target_salary                            Compare total with target
        bl    deserves_increase

```

## Putting it all together

### Assembler source code for SALARY program (3)

```

        WTO      'Employee has matched or exceeded target salary'
        b        resume_code
deserves_increase    dc      0h
        WTO      'Employee deserves a pay increase'
*
* Return to the caller of the program
*
resume_code          dc      0h
        LM      14,12,12(13)
*
        XR      15,15
        BR      14
* *****
* END OF PROGRAM - DATA FOLLOWS
* *****
*
WTO_BUF              DC H'0'
WTO_TEXT             DS CL256
PROCESS_TEXT        DC C'Processing employee '
* SALARY SCHEME DATA
TARGET_SALARY       DC F'24000'          TARGET SALARY FOR COMPANY
* EMPLOYEE RECORD STRUCTURE
EMPLOYEE            DC 0F
EMPLOYEE_NAME       DS CL40              EMPLOYEE'S NAME
EMPLOYEE_MONTHLY_PAY DS F                VALUE OF MONTHLY PAY
EMPLOYEE_BONUS      DS H                  YEARLY BONUS AMOUNT
EMPLOYEE_BENEFITS   DS F                  MONTHLY BENEFITS
EMPLOYEE_REC_LEN    EQU *-EMPLOYEE      SIZE OF EMPLOYEE RECORD

```

## Putting it all together

# Assembler source code for SALARY program (4)

```
* EMPLOYEE EXAMPLE DATA
EMPLOYEE_ID_1      DC      0F
                   DC CL40 'BOB SMITH'
                   DC F'2000'
                   DC H'1000'
                   DC F'50'

                LTORG ,
                END
```

## Putting it all together

### Job output for assembling, binding and running SALARY program (1)

```

1          J E S 2  J O B  L O G  --  S Y S T E M  M V 3 3  --  N O D E  W I N M V S 3 3
0
21.20.35 JOB47759 ---- TUESDAY, 29 APR 2014 ----
21.20.35 JOB47759 IRR010I USERID XXXXXXXX IS ASSIGNED TO THIS JOB.
21.20.35 JOB47759 IEF677I WARNING MESSAGE(S) FOR JOB XXXXXXXX ISSUED
21.20.35 JOB47759 ICH70001I XXXXXXXX LAST ACCESS AT 21:00:58 ON TUESDAY, APRIL 29, 2014
21.20.35 JOB47759 $HASP373 XXXXXXXX STARTED - INIT 1 - CLASS A - SYS MV33
21.20.35 JOB47759 IEF403I XXXXXXXX - STARTED
21.20.35 JOB47759 -
21.20.35 JOB47759 - -TIMINGS (MINS.)--
21.20.35 JOB47759 -JOBNAME STEPNAME PROCSTEP RC EXCP CPU SRB CLOCK SERV PG PAGE SWAP VIO SWAPS STEPNO
21.20.35 JOB47759 -XXXXXXX S1 00 95 .00 .00 .00 615 0 0 0 0 0 1
21.20.35 JOB47759 -XXXXXXX S2 LKED 00 28 .00 .00 .00 191 0 0 0 0 0 2
21.20.35 JOB47759 +Processing employee BOB SMITH
21.20.35 JOB47759 +Employee has matched or exceeded target salary
21.20.35 JOB47759 -XXXXXXX S2 GO 00 4 .00 .00 .00 51 0 0 0 0 0 3
21.20.35 JOB47759 IEF404I XXXXXXXX - ENDED
21.20.35 JOB47759 -XXXXXXX ENDED. NAME-
21.20.35 JOB47759 $HASP395 XXXXXXXX ENDED
TOTAL CPU TIME= .00 TOTAL ELAPSED TIME= .00
0----- JES2 JOB STATISTICS -----
- 29 APR 2014 JOB EXECUTION DATE
- 24 CARDS READ
- 669 SYSOUT PRINT RECORDS
- 0 SYSOUT PUNCH RECORDS
- 37 SYSOUT SPOOL KBYTES
- 0.00 MINUTES EXECUTION TIME
1 //XXXXXXX JOB NOTIFY=&SYSUID JOB47759
IEFC653I SUBSTITUTION JCL - NOTIFY=XXXXXXX
2 //S1 EXEC PGM=ASMA90
/* *****
/* CHANGE THE FOLLOWING LINE TO REFLECT THE PROGRAM NAME
/* *****

