



Configuring LPARs for Performance

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Agenda

Overview of Terms and Partitioning Controls

- Per CP Share
- Short CPs

Managing Capacity

- Intelligent Resource Director
- Initial Capping
- Soft Capping
- Group Capacity



Important Terms to Understand

- LPAR weight and per CP share
- Effective Dispatch Time
- Partition Dispatch Time
- Short CPs

Important Concepts to Understand

- LPAR weights become important only when the processor is very busy or capped
- There are two dispatchers involved in making resource allocations
 - PR/SM
 - Operating System



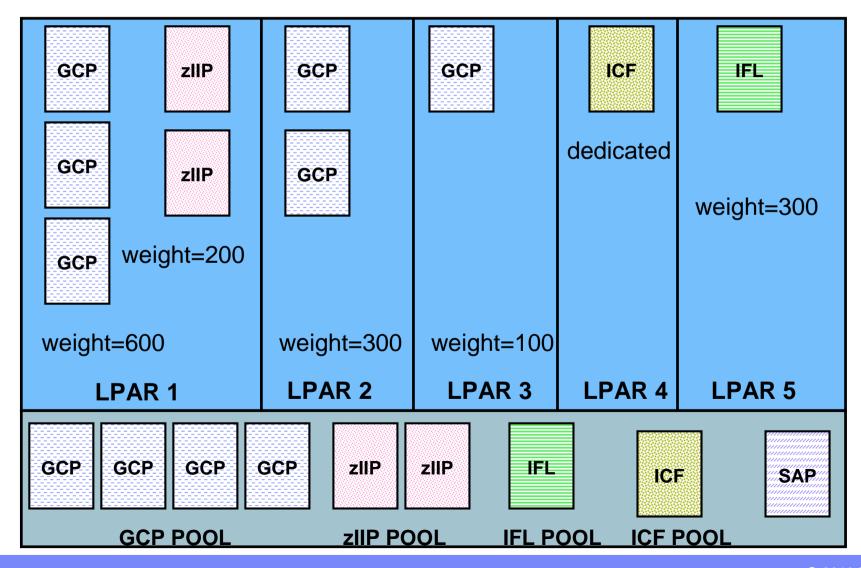
Partitioning Controls

- Number of partitions, their relative weights, and CP mode (dedicated or shared)
- Number of logical CPs defined to the partitions
- Horizontal or Vertical CP Management (Hiperdispatch)
- Capping Controls
 - Initial Capping (Hard Caps)
 —Group Capacity Controls
 - Defined Capacity (Soft Capping)
 Absolute Capping (NEW!)
- Ratio of logical CPs to physical CPs
- CP usage; either general purpose, or specialty CP (IFL / ICF / zAAP / zIIP) CPs
- Type of system control program (z/OS, z/VM, Linux, etc.)



Partitioning Controls

 A Partitions weight is relative to the summed weights of all of the partitions in their respective pools



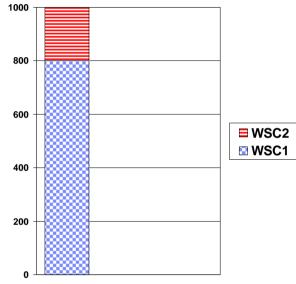


System z Virtualization

- 1 to 60 LPARs per CEC
- Number of CPs is CEC dependent
 - 1 to 64 for the 2097-E64 (z10)
 - 1 to 80 for the 2817-M80 (z196)
 - 1 to 101 for the 2827-HA1 (zEC12)
- Number of logical CPs is operating system dependent
- Operating System doesn't know it is not running on the hardware
 - More integration is happening over time i.e. hiperdispatch
- Dispatching can be done event driven (typical) or time sliced
- Dispatch interval is based on a heuristic method which depends upon the logical to physical ratio



Calculate LPAR Share



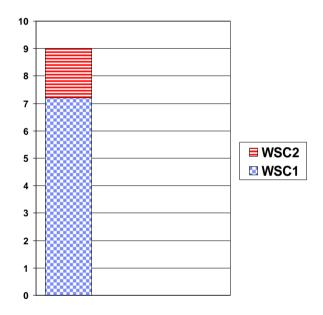
WSC1 Share: 800 / 1000 = 80%

WSC2 Share: 200 / 1000 = 20%



 All active LPARs are used even if an SCP is not IPL'ed

 Only LPARs with shared CPs are used in the calculation Processor guarantee = # of General Purpose Physical (GCP) * LPAR Share



