



I/O Synergy – The Whole is Greater than the Parts (I/O Performance)

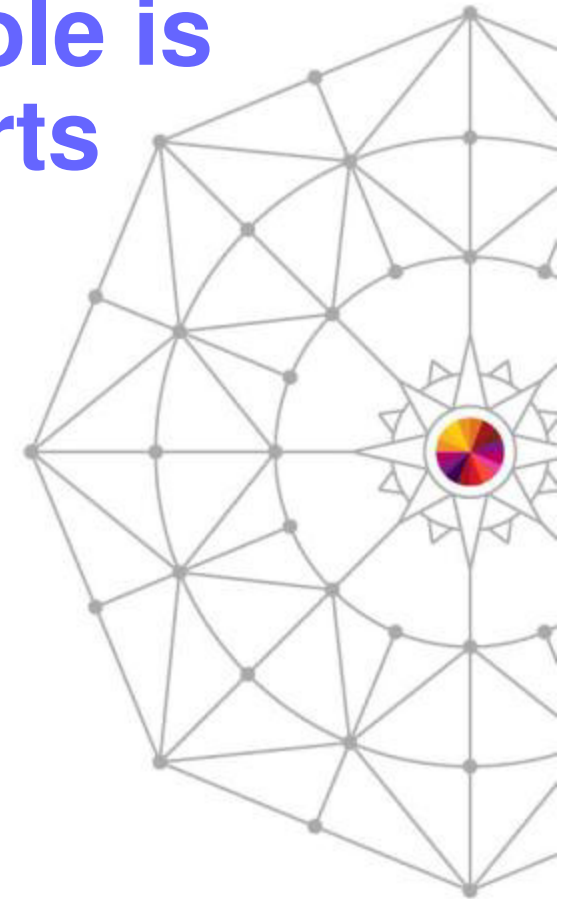
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Session 15081



Introduction

- Synergy focus areas
- I/O Performance Related Topics
- For each topic:
 - Overview
 - Pointers/Tips
 - Triva/Interesting Facts

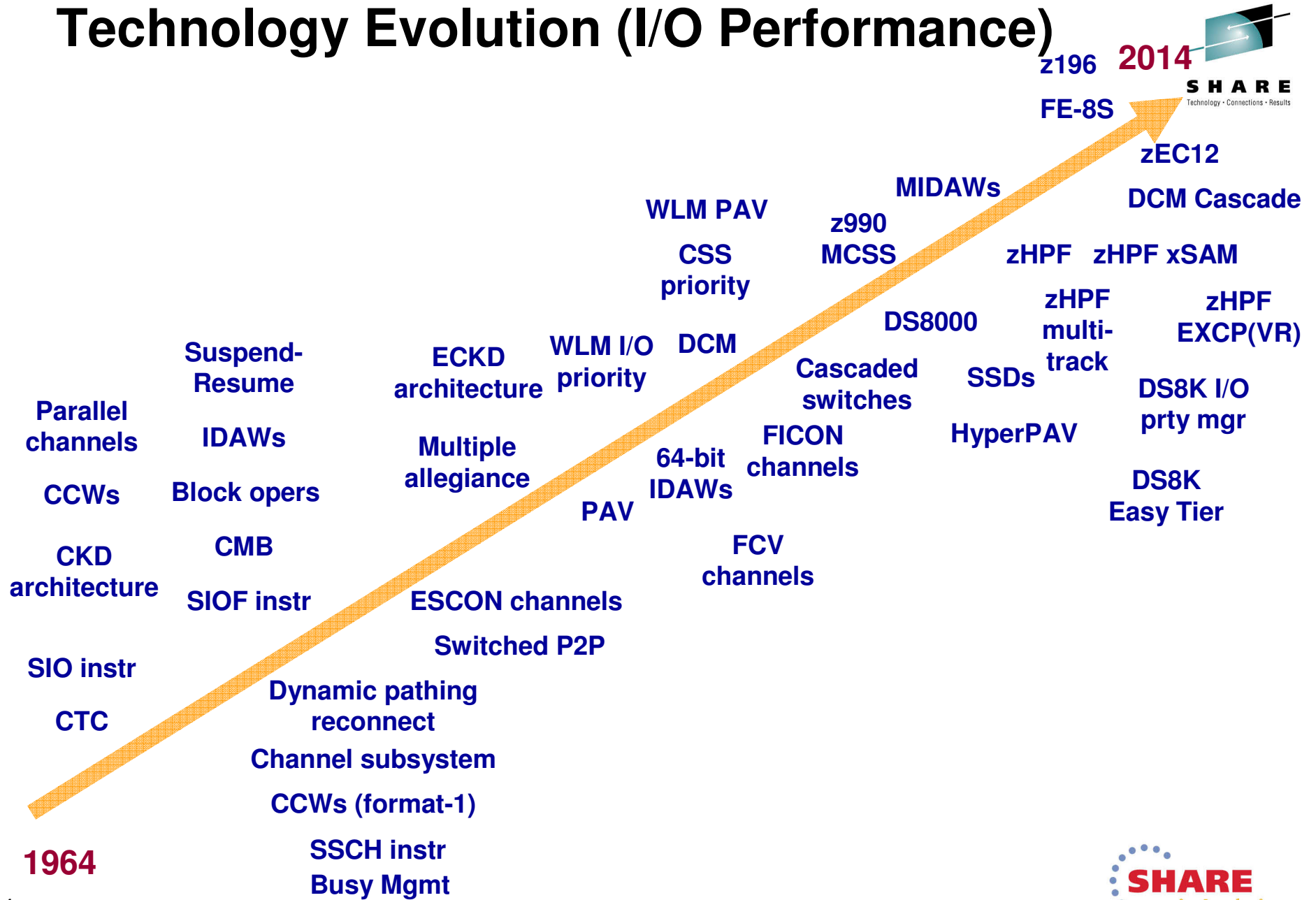
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I/O and Storage Constraints



Focus Area	Description	Client Impact
I/O Rates and Bandwidth	Ability to provide a balanced system for MIPs and I/O requests within the existing footprint and power envelope	Inability to achieve high CPU utilization because of I/O delays and concurrency limits. Inability to consolidate more physical footprints. Inability to grow the workloads on z/OS.
Latency	Fit within the existing batch window, ability to handle workload spikes	Bottlenecks will force clients to split workloads and spend resources managing around the I/O limits
Scalability	Number of devices, channels, devices per channel, aliases, I/O slots for accelerators, addressable storage capacity	Inability to add additional work and data to the existing z/OS infrastructure
Resiliency	Ability to recover from a channel path, device, storage controller, or site failure with minimal impact to the system and applications.	Application delays and outages
Complexity	I/O configuration planning, tuning, problem determination, capacity planning	Added cost and risk. Frequency of manual effort increases chance for errors, degrades resilience, and adds cost.

Technology Evolution (I/O Performance)

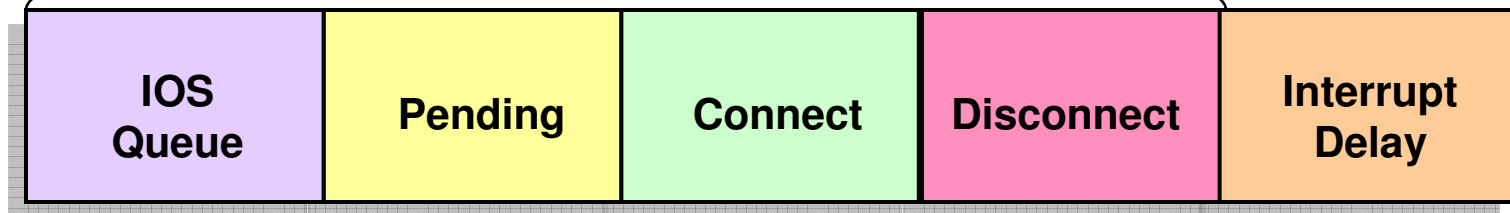


Technologies to Improve I/O Service Time



RMF Response Time

Response Time Components



Contributors

- High tran rates
- Long busy conditions
- Other I/O delays
- Device reserved
- CU recovery
- Fabric congestion
- Distance
- High port / HA utilization
- Extent conflict
- FICON channel multi-plexing
- Out of buffer credits
- Cache miss
- NVS destage
- PPRC delay
- Write pacing (XRC)
- LPAR weights
- Capping
- CPENABLE

Technologies

- HyperPAV
- WLM I/O priority
- zHPF
- CSS priority
- FICON DCM
- FICON express 8S
- zEC12 channel selection
- zHPF
- MIDAWs
- FICON express 8S
- CU cache alg
- Easy Tier
- SSDs
- PPRC FCP pre-deposit writes
- I/O prty mgr
- Hiper-dispatch



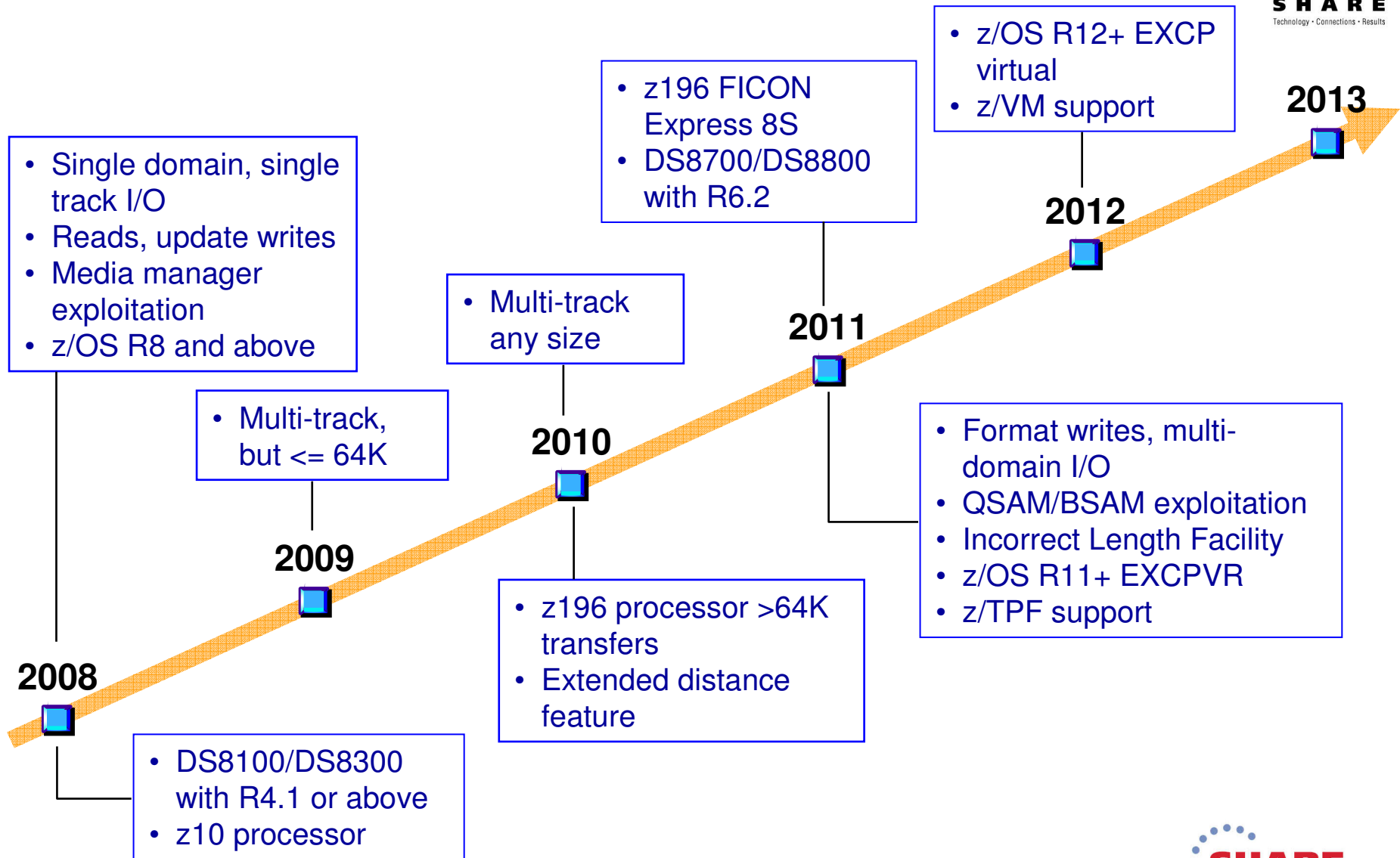
High Performance FICON for System z (zHPF)

zHPF is a critical element of the System z I/O strategy

How does zHPF improve I/O performance?

- Leverages optimizations made in the HBA for FCP
- Uses a protocol that has less chit-chat between the channel and device
- Packages entire command stream into one or two sequences
- Packages data more efficiently into frames (less buffer credits used)
- H/W data router on FICON express 8S provides additional acceleration

