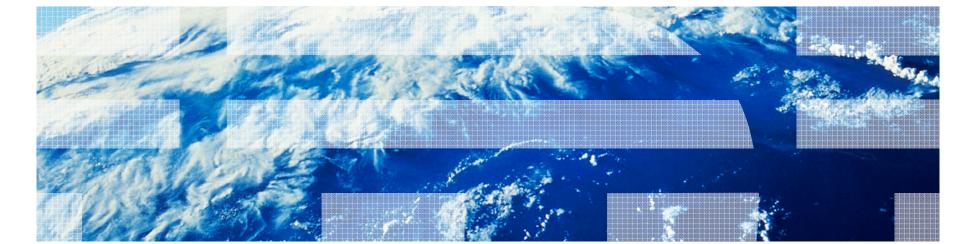
z/VM Performance Update 2013 Session 13505

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Agenda

- z/VM 6.3 Regression Performance
- Large Memory
- HiperDispatch
- Various other line items or small changes
- Monitor record changes
- Performance-related service
- z/VM Performance Toolkit



Regression Performance

- Ran our usual library of workloads -CMS interactive, various Apache configurations
- Results are within the usual 5% regression criteria
- Some workloads will see improvements:

 Constrained by reorder or demand scan
 A few heavy guests along with small virtual to logical processor ratio



Large Memory



Large Memory: Highlights

- Exploit a 1 TB central memory

 - Exploit larger real
 Allow larger total virtual
 - -XSTORE no longer required; best practice is not to use it
- Remove algorithms and techniques known not to scale
- Improve SET RESERVED in a couple key ways
 - -Make it more effective
 - -Extend it to NSS and DCSS
- Overhaul CP Monitor records appropriately
- New or changed z/VM Performance Toolkit screens

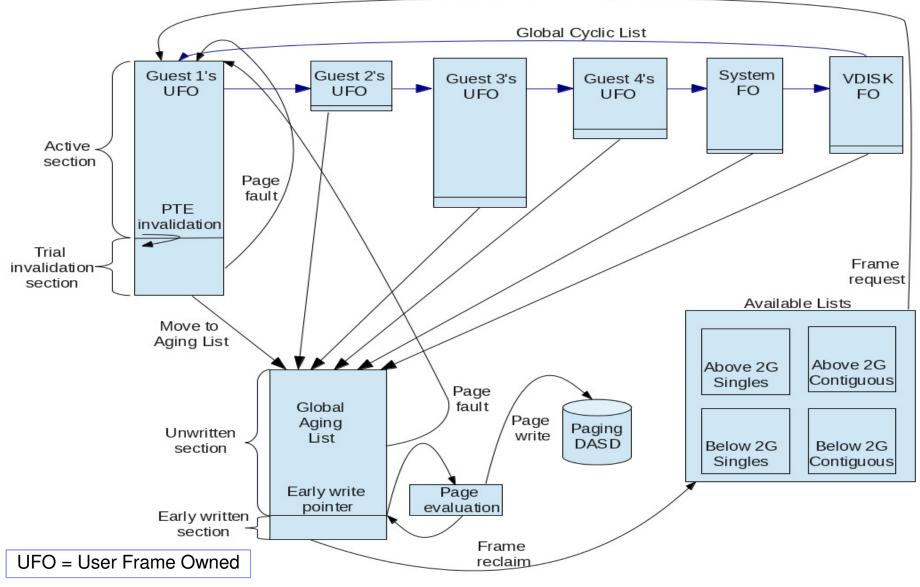


Large Memory, Scaling and Other Problems Removed

- Reorder processing is gone
- Demand scan no longer searches frame lists
- No more use of RRBE instruction (expensive instruction)
- DASD channel program now does scatter-gather I/O
- Unchanged pages usually not rewritten
- Remove use of scheduler lists
- Decrease priority of VDISK pages
- Better balance of memory used below 2 GB

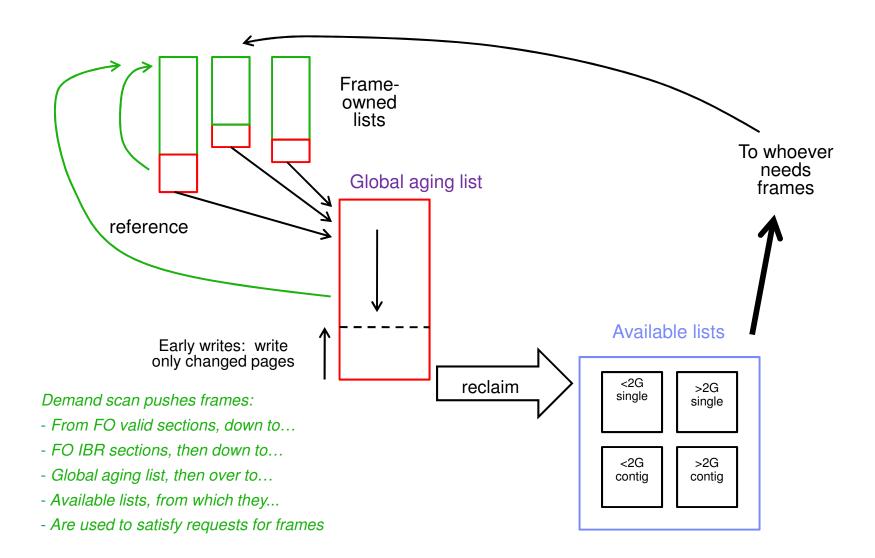


Memory Management Algorithm Visualization



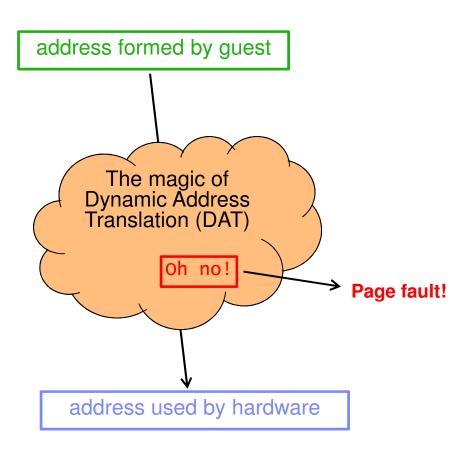


New Approach: The Big State Diagram





New Approach: Trial Invalidation



- Page table entry (PTE) contains an "invalid" bit
- What if we:
 - Keep the PTE intact but set the "invalid" bit
 - Leave the frame contents intact
 - Wait for the guest to touch the page
- A touch will cause a page fault, but...
- On a fault, there is nothing really to do except:
 - Clear the "invalid" bit
 - Move the frame to the front of the frame list to show that it was recently referenced
- We call this **trial invalidation**.



The "Sweet Spot" Workload

Our synthetic workload called *Sweet Spot* imitates behaviors we have seen in customer-supplied MONWRITE data.

	z/VM 6.2	z/VM 6.3	Delta	Pct. Delta
Cstore	256	384	128	
Xstore	128	0	-128	
External Throughput (ETR)	0.0746	0.0968	0.0222	29.8%
Internal Throughput (ITR)	77.77	105.60	27.83	35.8%
System Util/Proc	31.4	4.7	-26.7	-85.0%
T/V Ratio	1.51	1.08	-0.43	-28.5

By getting rid of both reorders and spin lock contention, we achieved huge drops in %CPU and T/V.



Large Memory Scaling Measurements

 VIRSTOR – Test case system started with CMS boot strap with controls over memory reference patterns and processor usage.