WSC1 Capacity: 9 * .80 = 7.2 CPs

WSC2 Capacity: 9 * .20 = 1.8 CPs

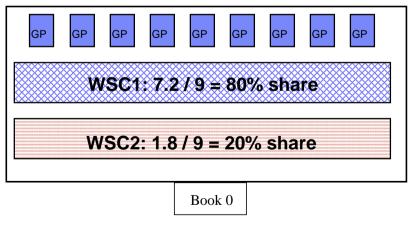


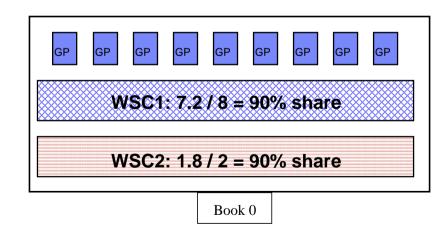
 The processor guarantee is used to offer protection to one LPAR over other busy LPARs demaning service



Determine Per CP Share - Horizontal CP Management

- PR/SM guarantees an amount of CPU service to a partition based on weights
- PR/SM distributes a partition's share <u>evenly across</u> the logical processors
- Additional logicals are required to receive extra service which is left by other partitions. The extra service is also distributed <u>evenly across</u> the logicals
- The OS must run on all logicals to gather all its share [z/OS Alternate Wait Management]





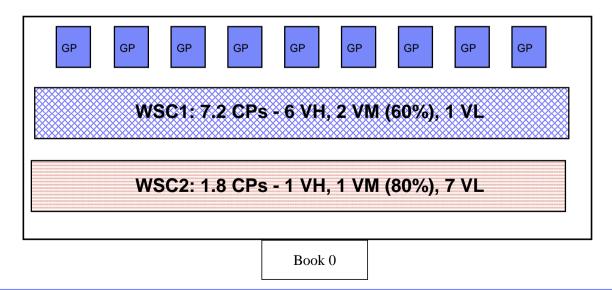


Biggest Per CP Share possible is best when processor is busy



Determine Per CP Share - Vertical CP Management

- Logical processors are classified as vertical high, medium or low
- PR/SM quasi-dedicates vertical high logicals to physical processors
- The remainder of the share is distributed to the vertical medium processors
- Vertical low processors are only given service when other partitions do not use their entire share and there is demand in the partition
- Vertical low processors are parked by z/OS when no extra service is available





System z Partitioning Controls

Access to resources is relative to other partitions on the CEC
 2827-704

Pool	LPAR Name	Weight	Logicals Defined	Logicals by Weight	Logical to Physical Ratio	HD=YES
GCP	LPAR1	600	3	2.4		1 VH, 2 VM
GCP	LPAR2	300	2	1.2		2 VM, 60% share
GCP	LPAR3	100	1	0.4		1 VM, 40% share
		1000	6		1.5 : 1	
zIIP	LPAR1	200	2	2		2 VH
		200	2		1:1	
IFL	LPAR5	300	1	1		
		300	1		1:1	
ICF	LPAR4	DED	1	1		
			1	1	1:1	



Logical Processor Utilization

- Measurement which states the busy of the logical CPs
 - Independent measure of capacity
 - Can run out of logical CP capacity before the processor is 100% busy
 - More logical CPs than weight means the utilization is artificially low

Physical Processor Utilization

- Differs from effective time when the number of logical CPs defined to the partition does not match the number of GCPs
- It is this metric which is used in Capacity Planning exercises

	- PAR	TITION	I DATA				L	OGICAL	PARTITION PROC	ESSOR DATA	AVERAGE	PROCESSO	R UTILIZATI	ON PERCENTA	AGES
			MS	U	-CAF	PING	PROC	ESSOR-	DISPATCH	TIME DATA	LOGICAL PRO	CESSORS	PHYSIC	AL PROCESSO	ORS
NAME	S	WGT	DEF	ACT	DEF	WLM%	NUM	TYPE	EFFECTIVE	TOTAL	EFFECTIVE	TOTAL	LPAR MGMT	EFFECTIVE	TOTAL
WSC1	A	370	0	700	NO	0.0	15.0	CP	01.45.57.466	01.46.19.021	47.09	47.25	0.10	28.26	28.35
WSC2	A	315	0	288	NO	0.0	15.0	CP	00.43.23.443	00.43.46.035	19.28	19.45	0.10	11.57	11.67
WSC3	A	315	0	178	NO	0.0	15.0	CP	00.26.39.732	00.27.00.535	11.85	12.00	0.09	7.11	7.20
WSC4	A	25	45	4	NO	0.0	2.0	CP	00.00.32.779	00.00.34.362	1.82	1.91	0.01	0.15	0.15
PHYSICAI	L *									00.01.05.674			0.29		0.29
TOTAL									02.56.33.422	02.58.45.630			0.59	47.08	47.67



RMF Partition Report

MVS PARTITION NAME	WSC1	NUMBER OF PHYSICAL PROCESSORS	31
IMAGE CAPACITY	2469	CP	25
NUMBER OF CONFIGURED PARTITIONS	17	IFL	1
WAIT COMPLETION	NO	ICF	2
DISPATCH INTERVAL	DYNAMIC	IIP	3