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (2)

```

3 // SET PRGNM=SALARY
  // * *****
  // * *****
  // * *****
4 // SET      SRCLIB=&SYSUID..STAGE1.ASM0.ANSWERS
  IEFC653I SUBSTITUTION JCL - SRCLIB=XXXXXXX.STAGE1.ASM0.ANSWERS
5 //SYSPRINT DD SYSOUT=*
6 //SYSLIN   DD DSN=&&TEMP, DISP=(, PASS), SPACE=(CYL,1)
7 //SYSIN    DD DSN=&SRCLIB(&PRGNM), DISP=SHR
  IEFC653I SUBSTITUTION JCL - DSN=XXXXXXX.STAGE1.ASM0.ANSWERS(SALARY), DISP=SHR
8 //SYSLIB   DD DSN=SYS1.MACLIB, DISP=SHR
9 //         DD DSN=PP.HLASM.ZOS201.SASMMAC1, DISP=SHR
10 //        DD DSN=PP.HLASM.ZOS201.SASMMAC2, DISP=SHR
11 //S2      EXEC LKEDG, COND=(8, LE),
  //         PARM.LKED='XREF, LIST, NCAL, MAP'
12 XXLKED   EXEC PGM=HEWLH096, PARM='SIZE=(384K, 96K), XREF, LIST, NCAL',      00033302
  XX        REGION=512K                                                    00066600
13 //SYSPRINT DD SYSOUT=*
  X/SYSPRINT DD  SYSOUT=A                                                    00100000
14 //SYSLIN   DD DISP=OLD, DSN=*.S1.SYSLIN
  X/SYSLIN   DD DDNAME=SYSIN                                                00150000
15 XXSYSLMOD DD DSN=&&GOSET(GO), SPACE=(1024, (50, 20, 1)), *00200000
  XX        UNIT=SYSDA, DISP=(MOD, PASS)                                    00250000
16 XXSYSUT1  DD UNIT=(SYSDA, SEP=(SYSLMOD, SYSLIN)), *00300000
  XX        SPACE=(1024, (200, 20))                                        00400000
17 XXGO      EXEC PGM=*.LKED.SYSLMOD, COND=(4, LT, LKED)                    00450000
18 //GO.SYSUDUMP DD SYSOUT=*
  // * *****
  // * *****
  // * *****
STMT NO. MESSAGE

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (3)

```

11 IEF001I PROCEDURE LKEDG WAS EXPANDED USING SYSTEM LIBRARY SYS1.PROCLIB
14 IEF648I INVALID DISP FIELD- PASS SUBSTITUTED
ICH70001I XXXXXXXX LAST ACCESS AT 21:00:58 ON TUESDAY, APRIL 29, 2014
IEF236I ALLOC. FOR XXXXXXXX S1
IEF237I JES2 ALLOCATED TO SYSPRINT
IGD101I SMS ALLOCATED TO DDNAME (SYSLIN )
        DSN (SYS14119.T212035.RA000.XXXXXXX.TEMP.H01 )
        STORCLAS (STANDARD) MGMTCLAS ( ) DATACLAS ( )
        VOL SER NOS= VIO
IGD103I SMS ALLOCATED TO DDNAME SYSIN
IEF237I ADA1 ALLOCATED TO SYSLIB
IEF237I ADA1 ALLOCATED TO
IEF237I ADA1 ALLOCATED TO
IEF142I XXXXXXXX S1 - STEP WAS EXECUTED - COND CODE 0000
IEF285I XXXXXXXX.XXXXXXX.JOB47759.D0000101.?          SYSOUT
IGD106I SYS14119.T212035.RA000.XXXXXXX.TEMP.H01      PASSED,      DDNAME=SYSLIN
IGD104I XXXXXXXX.STAGE1.ASM0.ANSWERS                 RETAINED,    DDNAME=SYSIN
IEF285I SYS1.MACLIB                                   KEPT
IEF285I VOL SER NOS= 33SY02.
IEF285I PP.HLASM.ZOS201.SASMMAC1                     KEPT
IEF285I VOL SER NOS= 33SY02.
IEF285I PP.HLASM.ZOS201.SASMMAC2                     KEPT
IEF285I VOL SER NOS= 33SY02.
IEF373I STEP/S1 /START 2014119.2120
IEF032I STEP/S1 /STOP 2014119.2120
CPU:      0 HR 00 MIN 00.01 SEC      SRB:      0 HR 00 MIN 00.00 SEC
VIRT:    220K SYS: 264K EXT: 65536K  SYS:    10404K
ATB- REAL:      36K SLOTS:      0K
        VIRT- ALLOC:      6M SHRD:      0M
IEF236I ALLOC. FOR XXXXXXXX LKED S2
IEF237I JES2 ALLOCATED TO SYSPRINT

```



## Putting it all together

### Job output for assembling, binding and running SALARY program (4)