zHPF Evolution

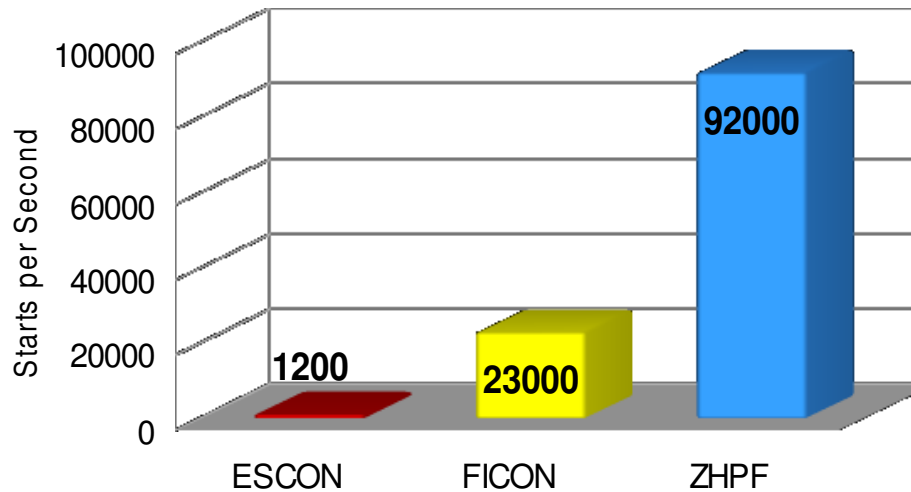


zHPF/FICON/ESCON Comparison



Start Rate Comparison

4K Block Size, 100% Utilization

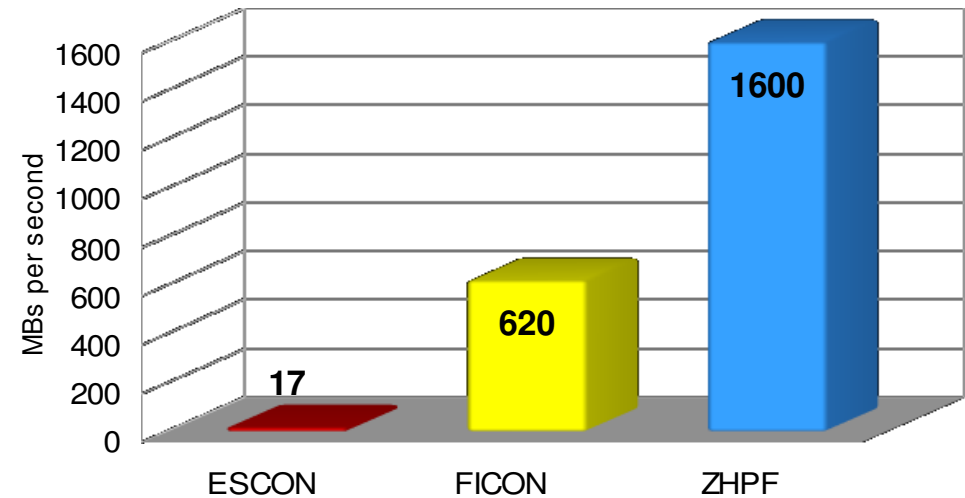


4X the start rate of FICON

2.5X the throughput of FICON

Throughput Comparison

Full duplex - large read/write mix

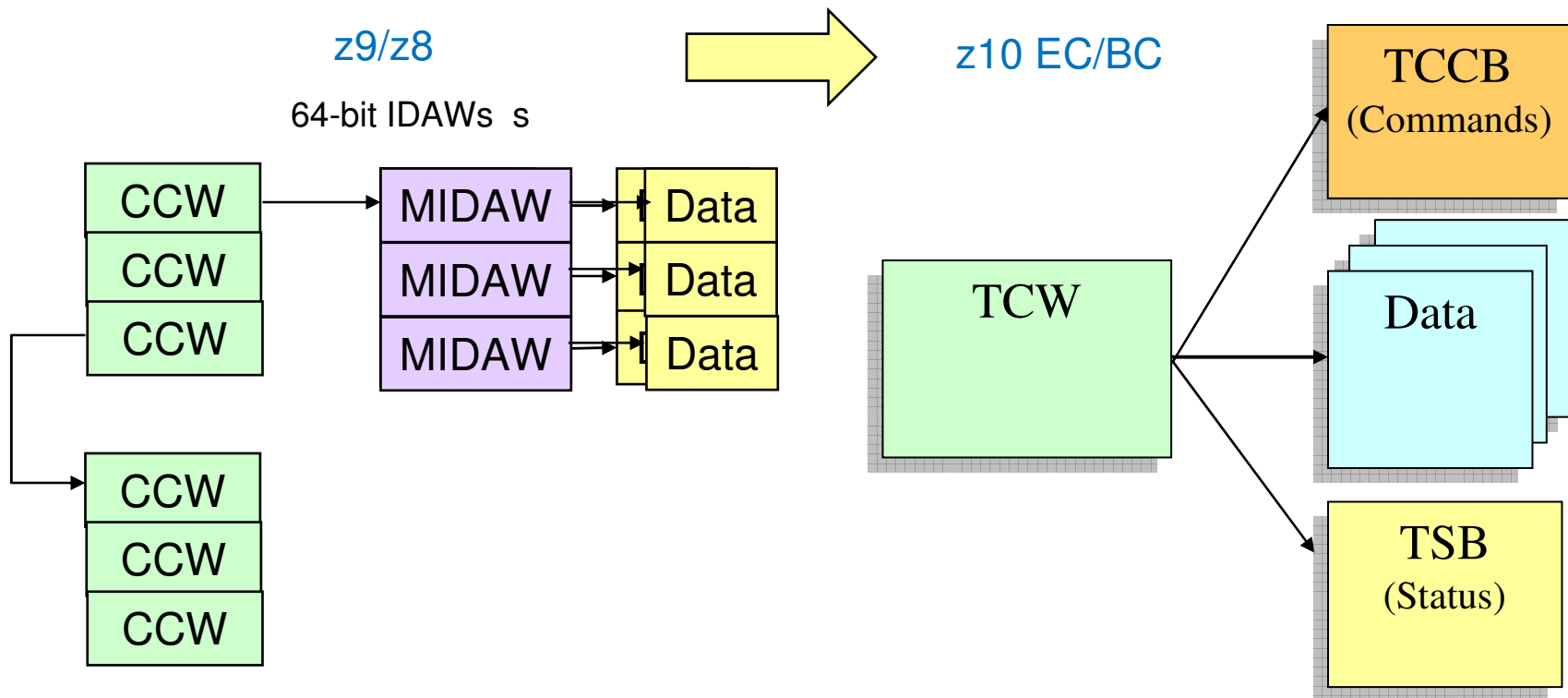


★ ESCON only supports half duplex

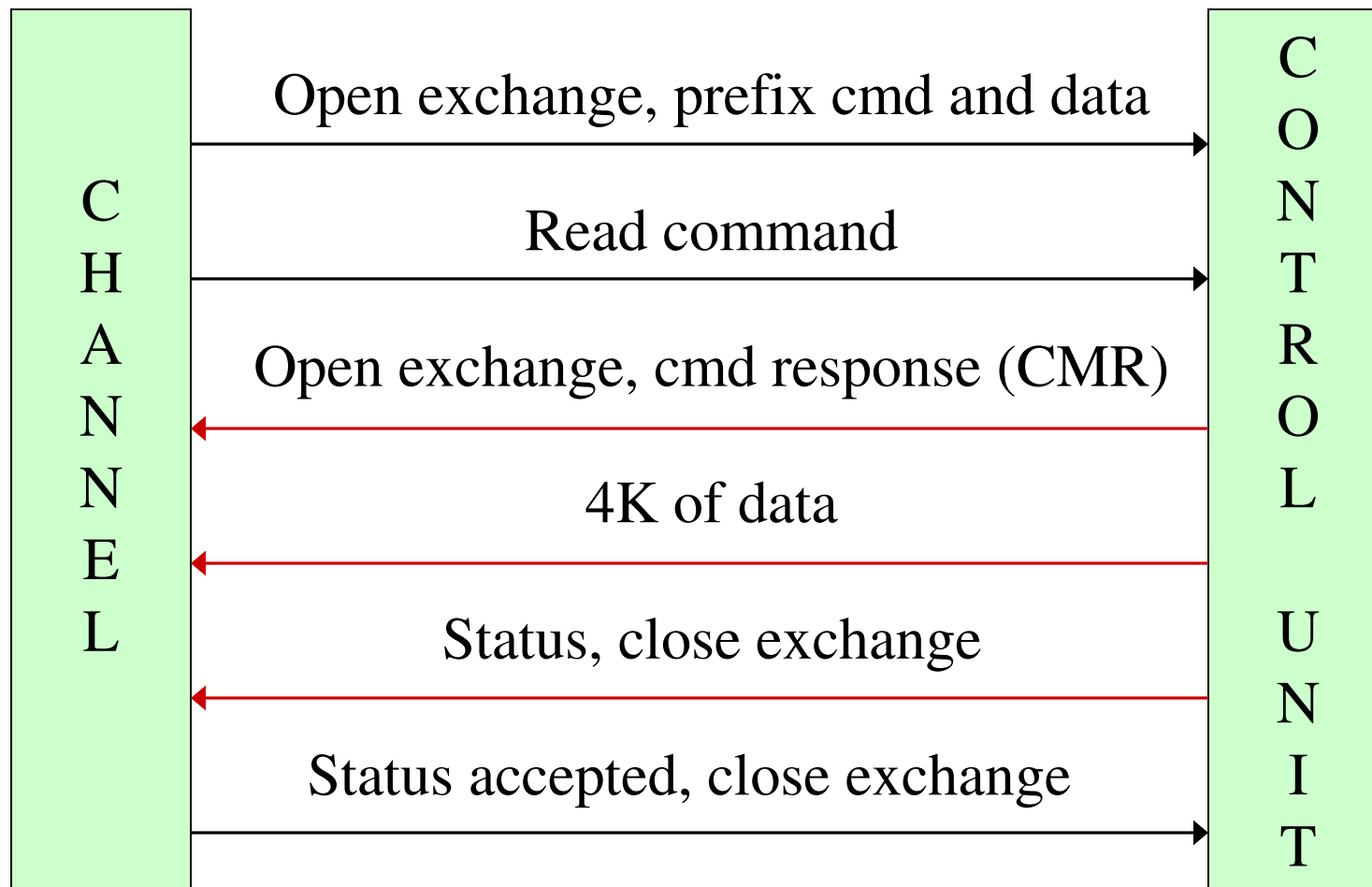


zHPF Trivia

zHPF is the first time we have dramatically changed the format of a channel program since System 360. Changes have been made to basic CCW channel program (e.g., format-1 CCWs, MIDAWs), and major changes have been made to the underlying architecture (XA, ESCON, FICON), but the channel program format was not radically changed.

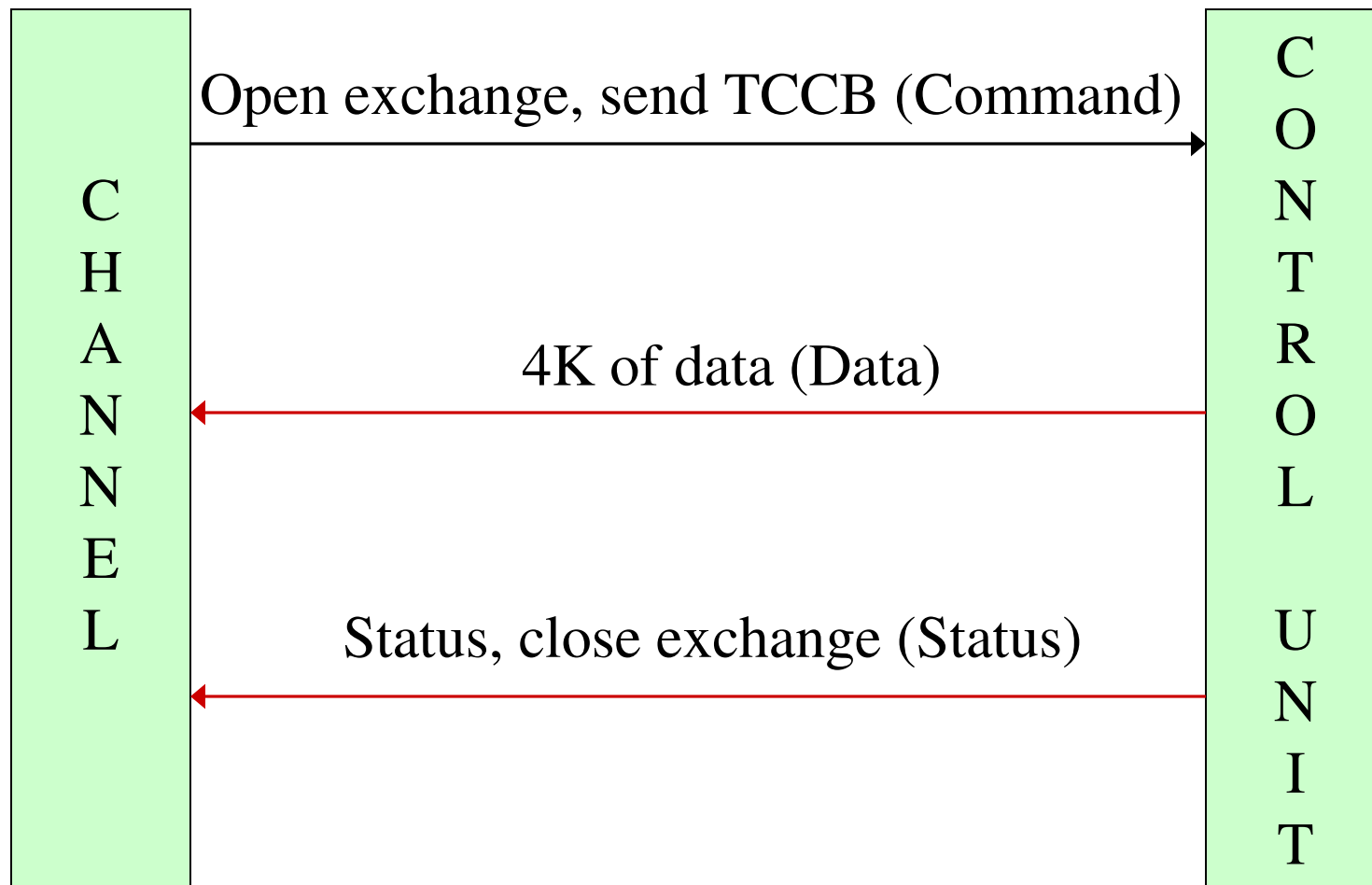


FC Link Protocol for FICON



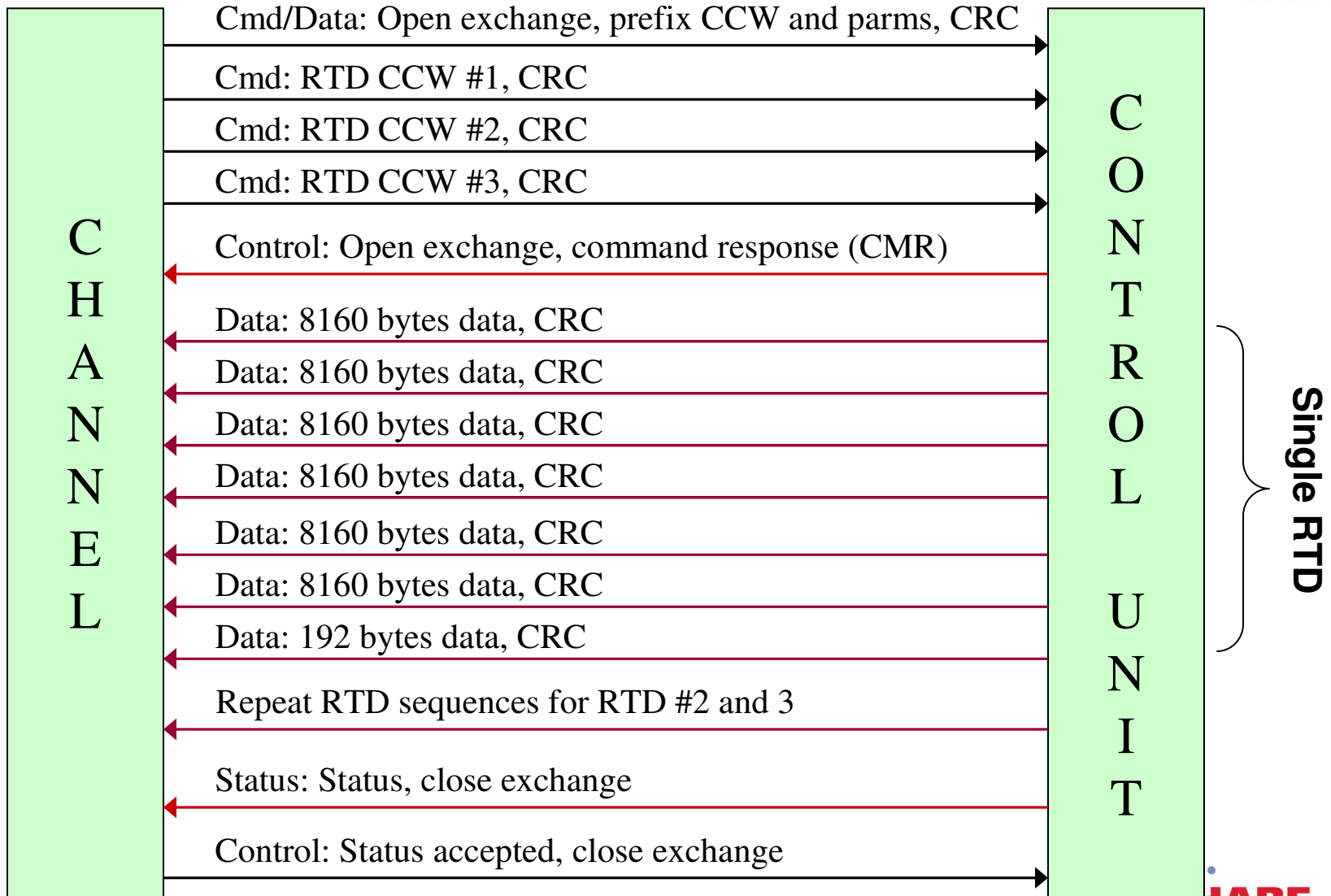
Example of a 4K read I/O request (6 sequences, 8 frames)

FC Link Protocol for zHPF (and FCP)



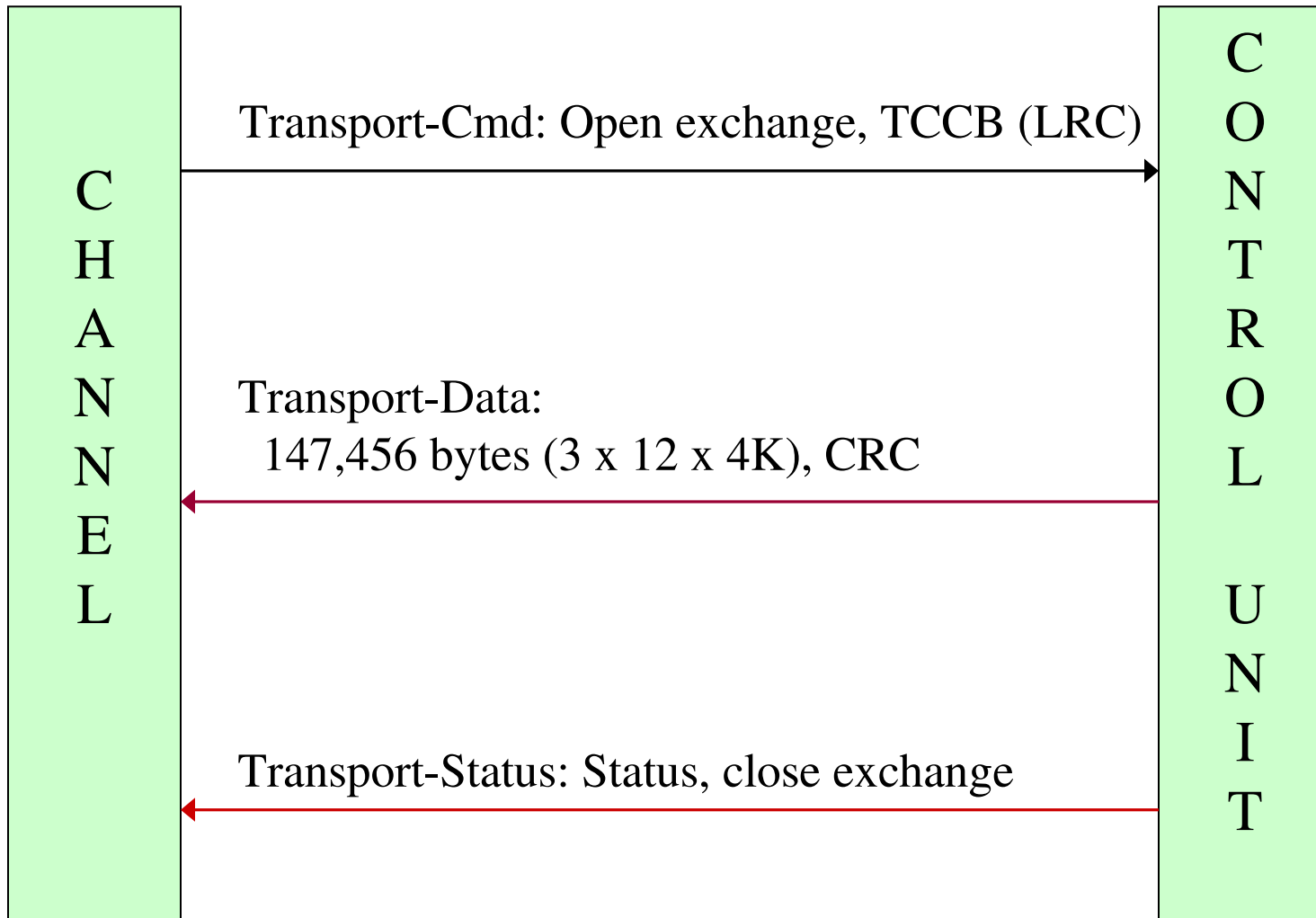
Example of a 4K read I/O request (3 sequences, 5 frames)