	PART	rition	DATA				LO	OGICAL
			MS	U	-CAP	PING	PROCE	ESSOR-
NAME	S	WGT	DEF	ACT	DEF	WLM%	NUM	TYPE
WSC1	A	370	0	700	NO	0.0	15.0	CP
WSC2	A	315	0	288	NO	0.0	15.0	CP
WSC3	A	315	0	178	NO	0.0	15.0	CP
WSC4	A	25	45	4	NO	0.0	2.0	CP
*PHYSICAL	k							
TOTAL								
CF01	A	DED					2	ICF
PHYSICAL	k							
TOTAL								
WSC1	А	10					3	IIP
WSC2	A	10					3	IIP
WSC3	А	10					3	IIP
WSC4	А	10					1	IIP
PHYSICAL	k							

Processor Running Time

- Default is limited to a range of 12.5-25 ms
- Dynamically calculated

25 ms * (Number of Physical Shared CPs) Total # of Logical CPs for all LPARs

- Vertical Highs get run time of 100 ms
- Recalculated when LPARs are stopped or started or CPs are Configured on/off
- When a logical CP does not go into a wait state during its run time, it loses the physical CP when it reaches the end of its run time



Managing Capacity on System z and z/OS

- Intelligent Resource Director
- PR/SM Initial Capping Hard Capping
- Defined Capacity Soft Capping
- Group Capacity
- Other Methods of Changing Capacity
 - Absolute Capping
 - WLM Resource Groups
 - Discretionary Goal Management
 - Config CPU Command
 - Customer Initiated Power Save Mode
 - OOCoD



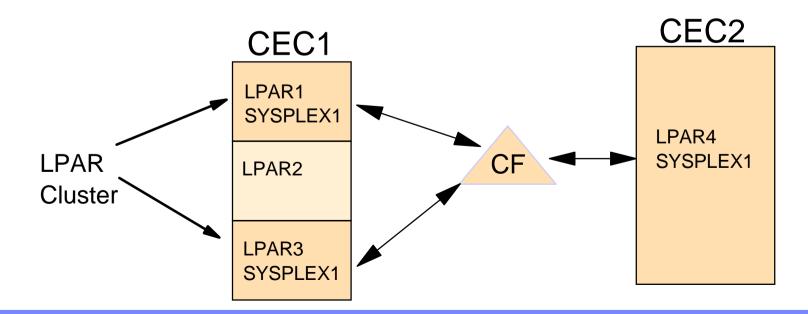
Intelligent Resource Director

WLM Function which:

- Manages LPAR Weights
- Varies logical CPs On and Off Disabled and replaced with Hiperdispatch=YES
- Manages CHPIDs
- Manages I/O Priorities

Scope is an LPAR Cluster

All MVS images on the same physical processor, in the same sysplex





IRD Management

WLM manages physical CPU resource across z/OS images within an LPAR cluster based on service class goals

- LPAR Weight Management
 - Dynamic changes to the LPAR weights
 - Sum of LPAR weights can be redistributed within the cluster
 - Partition(s) outside of the cluster are not affected
 - Moves CP resource to the partition which requires it
 - Reduces human intervention

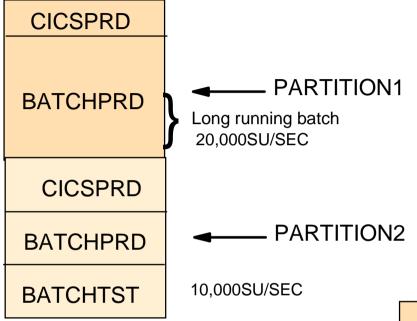
LPAR Vary CP Management

- Dynamic management of online CPs to each partition in the cluster
- Optimizes the number of CPs for the partition's current weight
- Prevents 'short' engines
- Maximizes the effectiveness of the MVS dispatcher
- Has an IEAOPTxx option (VARYCPUMIN) to set minimum number of CPs regardless of LPAR's weight
- Reduces human intervention
- Replaced by Hiperdispatch=yes

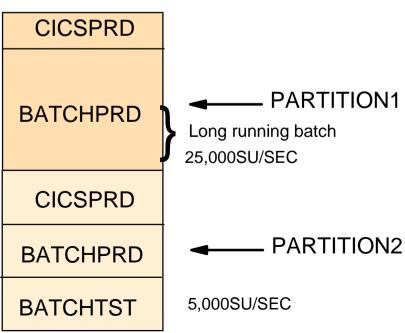


Benefit of LPAR IRD Management

This



Becomes





LPAR Weight Management

- Enabled using PR/SM server processor definition panels
 - WLM managed
 - Initial processing weight
 - Minimum processing weight
 - Maximum processing weight
- Weights should be 2 or 3 digit values to provide optimal results
 - Weight is increased by 5% of average LPAR weight value