```

IGD103I SMS ALLOCATED TO DDNAME SYSLIN
IGD101I SMS ALLOCATED TO DDNAME (SYSLMOD )
        DSN (SYS14119.T212035.RA000.XXXXXXX.GOSET.H01 )
        STORCLAS (STANDARD) MGMTCLAS ( ) DATACLAS ( )
        VOL SER NOS= VIO
IGD101I SMS ALLOCATED TO DDNAME (SYSUT1 )
        DSN (SYS14119.T212035.RA000.XXXXXXX.R0111176 )
        STORCLAS (STANDARD) MGMTCLAS ( ) DATACLAS ( )
        VOL SER NOS= VIO
IEF142I XXXXXXX LKED S2 - STEP WAS EXECUTED - COND CODE 0000
IEF285I XXXXXXX.XXXXXX.JOB47759.D0000102.?          SYSOUT
IGD106I SYS14119.T212035.RA000.XXXXXX.TEMP.H01      PASSED,      DDNAME=SYSLIN
IGD106I SYS14119.T212035.RA000.XXXXXX.GOSET.H01    PASSED,      DDNAME=SYSLMOD
IGD105I SYS14119.T212035.RA000.XXXXXX.R0111176    DELETED,     DDNAME=SYSUT1
IEF373I STEP/LKED      /START 2014119.2120
IEF032I STEP/LKED      /STOP 2014119.2120
        CPU:      0 HR 00 MIN 00.00 SEC   SRB:      0 HR 00 MIN 00.00 SEC
        VIRT:    100K SYS: 268K EXT:      1772K SYS: 10376K
        ATB- REAL:      0K SLOTS:      0K
        VIRT- ALLOC:      0M SHRD:      0M
IEF236I ALLOC. FOR XXXXXXX GO S2
IGD103I SMS ALLOCATED TO DDNAME PGM=*.DD
IEF237I JES2 ALLOCATED TO SYSUDUMP
Processing employee BOB SMITH
Employee has matched or exceeded target salary
IEF142I XXXXXXX GO S2 - STEP WAS EXECUTED - COND CODE 0000
IGD104I SYS14119.T212035.RA000.XXXXXX.GOSET.H01    RETAINED,    DDNAME=PGM=*.DD
IEF285I XXXXXXX.XXXXXX.JOB47759.D0000103.?          SYSOUT
IEF373I STEP/GO      /START 2014119.2120
IEF032I STEP/GO      /STOP 2014119.2120
        CPU:      0 HR 00 MIN 00.00 SEC   SRB:      0 HR 00 MIN 00.00 SEC

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (5)

```

VIRT:      8K  SYS:   252K  EXT:      0K  SYS:   10440K
ATB- REAL:                0K  SLOTS:                0K
VIRT- ALLOC:      0M  SHRD:      0M
IGD105I SYS14119.T212035.RA000.XXXXXXX.TEMP.H01      DELETED,   DDNAME=SYSLIN
IGD105I SYS14119.T212035.RA000.XXXXXXX.GOSET.H01      DELETED,   DDNAME=SYSLMOD
IEF375I JOB/XXXXXXX /START 2014119.2120
IEF033I JOB/XXXXXXX /STOP 2014119.2120
CPU:      0 HR  00 MIN  00.01 SEC      SRB:      0 HR  00 MIN  00.00 SEC

```

```

1 High Level Assembler Option Summary
-

```

```

(PTF UI11676) Page 1
HLASM R6.0 2014/04/29 21.20

```

```

0 No Overriding ASMAOPT Parameters
No Overriding Parameters
No Process Statements

```

```
Options for this Assembly
```

```

0 NOADATA
ALIGN
NOASA
NOBATCH
CODEPAGE(047C)
NOCOMPAT
NOBCS
NODECK
DXREF
ESD
NOEXIT
FLAG(0,ALIGN,NOCONT,EXLITW,NOIMPLEN,NOPAGE0,PUSH,RECORD,NOSUBSTR,USING0)
NOFOLD
NOGOFF
NOINFO

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (6)

```

LANGUAGE (EN)
NOLIBMAC
LINECOUNT (60)
LIST (121)
MACHINE (, NOLIST)
MXREF (SOURCE)
OBJECT
OPTABLE (UNI, NOLIST)
NOPCONTROL
NOPESTOP
NOPROFILE
NORA2
NORENT
RLD
RXREF
SECTALGN (8)
SIZE (MAX)
NOSUPRWARN
SYSPARM ()
NOTERM
NOTEST
THREAD
NOTRANSLATE
TYPECHECK (MAGNITUDE, REGISTER)
USING (NOLIMIT, MAP, NOWARN)
NOWORKFILE
XREF (SHORT, UNREFS)

```

No Overriding DD Names

```

1
-Symbol  Type  Id      Address  Length  External Symbol Dictionary
Owner Id  Flags  Alias-of

```

Page 2  
HLASM R6.0 2014/04/29 21.20

## Putting it all together

### Job output for assembling, binding and running SALARY program (7)

```

0          PC 00000001 00000000 00000000          00
SALARY    SD 00000002 00000000 000002C4          02
1
1  Active Usings: None
0  Loc  Object Code  Addr1 Addr2  Stmt  Source Statement                                     HLASM R6.0  2014/04/29 21.20
0
1 *****
2 * PUTTING IT ALL TOGETHER PROGRAM
3 *****
4 *
5 * The purpose of this small demo program is to demonstrate some small
6 * parts of assembler programming.
7 * The demo pretends that it has been passed an employee record via
8 * register 1.
9 * It will copy this record to some working storage and then proceed
10 * to determine whether or not the employee is eligible for a pay
11 * increase by comparing the employee's annual salary to the target
12 * salary.
13 *
14 * The employee's annual salary is calculated as:
15 *     12 x (MONTHLY_PAY-BENEFITS) + BONUS
16 *
17 *
18          ASMDREG ,
19+         PUSH PRINT                                     01-ASMDR
119+        POP PRINT                                     01-ASMDR
000000    000000 002C4 120 SALARY CSECT
121 SALARY AMODE 31
122 SALARY RMODE 24
123 * USUAL PROGRAM SETUP
000000 90EC D00C    0000C 124          STM 14,12,12(13)
000004 05C0       125          BALR 12,0

