



Managing z/VM and Linux Performance Best Practices

Mike Sine IBM

March 13, 2014 Session Number 14637









Special Notices and Trademarks

Special Notices

This presentation reflects the IBM Advanced Technical Skills organizations' understanding of the technical topic. It was produced and reviewed by the members of the IBM Advanced Technical Skills organization. This document is presented "As-Is" and IBM does not assume responsibility for the statements expressed herein. It reflects the opinions of the IBM Advanced Technical Skills organization. These opinions are based on the author's experiences. If you have questions about the contents of this document, please contact the author at sine@us.ibm.com.

Trademarks

The following are trademarks or registered trademarks of International Business Machines Corporation in the United States, other countries, or both.

IBM, the IBM logo, Candle, DB2, developerWorks, iSeries, Passport Advantage, pSeries, Redbooks, Tivoli Enterprise Console, WebSphere, z/OS, xSeries, zSeries, System z, z/VM.

A full list of U.S. trademarks owned by IBM may be found at http://www.ibm.com/legal/copytrade.shtml.

NetView, Tivoli and TME are registered trademarks and TME Enterprise is a trademark of Tivoli Systems, Inc. in the United States and/or other countries.

Microsoft, Windows, Windows NT, Internet Explorer, and the Windows logo are registered trademarks of Microsoft Corporation in the United States and/or other countries.

Java and all Java-based trademarks and logos are trademarks or registered trademarks of Sun Microsystems, Inc. in the United States, other countries, or both.

Linux is a trademark of Linus Torvalds in the United States, other countries, or both.

UNIX is a registered trademark in the United States and other countries licensed exclusively through The Open Group. Intel and Pentium are registered trademarks and MMX, Pentium II Xeon and Pentium III Xeon are trademarks of Intel Corporation in the United States and/or other countries.

Other company, product and service names may be trademarks or service marks of others.





AGENDA

- Introduction
- Monitoring Requirements
 - Virtual Linux and z/VM performance considerations
 - Don't forget the hardware
 - Integration from hardware systems applications Persistent historical views
- Enterprise Management
- Operational Requirements
 - Centralized Control
 - Including all Enterprise Virtual Machines
- Integrating Monitoring and Operations
- Bringing it all together





Virtual Linux servers have unique challenges versus running on physical machines.

- z/VM System Programmers and Linux Administrators may not be in the same organization.
- We find that it is easy to over allocate resources; therefore, our monitoring examines resource usage of hardware, hypervisor, as well as the virtual machine. Realtime and historical metrics demonstrate peaks periods as well as average runtimes.
- Cross-platform virtualization increases these challenges









AGENDA

- Introduction
- Monitoring Requirements
 - Virtual Linux and z/VM performance considerations
 - Don't forget the hardware
 - Integration from hardware systems applications Persistent historical views
- Enterprise Management
- Operational Requirements
 - Centralized Control
 - Including all Enterprise Virtual Machines
- Integrating Monitoring and Operations
- Bringing it all together





OMEGAMON XE on z/VM and Linux agents with ITM example





An Integrated Monitoring Approach

- Provides performance monitoring for z/VM and Linux guests
- Executes automated actions in response to defined events or situations (monitoring without automation is overhead)
- Integrates and Scales well across Enterprise for central control and trending:
 - Specifically focused on z/VM and Linux guests
 - Able to integrate z/VM and Linux into Enterprise Solution
 - Data warehousing for trend analysis (interactive and batch)



Workspaces to Manage z/VM and Linux



<u>z/VM</u>

- Processors
- SYSTEM Utilization, spinlocks
- Workload
 - Linux Appldata
 - Scaled & total CPU values
- LPAR Utilization
- PAGING and SPOOLING Utilization
- DASD
- Minidisk Cache
- Virtual Disks
- Channels
- CCW Translation
- REAL STORAGE Utilization
- NETWORK Utilization (Hiper Socket and Virtual Switch)
- TCPIP Utilization Server
- TCPIP Utilization Users
- Resource Constraint (Wait states)
- System Health

<u>Linux</u>

- Linux OS
- System Information
 - CPU aggregation
 - Virtual Memory Statistics
- Process
- Users
- Disk Usage
- File Information
- Network





Have I allocated enough Virtual CPUs to my guest?

- Do not define more virtual CPUs for a Linux guest than are needed.
 - The use of more than one processor requires software locks so that data or control blocks are not updated by more than one processor at a time.
- Linux makes use of a global lock, and when that lock is held, if another processor requires that lock, it spins.
 - Set the number of virtual processors based on need and not simply match the number of real that are available.
 - Careful when cloning as some Linux guests require more Virtual CPUs (ex: Running Websphere, Oracle) than others.



Aggregate monitoring of Virtual CPUs





•••• in Anaheim



z/VM Processor Utilization

- **Total Processor Utilization** This is the processor utilization from the VM perspective and includes CP, VM System, and Virtual CPU time.
- **System Time:** This is the processor time used by the VM control program for system functions that are not directly related to any one virtual machine. This should be less than 10% of the total.
- **CP Processor Time:** This is the processor time used by the VM control program in support of individual virtual machines.
- Virtual Processor Time: (Emulation Time): This is processor time consumed by the virtual machine and the applications within it.
- **Total to Virtual Ratio** The ratio of total processor time to virtual processor time is often used as an indicator of z/VM efficiency or overhead. The closer to 1.0, the better the z/VM efficiency. RoT: Should explore causes of a ratio over 1.30.





System Processor Utilization Workspace



Complete your session evaluations online at www.SHARE.org/AnaheimEval

in Anaheim

z/VM Workload Workspace









Spin Lock Wait

• Time Spinning on Locks Percent:

- The percentage of time processors spend spinning on formal spin locks. RoT: Should be less than 10%.
- Increases as number of logical processors increases.