FICON – Read 3 Contiguous Tracks (12x4K Records)



Total = 28 sequences, 82 frames

zHPF – Read 3 Contiguous Tracks (12x4K Records)

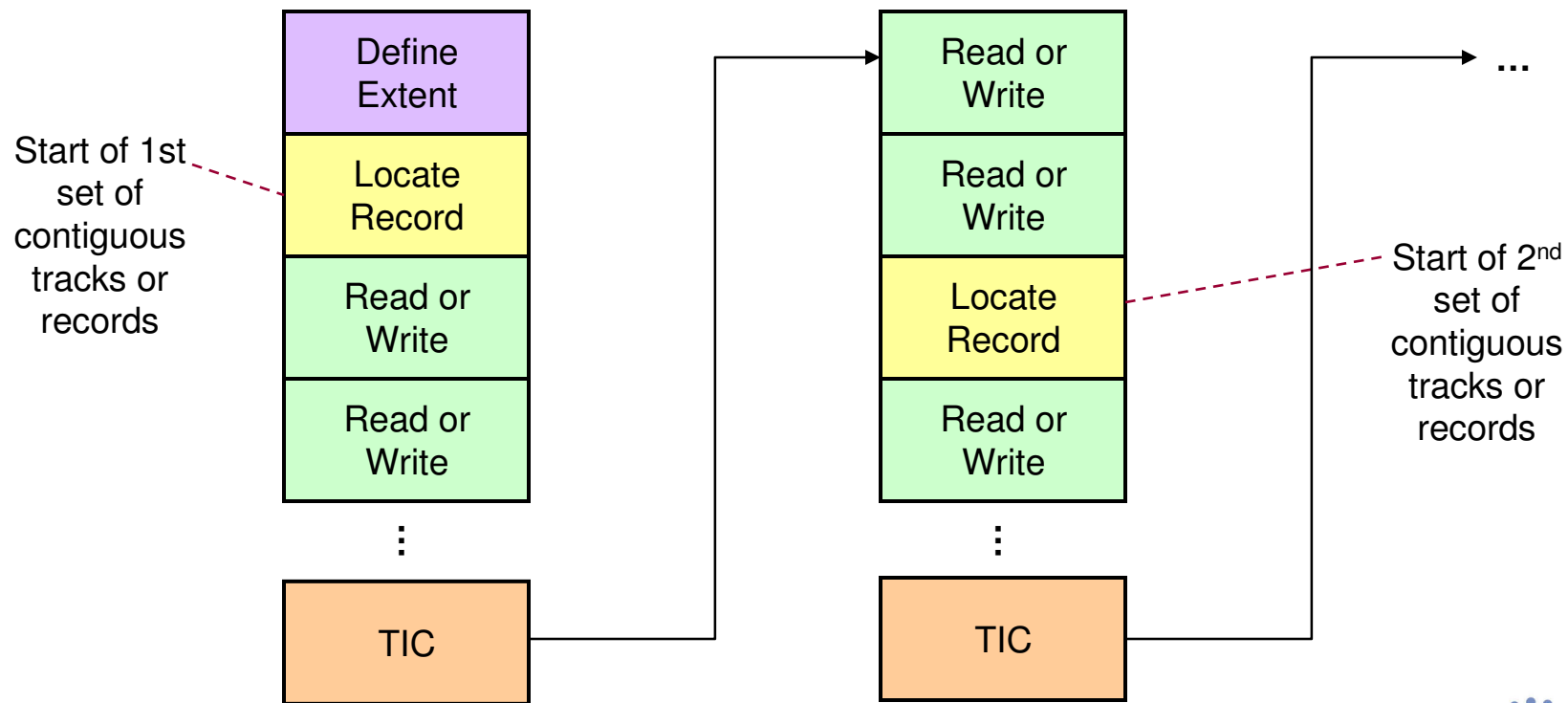


Total = 3 sequences, 75 frames



CCW Channel Programs – Virtually Unlimited in Size

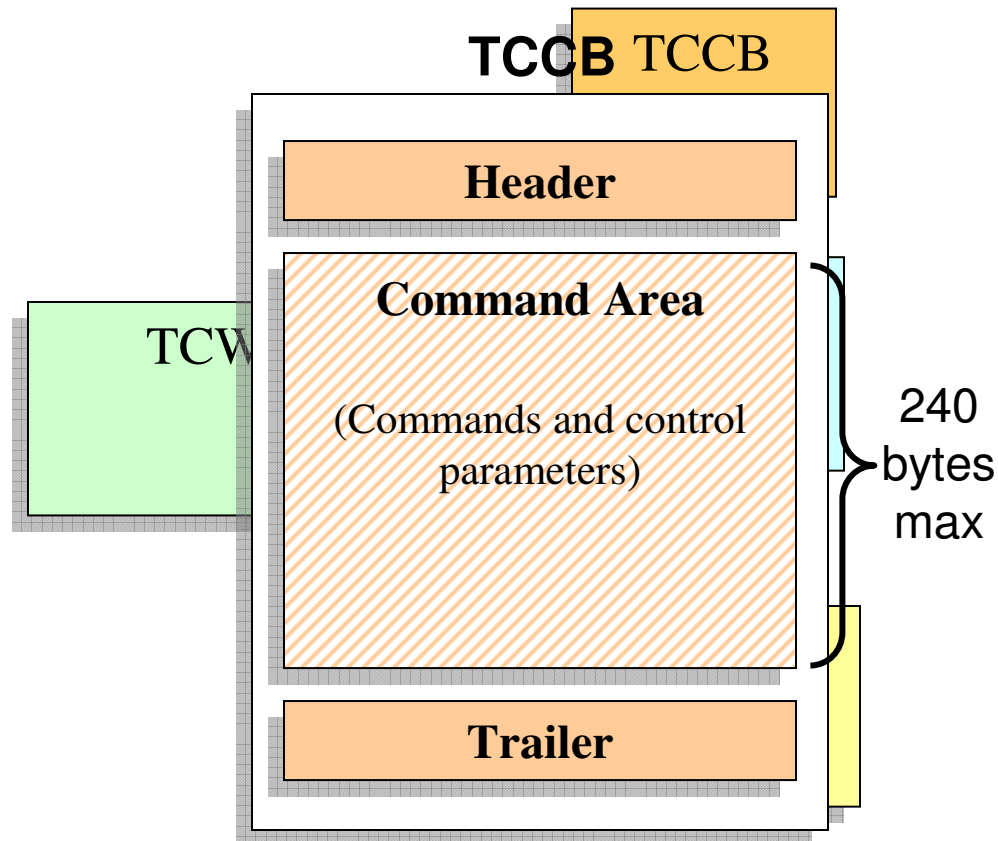
A single CCW channel program is virtually unlimited in size since CCWs can be chained together implicitly or explicitly. The size of a channel program is only limited by the amount of 31-bit virtual and real storage available.



zHPF Channel Programs – Limited by Architecture

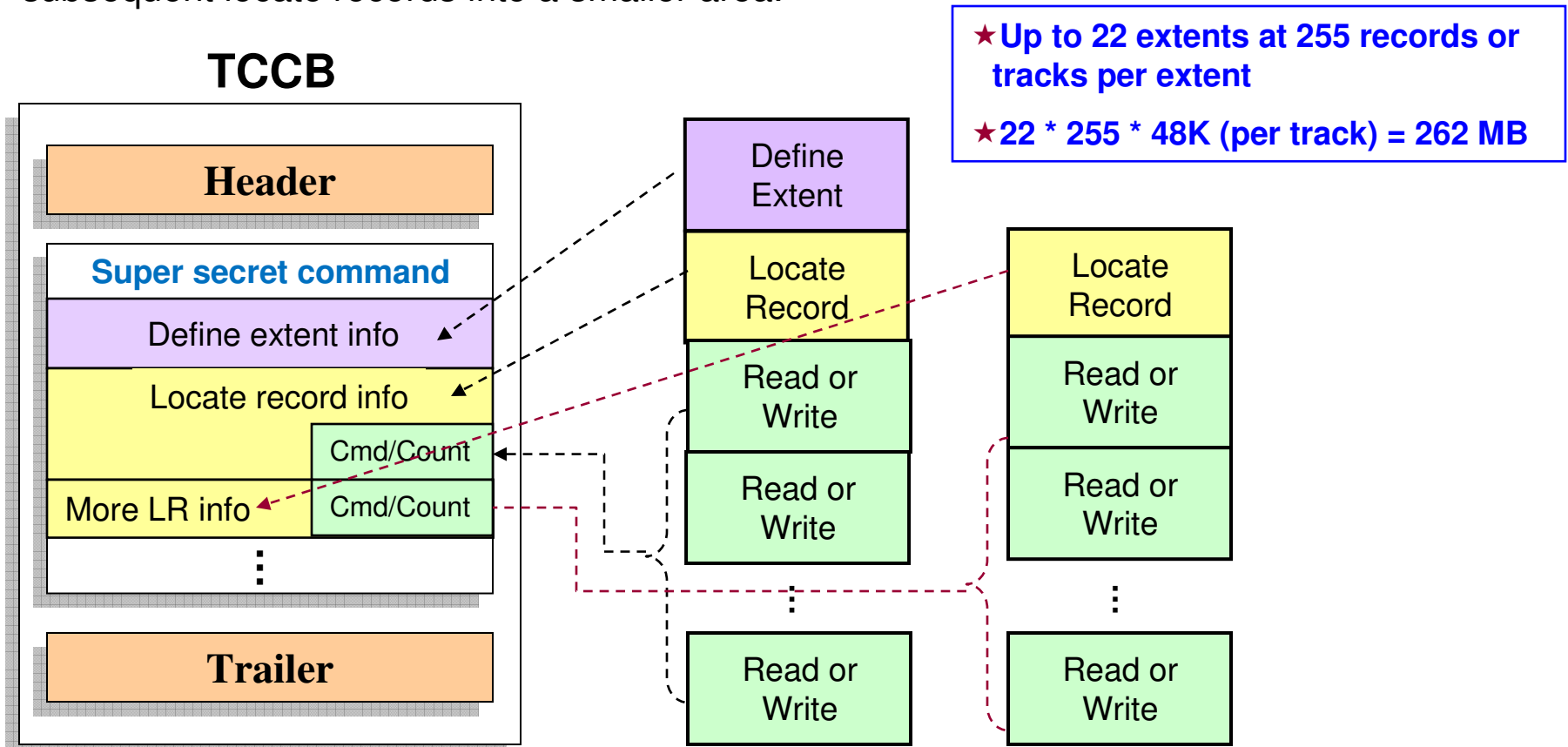


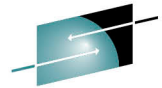
The command block for zHPF is basically a SCSI CDB, which limits the size to 240 bytes. How can we pack enough information into a channel program so that it can be used for things like DB2 list prefetch and QSAM/BSAM I/O?



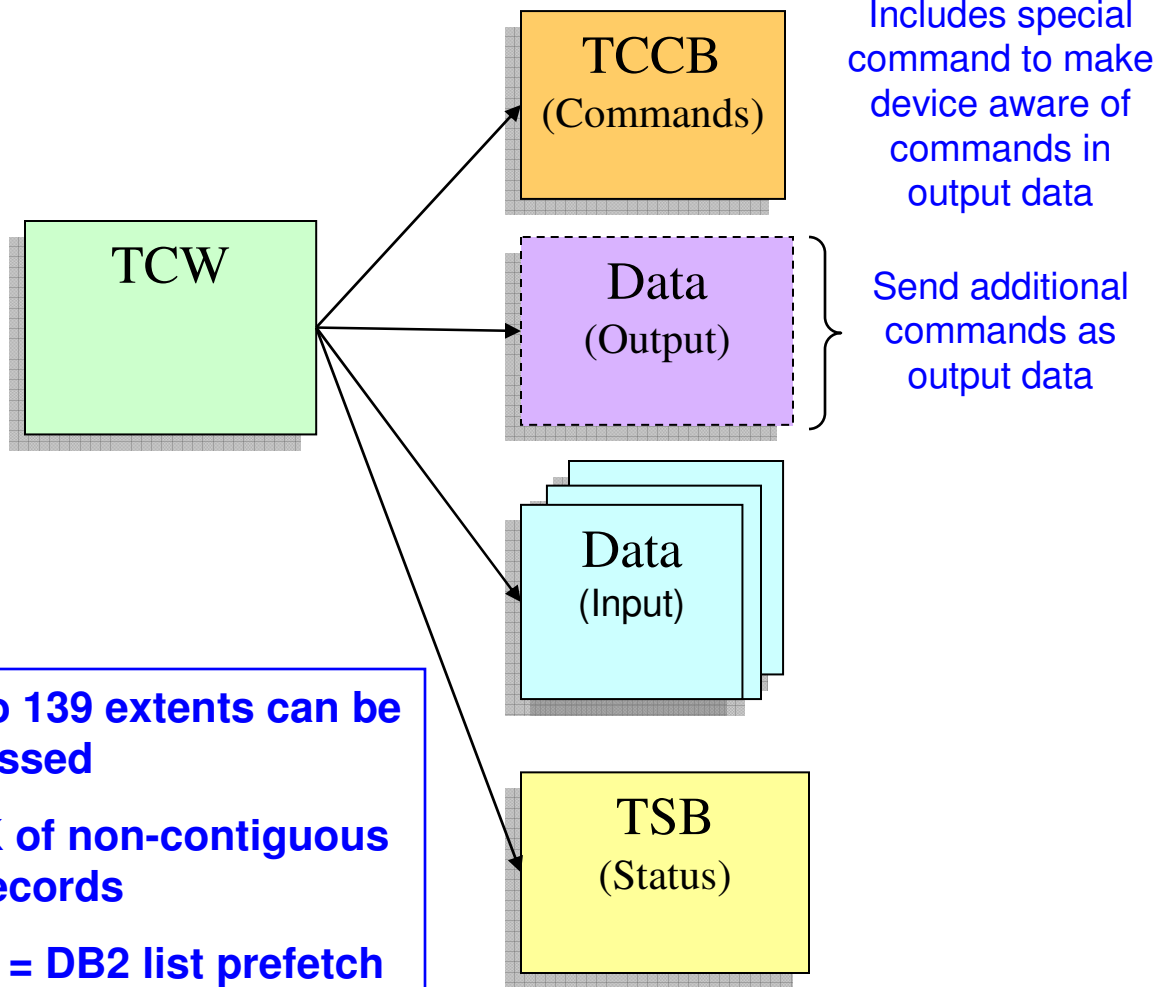
Step 1 – Collapse Lots of CCWs

For CCW channel programs, z/OS collapses the define extent and locate record CCWs into a single CCW. For zHPF, we can also collapse the read/write CCWs as well as the subsequent locate records into a smaller area.