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (8)

```

R:C 00006      126      USING *,12
000006 4110 C286      0028C 127 * Point register 1 at the first employee to process...
128      la    r1,employee_id_1
129 *
130 * The data for the employee record is passed in from register 1.
131 * This small program will determine whether or not the employee
132 * is eligible for a pay rise or not.
133 *
00000A D733 C252 C252 00258 00258 134      xc    employee(employee_rec_len),employee      Clear WS
000010 D233 C252 1000 00258 00000 135      mvc   employee(employee_rec_len),0(r1)      Copy record to WS
136 *
137 * Output the name of the employee that is being processed
138 *
000016 D7FF C13A C13A 00140 00140 139      xc    wto_text,wto_text      Clear text buffer
00001C D213 C13A C23A 00140 00240 140      mvc   wto_text(1'process_text),process_text
000022 D227 C14E C252 00154 00258 141      mvc   wto_text+1'process_text(1'employee_name),employee_name
142 * Calculate the amount of text to output
000028 A758 003C      0003C 143      lhi   r5,1'process_text+1'employee_name
00002C 4050 C138      0013E 144      sth   r5,wto_buf      Store length in buffer
000030 4150 C138      0013E 145      la    r5,wto_buf      Load address of bufffer
146      WTO   TEXT=(5)      Output text
000034      148+     CNOP  0,4
000034 A715 003A      000A8 149+     BRAS  1,IHB0002A      BRANCH AROUND MESSAGE      @LCC 01-WTO
000038 0008      150+     DC    AL2(8)      TEXT LENGTH      @YA17152 01-WTO
00003A 0010      151+     DC    B'0000000000010000'      MCSFLAGS      01-WTO
00003C 00000000      152+     DC    AL4(0)      MESSAGE TEXT ADDRESS      @L5A 01-WTO
000040 02      153+     DC    AL1(2)      VERSION LEVEL      @PJC 01-WTO
000041 00      154+     DC    B'000000000'      MISCELLANEOUS FLAGS      @L2A 01-WTO
000042 00      155+     DC    AL1(0)      REPLY LENGTH      @L2A 01-WTO
1
Active Usings: SALARY+X'6',R12

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (9)

```

0 Loc Object Code Addr1 Addr2 Stmt Source Statement HLASM R6.0 2014/04/29 21.20
0000043 68 156+ DC AL1(104) LENGTH OF WPX @L5C 01-WTO
000044 0080 157+ DC B'0000000010000000' EXTENDED MCS FLAGS @L2A 01-WTO
000046 0000 158+ DC AL2(0) RESERVED @L2A 01-WTO
000048 00000000 159+ DC AL4(0) REPLY BUFFER ADDRESS @P7C 01-WTO
00004C 00000000 160+ DC AL4(0) REPLY ECB ADDRESS @P7C 01-WTO
000050 00000000 161+ DC AL4(0) CONNECT ID @01C 01-WTO
000054 0000 162+ DC B'0000000000000000' DESCRIPTOR CODES @L2A 01-WTO
000056 0000 163+ DC AL2(0) RESERVED @L2A 01-WTO
000058 0000000000000000 164+ DC XL16'00000000000000000000000000000000' X01-WTO
000060 0000000000000000 + EXTENDED ROUTING CODES @L2A
000068 0000 165+ DC B'0000000000000000' MESSAGE TYPE @L2A 01-WTO
00006A 0000 166+ DC AL2(0) MESSAGE'S PRIORITY @L2A 01-WTO
00006C 4040404040404040 167+ DC CL8' ' JOB ID @L2A 01-WTO
000074 4040404040404040 168+ DC CL8' ' JOB NAME @L2A 01-WTO
00007C 4040404040404040 169+ DC CL8' ' RETRIEVAL KEY @L2A 01-WTO
000084 00000000 170+ DC AL4(0) TOKEN FOR DOM @P1C 01-WTO
000088 00000000 171+ DC AL4(0) CONSOLE ID @P1C 01-WTO
00008C 4040404040404040 172+ DC CL8' ' SYSTEM NAME @L2A 01-WTO
000094 4040404040404040 173+ DC CL8' ' CONSOLE NAME @L3A 01-WTO
00009C 00000000 174+ DC AL4(0) REPLY CONSOLE NAME/ID ADDR @L3A 01-WTO
0000A0 00000000 175+ DC AL4(0) CART ADDRESS @L4C 01-WTO
0000A4 00000000 176+ DC AL4(0) WSPARM ADDRESS @L6C 01-WTO
0000A8 177+IHB0002A DS 0H 01-WTO
0000A8 18E1 178+ LR 14,1 FIRST BYTE OF PARM LIST @L2A 01-WTO
0000AA 1BFF 179+ SR 15,15 CLEAR REGISTER 15 @L2A 01-WTO
0000AC 4AF1 0000 00000 180+ AH 15,0(1,0) ADD LENGTH OF TEXT + 4 @L2A 01-WTO
0000B0 1AEF 181+ AR 14,15 FIRST BYTE AFTER TEXT @L2A 01-WTO
0000B2 5050 1004 00004 182+ ST 5,4(0,1) STORE TEXT ADDR INTO PLIST @L5A 01-WTO
0000B6 0A23 183+ SVC 35 ISSUE SVC 35 @L6A 01-WTO
184 *