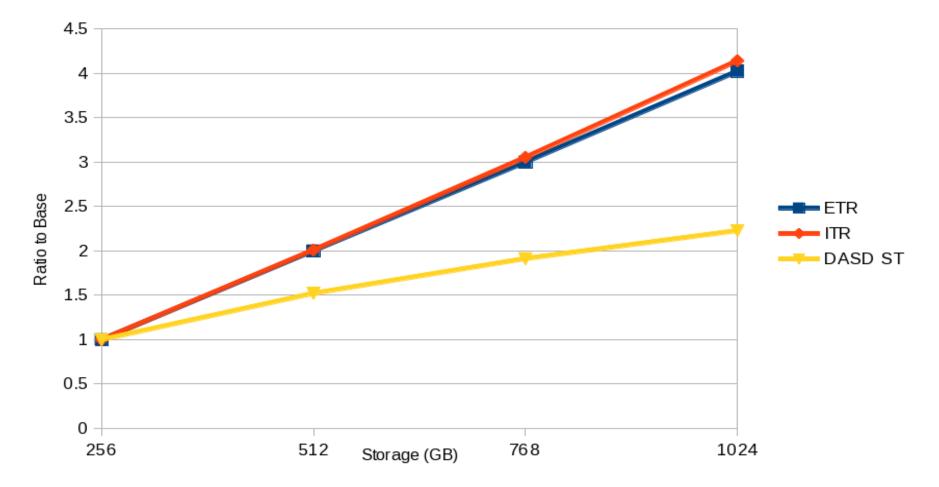
- Create workload similar to resource usage from customer Monwrite data

- Linux Apache Static Web serving
- Measure and test levels of servers at peak usage for 256 GB in an overcommitted environment
- Scale up from there to 1 TB

 All resources scaled up, though note that while additional DASD space was provided, it was on the same storage server.



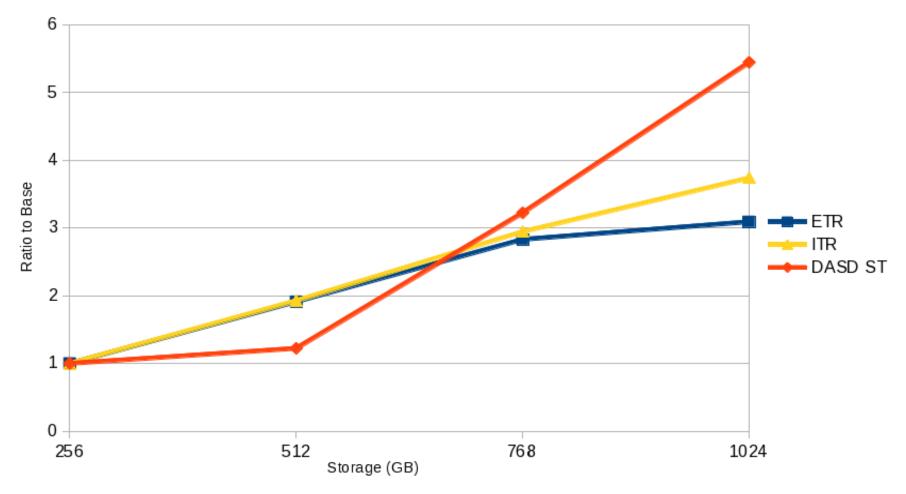
VIRSTOR Workload in Overcommitted Environment



ETR = External Throughput; ITR = Internal Throughput; DASD ST = DASD Service Time



Apache Workload in Overcommitted Environment



ETR = External Throughput; ITR = Internal Throughput; DASD ST = DASD Service Time

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Planning for Large Memory – Expanded Storage

- Convert Expanded Storage to Central Storage
 - Previously Expanded Storage helped mitigate the expense of stealing the lest than ideal page.
 - The Global Aging List and trial invalidation effectively removes the need for this expanded storage mitigation
 - The Early Writes also reduces the need for expanded storage
- Expanded Storage continues to be supported
 - -Attach to guests for testing and other uses.
- Statement of Direction: Support for Expanded Storage will be dropped in a future release



Planning for Large Memory – DASD Paging Space

- Please consult CP Planning and Administration book for new formula.
- Less over provisioning of paging DASD space is required in z/VM 6.3

 I/O Programs changed so no longer need contiguous slots on DASD for block sets
- Early Writes may increase the pages written
 - If your workload is not overcommitted at all, the consider turnig off Early Writes with SET AGELIST command
- PGT004 Abends still exist z/VM abends when all page and spool space has been used
- Set alerts to track and warn as appropriate
 - Operations Manager for z/VM



Planning for Large Memory – Reserving Memory

- Check or add SET RESERVED settings
- SET RESERVED DCSS MONDCSS 64M
 Reserve Memory for the Monitor DCSS
- Note, Live Guest Relocation does not preserve the reserved frames setting



Planning for Large Memory – Dump

- Plan enough dump space
 - http://www.vm.ibm.com/service/zvmpladm.pdf
- Recommend having dedicated dump space
- Learn how to use DUMPLD2 to process dumps as multiple files



Large Memory CP Monitor Changes

Domain	Record	Name	Туре	Title	Fields, N / D / C
D0	R3	MRSYTRSG	sample	Real Storage Data (Global)	DC
D0	R4	MRSYTRSP	sample	Real Storage Data (Per Processor)	D
D0	R6	MRSYTASG	sample	Auxiliary Storage (Global)	NC
D0	R7	MRSYTSHS	sample	Shared Storage Data	D
D0	R23	MRSYTLCK	sample	Formal Spin Lock Data	NC
D1	R7	MRMTRMEM	config	Memory Configuration Data	Ν
D1	R15	MRMTRUSR	config	Logged on User	С
D2	R4	MRSCLADL	event	Add User to Dispatch List	DC
D2	R5	MRSCLDDL	event	Drop User from Dispatch List	DC
D2	R6	MRSCLAEL	event	Add User to Eligible List	С
D2	R8	MRSCLSTP	event	System Timer Pop	D
D3	R1	MRSTORSG	sample	Real Storage Management (Global)	NDC
D3	R2	MRSTORSP	sample	Real Storage Activity (Per Processor)	D
D3	R3	MRSTOSHR	sample	Shared Storage Management	NC
D3	R14	MRSTOASI	sample	Address Space Information Record	NC
D3	R15	MRSTOSHL	event	NSS/DCSS/SSP Loaded into Storage	Ν
D3	R16	MRSTOSHD	event	NSS/DCSS/SSP Removed From Storage	NC
D4	R2	MRUSELOF	event	User Logoff Data	NDC
D4	R3	MRUSEACT	sample	User Activity Data	NDC
D4	R9	MRUSEATE	event	User Activity Data at Transaction End	DC



z/VM Performance Toolkit: Highlights

Changed screens:

- FCX102 SYSTEM, Some Internal System Counters
- FCX103 STORAGE, General Storage Utilization
- FCX133 NSS, NSS and DCSS Utilization and Paging Activity
- FCX146 AUXLOG, Auxiliary Storage Utilization, by Time
- FCX147 VDISKS, Virtual Disks in Storage
- FCX265 LOCKLOG, Spin Lock Log, by Time

Deleted screens:

- FCX254 AVAILLOG, Available List Management, by Time
- FCX259 DEMNDLOG, Demand Scan Details, by Time

New screens:

- FCX290 UPGACT, User Page Activity
- FCX291 UPGACTLG, User Page Activity (benchmarks a user)
- FCX292 UPGUTL, User Page Utilization Data
- FCX293 UPGUTLLG, User Page Utilization Data (benchmarks a user)
- FCX294 AVLB2GLG, Available List Data Below 2G, by Time
- FCX295 AVLA2GLG, Available List Data Above 2G, by Time
- FCX296 STEALLOG, Steal Statistics, by Time
- FCX297 AGELLOG, Age List Log, by Time

page state transition rates

page residency counts

available list counts

steal algorithm activity global aging list activity



z/VM Performance Toolkit: New Columns and Concepts

New Field	What this means
Inst	Instantiations: the rate at which valid memory is being created Instantiated: the amount of valid memory
Relse	Releases: the rate at which memory is being released
Inval	<i>Invalidations:</i> the rate at which demand scan is marking memory invalid as a way to determine whether it is being touched
Reval	<i>Revalidations</i> : the rate at which invalid pages are being made valid because somebody touched them
Ready	<i>Ready reclaims</i> or <i>ready steals:</i> the frame was found and selected for reclaim and had already been prewritten to auxiliary storage
Not Ready	<i>Notready reclaims</i> or <i>notready steals:</i> the frame was selected for reclaim but we had to wait for the auxiliary write (DASD) to finish before we could take it



z/VM Performance Toolkit: New Columns and Concepts

New Field	What this means
PNR	<i>Private, not referenced</i> : the page was read from aux as part of a block read, but it is still marked invalid because nobody has touched it yet
<i>x</i> <2G or <i>x</i> >2G	<i>Below 2 GB</i> or <i>Above 2 GB</i> : tells where the real backing frames are in real central
Sing	<i>Singles:</i> free frames surrounded by in-use frames (cannot coalesce)
Cont	Contigs: free frames in strings of two or more
Prot	<i>Protect threshold:</i> number of frames a singles-obtain must leave on a contigs-list

z/VM Performance Toolkit: New Report FCX292 UPGUTL

FCX292 Run 2013/04/10 07:38:36	UPGUTL User Page Utilization Data	Page 103
From 2013/04/09 16:02:10 To 2013/04/09 16:13:10		SYSTEMID CPU 2817-744 SN A6D85
For 660 Secs 00:11:00	"This is a performance report for SYSTEM XYZ"	z/VM V.6.3.0 SLU 0000

										Storad	• •								
					<										>				
	Data									<	- Inva	lid Bu	t Resi	dent -	>			Base	
	Spaces				<	Total	>	<-Lock	ed>	< U	FO>	< Pl	NR>	<-Age	List->			Space	Nr of
Userid	Owned	WSS	Inst	Resvd	T_A]]	T<2G	T>2G	L<2G	L>2G	U<2G	U>2G	P<2G	P>2G	A<2G	A>2G	XSTOR	AUX	Size	Users
>>Mean>>	.0	5284M	6765M	5611	5286м	27M	5259м	1010	232K	6565	2238K	59588	26M	53080	107м	.0	1815M	7108M	73
User Clas	ss Data	:																	
CMS1_USE	.0	3320ĸ	19M	.0	484K	.0	484K	.0	4096	.0	69632	.0	244K	.0	344к	.0	19м	2047м	1
LCC_CLIE	.0	364M	485M	.0	365M	11264	365м	.0	208K	.0	325K	.0	2686K	.0	8177K	.0	164M	1024M	8
LXA_SERV	.0	7974м	10G	.0	7978M	41M	7937м	.0	206к	9984	3327к	90624	39м	80725	161M	.0	2719M	10240м	48
User Data	a:																		
DISKACNT	.0	4976ĸ	5156K	0	4к	0	4к	0	0	0	4K	0	0	0	0	0	5152K	32M	
DTCVSW1	.0	184K	11M	0	196K	8к	188K	8к	4к	0	4K	0	0	0	168K	0	11M	32M	
DTCVSW2	.0	180ĸ	11M	0	184K	0	184K	0	4к	0	4K	0	0	0	164K	0	10M	32M	
EREP	.0	4912ĸ	4944K	0	4к	0	4к	0	0	0	4K	0	0	0	0	0	4940ĸ	32M	
FTPSERVE	.0	84K	5764K	0	88K	0	88K	0	4к	0	4K	0	0	0	76K	0	5760K	32M	
GCSXA	.0	204к	208K	0	8к	0	8к	0	4к	0	4к	0	0	0	0	0	200к	16M	
LCC00001	.0	364M	488M	0	365M	0	365M	0	204K	0	228K	0	2884K	0	8660K	0	192м	1024м	
LCC00002	.0	369м	492м	0	371M	20K	371M	0	204K	0	224к	0	2312K	0	7736к	0	159м	1024м	
	-			-		-		-		-		-		-		-			

Look for the new concepts: Inst IBR UFO PNR AgeList

0 252K

0 228K

0 2852K

0 2724K

0 8372K

0 8512K

0 215M 1024M

0 185M 1024M

• Amounts are in bytes, suffixed. Not page counts!

0 204K

0 204K

• FCX113 UPAGE is still produced.

0 364M

LCC00003

LCC00004

.0 363M

.0 363M

484M

483M

0 364M

0 363M 16K 363M

z/VM Performance Toolkit: New Report FCX292 UPGUTL

<				Res	sident					>		
	<> Invalid But Resident>											
<	Total	>	<-Lock	ed>	< UF	0>	< PI	NR>	<-Agel	_ist->		
T_A]]	T<2G	T>2G	L<2G	L>2G	U<2G	U>2G	P<2G	P>2G	A<2G	A>2G	XSTOR	AUX
365м	0	365м	0	204к	0	228к	0	2884к	0	8660K	0	192м

- Look for the new concepts: Inst IBR UFO PNR AgeList
- Amounts are in bytes, suffixed. Not page counts!
- FCX113 UPAGE is still produced.

z/VM Performance Toolkit: New Report FCX290 UPGACT

FCX290	Run 2013/04/10 07:38:36	UPGACT	Page	102
		User Page Activity		
From 2	013/04/09 16:02:10	SYSTEMID		
то 2	013/04/09 16:13:10	CPU 2817-744	SN A	6D85
For	660 Secs 00:11:00	"This is a performance report for SYSTEM XYZ" Z/VM V.6.3	.0 SLU	0000

. . . . -----> Storage -----> <----- Movement/s -----> Stl <--- Transition/s ----> <-Steal/s-> <Migrate/s> Nr of Userid Wt Inst Relse Inval Reval Ready NoRdy PGIN PGOUT Reads Write MWrit Xrel Users >>Mean>> 1.0 143K 5142 849K 718K 999K .0 .0 .0 958K 761ĸ .0 .0 73 User Class Data: CMS1 USE 1.0 15515 15801 2377 1632 5145 .0 .0 .0 .0 1980 .0 .0 1 LCC_CLIE 1.0 658K 20875 488K 486K 60875 .0 .0 .0 54212 22869 .0 .0 8 1.0 108K 1095 1191K 994K 1506K .0 .0 1447K 1153K LXA SERV .0 .0 .0 48 User Data: DISKACNT 1.0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 3072 2855 0 0 0 DTCVSW1 1.0 0 0 0 0 0 0 0 1.0 0 0 3004 2780 0 0 0 0 0 0 DTCVSW2 0 EREP 1.0 1434 1434 0 FTPSERVE 1.0 0 0 GCSXA 1.0 0 0 0 0 0 0 0 0 0 0 LCC00001 601K 18686 501ĸ 498K 65139 0 0 0 49866 23670 0 0 1.0 LCC00002 1.0 657K 24955 487ĸ 486K 54725 0 0 0 44522 18991 0 0 LCC00003 565K 23012 485K 481K 64065 0 0 0 44783 19859 0 0 1.0 LCC00004 1.0 602K 24104 499к 495K 63178 0 0 0 48811 24588 0 0 1.0 717K 25675 0 0 66002 28753 0 LCC00005 500K 499K 65865 0 0