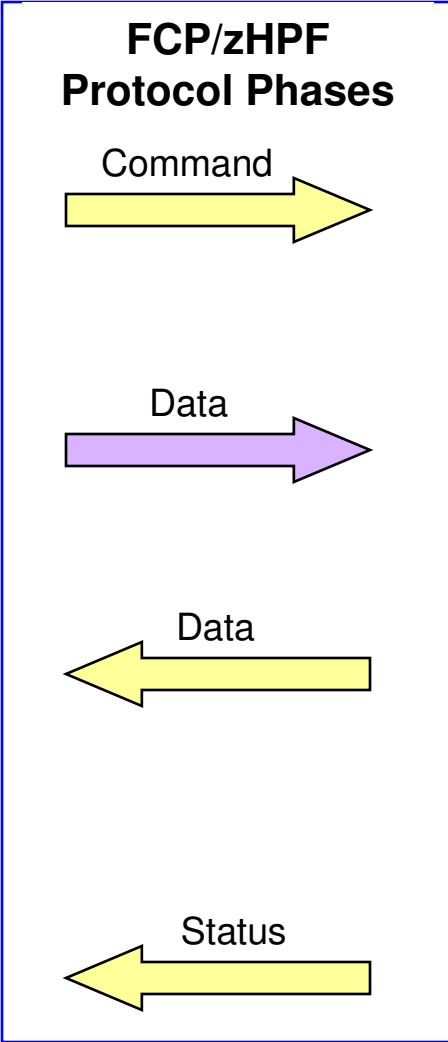




Step 2 – Extend Commands into Output Stream



- ★ Up to 139 extents can be accessed
- ★ 556K of non-contiguous 4K records
- ★ User = DB2 list prefetch



zHPF Tips – How to Tell if zHPF is Used



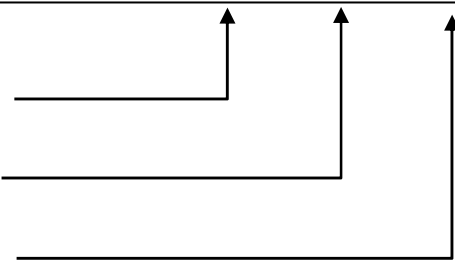
- Channel path activity report shows FICON vs zHPF activity
 - This is the only place where zHPF usage is recorded

CHANNEL PATH		UTILIZATION (%)				READ (MB/SEC)		WRITE (MB/SEC)		FICON OPERATIONS			ZHPF OPERATIONS			
ID	TYPE	G	SHR	PART	TOTAL	BUS	PART	TOTAL	PART	TOTAL	RATE	ACTIVE	DEFER	RATE	ACTIVE	DEFER
00	FC_S	5	Y	100.00	100.00	0.84	0.13	2.15	0.17	2.68	61.5	1.7	0.0	4.6	1.0	0.0
01	FC_S	5	Y	100.00	100.00	0.85	0.13	2.21	0.13	2.69	61.3	1.8	0.0	4.7	1.0	0.0
02	FC_S	4	Y	0.14	2.30	0.85	0.10	2.17	0.13	2.70	61.3	1.3	0.0	4.6	1.0	0.0
03	FC_S	4	Y	0.13	2.27	0.84	0.11	2.14	0.13	2.66	60.0	1.3	0.0	4.4	1.0	0.0
04	FC_S	5	Y	0.13	2.24	0.82	0.10	2.07	0.13	2.63	59.4	1.7	0.0	4.4	1.0	0.0
05	FC_S	5	Y	0.13	2.25	0.83	0.10	2.11	0.12	2.66	59.1	1.7	0.0	4.2	1.0	0.0
06	FC_S	4	Y	0.12	2.23	0.83	0.10	2.09	0.13	2.68	58.7	1.3	0.0	4.2	1.0	0.0

zHPF I/Os per second (physical channel)

Average number of active I/Os (open exchanges)

Average number of deferred I/Os (no available open exchanges)



zHPF Tips – Is zHPF Enabled?

```
D IOS, ZHPF
IOS630I 13.35.33 ZHPF FACILITY 742
HIGH PERFORMANCE FICON FACILITY IS ENABLED
```

```
D M=DEV(3100)
IEE174I 11.32.54 DISPLAY M 965
DEVICE 3100 STATUS=ONLINE
CHP                A4    B4    D0
ENTRY LINK ADDRESS 22    2323 B3
DEST LINK ADDRESS  42    23B2 42
PATH ONLINE        Y     Y     Y
...
FUNCTIONS ENABLED = MIDAW, ZHPF
```

```
D M=CU(3100)
IEE174I 11.08.13 DISPLAY M 226
CONTROL UNIT 3100
CHP                A4    B4    D0
ENTRY LINK ADDRESS 22    23    B3
...
ZHPF - CHPID       Y     Y     Y
ZHPF - CU INTERFACE Y     Y     Y
...
FUNCTIONS ENABLED = ZHPF
```

- Enables zHPF at a system level, default is NO
- Device could be disabled for any number of reasons including CU LIC feature
- Certain I/O errors can disable zHPF at the device or CU level (these are usually a sign of a microcode problem)
- zHPF CHPID/CU interface should always be Y unless it is connected to an ESCON channel/interface or FICON express
- If zHPF doesn't appear in functions enabled, then either LIC feature not installed or microcode problem

zHPF Tips – Is zHPF Enabled?

```
D SMS,OPTIONS
IGD002I 13:36:10 DISPLAY SMS 750
ACDS      = SYS1SMS.DR5.ACDS
COMMDS    = SYS1SMS.DR5.COMMDS
ACDS LEVEL = UNAVAIL
SMS PARMLIB MEMBER NAME = IGDSMSDR
INTERVAL = 15  DINTERVAL = 150
...
OAMTASK =          PDSE_SYSEVENT_DONTSWAP = NO
DB2SSID =          SAM_USE_HPF = YES
...
```

- IGDSMSxx parmlib member indicates whether QSAM/BSAM should use zHPF.

```
DEVSERV,QDASD,7000,RDFEATS
IEE459I 11.57.32 DEVSERV QDASD 848
  UNIT VOLSER SCUTYPE DEVTYPE          CYL  SSID...
07000 ----- 2107951 2107900          1113  7000...
  FEATURE CODES AT V020A0AF8
00070006001F00A6 DDDFFDE164400000 FD80000000000000 C3D000BB00000000
E785C00000000000 AEC0000000000000
```

- zHPF capability/feature bits (+x'28'). At a minimum, bits 0 & 2 should be on. Other bits may also need to be on for other exploitation (e.g, bit 5 for QSAM/BSAM). Contact IBM support for specific questions.

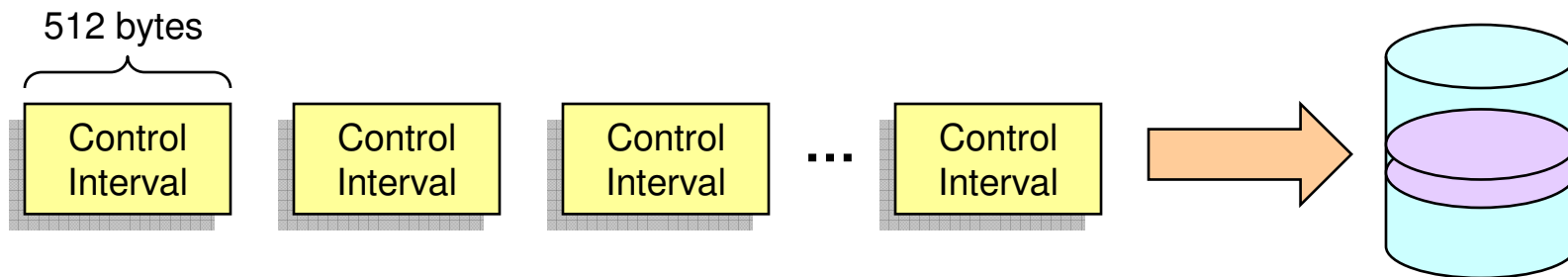
zHPF Trivia



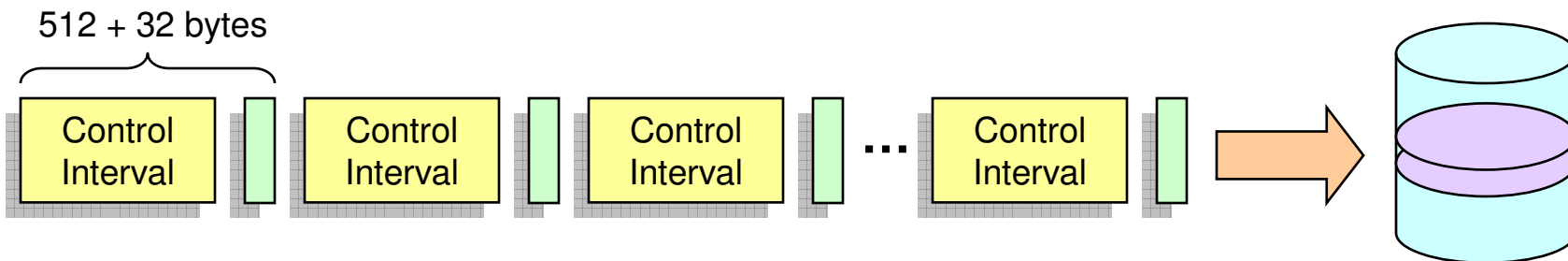
Although the channel supported a maximum transfer size of 2M in the initial zHPF support in the z10 processor, for performance reasons, the software limited the amount of data transferred to 64K. Where did this number come from?

Around 64K is where performance issues occurred.

512 bytes is the minimum VSAM control interval size. 128 of these fit into 64K (2 tracks are required).



Extended format data sets have a 32 byte suffix on each record.



☺ Therefore, we need to add $(64K / 512) * 32 = 4096$ to limit to allow 64K transfers for extended format data sets to eligible for zHPF.

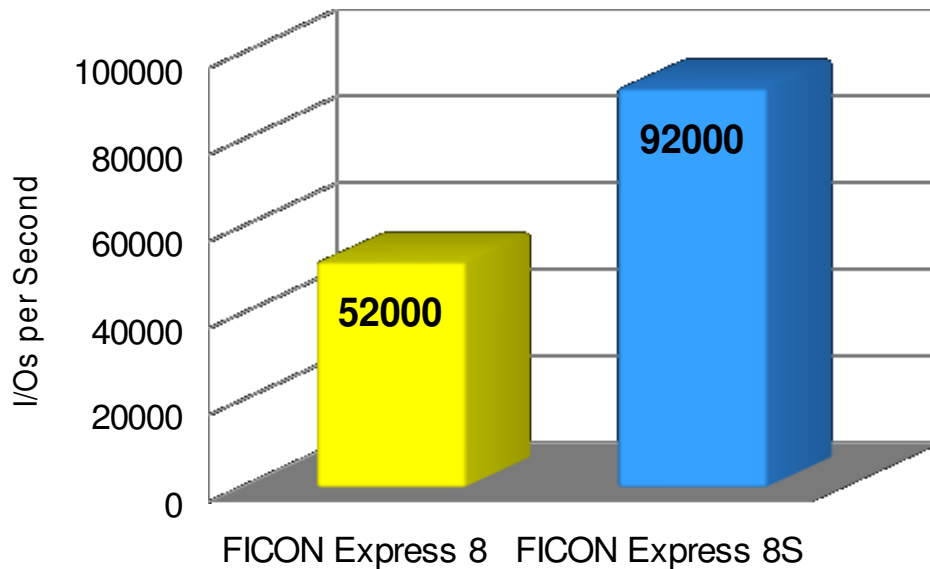


FICON Express 8S Channel

- zHPF protocol implemented in the hardware via a data router
- 76% higher start rates compared to zHPF on FICON express 8
- 107% higher throughput compared to zHPF on FICON express 8
- **100% of DB2 I/O can be converted to zHPF**