Spinlock Workspace



alony - Connections - Result 🐮 npmipsvt3.tivlab.raleigh.ibm.com - Remote Desktop - 🗆 × I OP 5 LOCKS per Second TOP 5 LOCKS by Duration 😋 view. [Physical] ШШ 1 шнц 0 1 Exclusive Spin Lock Rate per Second Exclusive Time Spinning on Locks in Microseconds - 🔲 **CP** Owned Devices * Shared Spin Lock Rate per Second Shared Time Spinning on Locks in Microseconds DASD LPAR 60 3.0 Network -Real Storage 50 System Rate M 2.0 40 TCPIP per 🔲 Workload 30 Second 🖻 – 🚂 vmlnx11:VL spuc 1.0 E- 1 z/VM Linux Systems 20 Channel 10 **CP** Owned Devices DASD 0.0 SRMSLOCK SRMSLOCK SRMATDLK HCPTRQLK RSAAVCLK RSA2GCLK SRMATDLK HCPTRQLK RSAAVCLK RSA2GCLK LPAR - Network Real Storage Ŧ Lock Name Lock Name 🕰 Physical Spin Lock Activity Top 5 Percent of Elapsed Time in Spin L. Total Exclusive **Exclusive Time** Exclusive Time Spinning Time Spinning Exclusive Time Spinning on Locks Percent LPAR. Spin Lock Lock Spin Lock Spinning Time Spinning v on Locks on Locks Shared Time Spinning on Locks Percent on Locks Name Name Calls Rate on Locks in Microseconds Percent per Second per Second in Microseconds Percent VIC SRMSLOCK 2.57 2.57 13.7 0.00 13.7 0.00 VIC SRMATDLK 51.7 1.77 0.00 51.7 1.77 0.00 VIC HCPTRQLK 5.9 0.59 0.00 5.9 0.59 0.00 Percent VIC SRMALOCK 0.0 0.00 0.00 0.0 0.00 0.00 VIC HCPPGDPL 0.0 0.00 0.00 0.0 0.00 0.00 VIC 0.0 0.00 0.00 0.0 0.00 HCPPGDSL 0.00 VIC HCPPGDTL 0.0 0.00 0.00 0.0 0.00 0.00 VIC HCPPGDAL 0.0 0.00 0.00 0.0 0.00 0.00 0 VIC RSAAVLLK 0.0 0.00 0.00 0.0 0.00 0.00 BUTDLKEY SRMATDLK HCPTMFLK SAAVOLK SA2GCLK VIC RSACALLK 0.0 0.00 0.00 0.0 0.00 0.00 VIC SYSDATLK 0.0 0.00 0.00 0.0 0.00 0.00

0.0

0.00

0.00

0.0



0.00

0.00

Complete your session evaluations online at www.SHARE.org/AnaheimEval

Lock Name

VIC

1

DCTLLOK



Is my Linux guest sized correctly?

- In general, do not define the Linux virtual machine larger than you need.
 - Excessive virtual machine sizes negatively impact performance.
 - Linux uses any extra storage for caching of data. For shared resources, this is an impact.
 - Reduce the size of the Linux guest until it starts to swap (use VDISK for swap).
 - Monitor swap usage.



Need breakdown of memory use



•••• in Anaheim



Page/Swap Attributes



•••• in Anaheim





VDISK

- What is it?
 - FBA (Fixed Block Architecture disk) device emulated inmemory
 - Translation: Very fast "device".
 - High performance paging device for Linux on z.
 - Memory is allocated by CP from the Dynamic Paging Area
 - Allocated only when referenced
 - Allocating a 10 MB device does NOT instantly consume 10 MB of pages.
 - Pages are allocated when needed.
 - Not recommended in a storage-constrained z/VM system.



VDISK Workspace



Technology - Connections - Recults VDISK - KYASH3 - SYSADMIN _ 8 × File Edit View Help 🖽 📧 🎠 🚸 🗿 🔽 🖶 🛑 🔘 🔾 🈂 🕋 🌐 🗞 📶 🙋 🙋 😂 🛄 📓 🗎 👰 🖓 🥑 🖅 🐚 💽 🔥 🎫 30 (1) ⟨¬ → ¬ 🚭 Navigator * 🗉 🖂 I Top 5 Paging Rates per Second / [] 8 🗆 × 📶 Top 5 Expanded Storage Paging Rate... 🥒 🔟 🖯 🗙 View: Physical æ ÷ 🌮 Pages Read from DASD per Second Pages to Central Storage per Second. * Pages to DASD per Second Pages Stolen per Second Pages Written to DASD per Second Pages from Central Storage per Second 🗄 📲 vmlnx11:VL 🔄 👫 z/VM Linux Systems Rate Rate per Second Channel per Second **CP** Owned Devices DASD ۲ LPAR 0 ACKERK Network ACKERK ANGELOM AVATAR BIGANG BRIANKT ANGELOM AVATAR BIGANG BRIANKT Real Storage System - 0700 - 0700 - OF00 TCPIP . 0298 7777 OFOC 0290 1111 Workload 0700 0700 -🕰 Physical VDISK Owner & Device Number VDISK Owner & Device Number / [] 8 🗆 × / ± 0 8 8 × II Top 5 Pages in Use 🔲 Virtual Disk Activity (A) Page: 1 of 2 Resident Pages Pai Pages Locked Pages LPAR Number Virtual I/O's System Device VDISK Owner VDISK Size Stolen Time fro Occupied Slots ID. Name Number of Links per Second per Second pe 🚖 XSTORE Pages 04/06/09 23:35:51 GDLVM7 GDLVM7 ACKERK 0299 100.000 0.00 0.00 1 . 04/06/09 23:35:51 GDLVM7 GDLVM7 ANGELOM 0700 7,000,000 1 0.00 0.00 1,600,000 Page 1.200.000 04/06/09 23:35:51 GDLVM7 GDLVM7 AVATAR 1111 4,000,000 1 0.00 0.00 800,000 Count 04/06/09 23:35:51 BIGANG GDLVM7 GDLVM7 0700 7.000.000 1 0.00 0.00 400,000 04/06/09 23:35:51 GDLVM7 1,440,000 GDLVM7 BRIANKT 0F00 1 0.00 0.00 04/06/09 23:35:51 GDLVM7 GDLVM7 CORAKR 05FF 10,000,000 0.00 0.06 1 EDLWRK14 - 05FA EDL/WRK8 EDLWRK22 EDLWRK23 EDLWRK2 04/06/09 23:35:51 GDLVM7 GDLVM7 CORAK2 05FF 20,000 1 0.00 0.00 04/06/09 23:35:51 GDLVM7 GDLVM7 CRASTDA 0999 4,000,000 1 0.00 0.01 0.00 0.00 04/06/09 23:35:51 GDLVM7 GDLVM7 DENISE 1111 4,000,000 1 - 05FA - 05FB 05FA 0.00 05FA 04/06/09 23:35:51 GDLVM7 GDLVM7 DENISE 020E 5,000,000 1 0.00 04/06/09 23:35:51 GDLVM7 GDLVM7 DENISE2 4,000,000 0.00 0.00 1111 1 VDISK Owner & Device Number ▶ ∓ Hub Time: Mon. 04/06/2009 11:38 PM 😳 Server Available VDISK - KYASH3 - SYSADMIN