Look for the new concepts: Inst Relse Inval Reval Ready NoRdy

z/VM Performance Toolkit: New Report FCX290 UPGACT

FCX290 Run 2013/04/10 07:38:36	UPGACT User Page Activity	Page 102
From 2013/04/09 16:02:10		SYSTEMID
то 2013/04/09 16:13:10		CPU 2817-744 SN A6D85
For 660 Secs 00:11:00	"This is a performance report for SYSTEM XYZ"	z/VM V.6.3.0 SLU 0000

	Stl -	< 1	ransit	tion/s	>	<-Stea	al/s->
Userid	Wt	Inst	Relse	Inval	Reval	Ready	NoRdy
>>Mean>>	1.0	143K	5142	849K	7 1 8K	999K	.0
User Class	5 Data	:					
CMS1_USE	1.0	15515	15801	2377	1632	5145	.0
LCC_CLIE	1.0	658K	20875	488K	486K	60875	.0
LXA_SERV	1.0	108K	1095	1191K	994K	1506K	.0
	0 0					•	
LCC00001 1.0 601K 1		498K 65139	0 0	0 49866 2		0	
LCC00002 1.0 657K 2 LCC00003 1.0 565K 2		486к 54725 481к 64065	0 0 0 0	0 44522 1 0 44783 1		0 0	
LCC00003 1.0 505K 2		495K 63178	0 0	0 48811 2		0	
LCC00005 1.0 717K 2		499K 65865	0 0	0 66002 2		0	

Look for the new concepts: Inst Relse Inval Reval Ready NoRdy

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z/VM Performance Toolkit: New Report FCX295 AVLA2GLG

FCX295 Run 2013/04/10 07:38:36	AVLA2GLG	Page 25
	Available List Data Above 2G, by Time	
From 2013/04/09 16:02:10		SYSTEMID
то 2013/04/09 16:13:10		CPU 2817-744 SN A6D85
For 660 Secs 00:11:00	"This is a performance report for SYSTEM XYZ"	z/VM V.6.3.0 SLU 0000

	<> Storage>							<-Fram	e Thre	sh>
Interval <avail< td=""><td>able></td><td><reque< td=""><td>sts/s></td><td><retur< td=""><td>rns/s></td><td><-Empt</td><td>y/s-></td><td>Sing</td><td><-Cont</td><td>igs-></td></retur<></td></reque<></td></avail<>	able>	<reque< td=""><td>sts/s></td><td><retur< td=""><td>rns/s></td><td><-Empt</td><td>y/s-></td><td>Sing</td><td><-Cont</td><td>igs-></td></retur<></td></reque<>	sts/s>	<retur< td=""><td>rns/s></td><td><-Empt</td><td>y/s-></td><td>Sing</td><td><-Cont</td><td>igs-></td></retur<>	rns/s>	<-Empt	y/s->	Sing	<-Cont	igs->
End Time Sing	Cont	Sing	Cont	Sing	Cont	Sing	Cont	Low	Low	Prot
>>Mean>> 23M	267M	47M	59м	47M	51M	.0	.0	1310	15	15
16:02:40 0	938м	32м	126м	502K	30310	.0	.0	1332	15	15
16:03:10 152K	4556K	50м	89м	49м	59м	.0	.0	1168	15	15
16:03:40 400к	4824к	68M	82м	71M	79м	.0	.0	1321	15	15
16:04:10 0	5896K	49м	72м	52M	70M	.0	.0	2409	15	15
16:04:40 0	2124к	40M	60м	41M	59м	.0	.0	1308	15	15
16:05:10 876к	3488K	54M	52м	55M	51M	.0	.0	1118	15	15
16:05:40 0	3624K	53M	58M	54M	57M	.0	.0	1409	15	15
16:06:10 2016к	4464К	49м	57M	51M	56M	.0	.0	1273	15	15

- Look for the new concepts: Singles Contigs Prot
- Amounts are in bytes, suffixed. Not page counts!
- FCX254 AVAILLOG is no longer produced.



z/VM Performance Toolkit: New Report FCX295 AVLA2GLG

FCX295 Run 2013/04/10 07:38:36	AVLA2GLG
	Available List Data Above 2G, by Time
From 2013/04/09 16:02:10	
то 2013/04/09 16:13:10	
For 660 Secs 00:11:00	"This is a performance report for SYS

	<> Storage>					<times> <-Frame Thresh></times>					
Interval	<avail< td=""><td>able></td><td><reques< td=""><td>sts/s></td><td><retur< td=""><td>ns/s></td><td><-Empt</td><td>y/s-></td><td>Sing -</td><td><-Cont</td><td>igs-></td></retur<></td></reques<></td></avail<>	able>	<reques< td=""><td>sts/s></td><td><retur< td=""><td>ns/s></td><td><-Empt</td><td>y/s-></td><td>Sing -</td><td><-Cont</td><td>igs-></td></retur<></td></reques<>	sts/s>	<retur< td=""><td>ns/s></td><td><-Empt</td><td>y/s-></td><td>Sing -</td><td><-Cont</td><td>igs-></td></retur<>	ns/s>	<-Empt	y/s->	Sing -	<-Cont	igs->
End Time	Sing	Cont	Sing	Cont	Sing	Cont	Sing	Cont	Low	Low	Prot
>>Mean>>	23M	267м	47M	59м	47м	51M	.0	.0	1310	15	15
16:02:40	0	938M	32M	126м	502K	30310	.0	.0	1332	15	15
16:03:10	152K	4556к	50м	89м	49м	59м	.0	.0	1168	15	15

- Look for the new concepts: Singles Contigs Prot
- Amounts are in bytes, suffixed. Not page counts!
- FCX254 AVAILLOG is no longer produced.



HiperDispatch

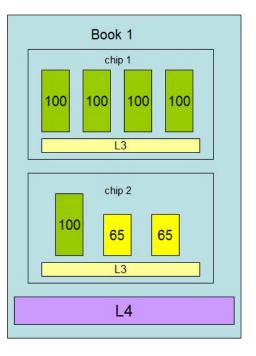


HiperDispatch: Highlights

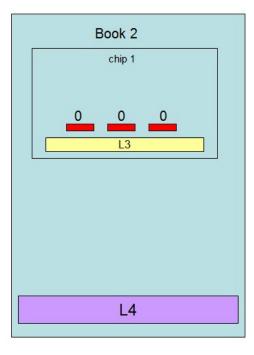
- Use of vertical mode partitions
 - -Running widely (larger n-way) if power is there
 - -Automatic reduction of MP level
- Topology-aware dispatching
- New or changed CP Monitor records
- New or changed z/VM Performance Toolkit screens
- How to plan for HiperDispatch



Vertical-Mode Partitions



Partition Topology



Features:

Concentrated entitlement Durable placement

Advantages:

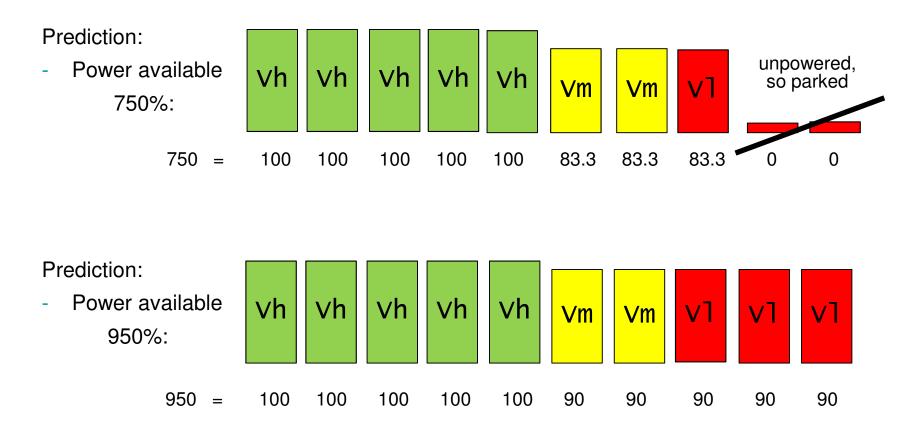
More secluded place to run

Opportunity to reduce MP level



Running Widely When The Power Is There

Entitlement: 630% via 5 Vh @ 100, 2 Vm @ 65, 3 VI @ 0



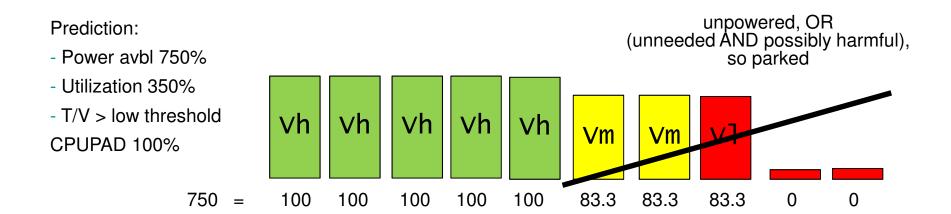
Vh = Vertical High; Vm = Vertical Medium; VI = Vertical Low

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Parking to Try To Reduce z/VM Overhead

Entitlement: 630% via 5 Vh @ 100, 2 Vm @ 65, 3 VI @ 0



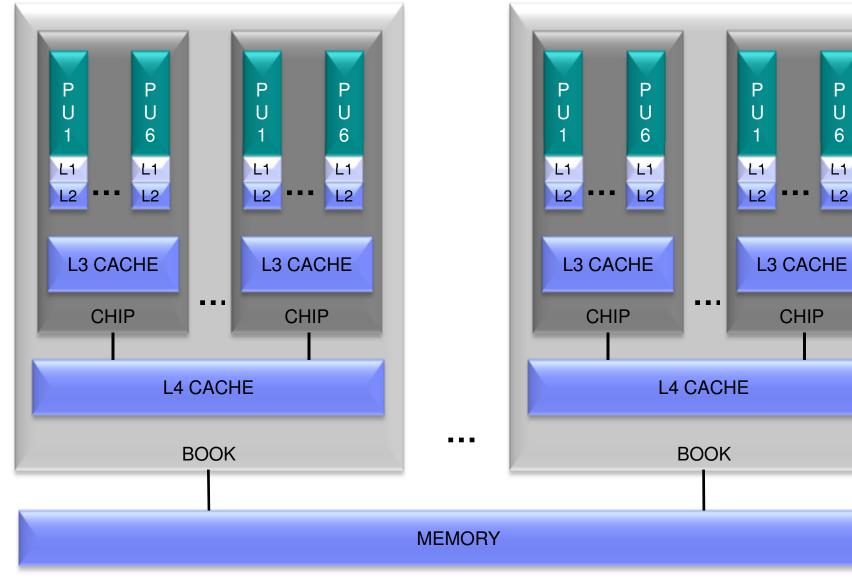
z/VM parks apparently unneeded processors, but only if T/V is projected high and load is projected below capacity. Safety margin is controlled via CP SET SRM CPUPAD.

Vh = Vertical High; Vm = Vertical Medium; VI = Vertical Low

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HiperDispatch- Dispatching Affinity



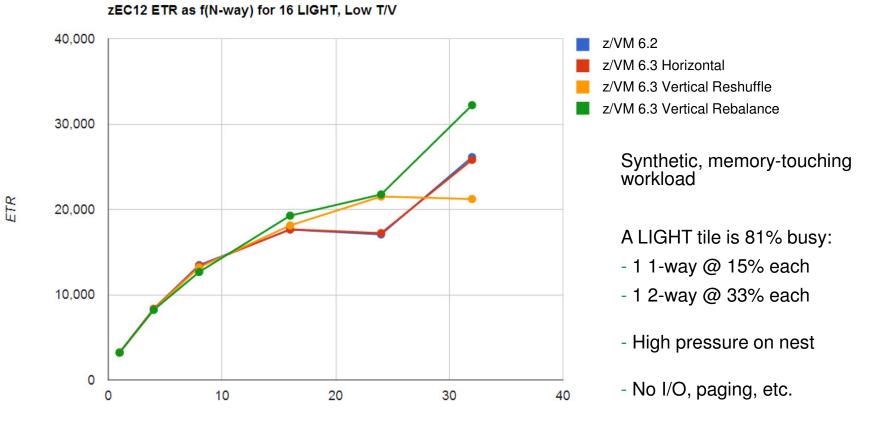


Topology-Aware Dispatching

- If can't send home, send close to home –Same chip? Same book?
- Try to keep a virtual MP's VCPUs together
- Try not to do long-drag steals —Cross-book, cross-chip
- Be smart about which real CPU we wake up —Same chip as stacked work? Same book?
- Rebalance:
 - -Optional replacement algorithm for Reshuffle
 - SET SRM DSPWDMethod
 - -only certain workloads are suitable
 - A few heavy users
 - Low Virtual:Logical PU ratio
 - Clearly distinguishable %CPU



Memory-Touching Workload, Light Edition

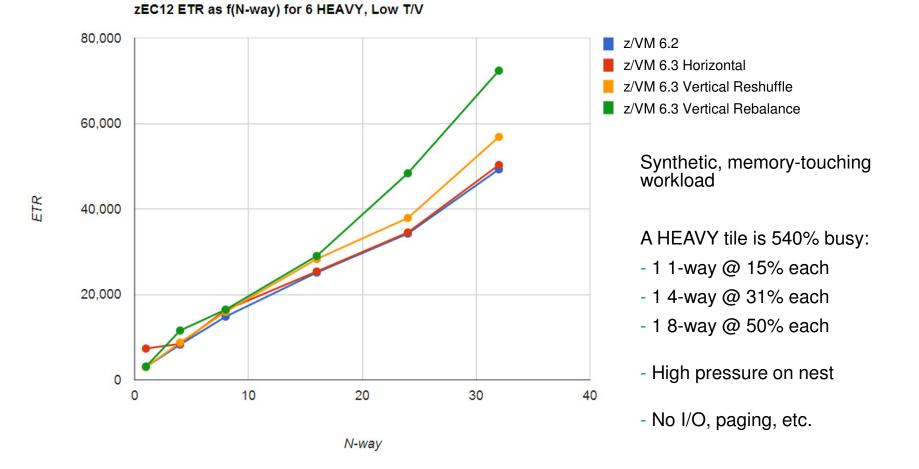


N-way

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Memory-Touching Workload, Heavy Edition



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Planning for HiperDispatch

Before:

- Decide what "success" looks like: metrics and values
- -Measure: transaction rates, MONWRITE data
- Ensure Global Performance Data Control was not disabled
 - -Logical partition activation profile
 - Default is enabled
- After:
 - -We think vertical and reshuffle are probably the right choice for you
 - Do same measurements
 - -Compare
- The way out:
 - CP SET SRM POLARIZATION HORIZONTAL
 - -SRM statement in system configuration file