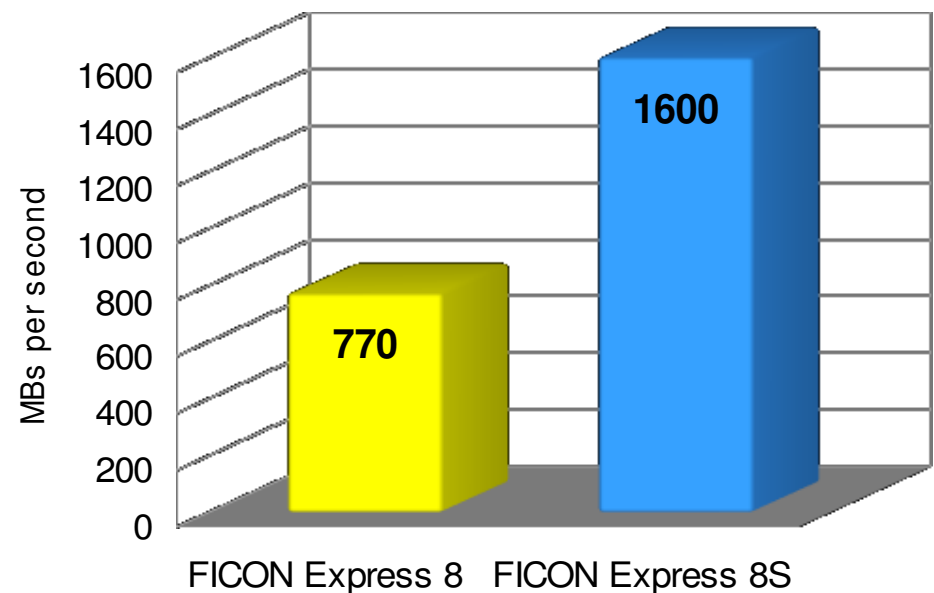
zHPF Start Rate Comparison

4K Block Size, 100% Utilization



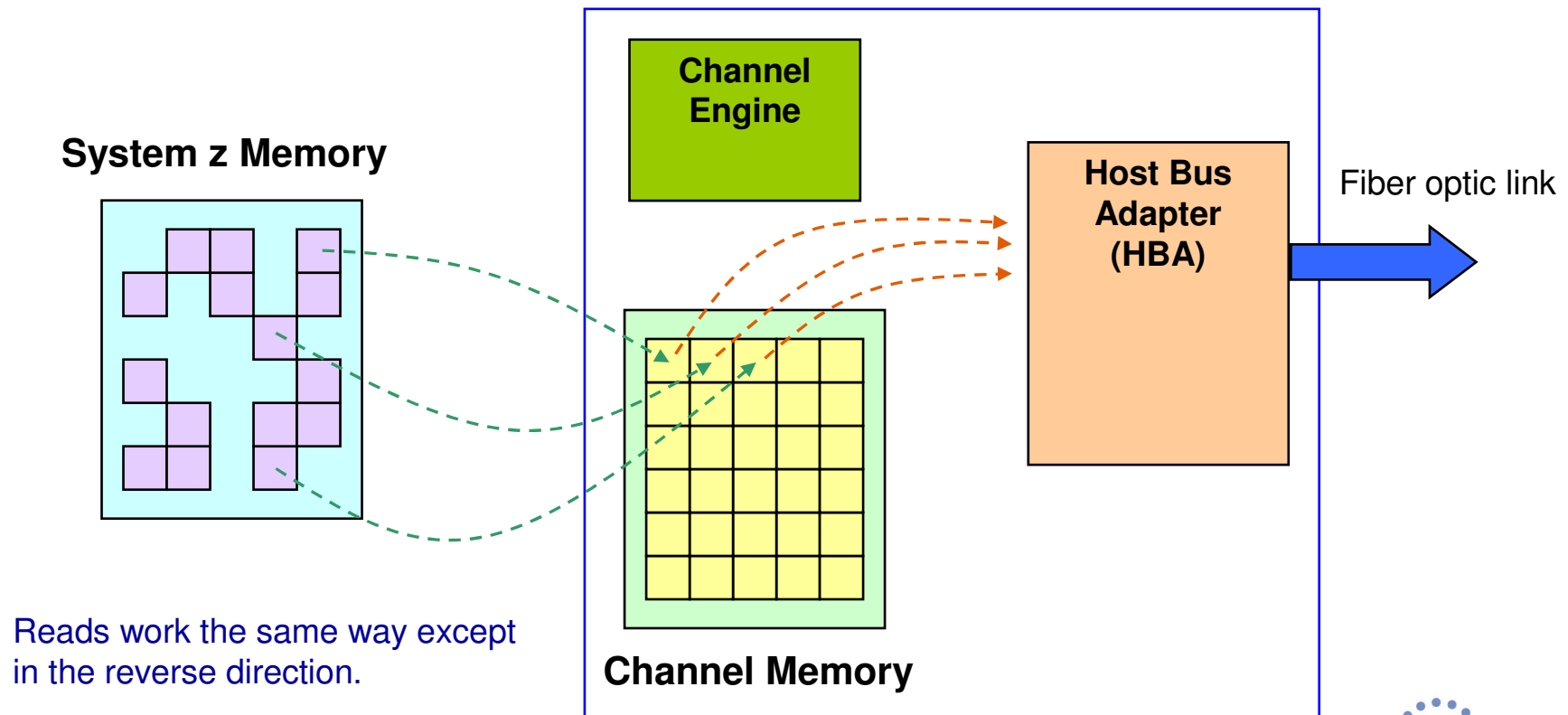
zHPF Throughput Comparison

Full duplex - large read/write mix



Older FICON Channels – Store and Forward

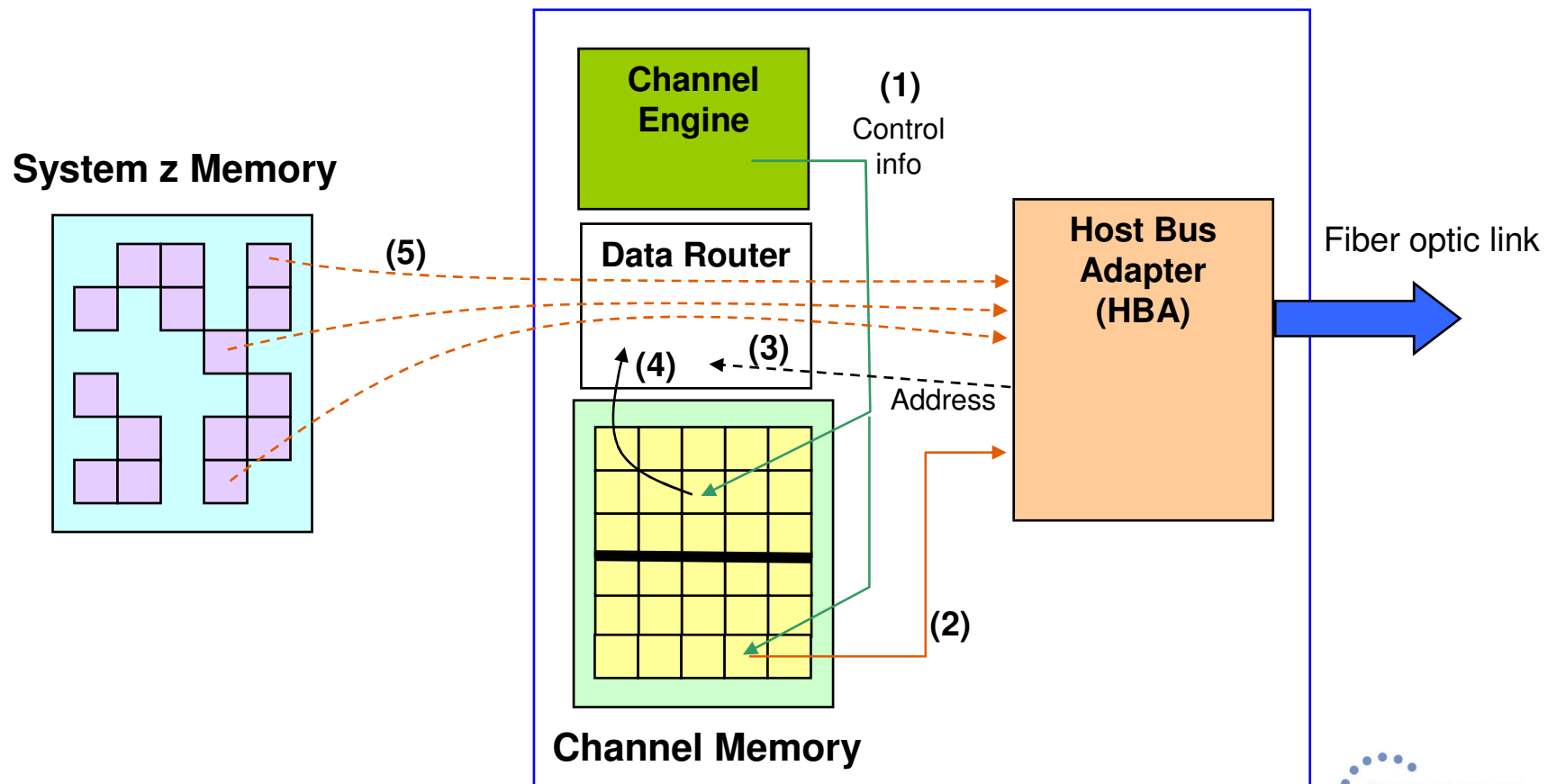
- ★ In older channels, the HBA did not access System z memory directly. A store and forward technique was used where the channel firmware would fetch information from System z memory into channel memory, which would then be accessed by the HBA.
- ★ A pipeline is established which allowed the channel to work in parallel with the HBA.



FICON Express 8S – H/W Data Router

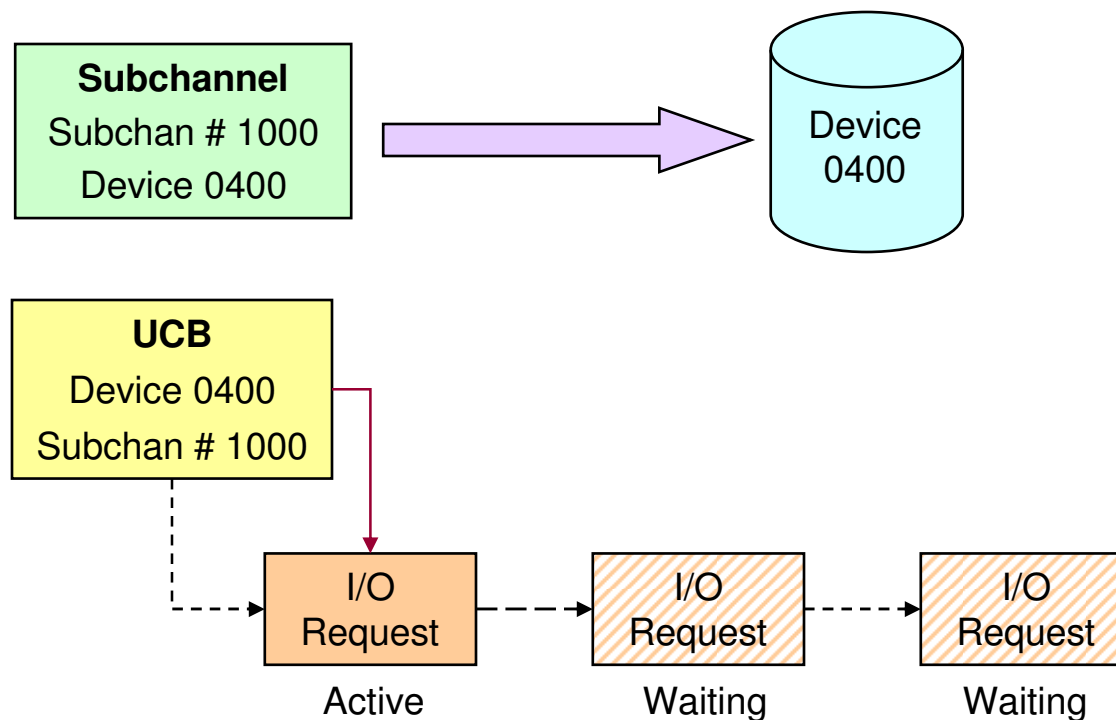


- ★ With the FICON express 8S channel, the HBA accesses System z memory directly through and under the control of a H/W data router.
- ★ The channel firmware is involved only in setting up the storage areas that need to be accessed in channel memory.



Parallel Access Volumes (PAV) – Before PAV

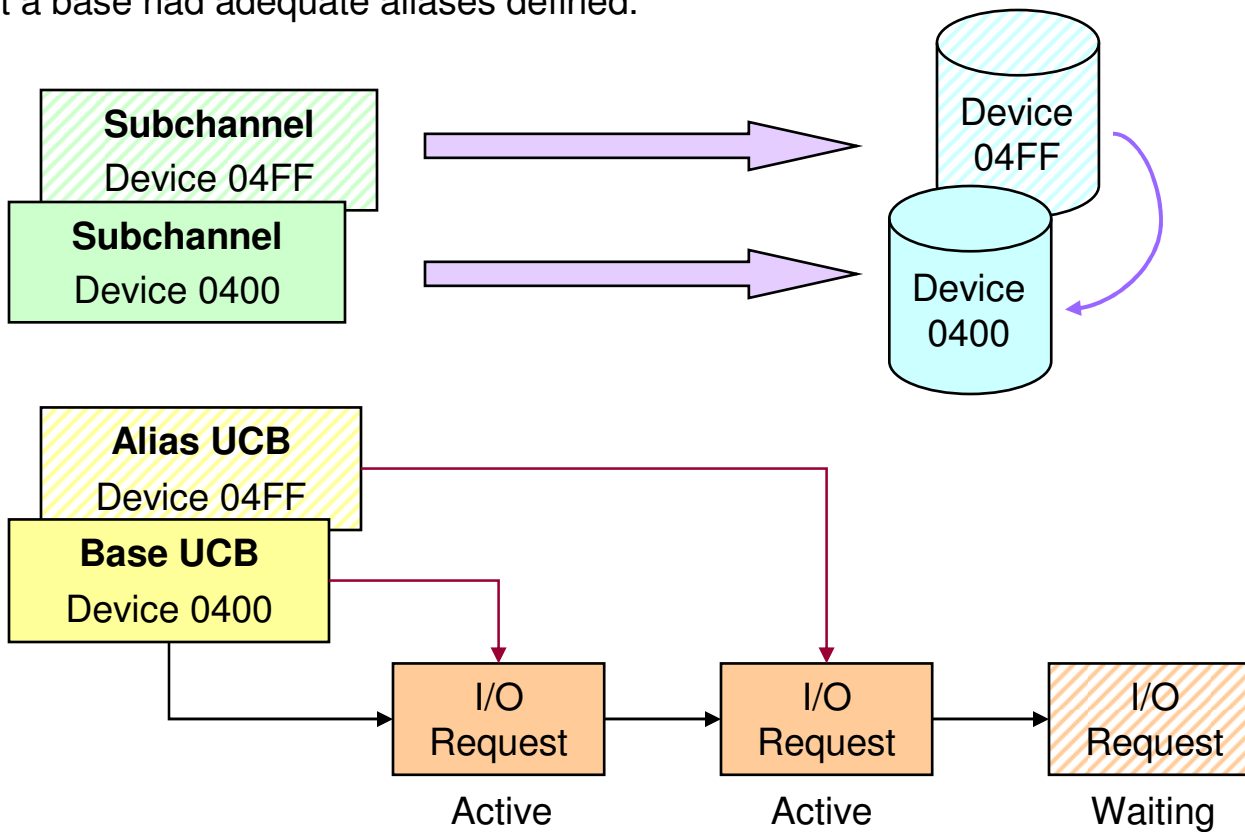
Although, storage controllers allow multiple I/Os to be issued to a single device from multiple systems (multiple allegiance), System z allows only one I/O to be issued to a device (subchannel in HSA) at a time for an LPAR.



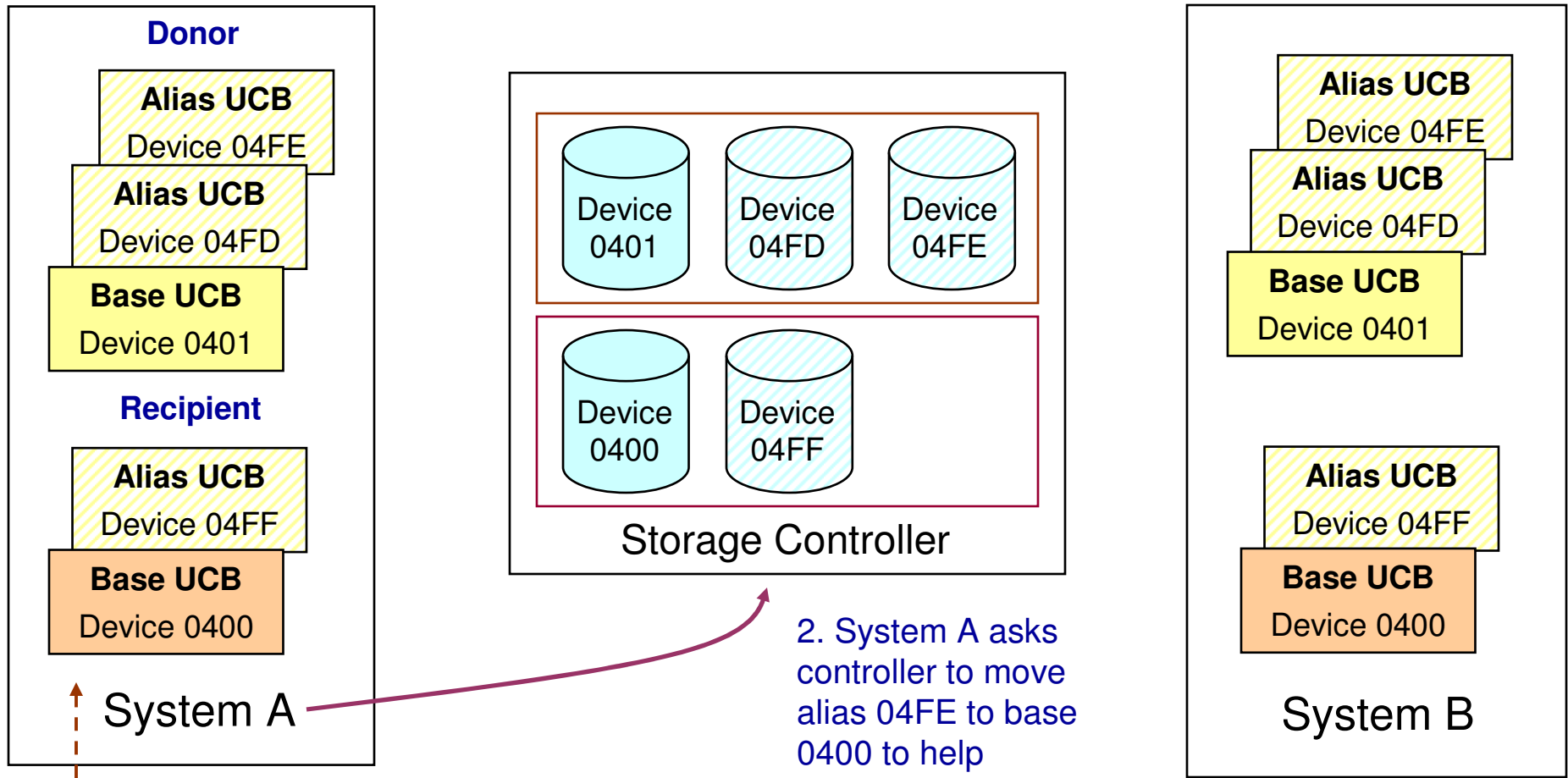
If multiple I/Os arrive for a device at the same time, the remaining I/O's will be queued internally in z/OS while waiting for the prior I/Os to complete (IOSQ time).

Static PAV

Static PAV allowed multiple *alias* devices to be defined for a *base* device. This allowed multiple I/Os to be done to a particular volume in parallel since a separate subchannel was assigned for each alias. However, the alias to base assignment was static, so this required careful planning to ensure that a base had adequate aliases defined.