Memory Configuration

- Plan on a virtual to real (V:R) memory ratio in the range of 1.5:1 to 3:1.
- z/VM's architecture still benefits from expanded storage:
 - Serves as high speed cache.
 - Increases consistency of response time.
 - See http://www.vm.ibm.com/perf/tips/storconf.html for the gory details.
- Rule of Thumb start with 20-25% of memory configured as expanded:
 - The lower the paging rate, the lower the amount of expanded storage required.
 - The greater the number of page frames available in central storage above 2GB, the higher the amount of expanded storage required.
 - Some workloads 2–4GB of expanded storage is sufficient, 1GB minimum. However, more and more Linux systems are running heavy workloads and the 20-25% rule still applies.
 - XSTORE obsolete for z/VM 6.3



OMEGAMON Memory Configuration









Paging Subsystem

- Plan for DASD page space utilization < 50%:
 - Page space tends to get fragmented over time.
 - Large contiguous free space allows for greater paging efficiency.
 - Monitor usage with OMEGAMON XE or Q ALLOC PAGE command.
- Do not mix page space with any other space on a volume.
- Recommend using devices of the same size/geometry.
- Calculation guidelines are located in the CP Planning and Administration Manual.
- z/VM 6.3 block paging changes eliminates the benefits of 50% page space, making monitoring even more important.



Planning for Large Memory

- Normal best practices for migrating from an earlier release still apply.
- Change your paging XSTORE into central
 - XSTORE provided an aging function. It helped catch reclaim selection "mistakes".
 - The new *IBR* concept and *global aging list* provide the same function but do so more efficiently in central.
- Plan enough DASD paging space
 - The system now pre-writes pages to DASD.
 - See z/VM manuals for detail recommendations.
- Plan a robust paging DASD configuration
 - Use plenty of paging volumes
 - Make the volumes all the same size
 - Put only paging space on the volumes you use for paging
 - Spread the paging volumes through your logical control units
 - Avoid logical control units that you know are hot on application I/O
 - Use plenty of chpids
 - Do not use ESCON chpids
 - Do not mix ECKD paging and SCSI paging
 - Leave reserved slots in the CP-owned list





OMEGAMON CP Owned Devices – Paging Subsystem



SHARE Technology - Connections - Results



z/VM Page Attributes





Minidisk Cache

- z/VM minidisk cache is a write-through cache:
 - Improves read I/O performance.
 - But it's not free.
- Not recommended for:
 - Memory constrained systems.
 - Linux swap file disks.
 - Flashcopy targets (see next chart)
- Default system settings are less than optimal.
- Recommended settings:
 - Eliminate MDC in expanded storage.
 - SET MDC XSTORE 0M 0M
 - Limit MDC in central storage 10% is a good starting point.
 SET MDC STORE 0M 256M
 - Monitor with product like OMEGAMON XE and/or the Q MDC command.







MDC and FlashCopy Interaction

- FlashCopy requests require z/VM to flush MDC for the entire minidisk.
- MDC Flush processing is very expensive even when there is no data in MDC to flush
 - System Time becomes very high.
- z/OS DFSMS and other utilities can make extensive use of FlashCopy for functions such as defragmentation
- Mitigations
 - Turn off MDC for minidisks that are FlashCopy targets



OMEGAMON MDISK Cache Allocations



Technology - Connections - Results Minidisk Cache - KYASH3 - SYSADMIN _ 8 × File Edit View Help 🌐 😡 📶 🙋 🔁 🔩 🛄 📓 🗎 👰 🖓 🔮 🖅 🐚 🔘 🔥 🎫 🎒 🖽 📧 光 🚸 🗿 🖻 🖶 0 <₽ • ⇒ 🔘 🈂 🌰 ٩ 30 -± II 🖂 Main Storage Frames Ravigator Cache Requests / [] 🖯 🗆 × / [] 8 🗆 × View: Physical -凶 Minidisk Cache Read Rate Actual Frames Below 2G * Invalid Request Rate Ideal Frames E- 2/VM Systems ⊕ ____perfkit5:VL Block Invalid Rate Actual Frames Above 2G 🖻 🖟 vmlnx11:VL 🔄 🖳 z/VM Linux Systems 1.0 14,000 Channel 12,000 CP Owned Devices DASD 10,000-Rate per Sec LPAR Fra 8,000 Network mes Real Storage 6,000 System 4,000 . TCPIP 🔲 Workload 2,000 🝓 Physical / ¥ 🛛 🖯 🗆 × / [] 🖯 🗆 × 🔲 Minidisk Cache Activity II Cache Age I Cache Expanded Storage Partition Max Cache Actual Cache Fair Share Min Block Life Ideal XSTORE System Size Size Size Size Avg XSTORE Age Actual XSTORE ID in Blocks in Blocks in Blocks in Blocks р 5,000 -WLAVMXA. 14336 4096 2660 2048 4,000 3,000 Seconds Blocks 2.000-1,000-. Þ 🕒 Hub Time: Thu, 07/23/2009 06:58 PM Minidisk Cache - KYASH3 - SYSADMIN 😲 Server Available