LPAR	Weight		
Receiver	400		
Donor	600		

Weight Donation

LPAR	Weight
Receiver	425
Donor	575



LPAR Weight Mgmt Requirements

- LPARs have shared CPs
- LPARs do not use Initial capping (Hard Caps)
- WLM LPAR CF structure is defined and connected
- LPAR CPU management is enabled on LPAR panels or Hiperdispatch=YES
- System must be running z/OS
- Processor must be a zSeries or System z processor



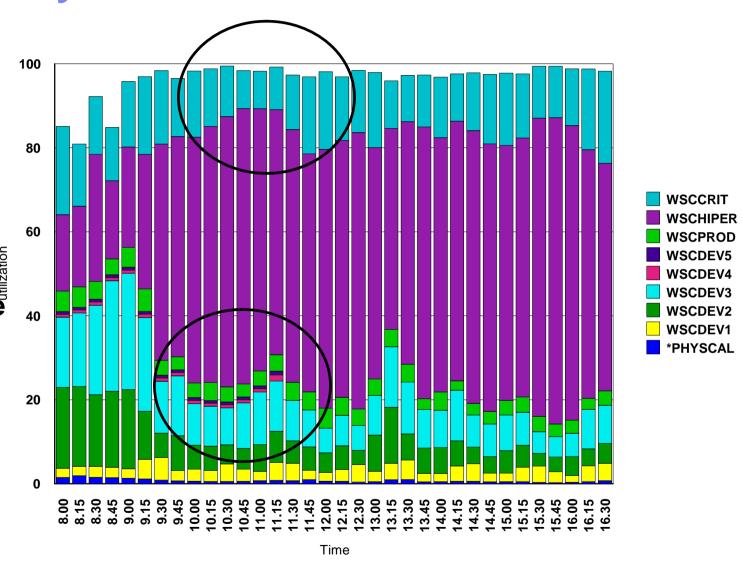
IRD Performance Management

- Need to have a multi system perspective when looking at overall throughput in an LPAR Cluster
 - WLM Policy and Goal Attainement
- Need to examine CEC demand within and outside the cluster
 - Whitespace



Overall CEC Busy

- IRD is active
- WSCCRIT,
 WSCHIPER and
 WSCPROD are in an LPAR Cluster called WSCPLEX
- When WSCDEV4 and WSCDEV5 are there the LPAR Cluster gets 82% of the CEC, when they are stopped the LPAR Cluster gets 89%
- CEC is very busy

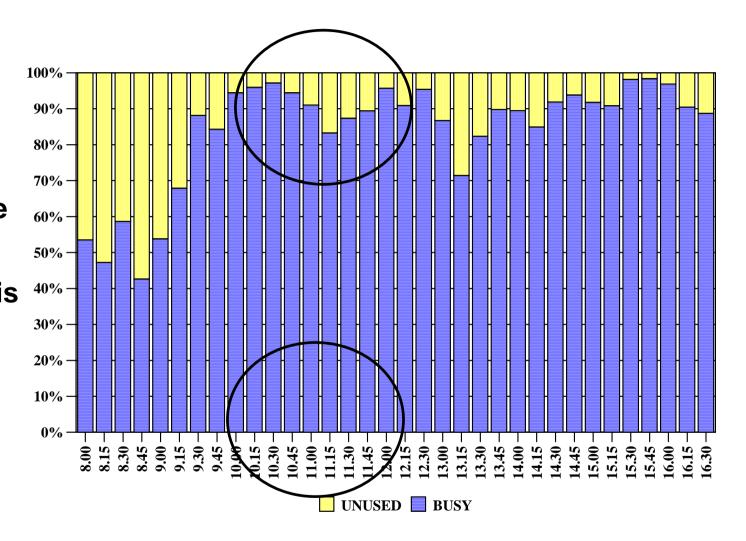


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WSCPLEX Cluster View of Capacity

- This chart represents the fair share of the CEC that the WSCPLEX cluster should have access to
- WSCPLEX Cluster is not using all of its capacity so is donating white space to the other LPARs





WSCCRIT Performance Degradation

HOUR	SCLASS	IMP	СР	PINDX
10	SYSTEM	0	10.5	1
10	SYSSTC	0	17.9	1
10	TSOL1	1	1.5	0.5
10	TRANONE	1	2.6	0.1
10	DMGMT	1	4	1.4
10	SERVERS	1	23.3	25.5
10	CICSL2	1	0.1	3.4
10	STCHI	1	1.2	2
10	TSOHI	1	1.3	0.6
10	TRANTWO	1	8.9	0.1
10	TRNMULT	1	6.5	0.3
10	TRNMULT	2	28.6	1.4
10	STC2	2	46.7	1.1
10	TRANFIVE	2	3.5	3.4
10	TSOHI	2	0.2	1.9
10	TSOI1	2	1.4	13.3
10	TRANONE	2	12.5	3.6
10	HOTPROD	2	0	0.1
10	DBASE	2	0	0.1
10	BATCHL1	3	52.2	2
10	TSOL1	3	1.6	1.3
10	DBASE	3	0	1.2
10	BATCHL2	5	2.7	3.7