Comments on Workloads

- Amenable workloads for z/VM HiperDispatch:
 - High-CPU, CPU-constrained workloads
 - Improving cache behavior stands to improve performance
 - Active VCPU : LCPU ratio isn't too large
 - High ratio has too much context switching to feel much effect
 - Runs in a partition having multiple topology containers
 - Gives z/VM an opportunity to separate guests from one another
- Compare those statements to IBM's statements about PR/SM and partitions



Comments on Workloads

- Indifferent workloads for z/VM HiperDispatch
 - Constrained by something else, such as I/O
 - Memory-overcommitted
 - High Virtual:Logical Processor ratio with every virtual CPU active just a little bit
 - Workloads with bad memory access habits
- Remember that vertical mode also keeps your partition away from the other partitions



HiperDispatch CP Monitor Changes

Domain	Record	Name	Туре	Description of Change
D0	R2	MRSYTPRP	sample	Added polarity, entitlement, and park-time accumulator
D0	R16	MRSYTCUP	sample	Added partition current weight
D0	R23	MRSYTLCK	sample	Added the HCPDSVTL topology lock
D1	R4	MRMTRSYS	config	Added bit indicating whether system is horizontal or vertical
D1	R5	MRMTRPRP	config	Added park state, polarization, entitlement, and topological location
D1	R16	MRMTRSCH	config	Added h/v bit, CPUPAD settings, and EXCESSUSE settings
D2	R7	MRSCLSRM	event	Added h/v bit, CPUPAD settings, and EXCESSUSE settings
D4	R2	MRUSELOF	event	Added rebalance results and steal results
D4	R3	MRUSEACT	sample	Added rebalance results and steal results
D5	R2	MRPRCVOF	event	Added park/unpark failure as reason varied off
D5	R3	MRPRCVON	event	Added parked as a state; use iff neither D5 R17 nor D5 R18 are seen
D5	R15 (new)	MRPRCDSV	event	Records assignment of processors to dispatch vectors
D5	R16 (new)	MRPRCPUP	event	Records park/unpark decision
D5	R17 (new)	MRPRCRCD	sample	Records processor's VMDBK steal behavior
D5	R18 (new)	MRPRCDHF	sample	Records PLDV population trends



z/VM Performance Toolkit

• Themes in the changes in existing Perfkit screens

- CPU entitlement appears in sensible places, e.g. FCX100 CPU
- Percent-parked appears in sensible places, e.g. FCX100 CPU
- Parked time is correctly accounted for, e.g. FCX126 LPAR %Susp
- SRM settings are reported where they ought to be, e.g. FCX154 SYSSET
- Interesting events are reported in FCX180 SYSCONF as they should
- Number of unparked CPUs appears in sensible places, e.g. FCX225 SYSSUMLG
- Counts of new monitor records appear in FCX155 MONDATA as they should
- Obsolete data is compatibly deleted in certain places, e.g. FCX144 PROCLOG

New reports sure to attract interest:

- FCX287 TOPOLOG shows a log of partition topology, container-major
- FCX298 PUORGLOG shows a log of partition topology, CPU-major
- FCX299 PUCFGLOG shows a log of the park/unpark state
- FCX301 DSVBKACT replaces the PLDV emptiness columns on FCX144 PROCLOG
- FCX302 PHYSLOG shows a physical CPU utilization log of the CEC by type pool
- FCX303 DSVSLOG replaces the PLDV steal columns on FCX144 PROCLOG
- FCX304 PRCLOG is where you should now look instead of FCX144 PROCLOG
- FCX306 LSHARACT reports the partitions' entitlements vs. logical CPU counts
- Obsolete reports
 - FCX144 PROCLOG is still there for now, but start using FCX304 PRCLOG instead



New Report PUORGLOG

1FCX298 Run 2013/05/20 10:39:48	PUORGLOG Processor Unit organization log
From 2013/05/19 03:39:31	
то 2013/05/19 03:41:31	
For 120 Secs 00:02:00	Result of GF003855 Run

Logical PU organization for Partition PPRF1 (GDLBOFVM)

Date	Time	CPU	туре	PPD	Ent.	Locat	ion
05/19	03:39:31	0	СР	VhD	100	1:6	
05/19	03:39:31	1	СР	VhD	100	1:6	
05/19	03:39:31	2	СР	VhD	100	1:5	Not
05/19	03:39:31	3	СР	VhD	100	1:5	
05/19	03:39:31	4	СР	VhD	100	1:5	Vh
05/19	03:39:31	5	СР	VhD	100	1:5	Vm
05/19	03:39:31	6	СР	VhD	100	1:5	VI
05/19	03:39:31	7	СР	VhD	100	1:4	
05/19	03:39:31	8	СР	VhD	100	1:4	Vhl
05/19	03:39:31	9	СР	VhD	100	1:4	Ent
05/19	03:39:31	10	СР	VhD	100	1:4	Loc
05/19	03:39:31	11	СР	VhD	100	1:2	
05/19	03:39:31	12	СР	VhD	100	1:2	
05/19	03:39:31	13	СР	VhD	100	1:2	
05/19	03:39:31	14	СР	VhD	100	1:2	

Notes:	
Vh	vertical high
Vm	vertical medium
VI	vertical low
VhD	vertical high, dedicated partition
Ent	entitlement (100% = 1 CPU's worth)
Location	book:chip (z10: book)

... truncated ...



New Report LSHARACT

1FCX306 Run 2013/06/24 09:36:54

LSHARACT Logical Partition Share

From 2013/02/19 11:49:58 To 2013/02/19 11:56:10 For 372 Secs 00:06:12

Result of GFCM0107 Run

LPAR Data, Collected in Partition RPRF2

Physical PUs, Shared: CP- 40 ZAAP- 2 IFL- 16 ICF- 1 ZIIP- 3 Dedicated: CP- 4 ZAAP- 0 IFL- 0 ICF- 0 ZIIP- 0

	- :			-	•	· - · · ·	•	•
-	c Partition	LPU	LPAR			<lpu td="" to<=""><td></td><td></td></lpu>		
	e Name			Entlment			Excess	
СР	RCPX4	10	10	59.3	•••	3.0	.0	
СР	RCTS1	5	10	59.3		311.9	252.6	-
СР	RCTS2	5	30	177.8		1.0	.0	-
CP	RCT1	20	30	177.8		111.3	.0	-
CP	RCT2	10	10	59.3		11.2	.0	
CP	REXT1	5	10	59.3		.0	.0	
CP	rext2	4	10	59.3		.0	.0	-
CP	RINS	10	10	59.3		.0	.0	0
CP	RPRF1	4	DED					
CP	RPRF2	24	335	1985.2		1548.4	.0	
CP	RSPX1	6	40	237.0		481.3	244.3	
CP	rspx2	6	40	237.0		499.7	262.7	0
CP	RSPX5	6	40	237.0		126.5	.0	0
СР	RST1	10	10	59.3		16.2	.0	
СР	RST1X	6	10	59.3		102.5	43.2	0
СР	RST2	6	50	296.3		.9	.0	0
CP	rst3	3	30	177.8		1.2	.0	0
ICF	RCTS2	1	10	25.0		.0	.0	-
ICF	RCT1	1	30	75.0		.0	.0	-
IFL	RCTS2	2	10	188.2		.0	.0	-
IFL	RCT1	2	30	564.7		.0	.0	u
IFL	RSTL1	16	45	847.1		449.2	.0	0
ZAA	P RCPX4	1	10	40.0		.1	.0	-
ZAA	P RCTS2	1	10	40.0		.0	.0	-
ZAA	P RCT1	1	30	120.0		.0	.0	u
ZII	P RCPX4	1	10	60.0		.3	.0	-
ZII	P RCTS2	1	10	60.0		.0	.0	-
	P RCT1	1	30	180.0		.0	.0	u

Features:

- Reports by partition and CPU type

You now have an easy way to see

the entitlements of your partitions.