OMEGAMON MDISK Cache Allocations – p. 2

📑 Minid	🖷 🛉 Minidisk Cache - KYASH3 - SYSADMIN														
<u>F</u> ile <u>E</u> d	<u>F</u> ile <u>E</u> dit <u>V</u> iew <u>H</u> elp														
< → → → ̄□ 🔚 🖽 🐻 🎠 ◈ ╣ 🗊 🛱 │ ● Ο 🍣 番 ﴿ │ 🎟 🗞 🖬 🖉 🖍 🖾 🖾 🖾 🖾 🗟 🗐 두 🥏 ⋮ ⊅ 🖆 ● 🔥 ≡ 🛎 │															
III Minidisk Cache Activity															
Block /alidates r Second	Full Read Hit Percent	ldeal Frames	Actual Frames Below 2G	Actual Frames Above 2G	Minimum Storage Frames	Maximum Storage Frames	Pages Deleted per Second	Steal Invoked per Second	MDC Bias	ldeal XSTORE in Blocks	Actual XSTORE in Blocks	Minimum XSTORE in Blocks	Maximum XSTORE in Blocks	XSTORE Page: Deleted per Second	s XS ⁻ pt
0.00	100.00	12288	5057	6306	2048	12288	0.00	0.00	1.00	4096	3928	1024	4096	0.0	0
	\smile				-							-			





Direct Access Storage Devices (DASD)

- Avg Pending Time for DASD
 - Average pending time for real DASD I/Os. RoT: Should be less than 1 millisecond.
- Items worth keeping an eye on:
 - Number of I/O's per Second, Percent Busy
 - Avg Service Time Average service time for real DASD devices (sum of the pending, connect, and disconnect times).
 - **DASD I/O Rate** Rate of traditional real I/Os per second to real DASD devices. Worth monitoring.



DASD I/O Workspace



Technology - Connections - Results 🖬 DASD - KYASH3 - SYSADMIN _ 8 × File Edit View Help 🌐 😡 📶 🙋 🔁 🔩 🛄 📓 🗎 👰 🖓 🔮 🖅 🐚 🔘 🔥 🎫 🎒 1 🔲 🖽 📧 光 🚸 🗿 🖻 🖶 0 0 😂 🍈 1 30 ← → → → Revigator 1 II II / [] 🖯 🗆 × / [] 🖯 🗆 × II Top 5 Device Busy II Top 5 I/O Rate View: Physical -凶 Percent Busy Number IO per Seconds -E- a z/VM Systems 🗄 --- 🚂 perfkit5:VL 100 -🗄 --- 🚂 vmlnx11:VL 🔄 📲 z/VM Linux Systems 80-Channel CP Owned Devices 60 DASD 40 LPAR Network Real Storage 20 System TCPIP n -530PRD -530PRD TVM605 **TVM806** TVM605 **TVM806** VMCDRF VMCDRF VMCDRB VMCDRB Workload -🝓 Physical 📶 Top 5 Servi... 🖉 🔟 🖯 🗙 / \$ 00 8 8 × 📶 Top 5 1/0... 🖉 🔟 🖯 🗶 DASD I/O Activity 8 Page: 1 of 2 Connection Time Average Queued IO Volume Average Average Number Average Average Disconnect Time Connection Percent Device Device Serial Queued V Service Disconnect F 10 Average Pending Time Address Type Time Busy Number 10 Time per Second Time * VM54SP 5A1A 3390 0.60 0.00 0.90 0.00 0 3 . VM54RS 5AE9 3390 0.50 0 0.00 0.80 0 0.00 VMSL51 5A57 3390 0.40 0 0.00 0.70 0 0.00 0 0.30 0.00 0.70 0.00 VMSL54 5A5A 3390 0 VMSL50 3390 0.30 0 0.00 0.70 0 0.00 5A56 VM53PA 5A08 3390 0.40 0 0.00 0.70 0 0.00 0 VMCD02 5A04 3390 0.40 0.00 0.70 0 0.00 0.40 (VMCD02) 0.30 0 0.70 VMSL53 5A59 3390 0.00 0 0.00 0.30 0 VMCD05 5A3A 3390 0.00 0.60 0 0.00 0.0 -ZVT105 - TVM805 - TVM806 -530PRD VM54SP VMCDRB VMSLBA VMCDRP VM54R9 VMCD02 VMSLHC 5A39 3390 0.30 0 0.00 0.60 0 0.00 VM54GS 5A35 3390 0.30 0 0.00 0.60 0 0.00 4 ▶ **∓** Hub Time: Fri, 07/24/2009 12:06 PM 😳 Server Available DASD - KYASH3 - SYSADMIN



System Dump & Spool Space



- Dump Space
 - Ensure there is sufficient dump space defined to the system.
 - Dump space requirements vary according to memory usage.
 - Q DUMP identifies allocated dump space.
 - Calculation guidelines are located in CP Planning and Administration Manual.
- Spool Space
 - Various uses:
 - User printer, punch, reader files (console logs)
 - DCSS, NSS
 - System files
 - Page space overflow
 - Spool Management:
 - Monitor with OMEGAMON, Operations Manager, Q ALLOC SPOOL cmd
 - SFPURGER utility:
 - Rule based tool to clean up spool space.
 - Included in the no charge CMS Utilities Feature (CUF).