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (10)

```

185 * Calculate employee's yearly pay as 12*(MONTHLY_PAY-BENEFITS)+BONUS
186 * We will use register 3 as a work register
187 *
0000B8 5830 C27A          00280 188      l      r3,employee_monthly_pay
0000BC 5B30 C282          00288 189      s      r3,employee_benefits      MONTHLY-BENEFITS
0000C0 5C20 C2BA          002C0 190      m      r2,=f'12'                  Multiply by 12
0000C4 4A30 C27E          00284 191      ah     r3,employee_bonus        Add yearly bonus
0000C8 5930 C24E          00254 192      c      r3,target_salary           Compare total with target
0000CC 4740 C106          0010C 193      bl     deserves_increase
194      WTO      'Employee has matched or exceeded target salary'
0000D0          196+    CNOP  0,4                                01-WTO
0000D0 A715 001B          00106 197+    BRAS  1,IHB0004A                BRANCH AROUND MESSAGE      @LCC 01-WTO
0000D4 0032          198+    DC    AL2(50)                   TEXT LENGTH                @YA17152 01-WTO
0000D6 0000          199+    DC    B'0000000000000000'      MCSFLAGS                   01-WTO
0000D8 C594979396A88585 200+    DC    C'Employee has matched or exceeded target salary' X01-WTO
0000E0 408881A2409481A3 +      MESSAGE TEXT                @L6C
000106          201+IHB0004A DS  0H                                01-WTO
000106 0A23          202+    SVC  35                        ISSUE SVC 35                @L6A 01-WTO
000108 47F0 C130          00136 203      b      resume_code
00010C          204 deserves_increase dc 0h
205      WTO      'Employee deserves a pay increase'
00010C          207+    CNOP  0,4                                01-WTO
00010C A715 0014          00134 208+    BRAS  1,IHB0006A                BRANCH AROUND MESSAGE      @LCC 01-WTO
000110 0024          209+    DC    AL2(36)                   TEXT LENGTH                @YA17152 01-WTO
000112 0000          210+    DC    B'0000000000000000'      MCSFLAGS                   01-WTO
1
Active Usings: SALARY+X'6',R12
0 Loc Object Code Addr1 Addr2 Stmt Source Statement HLASM R6.0 2014/04/29 21.20
0000114 C594979396A88585 211+ DC C'Employee deserves a pay increase' X01-WTO
00011C 408485A28599A585 + MESSAGE TEXT @L6C
000134 212+IHB0006A DS 0H 01-WTO

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (11)

```

000134 0A23          213+      SVC   35          ISSUE SVC 35          @L6A 01-WTO
                   214 *
                   215 * Return to the caller of the program
                   216 *
000136          217 resume_code      dc      0h
000136 98EC D00C      0000C      218          LM      14,12,12(13)
                   219 *
00013A 17FF          220          XR      15,15
00013C 07FE          221          BR      14
                   222 * *****
                   223 * END OF PROGRAM - DATA FOLLOWS
                   224 * *****
                   225 *
00013E 0000          226 WTO_BUF          DC H'0'
000140          227 WTO_TEXT         DS CL256
000240 D799968385A2A289 228 PROCESS_TEXT    DC C'Processing employee '
                   229 * SALARY SCHEME DATA
000254 00005DC0      230 TARGET_SALARY    DC F'24000'          TARGET SALARY FOR COMPANY
                   231 * EMPLOYEE RECORD STRUCTURE
000258          232 EMPLOYEE         DC 0F
000258          233 EMPLOYEE_NAME    DS CL40          EMPLOYEE'S NAME
000280          234 EMPLOYEE_MONTHLY_PAY DS F          VALUE OF MONTHLY PAY
000284          235 EMPLOYEE_BONUS   DS H          YEARLY BONUS AMOUNT
000288          236 EMPLOYEE_BENEFITS DS F          MONTHLY BENEFITS
                   237 EMPLOYEE_REC_LEN  EQU *-EMPLOYEE    SIZE OF EMPLOYEE RECORD
                   238 * EMPLOYEE EXAMPLE DATA
00028C          239 EMPLOYEE_ID_1     DC 0F
00028C C2D6C240E2D4C9E3 240          DC CL40 'BOB SMITH'
0002B4 000007D0      241          DC F'2000'
0002B8 03E8          242          DC H'1000'
0002BA 0000

```



# Putting it all together

## Job output for assembling, binding and running SALARY program (12)

```

0002BC 00000032          243          DC F'50'
0002C0          244          LTORG ,
0002C0 0000000C          245          =f'12'
          246          END
    