Dynamic (WLM Managed) PAV

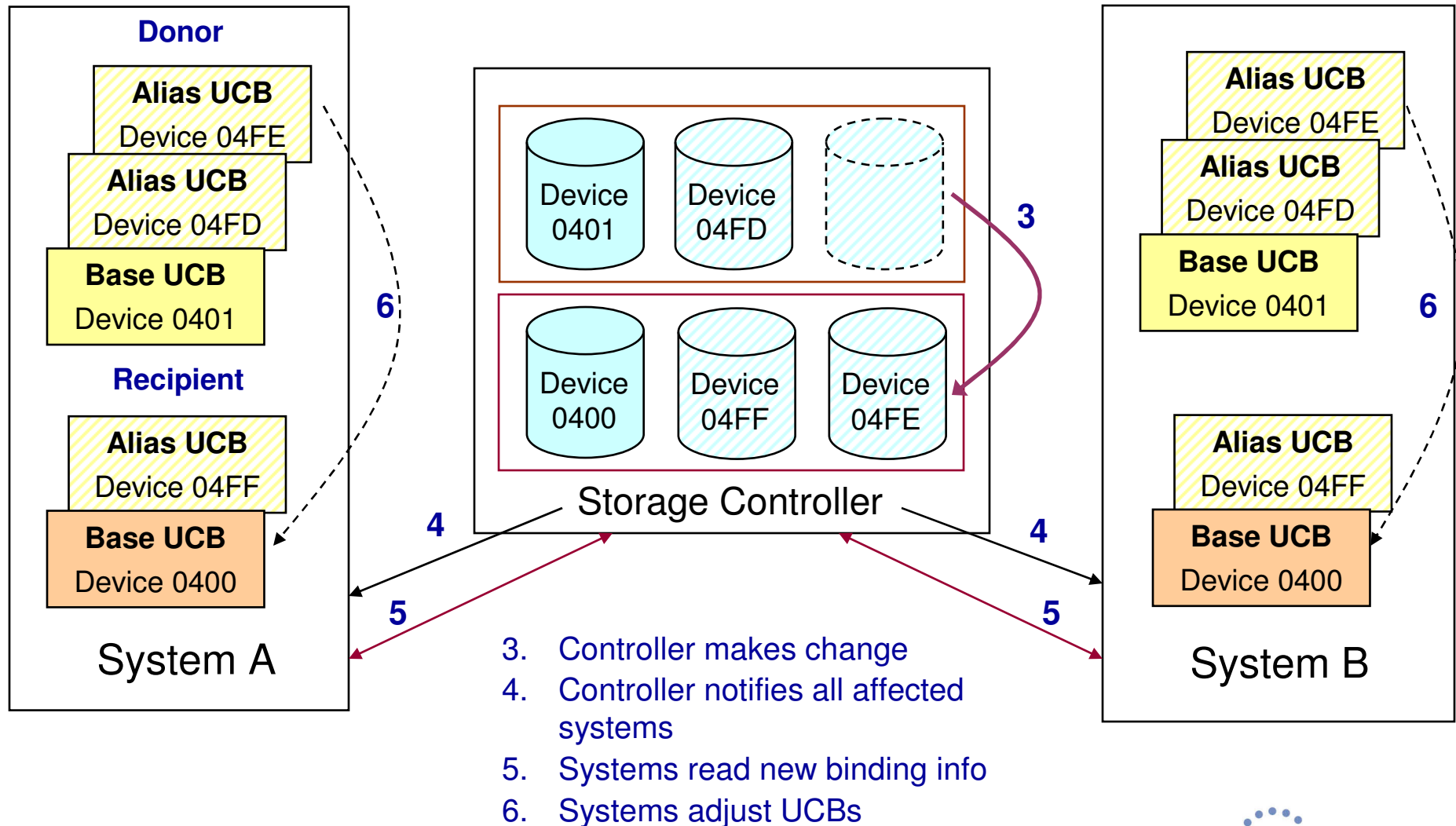


1. WLM on system A notices high IOSQ time on 0400 based on sysplex wide performance information

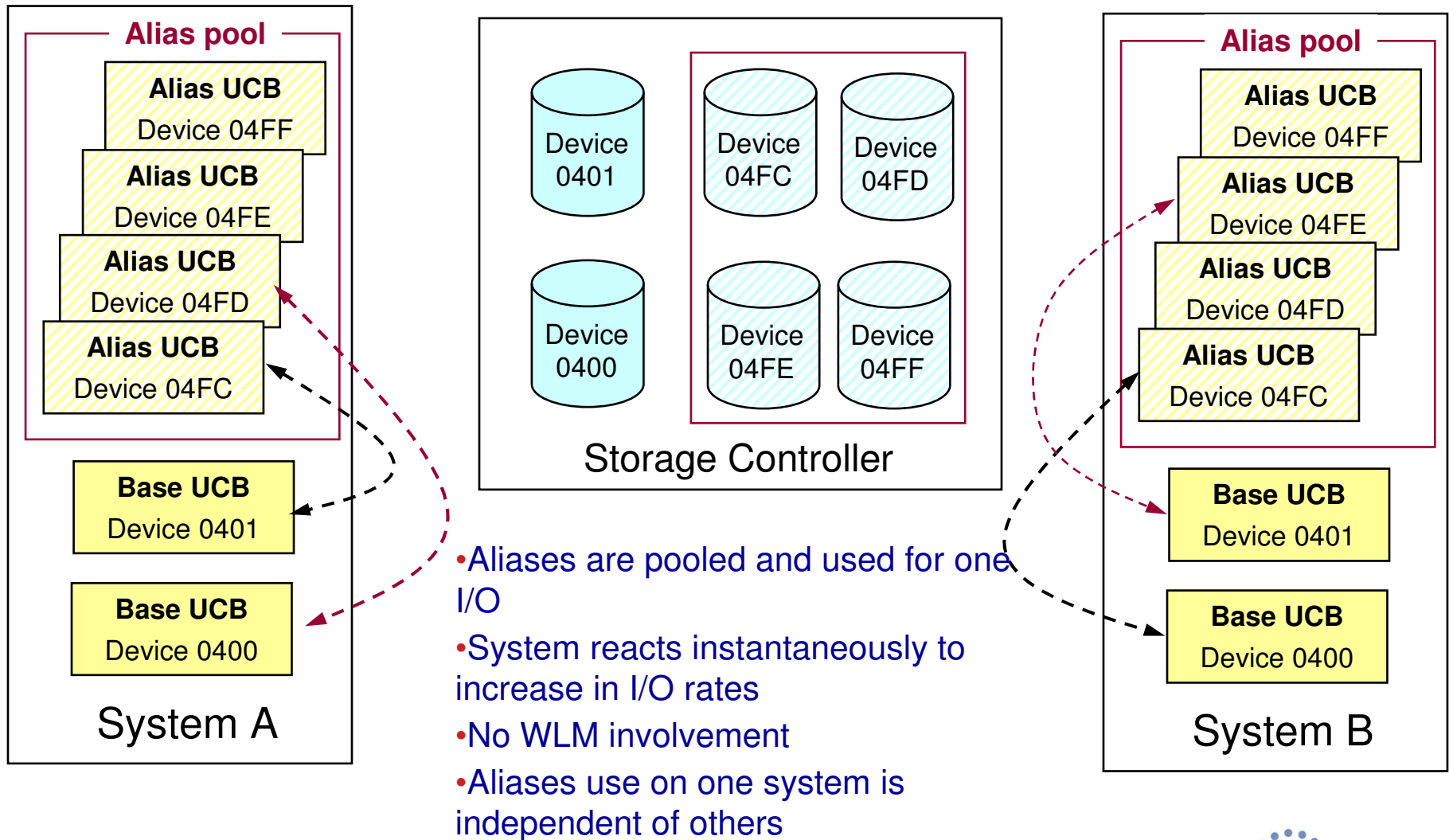
2. System A asks controller to move alias 04FE to base 0400 to help



Dynamic (WLM Managed) PAV



HyperPAV



HyperPAV Tips – How Many Aliases?

- Little's Law:
 - Number of active requests = I/O rate * I/O service time
- Applied to HyperPAV:
 - No. of aliases = Peak avg. I/O rate * service time * multiplier
- Multiplier added since peak average service time doesn't account for workload spikes
 - Multiplier should be at least 2 since peak concurrent I/Os will be higher than the average
 - Monitor for increases in IOSQ time or HPAV wait delay (see next page)

HyperPAV Tips - Monitoring



D I R E C T A C C E S S D E V I C E A C T I V I T Y															
z/OS V1R13				SYSTEM ID TRX2				DATE 11/23/2011							
				CONVERTED TO z/OS V2R1 RMF				TIME 12.30.00							
STORAGE GROUP	DEV NUM	DEVICE TYPE	NUMBER OF CYL	VOLUME SERIAL	PAV	LCU	ACTIVITY RATE	RESP TIME	IOSQ TIME	CMR DLY	DB DLY	INT DLY	PEND TIME	DISC TIME	CONN TIME
XTEST	2209	33903	3339	TRXSXA	1	0032	24.901	6.54	2.27	.202	.000	.020	.212	2.96	1.10
	220A	33909	10017	TRXT01	1	0032	20.085	6.12	2.06	.207	.000	.019	.217	2.80	1.04
	220B	33909	10017	TRXT02	1	0032	30.702	7.18	2.28	.210	.000	.013	.220	3.73	0.95
	220C	33909	10017	TRXT03	1	0032	27.103	6.93	2.03	.201	.000	.016	.211	3.67	1.02
	220D	33909	10017	TRXT04	1	0032	31.040	6.30	2.25	.208	.000	.022	.218	2.85	0.98
	220E	33909	10017	TRXT05	1	0032	33.010	6.62	2.31	.203	.000	.017	.213	3.08	1.01

First sign that I may not have enough aliases defined for this LCU is high IOSQ time.

Note that there may be other performance problems that are causing high IOSQ queue. For example, high disconnect time due to staging data into cache. Adding aliases won't necessarily help in these cases.



HyperPAV Tips – Monitoring



I/O QUEUING ACTIVITY													
z/OS V1R13				SYSTEM ID TEST1				START ...					
TOTAL SAMPLES = 901 IODF = 72				RPT VERSION V1R13 RMF				END ...					
				CR-DATE: 11/01/2013				CR-TIME: ...					
LCU	CU	DCM	GROUP	CHAN	CHPID	% DP	% CU	AVG CUB	AVG CMR	...	AVG CSS	HPAV	
		MIN	MAX	DEF	TAKEN	BUSY	BUSY	DLY	DLY		DLY	WAIT	MAX
00A0	2600				0C	0.003	0.00	0.00	0.0	0.1			
					0F	0.007	0.00	0.00	0.0	0.0			
					42	0.007	0.00	0.00	0.0	0.2			
					45	0.007	0.00	0.00	0.0	0.6			
					00	0.007	0.00	0.00	0.0	1.2			
					03	0.007	0.00	0.00	0.0	0.5			
					06	0.003	0.00	0.00	0.0	1.3			
					09	0.003	0.00	0.00	0.0	2.2			
					*	0.043	0.00	0.00	0.0	0.6	0.1	0.000	10

Average number of I/O requests that had to wait because no aliases available

Maximum number of aliases in use during this interval. **Be careful with this number because short spike may cause it to increase!**



HyperPAV Tips - Monitoring



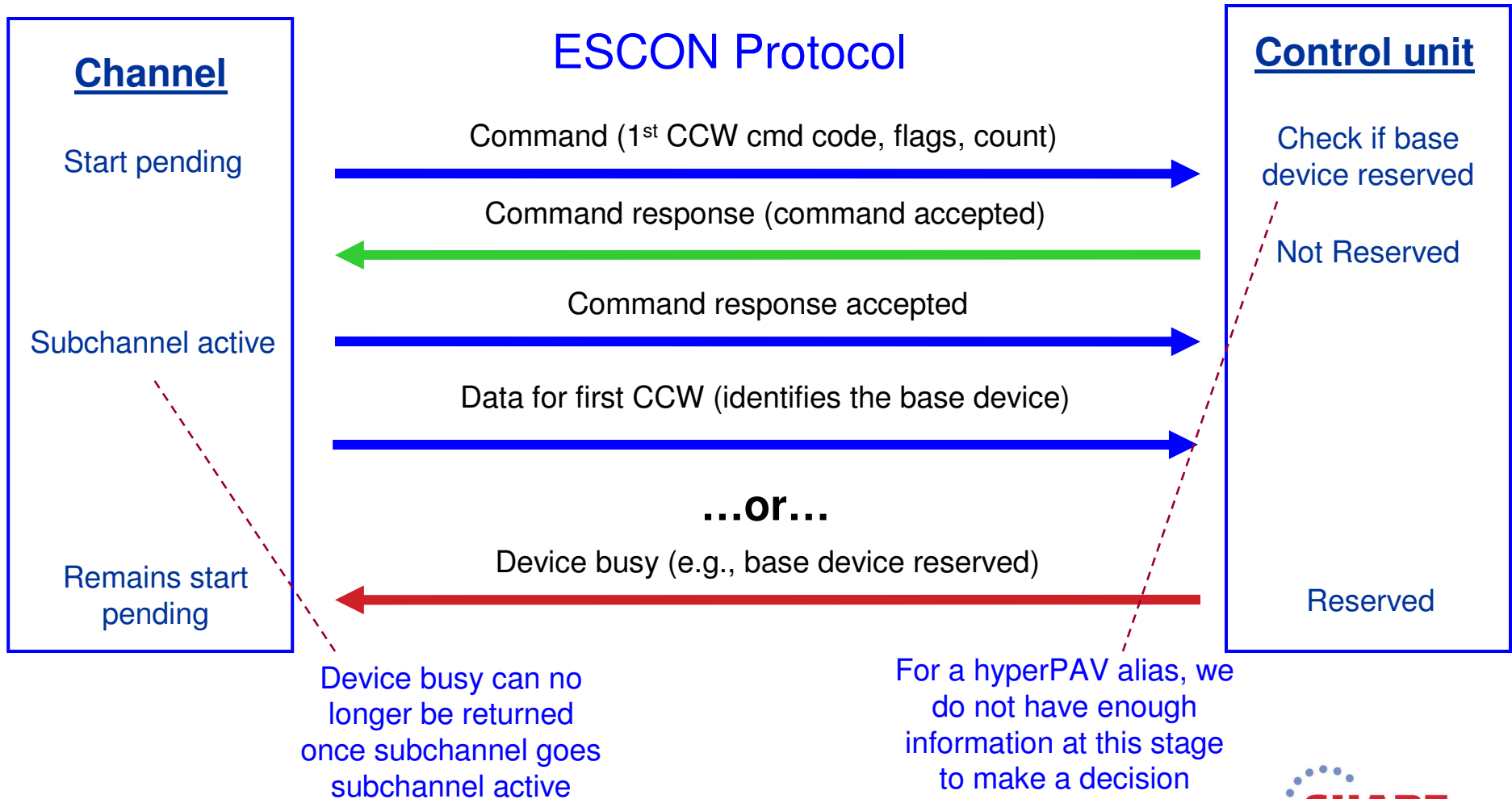
```
SY1  d m=dev(0710)
SY1  IEE174I 23.35.49 DISPLAY M 835
DEVICE 0710    STATUS=ONLINE
CHP                10    20    30    40
DEST LINK ADDRESS  10    20    30    40
PATH ONLINE       Y     Y     Y     Y
CHP PHYSICALLY ONLINE Y   Y   Y   Y
PATH OPERATIONAL  Y     Y     Y     Y
MANAGED           N     N     N     N
CU NUMBER         0700 0700 0700 0700
MAXIMUM MANAGED CHPID(S) ALLOWED:  0
DESTINATION CU LOGICAL ADDRESS = 07
SCP CU ND         = 002107.000.IBM.TC.03069A000007.00FF
SCP TOKEN NED     = 002107.900.IBM.TC.03069A000007.0700
SCP DEVICE NED    = 002107.900.IBM.TC.03069A000007.0710
HYPERPAV ALIASES IN POOL 32
```

Sometimes an I/O configuration or z/OS issue will cause aliases not to be available for use. You should compare this number with what you have defined for your LCU.

PAV Trivia



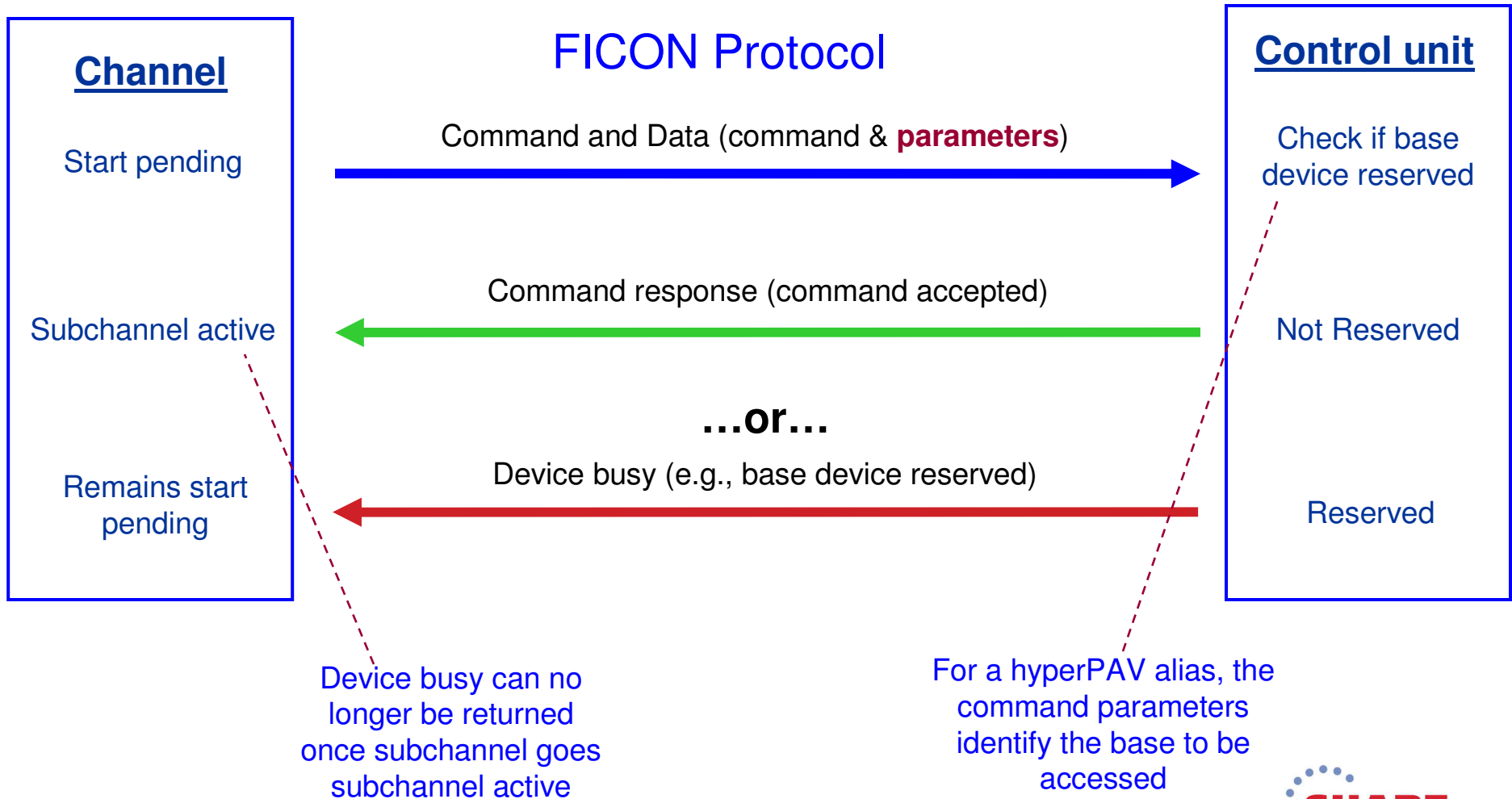
Why is HyperPAV supported only for FICON channels?