HOUR	SCLASS	IMP	СР	PINDX
11.15	SYSTEM	0	14	1
11.15	SYSSTC	0	21.7	1
11.15	TSOL1	1	1.4	1.3
11.15	TSO1	1	1	0.3
11.15	TRANONE	1	5.1	1.1
11.15	SERVERS	1	17.5	2.9
11.15	TRANTWO	1	1.1	3.4
11.15	CICSL2	1	0	72.1
11.15	DMGMT	1	3	3.9
11.15	STCHI1	1	3.7	5.3
11.15	TRNMULT	1	9.8	1.7
11.15	STC2	2	25.2	12.7
11.15	TRNMULT	2	24.7	107.2
11.15	TRANTWO	2	2.7	236.5
11.15	TSOL1	2	1	4.2
11.15	TSOHI	2	0.8	30.9
11.15	TRANONE	2	9.7	198.9
11.15	DBASE	2	0	14.5
11.15	BATCHL1	3	0	5138.8
11.15	DBASE	3	0	160.3



WSCPLEX LPAR Share Within the Cluster

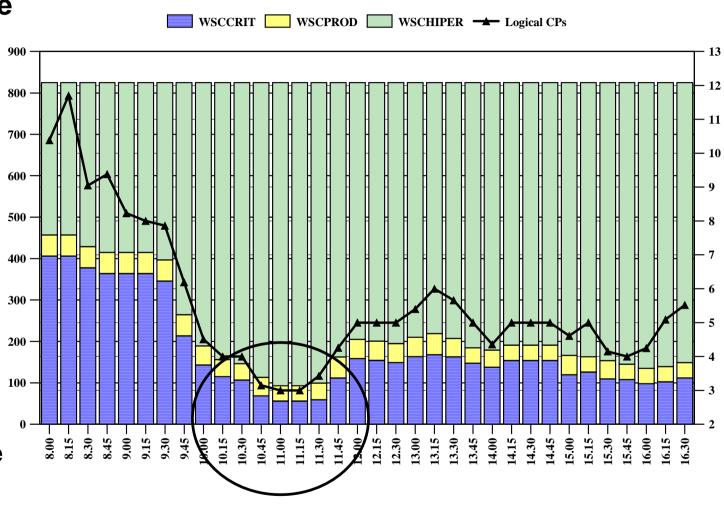
 This chart shows the change in LPAR weight over time

WSCCRIT is losing weights and logical CPs

Two issues to examine:

Why did we donate white space?

WSCCRIT suffers
 performance
 problems but did the
 benefit to
 WSCHIPER
 outweigh the costs?





WSCHIPER Gets Additional Weight

LPAR Weight

	10.00	11.15
WSCHIPER	636	732
WSCCRIT	143	56

LPAR Busy

	10.00	11.15
WSCHIPER	58.5	62.5
WSCCRIT	15.7	8.93

- WSCHIPER gets more weight but doesn't do more work
- High PIs makes IRD hesitant to move weight back
- High CEC Busy means no additional logicals can be added to WSCCRIT
- Low number of logical CPs means WSCCRIT can't schedule the work and hence the whitespace is donated

HOUR	SCLASS	IMP	СР	PINDX
10	SYSSTC	0	34.1	1
10	SYSTEM	0	20.2	1
10	CICSL1	1	307.9	1
10	CICSL2	1	182.4	1.1
10	CICSL3	1	81.6	1.2
10	SERVERS	1	59.6	1.3
10	STCHI	1	12.7	1.4
10	OMVS	2	0.1	0.4
10	STC2	2	33.9	1.3
10	BATCHL1	3	135.2	1.4
10	STCLO	3	1.3	2.4
10	TSOL1	3	0.2	0
10	BATCHL2	5	5	2.2
11.15	SYSSTC	0	35.9	1
11.15	SYSTEM	0	31.3	1
11.15	CICSL1	1	315.8	1
11.15	CICSL2	1	193.7	1
11.15	CICSL3	1	78.2	1.1
11.15	SERVERS	1	53.4	1.3
11.15	STCHI	1	20.7	1.2
11.15	OMVS	2	0.8	0.3
11.15	STC2	2	5	1.1
11.15	BATCHL1	3	118.3	1.5
11.15	STCLO	3	1.4	1.5
11.15	TSOL1	3	0.1	0
11.15	BATCHL2	5	9.4	1.2



What are the Tuning Options

- If Hiperdispatch = No then use VaryCPUMIN to keep sufficient logicals available
- Update the WLM policy so the goals are more reasonable
- Provide protection with IRD Minimum values