```

1 Ordinary Symbol and Literal Cross Reference										Page	6
-Symbol	Length	Value	Id	R	Type	Asm	Program	Defn	References	HLASM R6.0	2014/04/29 21.20
Odeserves_increase											
	2	0000010C	00000002		H	H		204	193B		
EMPLOYEE	4	00000258	00000002		F	F		232	134M 134 135M 237		
EMPLOYEE_BENEFITS											
	4	00000288	00000002		F	F		236	189		
EMPLOYEE_BONUS											
	2	00000284	00000002		H	H		235	191		
EMPLOYEE_ID_1											
	4	0000028C	00000002		F	F		239	128		
EMPLOYEE_MONTHLY_PAY											
	4	00000280	00000002		F	F		234	188		
EMPLOYEE_NAME											
	40	00000258	00000002		C	C		233	141 141 143		
EMPLOYEE_REC_LEN											
	1	00000034	00000002	A	U			237	134 135		
IHB0002A	2	000000A8	00000002		H	H		177	149B		
IHB0004A	2	00000106	00000002		H	H		201	197B		
IHB0006A	2	00000134	00000002		H	H		212	208B		
PROCESS_TEXT											
	20	00000240	00000002		C	C		228	140 140 141M 143		
resume_code											
	2	00000136	00000002		H	H		217	203B		
R1	1	00000001	00000001	A	U			28	128M 135		
R2	1	00000002	00000001	A	U			29	190M		
R3	1	00000003	00000001	A	U			30	188M 189M 191M 192		

# Putting it all together

## Job output for assembling, binding and running SALARY program (13)

```

0002BC 00000032          243          DC F'50'
0002C0          244          LTORG ,
0002C0 0000000C          245          =f'12'
          246          END
    
```

1 Ordinary Symbol and Literal Cross Reference										Page	6
-Symbol	Length	Value	Id	R	Type	Asm	Program	Defn	References	HLASM R6.0	2014/04/29 21.20
Odeserves_increase											
	2	0000010C	00000002		H	H		204	193B		
EMPLOYEE	4	00000258	00000002		F	F		232	134M 134	135M	237
EMPLOYEE_BENEFITS											
	4	00000288	00000002		F	F		236	189		
EMPLOYEE_BONUS											
	2	00000284	00000002		H	H		235	191		
EMPLOYEE_ID_1											
	4	0000028C	00000002		F	F		239	128		
EMPLOYEE_MONTHLY_PAY											
	4	00000280	00000002		F	F		234	188		
EMPLOYEE_NAME											
	40	00000258	00000002		C	C		233	141 141	143	
EMPLOYEE_REC_LEN											
	1	00000034	00000002	A	U			237	134 135		
IHB0002A	2	000000A8	00000002		H	H		177	149B		
IHB0004A	2	00000106	00000002		H	H		201	197B		
IHB0006A	2	00000134	00000002		H	H		212	208B		
PROCESS_TEXT											
	20	00000240	00000002		C	C		228	140 140	141M	143
resume_code											
	2	00000136	00000002		H	H		217	203B		
R1	1	00000001	00000001	A	U			28	128M 135		
R2	1	00000002	00000001	A	U			29	190M		
R3	1	00000003	00000001	A	U			30	188M 189M	191M	192

## Putting it all together

### Job output for assembling, binding and running SALARY program (14)

```

R5          1 00000005 00000001 A  U          32 143M 144 145M
SALARY      1 00000000 00000002      J          120 121 122
TARGET_SALARY
          4 00000254 00000002      F  F          230 192
WTO_BUF     2 0000013E 00000002      H  H          226 144M 145
WTO_TEXT    256 00000140 00000002      C  C          227 139M 139 140M 141M
=f'12'      4 000002C0 00000002      F          245 190

```

```
1 Unreferenced Symbols Defined in CSECTs
```

```

- Defn Symbol
0 103 AR0
  104 AR1
  113 AR10
  114 AR11
  115 AR12
  116 AR13
  117 AR14
  118 AR15
  105 AR2
  106 AR3
  107 AR4
  108 AR5
  109 AR6
  110 AR7
  111 AR8
  112 AR9
   84 CR0
   85 CR1
   94 CR10
   95 CR11
   96 CR12
   97 CR13

```

```

Page 7
HLASM R6.0 2014/04/29 21.20

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (15)

```
98 CR14
99 CR15
86 CR2
87 CR3
88 CR4
89 CR5
90 CR6
91 CR7
92 CR8
93 CR9
46 FR0
47 FR1
56 FR10
57 FR11
58 FR12
59 FR13
60 FR14
61 FR15
48 FR2
49 FR3
50 FR4
51 FR5
52 FR6
53 FR7
54 FR8
55 FR9
27 R0
37 R10
38 R11
39 R12
40 R13
```

# Putting it all together

## Job output for assembling, binding and running SALARY program (16)

```

    41 R14
    42 R15
1
- Defn Symbol
0 31 R4
  33 R6
  34 R7
  35 R8
  36 R9
  65 VR0
  66 VR1
  75 VR10
  76 VR11
  77 VR12
  78 VR13
  79 VR14
  80 VR15
  67 VR2
  68 VR3
  69 VR4
  70 VR5
  71 VR6
  72 VR7
  73 VR8
  74 VR9

1
- Con Source
0 L1 SYS1.MACLIB
  L3 PP.HLASM.ZOS201.SASMMAC2

1
-