PAV Trivia



Why is HyperPAV supported only for FICON channels?



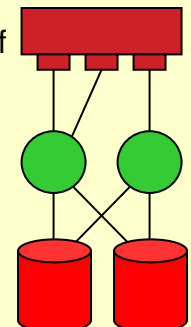
Dynamic Channel-Path Management (DCM)

- Goals
 - Simplify I/O configuration management
 - Improve performance management
 - Enhance RAS
- System manages a pool of channels that are dynamically added to control units as needed
 - Managed channel
 - Managed control unit
- System will add or move channels based on:
 - Channel utilization hot spots
 - Port utilization hot spots
 - Availability need – down to your last channel path

Dynamic Channel-Path Management (DCM)

Topology Discovery

- ★ Internal representation of connections between:
 - Channels
 - Switches
 - Control Units
- ★ Updated as changes to I/O topology occur



Data Collection

- ★ Collect performance data (e.g., channel and CU port utilization, connect time, bytes transferred)

DCM Algorithms

- ★ Check for hot spots and performance issues
- ★ Determine what changes can be made
- ★ Predict new channel utilizations
- ★ Select the best change
- ★ Change the I/O configuration via dynamic I/O
- ★ Vary the path online to all systems

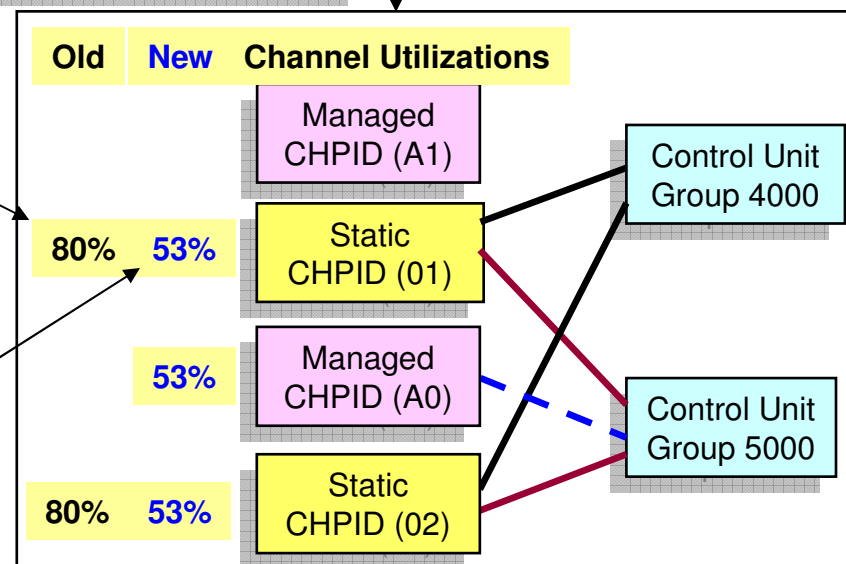
Possible changes

1. **Add A0 to 5000**
2. Add A1 to 5000
3. Add A0 to 4000
4. Add A1 to 4000

Dynamic I/O

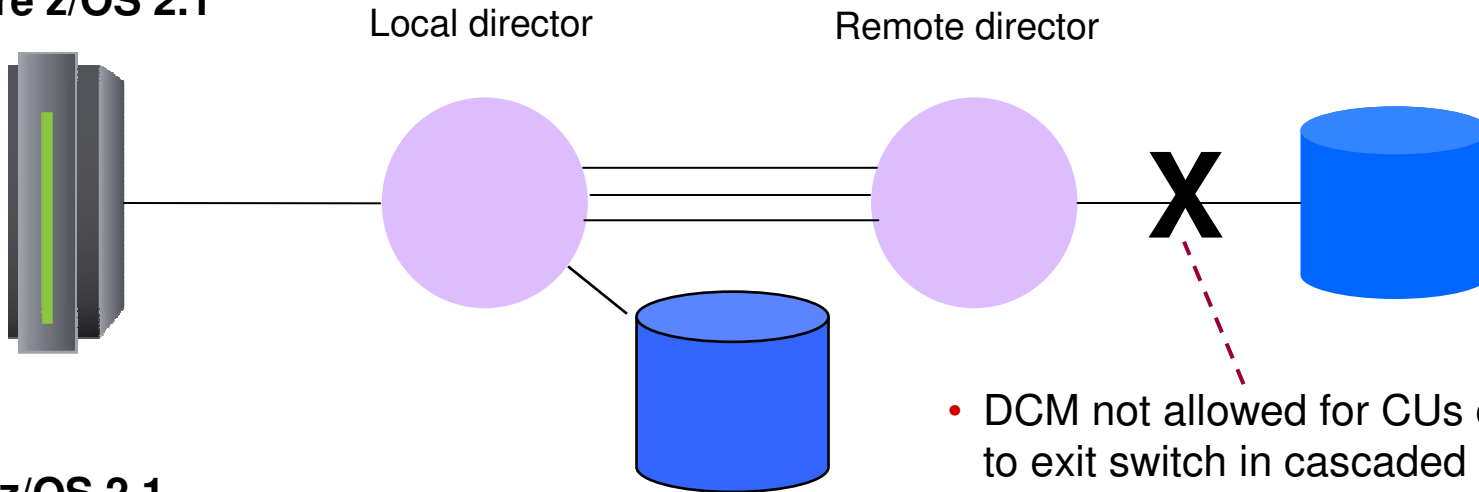
- ★ Add, replace or delete CHPID from control unit group

DCM Change



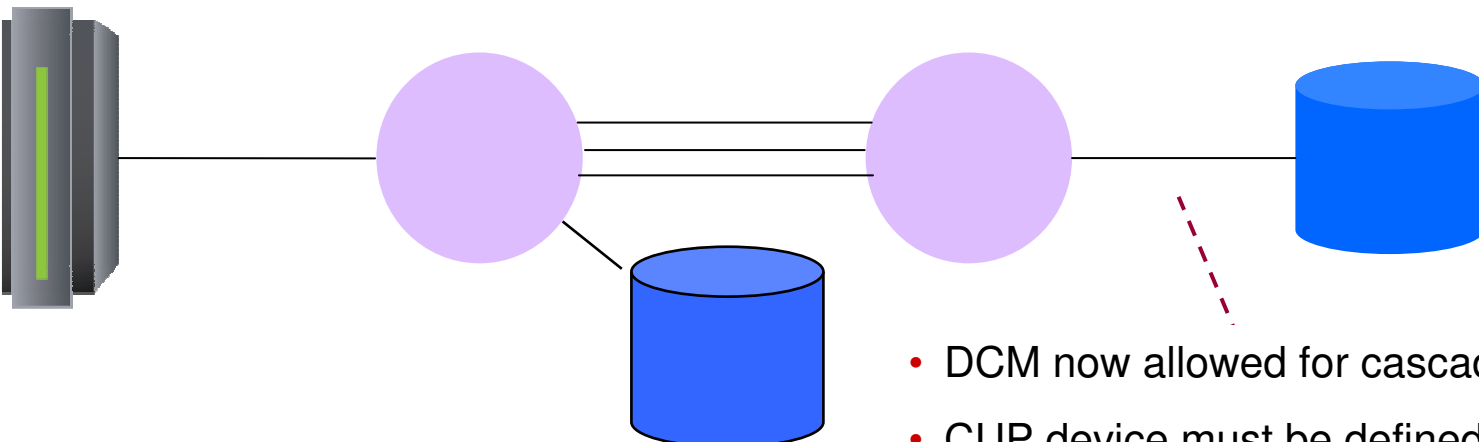
DCM Multi-Switch Cascade

Before z/OS 2.1



- DCM not allowed for CUs connected to exit switch in cascaded config

With z/OS 2.1



- DCM now allowed for cascaded
- CUP device must be defined on local and remote switches with 2 byte link addresses

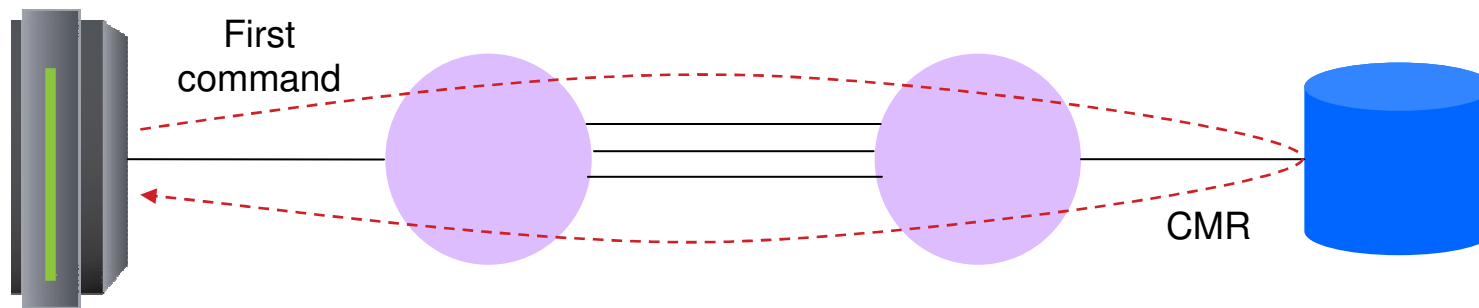
zEC12 Enhanced Channel Path Selection

- Overall I/O performance can be affected by poor performance on one channel path
 - See RMF I/O queuing report
- There are many possible causes:
 - Multi-system workload spikes
 - SAN congestion or destination port congestion
 - Control unit host adapter or port congestion
 - Firmware failures in the SAN, channel extenders, control units
 - Hardware failures (link speeds did not initialize correctly)
 - Mis-configuration and cabling errors
- Channel subsystem will now balance paths based on initial command response time

What is Initial Command Response Time?



- Initial command response (CMR) time is the amount of time from when the channel sends the first command until it gets a response from the control unit
 - One round trip through the fabric
 - Originally created to distinguish between IFCCs that occur because of congestion vs lost frames
 - Good for detecting fabric congestion and other problems on a path



zEC12 Enhanced Channel Path Selection



21:21:03 I= 40% DCM Group

Path	DCM	CTL	Units	MN	MX	DEF	LCU	Cont Rate	Del Lngth	Q	AVG CSS	CHPID Taken	%DP Busy	%CU Busy	AVG CUB	AVG CMR
5D		3000					010B					404.96	0.0	0.0	0.0	0.1
50		3000					010B					404.96	0.0	0.0	0.0	0.1
3D		3000					010B					404.95	0.0	0.0	0.0	0.0
C6		3000					010B					404.96	0.0	0.0	0.0	0.1
52		3000					010B					404.96	0.0	0.0	0.0	0.0
D1		3000					010B					404.96	0.0	0.0	0.0	6.5
95		3000					010B					404.96	0.0	0.0	0.0	6.1
C2		3000					010B					404.19	0.0	0.0	0.0	11
							010B	0.0	0.00	0.1		3238.9	0.0	0.0	0.0	2.9
5D		3100					010C					404.57	0.0	0.0	0.0	0.1
50		3100					010C					404.57	0.0	0.0	0.0	0.1
3D		3100					010C					404.55	0.0	0.0	0.0	0.0
C6		3100					010C					404.57	0.0	0.0	0.0	0.1
52		3100					010C					404.57	0.0	0.0	0.0	0.0
D1		3100					010C					404.56	0.0	0.0	0.0	6.5
95		3100					010C					404.56	0.0	0.0	0.0	6.1
C2		3100					010C					403.86	0.0	0.0	0.0	11
							010C	0.0	0.00	0.1		3235.8	0.0	0.0	0.0	2.9

Before

Substantial improvement in I/O rates

22:08:33 I= 57% DCM Group

Path	DCM	CTL	Units	MN	MX	DEF	LCU	Cont Rate	Del Lngth	Q	AVG CSS	CHPID Taken	%DP Busy	%CU Busy	AVG CUB	AVG CMR
5D		3000					010B					820.56	0.0	0.0	0.0	1.2
50		3000					010B					757.73	0.0	0.0	0.0	1.3
3D		3000					010B					998.91	0.0	0.0	0.0	1.1
C6		3000					010B					913.45	0.0	0.0	0.0	2.0
52		3000					010B					850.12	0.0	0.0	0.0	1.6
D1		3000					010B					370.59	0.0	0.0	0.0	1.6
95		3000					010B					576.98	0.0	0.0	0.0	1.2
C2		3000					010B					637.72	0.0	0.0	0.0	1.2
							010B	0.0	0.00	0.2		5926.1	0.0	0.0	0.0	1.4
5D		3100					010C					485.75	0.0	0.0	0.0	1.3
50		3100					010C					543.74	0.0	0.0	0.0	1.4
3D		3100					010C					1395.7	0.0	0.0	0.0	1.1
C6		3100					010C					1172.8	0.0	0.0	0.0	1.8
52		3100					010C					1237.1	0.0	0.0	0.0	1.4
D1		3100					010C					303.98	0.0	0.0	0.0	2.0
95		3100					010C					296.60	0.0	0.0	0.0	1.5
C2		3100					010C					246.21	0.0	0.0	0.0	1.3
							010C	0.0	0.00	0.3		5681.8	0.0	0.0	0.0	1.4

After



I/O Rate Health Check (OA40508)



```
CHECK (IBMIOS, IOS_IORATE_MONITOR)
START TIME: 04/22/2013 08:44:14.271360
CHECK DATE: 20120430 CHECK SEVERITY: MEDIUM
CHECK PARM: THRESHOLD(100), RATIO(2), XTYPE(), XCU()
```

IOSHC133I I/O Rate Report

The following control units show inconsistent I/O rates based on these parameters:

THRESHOLD = 100 ← I/Os per second to the CU must be at least this high
RATIO = 2 ← I/O rate ratio between the highest and lowest paths

```
I/O RATE EXCEPTION DETECTED AT: 04/22/2013 08:44:14.254730
CONTROL UNIT = 0500
ND = 002107.000.IBM.PK.000000000002
```

CHPID	ENTRY LINK	EXIT LINK	CU INTF	I/O RATE	AVG CMR	IOR EXC
14	B153	B177	0001	39.603	2.560	*
44	B353	B375	0012	101.38	2.112	
16	B055	B277	0013	98.019	2.134	
46	B154	B376	0104	50.693	2.048	

This is the exception path

Exception message appears in system log

•Medium Severity Exception *

```
IOSHC132E Analysis of I/O rates detected one or more control units with an exception.
```

Additional Reading

- *Evolution of the System z Channel* – Presentation given at Share August 2011 by Patty Driever and Howard Johnson
- *IBM zEnterprise 196 and IBM zEnterprise 114 I/O and FICON Express8S Channel Performance*
- *FICON DCM for System Programmers* - Whitepaper on Techdocs
- *IBM zEnterprise Storage I/O Advancements* - IBM Journal of Research and Development, Vol 56