Providing Protection in an IRD Environment

- Decisions across LPARs are based on WLM Policy
 - Ensure WLM definitions are well defined and accurate
 - Review Performance data at the LPAR Cluster level
- Protection comes from the use of MIN weights
 - Special protection for LPARs with high percentage of work which can be donated

Initial Weight = 6 CPs

IMP	CPs
0	
1	
2	
3	
4	1.3
5	0.8
6	0.4

Initial Weight = 3 CPs

Min Weight = 0.5 CPs

IMP	CPs						
0	0.4						
4	0.5						
5	1.3						
6	0.8						
2.6 CPs							

→

Initial Weight = 10 CPs

IMP	CPs
0	
1	
2	
3	
4	2.3
5	1.2
6	0.9

4.5 CPs

Min Weight = 5.5 CPs

2.5 CPs



PR/SM Initial Capping – Hard Cap

- HiperDispatch=No
 - The LPAR's relative weight per CP is the share for each logical CP and the goal of the LPAR dispatcher is to give each logical CP its share of the total relative weight
 - Capping is done on a logical CP basis
- Hiperdispatch=YES
 - Vertical High's will be capped at 100% of the logical
 - Vertical Mediums and Vertical Lows will share the allowed weight on a per CP basis



PR/SM – Weight Enforcement

Weight Enforcement Depends Upon LPAR definitions

- LPAR with Initial Capping
 - Enforces processing weights to within 3.6% of the LPAR's physical per CP share for logical CPs entitled to <u>1/10 or more</u> of one physical CP
- LPAR is Uncapped
 - Enforces the processing weights to within 3.6% of the LPAR's physical per CP share for logical CPs entitled to 1/2 or more of one physical CP
- LPAR Logical CP fails enforcement levels
 - Enforce the processing weights to within 3.6% of the total capacity of the shared physical CP resources
- Typically in most cases PR/SM will manage the processing weights to within 1% of the LPAR's physical per CP share



PR/SM Initial Capping – Weight Allocations

- An LPAR's Hard Capped Capacity is relative to the other LPARs
 - If an LPAR is started or stopped on a CEC with a hard cap a weight change must be done concurrently or the capped LPAR's allowed capacity will change
 - With Hiperdispatch you need to deactivate the LPAR so the VHs are reallocated correctly otherwise VLs will be used
 - WSC2 needs to go from 4 VH, 2 VM to 12 VH, 1 VM

2817-718

Name	Status	Weight	Capped	Weight in CPs	Status	Weight	Capped	Weight in CP
WSC1	Α	400	NO	7.2	D			
WSC2	Α	300	NO	5.4	Α	300	NO	9
WSC3	Α	300	YES	5.4	Α	300	YES	9

New Absolute Capping can help here



Defined Capacity – Soft Capping

Specified via LPAR definitions

- Provides sub-CEC pricing by allowing definition of LPAR capacity in MSUs
 - Allows a defined capacity smaller than the total capacity of the LPAR
 - Provides 1 MSU of granularity
- Only way to get a soft cap
- Initial Capping (PR/SM Hard Cap) and Defined Capacity cannot be defined for the same partition
 - Rule applies to GCPs only, Specialty CPs can be hard capped while the GCPs are soft capped
- LPAR must be using Shared CPs (Dedicated CPs are not allowed)

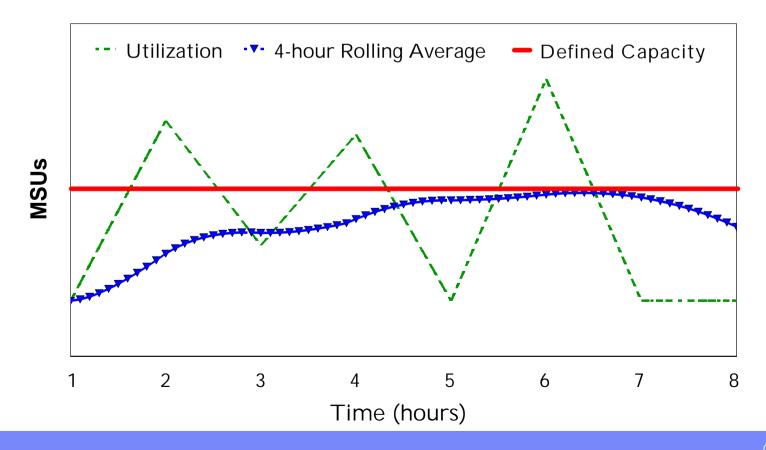
All sub-capacity software products in an LPAR have the same capacity

- LPAR partition capacity based pricing not product usage based pricing
- Regardless of actual CPU seconds used by the product