```

Unreferenced Symbols Defined in CSECTs

Page 8  
HLASM R6.0 2014/04/29 21.20

```

1
- Con Source
0 L1 SYS1.MACLIB
  L3 PP.HLASM.ZOS201.SASMMAC2

1
-

```

Macro and Copy Code Source Summary

Volume	Members	Page
33SY02	SYSSTATE WTO	9
33SY02	ASMDREG	10
	Using Map	10

Page 9  
HLASM R6.0 2014/04/29 21.20

Page 10  
HLASM R6.0 2014/04/29 21.20

# Putting it all together

## Job output for assembling, binding and running SALARY program (17)

```

    41 R14
    42 R15
1
- Defn Symbol
0 31 R4
  33 R6
  34 R7
  35 R8
  36 R9
  65 VR0
  66 VR1
  75 VR10
  76 VR11
  77 VR12
  78 VR13
  79 VR14
  80 VR15
  67 VR2
  68 VR3
  69 VR4
  70 VR5
  71 VR6
  72 VR7
  73 VR8
  74 VR9

1
- Con Source
0 L1 SYS1.MACLIB
  L3 PP.HLASM.ZOS201.SASMMAC2

1
-

```

Unreferenced Symbols Defined in CSECTs

Page 8  
HLASM R6.0 2014/04/29 21.20

```

1
- Con Source
0 L1 SYS1.MACLIB
  L3 PP.HLASM.ZOS201.SASMMAC2

1
-

```

Macro and Copy Code Source Summary

Con	Source	Volume	Members	Page
0	L1 SYS1.MACLIB	33SY02	SYSSTATE WTO	9
	L3 PP.HLASM.ZOS201.SASMMAC2	33SY02	ASMDREG	10

Using Map  
HLASM R6.0 2014/04/29 21.20

## Putting it all together

### Job output for assembling, binding and running SALARY program (18)

```

      Stmt  -----Location----- Action -----Using----- Reg Max      Last Label and Using Text
              Count      Id      Type      Value      Range      Id      Disp      Stmt
0   126   00000006   00000002 USING   ORDINARY   00000006 00001000 00000002 12 002BA   203 *,12
1
- Register References (M=modified, B=branch, U=USING, D=DROP, N=index)                                     Page 11
0   0(0)    124    218M
   1(1)    124    128M  135   149M  178   180N  182   197M  208M  218M
   2(2)    124    190M  218M
   3(3)    124    188M  189M  190M  191M  192   218M
   4(4)    124    218M
   5(5)    124    143M  144   145M  182   218M
   6(6)    124    218M
   7(7)    124    218M
   8(8)    124    218M
   9(9)    124    218M
  10(A)    124    218M
  11(B)    124    218M
  12(C)    124    125M  126U  218M
  13(D)    124    218
  14(E)    124    178M  181M  218M  221B
  15(F)    124    179M  179   180M  181   218M  220M  220
1
- Diagnostic Cross Reference and Assembler Summary                                     Page 12
0
0   No Statements Flagged in this Assembly
HIGH LEVEL ASSEMBLER, 5696-234, RELEASE 6.0, PTF UI11676
OSYSTEM: z/OS 02.01.00      JOBNAME: XXXXXXXX      STEPNAME: S1      PROCSTEP: (NOPROC)
0Data Sets Allocated for this Assembly
Con DDname      Data Set Name      Volume  Member
P1 SYSIN        XXXXXXXX.STAGE1.ASMO.ANSWERS  33P002  SALARY
L1 SYSLIB       SYS1.MACLIB        33SY02
L2              PP.HLASM.ZOS201.SASMMAC1      33SY02

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (19)

```

L3          PP.HLASM.ZOS201.SASMMAC2          33SY02
SYSLIN     SYS14119.T212035.RA000.XXXXXXX.TEMP.H01
SYSPRINT   XXXXXXXX.XXXXXXX.JOB47759.D0000101.?

64840K allocated to Buffer Pool      Storage required      384K
 91 Primary Input Records Read      3707 Library Records Read      0 Work File Reads
 0 ASMAOPT Records Read             384 Primary Print Records Written 0 Work File Writes
10 Object Records Written           0 ADATA Records Written

0Assembly Start Time: 21.20.35 Stop Time: 21.20.35 Processor Time: 00.00.00.0068
Return Code 000

1z/OS V2 R1 BINDER      21:20:35 TUESDAY APRIL 29, 2014
BATCH EMULATOR JOB(XXXXXXX ) STEP(S2 ) PGM= HEWLH096 PROCEDURE(LKED )
IEW2278I B352 INVOCATION PARAMETERS - XREF,LIST,NCAL,MAP
IEW2650I 5102 MODULE ENTRY NOT PROVIDED. ENTRY DEFAULTS TO SECTION SALARY.