- Reports entitlement in percent
- Reports percent-busy of the partition's CPUs of that type
- Reports whether the partition is consuming beyond its entitlement ("Excess")
- Reports LPU configuration wrt entitlement:
 - "o" overconfigured
 - "u" underconfigured
 - "-" apparently just right

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3 0

New Report LSHARACT

1FCX306 Run 2013/06/24 09:36:54 From 2013/02/19 11:49:58 To 2013/02/19 11:56:10 For 372 Secs 00:06:12

Logical Partition Share

LSHARACT

Result of GFCM0107 Run

LPA	LPAR Data, Collected in Partition RPRF2										
										_	
Pro	c Partition	LPU	LPAR				<lp< td=""><td>J Tota</td><td>al,%</td><td>⊳ LPU</td><td></td></lp<>	J Tota	al,%	⊳ LPU	
Тур	e Name	Count W	eight E	ntlment	Ty	peCap	В	usy Ex	kces	s Con	f
CP	RCPX4	10	10	59.3				3.0		0 O	
CP	RCTS1	5	10	59.3			31	1.9 2	252.	6 o	
CP	RCTS2	5	30	177.8			-	1.0		0 O	
CP	RCT1	20	30	177.8			11	1.3		0 O	
CP	RCT2	10	10	59.3			1	1.2		0 O	
CP	REXT1	5	10	59.3				.0		0 O	
CD		И	10	ICF RC	-1		30	75.0		<u> </u>	.0 -
	- "o" - overconf	igured		IFL RC		2	10	188.2		.0	.0 -
	"u" underee	afiaurad		IFL RC		2	30	564.7		.0	.0 u
	- "u" – undercoi	nigurea		IFL RS		16	45	847.1		449.2	.0 0
	- "-" - apparent	ly just right		ZAAP RCI ZAAP RCI		1 1	10 10	40.0 40.0		.1 .0	.0 - .0 -
		iy just right		ZAAP RC	-	1	30	120.0		-	.0 u
				ZIIP RC	PX4	1	10	60.0		.3	.0 -
				ZIIP RC		1	10	60.0			.0 -
				ZIIP RC	ΙL	T	30	180.0		.0	.0 u

New Report PUCFGLOG

1FCX299 Run 2013/06/24 09:36:54	PUCFGLOG Processor Unit Configuration log	Page	6
From 2013/02/19 11:49:52 To 2013/02/19 11:56:10	GFCM0107 CPU 2817-744	SN B6D8	85
For 378 Secs 00:06:18	Result of GFCM0107 Run z/VM v.6.3	.0 SLU 000)0

т	Type < Last	t> <	Next>
Date Time Type OnL Entitl	Cap CPUPAD EX Load XP	XPF T/V LCei XPF	T/V N NotVh UpCap LPU Unparked mask
02/19 11:49:54 CP 24 1985.2	100.0 70 2.2 1159.4	892.8 3.519 3.9 885.9 20	00.5 2 .0 200.0 00300000_0000000
02/19 11:49:56 CP 24 1985.2	100.0 70 .5 1153.3	888.1 256.0 1.7 883.4 20	01.3 2 .0 200.0 00300000_0000000
02/19 11:49:58 CP 24 1985.2	100.0 70 .5 1159.7	893.1 122.3 1.7 885.2 20	
02/19 11:50:00 CP 24 1985.2	100.0 70 .7 1136.7	875.4 53.45 1.7 857.7 17	
02/19 11:50:02 CP 24 1985.2	100.0 70 .9 1128.6	869.2 4.531 1.7 863.0 17	
02/19 11:50:04 CP 24 1985.2	100.0 70 1.3 1034.5	778.8 1.822 1.8 688.3 17	
02/19 11:50:06 CP 24 1985.2	100.0 70 .6 1157.1	891.1 38.57 1.8 856.4 16	
02/19 11:50:08 CP 24 1985.2	100.0 70 .5 1162.9	895.5 250.8 1.7 856.9 21	
02/19 11:50:10 CP 24 1985.2	100.0 70 44.8 1161.8	894.7 2.214 89.1 858.9 21	1.1 2 .0 200.0 00300000_0000000
02/19 11:50:12 * CPU Park/Unpark S	5		
02/19 11:50:12 CP 24 1985.2		881.9 1.517 354.6 858.5 19	97.6 5 .0 500.0 00300000_0000000
02/19 11:50:14 * CPU Park/Unpark S	5		
02/19 11:50:14 CP 24 1985.2	100.0 70 501.6 1155.6	890.0 1.009 803.5 858.3 19	97.5 10 .0 1000.0 013c0000_0000000
02/19 11:50:16 * CPU Park/Unpark S	5		
02/19 11:50:16 CP 24 1985.2		883.6 1.001 1497.6 857.9 14	6.5 16 .0 1600.0 0FFC0000_0000000
02/19 11:50:18 * CPU Park/Unpark S	5		
02/19 11:50:18 CP 24 1985.2		889.6 1.001 2199.1 857.7 13	30.3 23 100.0 2300.0 FFFF0000_0000000
02/19 11:50:20 * CPU Park/Unpark S	5		
02/19 11:50:20 CP 24 1985.2	100.0 70 2297.6 1179.7	908.5 1.001 2995.8 860.2 12	25.6 24 100.0 2400.0 FFFFE00_0000000
02/19 11:50:22 * CPU Park/Unpark S		981 4 1 005 2406 6 954 2 12	
02/19 11:50:22 CP 24 1985.2	100.0 70 2397.1 1144.5	881.4 1.005 2496.6 854.3 12	
02/19 11:50:24 CP 24 1985.2	100.0 70 2080.5 1181.8	910.1 1.002 2569.2 887.6 12	
02/19 11:50:26 CP 24 1985.2	100.0 70 1681.3 1140.0	878.0 1.002 2660.9 845.8 12 900.7 1.002 2684.7 886.2 1.	
02/19 11:50:28 CP 24 1985.2	100.0 70 1632.4 1169.6		
02/19 11:50:30 CP 24 1985.2 02/19 11:50:32 CP 24 1985.2	100.0 70 1587.7 1149.4 100.0 70 1878.3 1129.6	885.2 1.002 2635.4 869.6 1. 869.9 1.011 2560.8 854.7 1.	
.,		869.9 1.011 2560.8 854.7 1. 905.8 1.002 2425.8 884.3 1.	
02/19 11:50:34 CP 24 1985.2	100.0 70 1824.3 1176.2	903.0 1.002 2423.0 684.3 1.	007 24 100.0 2400.0 FFFFFF00_0000000

Look for: effect of high T/V, workload ramp-up, U' and XPF' values; power of a non-Vh

New Report DSVSLOG

1FCX303 Run 2013/05/20 10:32:38	DSVSLOG DSVBK Steals per logical CPU Log, by Time
From 2013/05/19 02:03:25	
то 2013/05/19 02:05:19	
For 114 Secs 00:01:54	Result of GF003820 Run