Rolling 4 hr Average & Defined Capacity

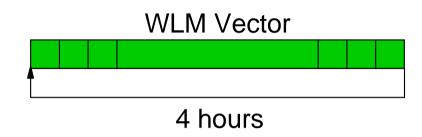
- Rolling 4-hour average tracked by Workload Manager
 - Rolling 4-hour average is not permitted to exceed defined capacity
 - May exceed during early capping intervals
 - If 4-hour average exceeds defined capacity, LPAR gets soft capped



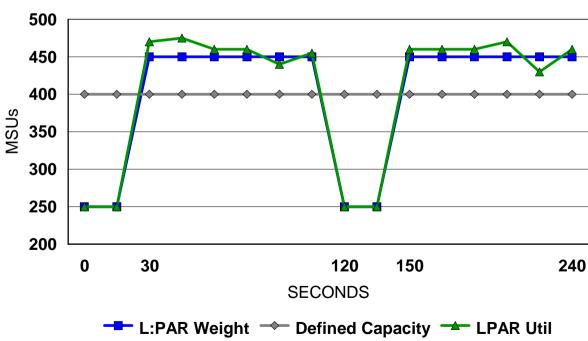


Managing to the Rolling 4hr Average

- LPAR Effective dispatch time for partition averaged in 5 min intervals
 - 48 entry vector
 - Every 10 seconds WLM issues Diagnose command to hardware to get effective time
 - Vector wraps after 4 hours
- Calculate a capping pattern
 - Control ratio of capped versus non-capped periods to keep partition usage at defined capacity
 - Capping state should change no more than once per minute
 - Limit partition to it's weight



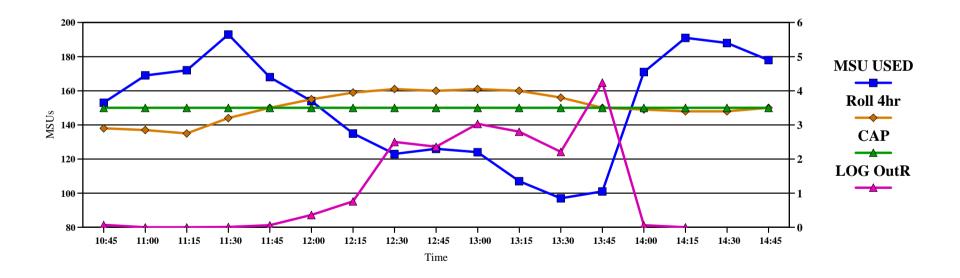






Managing to the Rolling 4hr Average

- When softcapped the LPAR is allowed to continually use the amount of capacity defined
- Work is not stopped to "make up" for time period when rolling 4hr average exceeds the defined capacity





LPAR Group Capacity Basics

Manage CPU for a group of z/OS LPARs on a single CEC

- Limit is set to total usage by all LPARs in group
 - Level of granularity is 1 MSU
 - Members which don't want their share will donate to other members
- Independent of sysplex scope and IRD LPAR cluster
- Works with defined capacity limits on an LPAR
 - Target share will not exceed defined capacity
- Works with IRD
- Can have more than one group on a CEC but an LPAR may only be a member of one group
- LPARs must share engines and specify WAIT COMPLETION = NO

Capacity groups are defined on the HMC Change LPAR Group Controls panels

- Specify group name, limit in MSUs, and LPARs in the group
- Members can be added or removed dynamically