1          *** M O D U L E M A P ***

-----
CLASS B_TEXT          LENGTH =      2C4  ATTRIBUTES = CAT,   LOAD, RMODE= 24
                     OFFSET =      0  IN SEGMENT 001   ALIGN = DBLWORD
-----

SECTION  CLASS
OFFSET  OFFSET  NAME          TYPE      LENGTH  DDNAME  SEQ  MEMBER
-----
          0  SALARY          CSECT     2C4    SYSLIN   01  **NULL**
1      ***  DATA SET SUMMARY  ***

DDNAME  CONCAT  FILE IDENTIFICATION

```



## Putting it all together

### Job output for assembling, binding and running SALARY program (20)

```
SYSLIN      01      SYS14119.T212035.RA000.XXXXXXX.TEMP.H01
```

```
*** E N D   O F   M O D U L E   M A P ***
```

```
1          C R O S S - R E F E R E N C E   T A B L E
```

```
TEXT CLASS = B_TEXT
```

```
----- R E F E R E N C E ----- T A R G E T -----
CLASS          ELEMENT |          ELEMENT |
OFFSET SECT/PART (ABBREV)  OFFSET  TYPE | SYMBOL (ABBREV)  SECTION (ABBREV)  OFFSET CLASS NAME |
*** NO ADDRESS CONSTANTS FOR THIS CLASS ***
*** E N D   O F   C R O S S   R E F E R E N C E ***
```

```
*** O P E R A T I O N   S U M M A R Y   R E P O R T ***
```

```
1PROCESSING OPTIONS:
```

```
ALIASES          NO
ALIGN2           NO
AMODE            UNSPECIFIED
CALL             NO
CASE             UPPER
COMPAT           UNSPECIFIED
COMPRESS         AUTO
DCBS             NO
```

## Putting it all together

### Job output for assembling, binding and running SALARY program (21)

```
DYNAM          NO
EXTATTR        UNSPECIFIED
EXITS:         NONE
FILL           NONE
GID            UNSPECIFIED
HOBSET         NO
INFO           NO
LET            04
LINECT        060
LIST           SUMMARY
LISTPRIV       NO
LONGPARM       NO
MAP            YES
MAXBLK         032760
MODMAP         NO
MSGLEVEL       00
OVLY           NO
PRINT          YES
RES            NO
REUSABILITY    UNSPECIFIED
RMODE          UNSPECIFIED
SIGN           NO
STORENX        NOREPLACE
STRIPCL        NO
STRIPSEC       NO
SYMTRACE
TERM           NO
TRAP           ON
UID            UNSPECIFIED
UPCASE         NO
WKSPACE        000000K,000000K
```

## Putting it all together

### Job output for assembling, binding and running SALARY program (22)

```
XCAL          NO
XREF          YES
***END OF OPTIONS***
```

#### 1SAVE OPERATION SUMMARY:

```
MEMBER NAME      GO
LOAD LIBRARY     SYS14119.T212035.RA000.XXXXXXX.GOSET.H01
PROGRAM TYPE     LOAD MODULE
VOLUME SERIAL
MAX BLOCK        32760
DISPOSITION      ADDED NEW
TIME OF SAVE     21.20.35  APR 29, 2014
```

#### 1SAVE MODULE ATTRIBUTES:

```
AC              000
AMODE           31
COMPRESSION     NONE
DC              NO
EDITABLE        YES
EXCEEDS 16MB   NO
EXECUTABLE      YES
LONGPARM        NO
MIGRATABLE     YES
OL              NO
OVLY           NO
```

## Putting it all together

### Job output for assembling, binding and running SALARY program (23)

```

PACK, PRIME          NO, NO
PAGE ALIGN          NO
REFR                NO
RENT                NO
REUS                NO
RMODE                24
SCTR                NO
SIGN                NO
SSI
SYM GENERATED      NO
TEST                NO
XPLINK              NO
MODULE SIZE (HEX)   000002C8

```

1 ENTRY POINT AND ALIAS SUMMARY:

```

NAME:                ENTRY TYPE AMODE C_OFFSET CLASS NAME          STATUS
SALARY                MAIN_EP      31 00000000 B_TEXT
*** E N D   O F   O P E R A T I O N   S U M M A R Y   R E P O R T ***

```

```

1z/OS V2 R1 BINDER      21:20:35 TUESDAY APRIL 29, 2014
BATCH EMULATOR JOB(XXXXXXX ) STEP(S2 ) PGM= HEWLH096 PROCEDURE(LKED )
IEW2008I 0F03 PROCESSING COMPLETED. RETURN CODE = 0.

```

## Putting it all together

### Job output for assembling, binding and running SALARY program (24)

```
1-----  
MESSAGE SUMMARY REPORT  
-----  
TERMINAL MESSAGES      (SEVERITY = 16)  
NONE  
  
SEVERE MESSAGES        (SEVERITY = 12)  
NONE  
  
ERROR MESSAGES         (SEVERITY = 08)  
NONE  
  
WARNING MESSAGES       (SEVERITY = 04)  
NONE  
  
INFORMATIONAL MESSAGES (SEVERITY = 00)  
2008  2278  2650  
  
**** END OF MESSAGE SUMMARY REPORT ****
```