VMDUMP Processing Concern



- VMDUMP is a very helpful command for problem determination.
- Some weaknesses:
 - Does not scale well, can take up to 40 minutes per GB.
 - It is not interruptible
 - APAR VM64548 is open to address this.
- Linux provides a disk dump utility which is much faster relative to VMDUMP.
 - It is disruptive
 - Does not include segments outside the normal virtual machine.
- See <u>http://www.vm.ibm.com/perf/tips/vmdump.html</u>
- If you increase central, make sure you also increase dump space
 - More guidance available on <u>www.vm.ibm.com/techinfo/</u>
 - Download updated "Allocating Space for CP Hard ABEND Dumps"



System Dump & Spool Space







Tips—Overall Health of Your System



SHARE Technology - Connections - Results

At a quick glance you can see the %CPU usage, what your overcommit ratio is, the number of users in a wait state, and paging rates of all your z/VM systems






30

Resource Constraint Analysis (Waits)







Do not ignore the hardware!

- Just because Linux resources are virtual, do not ignore the hardware!
 - Hardware is another potential layer of shared resources.
 - LPAR weight, CPU sharing, LPAR load, and other attributes need to be monitored for overall system performance.
 - The measurement should include the entire CEC and not just the LPAR hosting z/VM.



Processors

- Logical Processors
 - LPAR recommendation no greater than a 4:1 logical to real ratio.
 - z/VM 5.1 z/VM 5.2 support up to 24 processors.
 - z/VM 5.3 z/VM 6.x support up to 32 processors.





LPAR Utilization Workspace



Technology - Connections - Results LPAR - KYASH3 - SYSADMIN _ 8 × File Edit View Help 🌐 🗞 💷 🖉 🜊 🛄 📓 🗎 👰 🖓 🖅 🐚 💽 🙏 🎫 🎒 🖽 🖾 🏹 🧇 <mark>8</mark>1 0 🔁 ← → → → ? Ē 1 30 🝓 Navigator * -/ [] 🖯 🗆 × / [] 🖯 🗆 × IL LPAR Busy IL LPAR Load -凶 Ć, View: Physical LPAR Busy LPAR Load * E- 2/VM Systems Physical CPU Busy 😥 🕞 perfkit5:VL 801 🔄 🚡 vmlnx11:VL 100 -🗄 損 z/VM Linux Systems 60-80-Channel CP Owned Devices percent percent • 60 40-DASD LPAR 40 . Network 20-20 Real Storage System n TCPIP CANVM1 RALHCD RALHOD **CANVM1** CANSYSA RALNSBO CANSYSA RALNSBO TIVVMTO TIVVMTO. 🔲 Workload 🝓 Physical II LP... / □ ⊟ □ × 📶 Parti... 🖊 🔟 🖯 🗶 / ¥ 🛛 🖯 🗆 × 🛄 LPAR Utilization Total L LPAR LPAR LPAR Busy LPAR Processor LPAR LPAR LPAR LPAR LPAR Weight LPAR Suspended Time LPAR Busy Su Number Name Percent Weight Type Status Load CPU Capped Percent Т 120 2 NO 1 CANSYSA 9.40 100.00 CP 18.80 ACTIVE 2.70 6.25 114.00 CP 12.50 ACTIVE* 1.80 2 NO 2 CANVM1 100 5 RALHCD 0.00 0.00 0.00 INACTIVE 0.00 1 Unknown Unknown 80 3 RALNS60 100.00 DED IFL 500.00 ACTIVE 71.40 5 NO verage 1 NO 4 TIVVMT01 0.20 5.00 CP 0.20 ACTIVE 0.00 60-40-20 0.0 CANVM1 CANVMI RALHOD RALNS60 RALHOD TIVVMT01 CANSYS/ TIVVMTO: CANSYS/ 4 + Hub Time: Fri, 07/24/2009 11:05 AM 🙄 Server Available LPAR - KYASH3 - SYSADMIN





Processor by LPAR name workspace



Technology · Connections · Results





CPC workspace



📃 CPC - ZVMCVT04	- SYSADMIN	*ADMI	N MODE*															
<u>File Edit View H</u>	telp		3 0 0 1	ាក	<u> </u>		a II. 🕾 I	a 💿 🥅		Ch 🗑 🔗								
									Processor Busy by Count of Processors									
						busy'i creciik												
(inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8 (inux-koa8	EVL VL inannel P Owned Dev ASD PAR Stwork sal Storage rstem CPIP orkload /L inux System	s rices	nysiitiai				100 90 90 90 90 90 90 90 90 90 90 90 90 9		Scaled	I LPAR Busy P	eroent ad Percent			Available CPU Capacity	2,800 2,400 1,600 800 400 0		Available CPU Capacity	
Physical	Channel Channel				-	CP TP No CP NP Processor Type			다. 구 별 사용 다 우 우 우 Processor Type									
Processor Utiliza	tion by CPC a	and Proce	essor Type															/ ¥ 🛛 🖯 ×
Time	System ID	LPAR Name	Processor Type	Number of Processors	Total LPAR Weight	Dedicated Processor Count	LPAR Busy Percent	LPAR Overhead Percent	Overhead not Charged to an LPAR Percent	Overall LPAR Busy Percent	Scaled LPAR Busy Percent	Scaled LPAR Overhead Percent	Scaled Ov not Cha to an LF Perce	erhead rged PAR nt	Scaled Overall LPAR Busy percent	Available CPU Capacity	System Name	
07/27/12 08:45:12	GDLEST3	EST3	CP	42	210	12	1184	105	225	1514	28	2		5	36	2686	LXI01500.endicott.ibm.com:VL	
07/27/12 08:45:12	GDLEST3	EST3	zaap	2	30	0	0	0	0	0	0	0		0	0	200	LXI01500.endicott.ibm.com:VL	
07/27/12 08:45:12	GDLEST3	EST3	IFL	16	175	0	391	118	358	867	24	7		22	54	733	LXI01500.endicott.ibm.com:VL	
07/27/12 08:45:12	GDLEST3	EST3	ICF	1	50	0	0	0	27	27	0	0		27	27	73	LXI01500.endicott.ibm.com:VL	
07/27/12 08:45:12	GDLEST3	EST3	zIIP	3	50	0	0	0	1	1	0	0		0	0	299	LXI01500.endicott.ibm.com:VL	