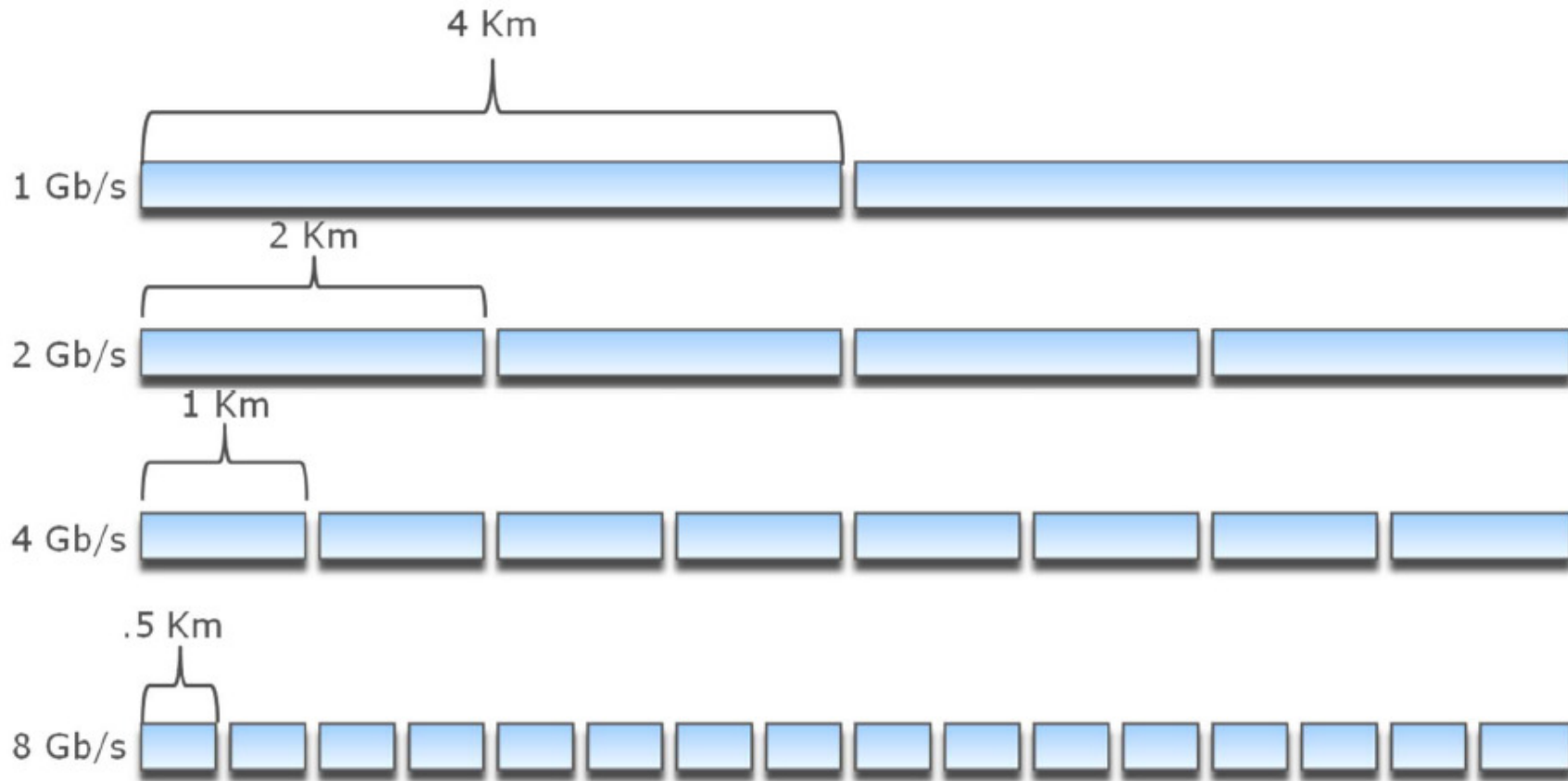
Thank you



Appendix

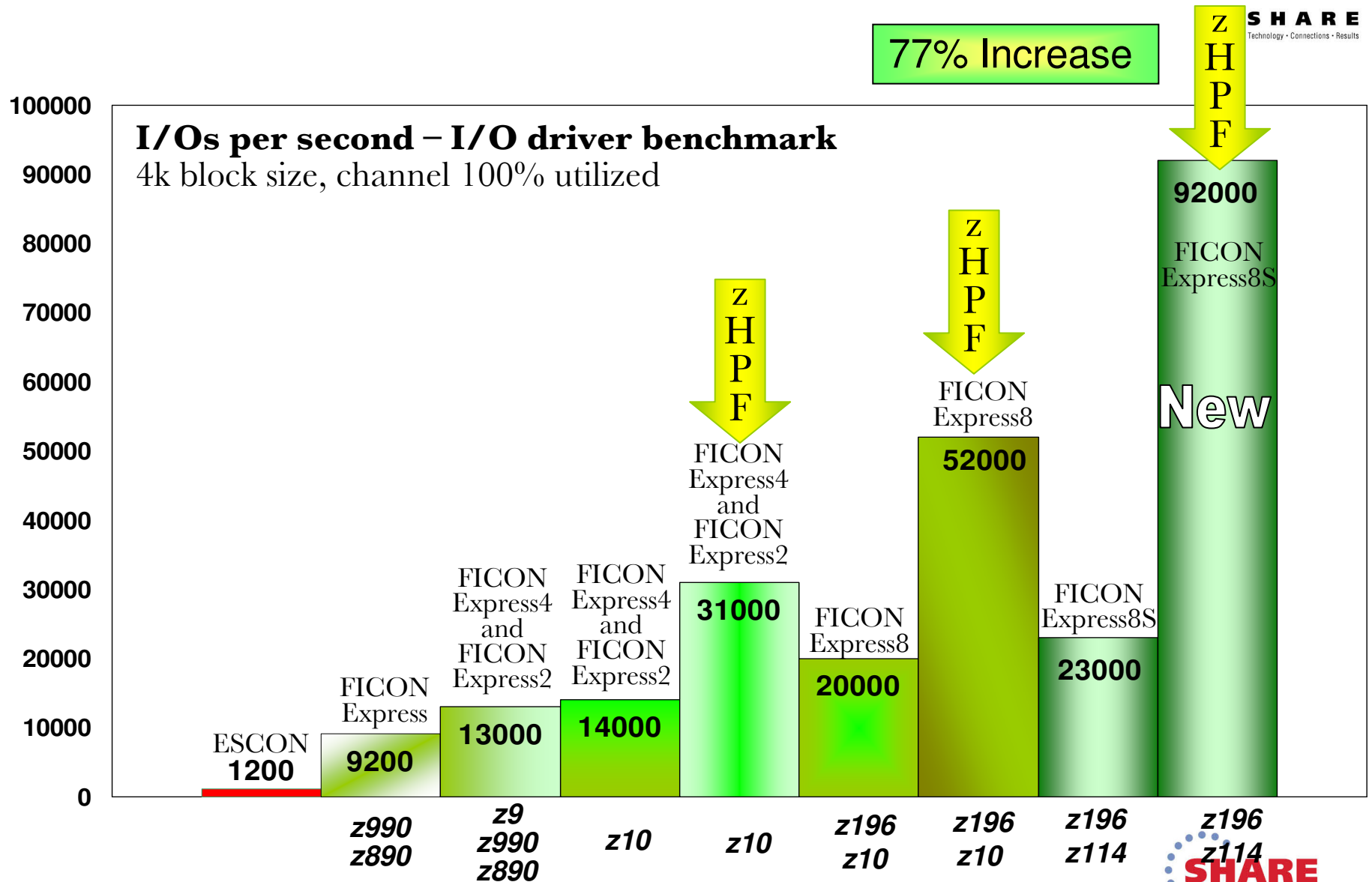
FICON Trivia

How long is a fiber channel frame?

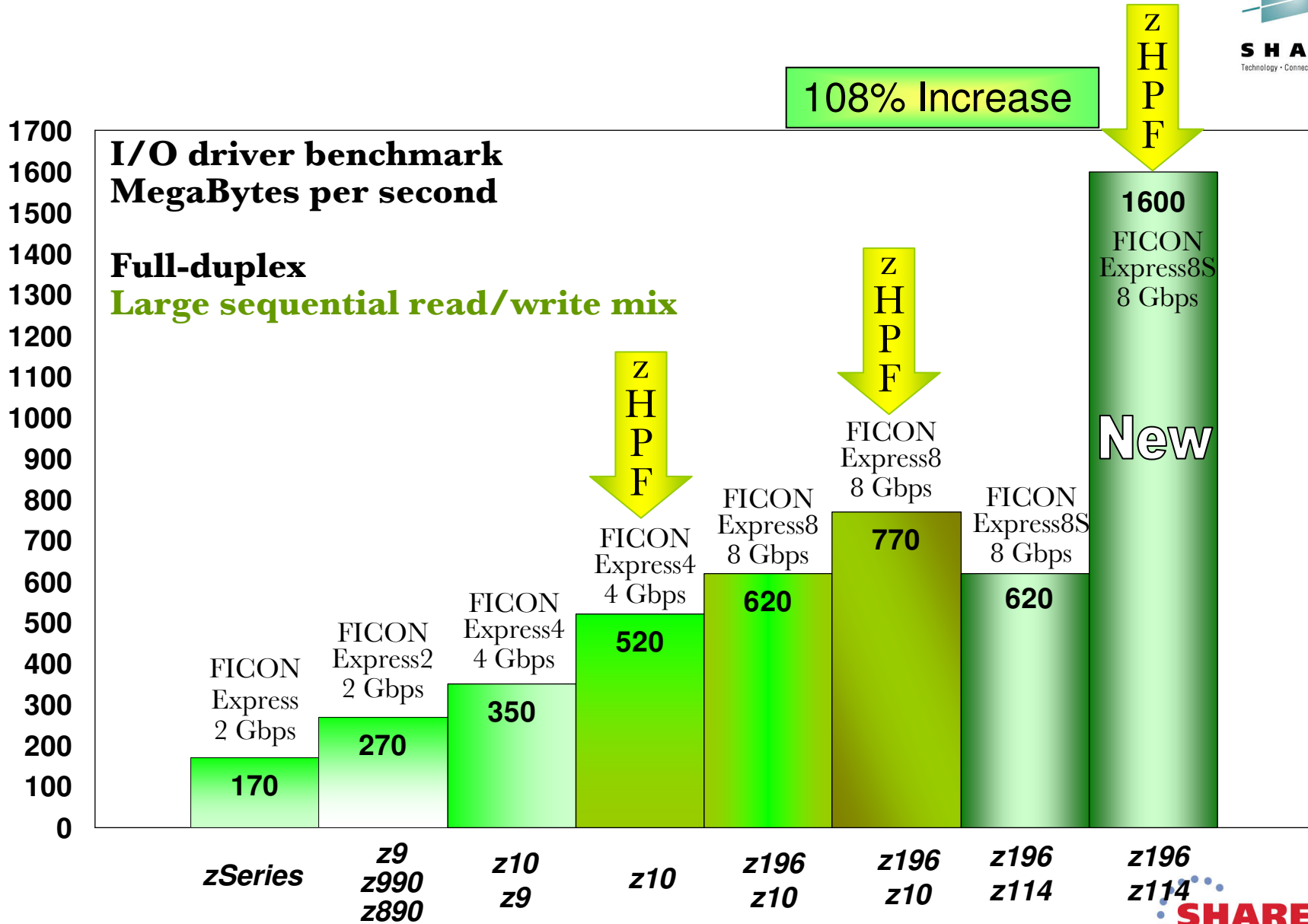


★ The fibre channel frame is the basic building block in FICON. A frame can have a payload up to 2112 bytes. Each byte is represented by 10 bits (8b/10b encoding).

zHPF - 4X More I/Os per Second than FICON I/O



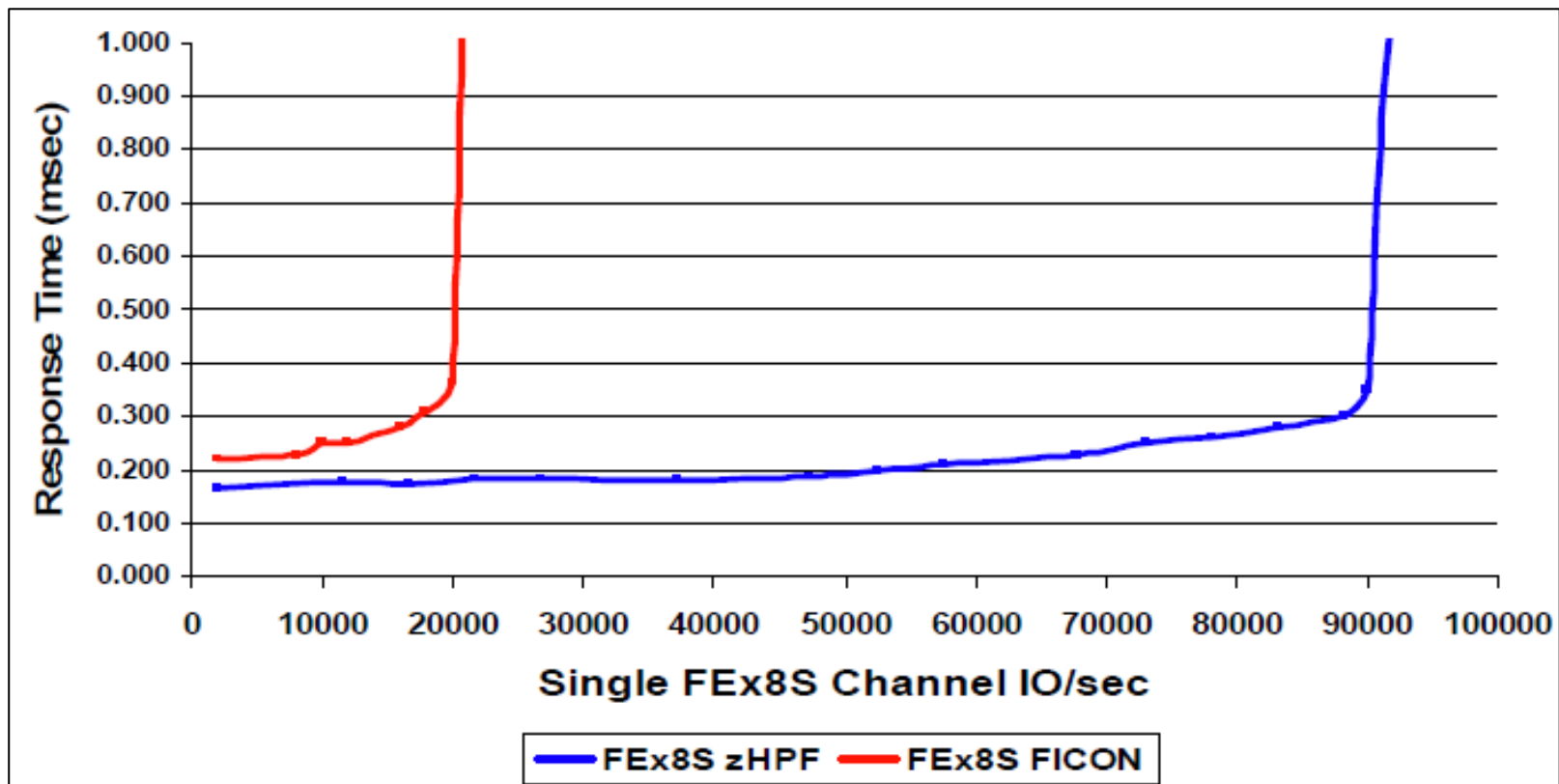
zHPF - More than 2X FICON Throughput



zHPF - Response Time Improves Too



Single FICON Express8S channel: zHPF vs FICON READ 4k bytes/IO
Total I/O Response Time vs IO/sec



CMR Health Check Report Example

```
CHECK (IBMIOS, IOS_CMRTIME_MONITOR)
START TIME: 12/10/2011 16:34:03.455536
CHECK DATE: 20100501 CHECK SEVERITY: MEDIUM
CHECK PARM: THRESHOLD (3) , RATIO (5) , XTYPE ( ) , XCU ( )
```

IOSHC113I Command Response Time Report

The following control units show inconsistent average command response (CMR) time based on these parameters:

THRESHOLD = 3

RATIO = 5

CMR TIME EXCEPTION DETECTED AT: 12/10/2011 16:29:24.212239

CONTROL UNIT = 25C0

ND = 002107.941.IBM.75.0000000WH391

CHPID	ENTRY LINK	EXIT LINK	CU INTF	I/O RATE	AVG CMR
81	2C51	2DC4	0030	72.330	9.21
22	3C1B	3DC2	0031	71.651	9.47
82	2C52	2DC0	0032	72.333	8.70
84	2C54	2DCC	0100	71.810	1.92
21	3C19	3DD2	0231	72.122	1.79

These are the exception paths

Exception message appears in system log

* Medium Severity Exception *

IOSHC112E Analysis of command response (CMR) time detected one or more control units with an exception.