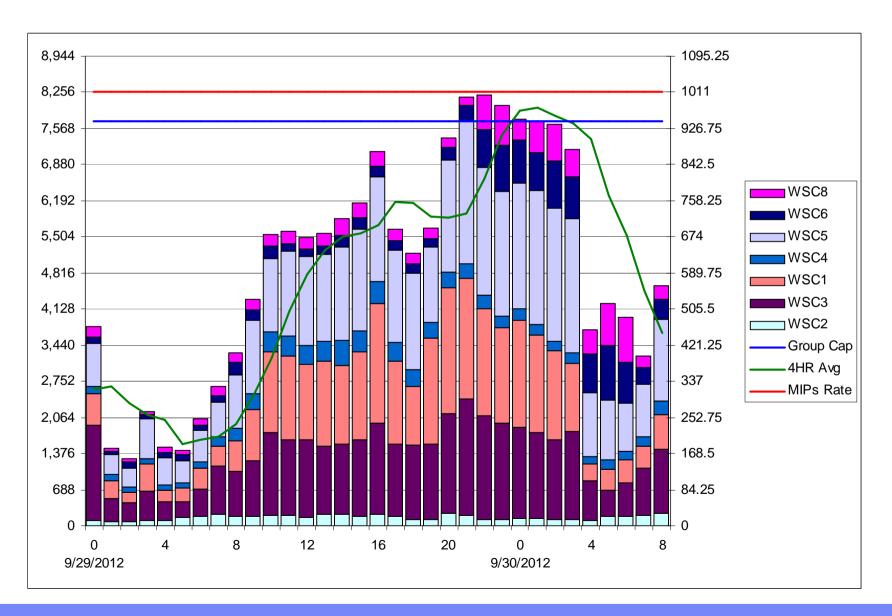
LPAR Group Capacity Basics

- Uses WLM rolling 4 hr avg in MSUs to manage the Group Capacity limit
 - Cap enforced by PR/SM if group rolling 4 hr avg exceeds limit
 - Each member is aware of other members' usage and determines its share based on its weight as a percentage of total weight for all members in group
 - NOTE: When using IRD the weights can change and therefore the target MSU value can change
 - The defined capacity limit, if also specified, is never exceeded
- Until members "learn" about the group and build a history, the cap is not enforced
 - May take up to 4 hours (48 measurements at 5 minute intervals are maintained for rolling 4 hour average) for capping to start
 - Similar to the bonus period with defined capacity
 - When new member joins the group, it has to build up its history and during this time the group usage may exceed the capacity limit
 - Capping is removed when the group rolling 4 hour average drops below group limit
- Example shows how many MSUs each LPAR would get if they all wanted their share. Target MSUs based on a group limit of 200. Total group weight is 500.

LPAR	WEIGHT	SYSPLEX	CAPACITY GROUP	TARGET MSU
LPAR1	150	PLEX1	GROUPA	60
LPAR2	300	PLEX2	GROUPA	120
LPAR3	500	PLEX1	n/a	n/a
LPAR4	50	PLEX1	GROUPA	20



Example of Group Capacity





RMF Group Capacity Reporting

- CAPPING WLM% (percentage of time WLM capped the partition) is insufficient when the partition is member of a capacity group:
 - WLM% only tells to what extent a partition is <u>subject to capping</u> but not whether the partition was actually capped
 - WLM% is more a matter of <u>how</u> WLM caps the partition instead of <u>how much</u> it is being capped
- CAPPING ACT% displays the percentage of time where the partition was actually capped
 - Users of Capacity Groups can determine the available (unused) capacity for their group and whether the partition was actually capped:

NUMBER OF PHYSICAL PROCESSORS						6			GROUP	NAME	ATS	
CP						5			LIMIT		141	
IIP						1	1 AVAILABLE 1				1	
GROUP-	CAPACITY	PARTITION	SYSTEM	MS	SU	WGT		CAPPING	G	- ENTITL	EMENT -	
NAME	LIMIT			DEF	ACT		DEF	WLM%	ACT%	MINIMUM	MAXIMUM	
ATS	141	WSC1	WSC1	0	0	25	NO	0.0	0.0	7	141	
		WSC2	WSC2	0	85	380	NO	87.5	13.1	119	141	
		WSC3	WSC3	0	24	25	NO	0.0	0.0	7	141	
		WSC4	WSC4	0	2	20	NO	0.0	0.0	6	20	
			TOTAL		111	450						

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Intersection of IRD and Group Capacity

- OA29314 DOC IRD and Group Capacity
- WLM only manages partitions in a Group Capacity which meet the following conditions:
 - Partition must not be defined with dedicated processors
 - Partition must run with Shared processors and Wait Complete=No must be set
 - Operating System must be z/OS 1.8 and above
 - z/OS cannot be running as a z/VM Guest
 - PR/SM Hard Capping is not allowed
- Any LPAR not meeting the conditions is removed from the Group and the remaining LPARs are managed to the Group Limit
- Group Capacity will function with IRD weight management as long as the partitions in the Group are not subject to capping
 - No Weight moves will take place as long as the Group is being capped



Enhanced SMF Recording

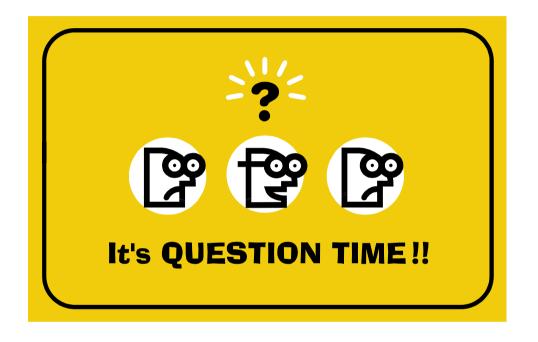
- It is recommended to turn on recording of SMF 99 subtype 11 when you start to exploit group capping
 - The collected data is small and only written every 5 minutes
 - Size is about 1300 bytes fixed + 240 bytes per LPAR on a CEC
 - Approximately 3k for a CEC with 8 partitions
 - The data is crucial for all analysis done by IBM therefore recommend the data be collected unconditionally



Summary

- LPAR controls are important in controlling capacity available to workloads
- IRD weight management is still valuable if you have the right environment
 - Measure and manage at the LPAR Cluster level
- Capping controls are inter-related and can be used to control overall CEC capacity
 - Be aware of the impacts on performance







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