History On-Demand with Persistent Historical Views



This makes it easier to see anomalies, or match spikes. Capturing performance data as a base line is a must:

- General history data business as usual.
- Detailed raw monitor data prior to and following any major changes.
- Ability to review attributes of a past incident through the enterprise view!
- On-Demand through the Portal or Batch



On-Demand: Persistent Historical Views



6 Mon, 02/09/2009 05:45 PM - Mon, 02/16/2009 06:30 PM

30

Historical Navigation Mode: Mon, 02/09/2009 05:45 PM - Mon, 02/16/2009 06:30 PM 🜔 Hub Time: Mon, 02/16/2009 08: 🌎 Server Available 🚽 System Information (Superseded) - hasle330.wsclab.washington.ibm.com - Mike Sine 🛛 *ADMIN MODE*

Pages Swapped In Pages Swapped in persec



/ 1 日

On-Demand: Persistent Historical Views





Max and Avg CPU example:





2011-06-30 2011-07-02 2011-07-04 2011-07-06 2011-07-08 2011-07-10 2011-07-12 2011-07-16 2011-07-18 2011-07-20 2011-07-22 2011-07-24 2011-07-26 2011-07-28

Legend:

Max_CPU_Percent: Avg_CPU_Percent: Mean of the Maximum: Mean of the Averages: AVG_Main_Memory_Util: AVG_Cache_Used: AVG_Page_Alloc_Rate: AVG_Swap_Used: Maximum CPU for the day as a percent of the number of virtual CPUs Average CPU for the day as a percent of virtual CPUs 30 day average for Maximum CPU percentages 30 day average for the average CPU percentages Average main memory utilization for the day as a percent Average size of memory used to cache buffers in megabytes Average number of pages obtained from available list in 4 kilobyte pages per second The percent of swap space used.



Avg Linux Memory breakdown example:









Tivoli Common Reporting (TCR)

- TCR reports available on the OPAL website
 - http://www-18.lotus.com/wps/portal/topal
- What is TCR?
 - Tivoli Common Reporting.
 - Consistent approach to viewing and administering reports.
 - Cognos based.
 - Flexible development environment (Eclipse based) for creating report definitions.
 - Five templates provided for download.
 - Taking suggestions for more





Sample Reports Available

- z/VM VM System CPU Utilization
- z/VM VM System Paging Utilization
- z/VM Linux System CPU Utilization
- z/VM VM System CP-Owned Device Utilization
- z/VM VM System TCP Server Statistics



SHARE Technology - Connections - Results

Tivoli

Ly VIII Dystelli Gi o otilizatioi	z/VM	System	CPU	Uti	lization
-----------------------------------	------	--------	-----	-----	----------

Report Period	All	Significant Resources Selected	5
Start Date	Dec 31, 1969 12:00 AM	End Date	Nov 30, 2007 11:59 PM
System ID	All	LPAR Name	All

IBM.

LPAR Busy



Available Summarization Time Periods:	
Hourly	
Daily	
Weekly	
Monthly	
Not Summarized Data	

LPAR Name	LPAR Busy	LPAR Load	LPAR Suspend	LPAR Overhead	Date/Time	
80770804307075782008	100000000000000000000000000000000000000	100000000000000000000000000000000000000	Time	Time	221212220022	



Complete your session evaluations online at www.SHARE.org/AnaheimEval



IBM.

2 / 18

Suctom - TIV	MV66				
LPAR Name	LPAR Busy	LPAR Load	LPAR Suspend Time	LPAR Overhead Time	Date/Time
RALNS31	100	4.2	0	.6	Nov 29, 2007 4:02 PM
RALNS32	100	4.2	0	.6	Nov 29, 2007 4:02 PM
RALNS61	100	4.2	0	.6	Nov 29, 2007 4:02 PM
TIVMV51	100	2.09	0	.6	Nov 29, 2007 4:02 PM
TIVMV510	100	2.09	0	.6	Nov 29, 2007 4:02 PM
RALNS31	100	4.2	0	.6	Nov 29, 2007 4:08 PM
RALNS32	100	4.2	0	.6	Nov 29, 2007 4:08 PM
RALNS61	100	4.2	0	.6	Nov 29, 2007 4:08 PM
TIVMV51	100	2.09	0	.6	Nov 29, 2007 4:08 PM
TIVMV510	100	2.09	0	.6	Nov 29, 2007 4:08 PM
RALNS31	100	4.2	0	.6	Nov 29, 2007 4:22 PM
RALNS32	100	4.2	0	.6	Nov 29, 2007 4:22 PM
RALNS61	100	4.2	0	.6	Nov 29, 2007 4:22 PM
TIVMVS1	100	2.09	0	.6	Nov 29, 2007 4:22 PM
TIVMV510	100	2.09	0	.6	Nov 29, 2007 4:22 PM
RALNS31	100	4.2	0	.6	Nov 29, 2007 4:37 PM
RALNS61	100	4.2	0	.6	Nov 29, 2007 4:37 PM
TIVMV51	100	2.09	0	.6	Nov 29, 2007 4:37 PM
TIVMVS10	100	2.09	0	.6	Nov 29, 2007 4:37 PM
RALNS32	99.8	4.2	0	.6	Nov 29, 2007 4:37 PM
RALNS31	100	4.2	0	.6	Nov 29, 2007 4:52 PM
RALNS32	100	4.2	0	.6	Nov 29, 2007 4:52 PM
TIVMV51	100	2.09	0	.6	Nov 29, 2007 4:52 PM
TIVMVS10	100	2.09	0	.6	Nov 29, 2007 4:52





Tivoli



AGENDA

- Introduction
- Monitoring Requirements
 - Virtual Linux and z/VM performance considerations
 - Don't forget the hardware
 - Integration from hardware systems applications Persistent historical views
- Enterprise Management
- Operational Requirements
 - Centralized Control
 - Including all Enterprise Virtual Machines
- Integrating Monitoring and Operations
- Bringing it all together



Addressing the trend – zEnterprise will enable management of diverse resources across diverse platforms as a single Workload

- A Platform Workload is a grouping mechanism and "management view" of virtual servers supporting a business application
- Provides the context within which associated platform resources are presented, monitored, reported, and managed
- Management policies are associated to Platform Workload
 - Currently supports Performance Policy





Looking at managing the zEnterprise aka "systems of systems" getting yourself organized..

Visibility

See your Business

Control

Manage service risk and compliance

Automation

Optimize business service delivery



Manage different Hypervisors as Centralized resource.



Monitoring and Managing the Enterprise – zEnterprise will enable the management of Resources across Virtual Servers







- Manage resources across virtual servers to achieve workload goals
 - Detect that a virtual server is part of Workload not achieving goals
 - Determine that the virtual server performance can be improved with additional resources
 - Project impact on all effected Workloads of moving resources to virtual server
 - If good trade-off based on policy, redistribute resources
 - Initially support CPU management



Business views across the zEnterprise





• ITM Infrastructure is shown separate to highlight components, however, each of these ITM components can reside on the zEnterprise.

• OMEGAMON agents can monitor z/OS system and subsystems, z/VM system and LPAR components, and Linux on z.

• ITM agents can monitor Linux on System z, Linux on System x, and AIX on Power7, and supported applications and databases.



Common Interface across the zEnterprise



Complete your session evaluations online at www.SHARE.org/AnaheimEval. Note: All statements regarding IBM's plans, directions, and intent are subject to change or withdrawal without notice, and represent goals and objectives only.



The future is ensembles and multiple hypervisors





AGENDA

- Introduction
- Monitoring Requirements
 - Virtual Linux and z/VM performance considerations
 - Don't forget the hardware
 - Integration from hardware systems applications Persistent historical views
- Enterprise Management
- Operational Requirements
 - Centralized Control
 - Including all Enterprise Virtual Machines
- Integrating Monitoring and Operations
- Bringing it all together





Operations Manager for z/VM

Increase productivity

- > Authorized users view and interact with monitored virtual machines without logging onto them
- Multiple users view/interact with a virtual machine simultaneously

Improve system availability

- > Monitor virtual machines and processes
- > Take automated actions based on console messages
- Reduce problems due to operator error





Features and Functions

- Monitor service machine consoles
- Monitor spool usage
- Monitor system events
- View and interact with monitored consoles from authorized user IDs
- Find and view spool files
- Schedule events/actions
- Dynamic configuration
- Separation of access control





Enterprise level console/syslog management:

Technology - Connections - Results



Complete your session evaluations online at www.SHARE.org/AnaheimEval. Note: All statements regarding IBMs plans, directions, and intent are subject to change or withdrawal without notice, and represent goals and objectives only.



Monitor Service Machines

- Define rules to
 - Scan console messages for text matching
 - Includes column, wildcard, and exclusion support
 - Optionally restrict to specific user ID(s)
 - Take actions based on matches
- Multiple rules can apply to one message
 - Rules processed in order of definition in the configuration file
 - FINAL option available to indicate no additional rules should be evaluated





View and Interact with Consoles

- Authorized users can view live consoles of monitored service machines and guests
 - Multiple users can view the same console simultaneously
 - No need to logon to the service machine to see its console
 - Test data and Linux syslog data treated as a "console"
 - Views can be defined to look at a group of consoles in one view
- Full screen mode
 - Scroll up and down to view and search historical data
 - Auto scroll (on or off) as new output is displayed on the console
 - From command line, issue commands back to the monitored console
- Amount of data that is visible depends on specified or default data space size
- Rules/actions may modify the view
 - Suppress messages from the console
 - Hold or highlight messages with color, blinking, etc.
- Authorized users can view the log file
 - Can also request a copy of the log file from today or a previous day





Monitor and View Spool Files

- Create spool monitors to trigger actions when
 - Percent of spool usage falls within a specified range
 - Percent of spool usage increases at a specified rate
- Actions triggered can be the same actions used by console monitoring
- Authorized users can
 - Display a list of spool files based on one or more attributes
 - Owner
 - Size
 - Date created
 - From the list the user can
 - · View the contents of an individual spool file
 - Transfer, change, or purge a spool file





Schedule Events and Actions

- Define schedules
 - Hourly, daily, weekly, monthly, or yearly, nth weekday of the month
 - Once on specified month, day, year, and time
 - At regular intervals
 - Every x hours and y minutes
 - Within a specified window of time
 - Specify start time
 - Specify conflicting schedules
 - Specify maximum time to defer this schedule
 - Within limits
 - Restrict to specific days of the week: Monday through Sunday plus holidays
 - Restrict to certain hours of the day
- Specify the action associated with the schedule
 - Actions specified are the same as those for console and spool monitoring





Respond to System Events

- Create monitors for z/VM system events (*VMEVENT) related to user IDs
 - Logon
 - Logoff
 - Failure condition (typically CP READ)
 - Logoff timeout started
 - Forced sleep started
 - Runnable state entered (VM READ)
 - Free storage limit exceeded
- Optionally restrict to specific user ID(s)
- Specify the action associated with the event
 - Actions specified are the same as those for schedules and console and spool monitors





Dynamic Configuration

- Initial configuration file loaded at startup
 - May imbed other configuration files
- Most configuration options can be updated while Operations Manager is running
 - Add, delete, or change:
 - Rules, actions, monitors, schedules, holidays, groups, user authorization
 - Suspend or resume rules, monitors, schedules
- Multiple methods
 - GOMCMD command interface
 - Load a new or updated configuration file
 - Commands in DEFACTN statements





Operations Manager







Summary

- Use Operations Manager to
 - Automate daily operations
 - Prevent problems rather than react to them
 - Automate reactions to problems when they can't be prevented
 - Improve problem determination procedures
 - Increase programmer and operator productivity





AGENDA

- Introduction
- Monitoring Requirements
 - Virtual Linux and z/VM performance considerations
 - Don't forget the hardware
 - Integration from hardware systems applications Persistent historical views
- Enterprise Management
- Operational Requirements
 - Centralized Control
 - Including all Enterprise Virtual Machines
- Integrating Monitoring and Operations
- Bringing it all together





Bring it all together

It is often that a unit of work is serviced by multiple applications and databases across multiple operating systems, including z/VM and Linux. Integrated views allow:

- Unit of work, or application tracking
- Business views
- Single skill sets to monitor dissimilar hardware, operating system, and application environments.




Application View: Scaling Scenario

- WebSphere MQ on Linux for System z receives "order requests" in the form of Queue messages, and places them on a queue.
- A WebSphere Application Server is invoked to periodically check the queue for messages and process them to a DB2 on z/OS database.
- The orders are coming too fast for the Websphere application to process.
- A second Linux server is started with another copy of Websphere application server to aid in the processing of requests.



Application View: Scaling Scenario



- Trigger: Queue Depth
- Options for triggering actions can be based on things such as:
 - The number of orders received but not yet processed (the number of messages on the queue)
 - The amount of time it is taking to process the orders
 - The response time of the web application
 - The CPU usage of the z/VM Guest
 - Other things I haven't given much thought to yet.



MQ Series Queue growth started



•••• in Anaheim

📲 Navigator 🌲 🗉 🖯 🗙	🔲 Order Pr	ocessing Ti	me - M 🗸	* 🛛 🖯 🗆 ×	🔲 Order Proces	sing Guests 🛛 🔻 🔟 🖯 🗙
🕐 🏀 View: P10rders 💌	ORDERS F	ROCSECS			Orde	er Processing Guests
P10rders P10rders_Guests P10rders_MQ P10rders_Web	8	189]		User ID 1	
∎ P10rders_zVM	🔲 Daily Par	🛄 Daily Parts Orders				/ ∓ 🛛 🖯 ×
	PART	ORDERS	QUANTITY	V AVGPROCSECS	MAXPROCSECS	LocalTimeStamp
	GIZMOS	1	679	247	247	2008/05/09 11:46:51 030
	FOOBARS	2	1149	197	216	2008/05/09 11:46:51 010
	THINGYS	2	918	190	231	2008/05/09 11:46:51 040
	GADGETS	2	1740	164	195	2008/05/09 11:46:51 020
	WIDGETS	1	793	163	163	2008/05/09 11:46:51 050
Physical Republication of the second						
💷 z/VM guest processor 🖉 🗉 😑 🗙 🖬	Applica 🥖		× 📶 Ore	lers R 🥒 🏝 🔟	8 🗆 × 🖬 B	ar Chart 🥒 🏝 🔟 🖻 🗮 🗙
z/VM guest processor usage 🛛 🔄	¹⁶ 1		3			
2.0 1.0 0.0 ESMTS105 ESMTS105		Current Dep	3 1 		8000 6000 4000 2000	
🔲 z/VM Resources allo 🖉 🐺 🖽 🖶 🗖 🗙	6				Requests	Avg.Resp. (ms)
Linux Guest ID Time c OOSP1A.ESMTS105:LZ 05/09/08 11:46:27	4 2 0 APPQC01		05/09/08 11:00:00	0 (Requests)(05/09/08 5 09/09/08 5 09/09/08 11:20:00 00 00 00 00 00 00 00 00 00	11:40:26) 0	05/09/08 11:45:00 05/09/08 11:30:00 05/09/08 11:30:00
				1	p	

-

•••• in Anaheim

Scaling Scenario





Adjusting Resources for a Linux Guest

- Virtual CPU consumption is high for a Linux guest
- Detect the alert
 - Automation receives the message
- Action is triggered by a rule in Operations Manager
- Operations Manager issues CP commands to tune the guest
 - SET QUICKDSP
 - SET SHARE
- Ability to monitor the output is key



Adjusting resources for a Linux guest







OMEGAMON Configuration

- Define a situation (alert) to detect high CPU consumption for Linux virtual machines.
- Define the automated "Take Action" to:
 - Direct a message to console monitored by Operations Manager.
 - Include in the message keywords to trigger Operations Manager rule.
 - Guest Name
 - Guest need CPU priority text
 - Any unique data desired for specific customer environment.





Е ults



Hindi



Traditional Chinese

Спасибо

Russian

Thank You

English



Korean

Gracias

Spanish

Obrigado

Brazilian Portuguese

Grazie

Arabic

Italian



Simplified Chinese

ありがとうございました

Japanese



Merci

French







Managing z/VM and Linux Performance Best Practices



Mike Sine IBM

March 13, 2014 Session Number 14637