	С					Pct						
Interval	Р					Park	<	DSVBK S	Steal /s			>
End Time	U	туре	PPD	Ent.	DVID	тime	Lv1-00	Lv1-01	Lv1-02	Lv1-03	Lv1-04	Lv1-05
>>Mean>>	0	СР	Vh	100	0000	0	4.404	4.088	.000			
>>Mean>>	1	СР	Vh	100	0001	0	2.456	2.561	.000			
>>Mean>>	2	СР	Vh	100	0002	0	6.877	.921	.000			
>>Mean>>	3	СР	Vh	100	0003	0	7.596	.930	.000			
>>Mean>>	4	СР	Vh	100	0004	0	4.500	.482	.000			
>>Mean>>	5	СР	Vh	100	0005	0	3.614	.228	.000			
>>Mean>>	6	СР	Vh	100	0006	0	4.518	.482	.000			
>>Mean>>	7	СР	Vh	100	0007	0	2.912	.386	.000			
>>Mean>>	8	СР	Vh	100	0008	0	1.412	.421	.000			
>>Mean>>	9	СР	Vh	100	0009	0	1.386	.184	.000			
>>Mean>>	10	СР	Vh	100	000A	0	2.070	.544	.000			
>>Mean>>	11	СР	Vh	100	000в	0	2.114	.149	.000			
>>Mean>>	12	СР	Vh	100	000C	0	5.886	1.623	.000			
>>Mean>>	13	СР	Vh	100	000d	0	3.772	.702	.000			
>>Mean>>	14	СР	Vh	100	000e	0	3.026	.675	.000			
>>Mean>>	15	СР	Vh	100	000F	0	2.658	.360	.000			
>>Total>	16	СР	Vh	1600	MIX	0	59.202	14.737	.000			

Reports VCPU steal behavior by the distance the steal dragged the VCPU.

- Lvl-00: you stole it from a CPU in your chip
- Lvl-01: you stole it from a CPU in your book

- (z10: ... in your book) (z10: ... in another book)
- LvI-02: you stole it from a CPU on another book
- (z10: ... not applicable)

New Report PHYSLOG

1FCX302 Run 2013/06/24 09:36:54 PHYSLOG Real CPU Utilization Log From 2013/02/19 11:49:58 To 2013/02/19 11:56:10 For 372 Secs 00:06:12 Result of GFCM0107 Run

			Ir	nterval	<p< th=""><th>U Num></th><th>Total</th><th></th><th></th><th></th></p<>	U Num>	Total			
End Time	туре	Conf	Ded	Weight	%LgclP	%Ovrhd	LpuT/L	%LPmgt	%Total	ТуреТ/L
>>Mean>>	СР	44	4	675	3387.1	27.947	1.008	31.870	3446.9	1.018
>>Mean>>	ZAAP	2	0	50	.093	.042	1.451	.424	.559	6.015
>>Mean>>	IFL	16	0	85	448.16	1.017	1.002	2.108	451.28	1.007
>>Mean>>	ICF	1	0	40	.004	.003	1.624	2.257	2.263	563.66
>>Mean>>	ZIIP	3	0	50	.193	.090	1.465	1.204	1.487	7.694
>>Mean>>	>Sum	66	4	900	3835.5	29.099	1.008	37.864	3902.5	1.017

Features:

- Tallied by CPU type (CP, IFL, ...)
- Reports all three ways CPU gets used:
 - By logical CPUs
 - By PR/SM, chargeable
 - By PR/SM, unchargeable

New concepts:

- LPU T/L: like "guest T/V"
- Type T/L: like "system T/V"



Other Changes in z/VM 6.3

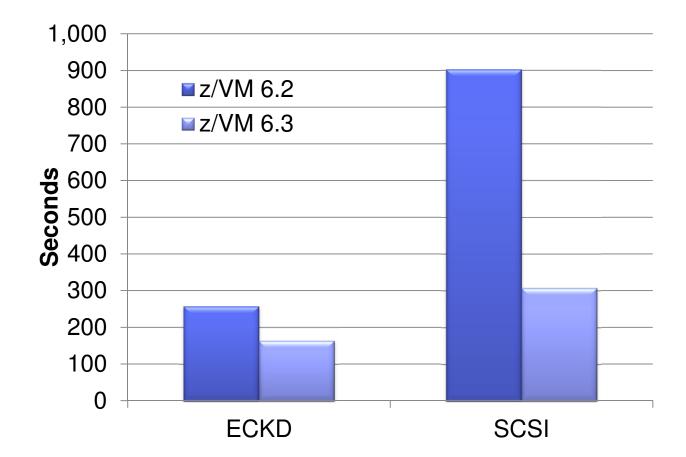


Large Memory Dump: Highlights

- Assures the system can produce a dump of a 1 TB system
- Includes changes to hard abend dump, SNAPDUMP, dump loader, and VM Dump Tool
- Changes to help improve the speed of hard abend dumps — Speed of dump-to-ECKD & dump-to-SCSI has been improved
- Can now stand-alone dump in hard-abend format to either ECKD or SCSI
- Emphasis on specifying DUMP operand on CP_OWNED statement
- Recovery of preallocated dump space after a SNAPDUMP
- SET DUMP can now list up to 30 DASD devices to receive the dump
- Monitor Domain 1 Record 17 new fields to describe frames constituting 2 MB buffer reserved for dump



Large Memory Dump: 256 GB System



SNAPDUMP Performance 256 GB System with approximately 4 GB of memory dumped.



z/VM 6.3: Other Performance Items

- FCP Data Router
 - Data movement assist for FCP card
 - Lets card move data from System z memory directly to its SCSI card
 - Available on z196 GA2 and later
- Local TLB Clearing Facility
 - Guests can now use IPTÉ or IDTE with the local-clearing-control (LC) bit
 - Ask your OS whether it does this
- Access-Exception Fetch/Store Indication Facility
 - More detail in Translation Exception Identifier when a storage access is denied
 - Ask your OS whether it can make use of this
- VSWITCH Recovery Stall Prevention

 Properly handle missing interrupt from Clear Subchannel (CSCH) related to uplink

 port
- CCW fast-trans and MDC both now allow Prefix-LRE CCW Important because some later Linux distros use Prefix-LRE
- VM65156 in base

-Missing path in guest mask no longer causes MDC bypass



Monitor Record Changes

- All the HiperDispatch changes
- All the Large Memory changes
- For FCP Data Router: D1 R19, D6 R25
- HiperSockets changes: D1 R19, D6 R25, D6 R26, D6 R27
- For Large Memory Dump: D1 R7, D3 R1
- VSWITCH Edge Port Aggregator: D6 R21, D6 R35
- VSWITCH Recovery Stall Prevention: D6 R22
- Additional debug: D0 R17, D0 R20, D3 R4, D3 R11, D5 R8, D5 R10, D6 R3, D6 R4, D6 R7, D6 R8, D6 R14, D6 R31, D9 R3



z/VM Performance Toolkit

- High Performance FICON changes

 SYSLOG, SYSTEM, DEVICE HPF, HPFLOG, SYSCONF, IOCHANGE, LCHANNEL all updated
- VSWITCH HiperSockets Bridge changes

 GVNIC, VNIC, GVSWITCH, VSWITCH, QDIO, IOCHANGE all updated
- LGR changes

 New reports LGRELOG and LGRDATA
- Large Memory Changes
 - 6 changed, 2 deleted, 8 new
- HiperDispatch Changes
 7 changed, 1 obsolete, 8 new



The CPU Measurement Facility

CPU Measurement Facility Counters

- CPU MF counters are a System z hardware facility that records the performance of the CPU and nest –Instructions, cycles, cache misses, … processor-ish stuff
- Available on zEC12, zBC12, z196, z114, and z10 EC & BC
- The CPU MF counter values:
 - Help IBM to understand how your workload stresses a CEC
 - -Help IBM to map your workload into the LSPR curves
 - Help IBM to understand your system when there is a processor performance problem
- Support in base of z/VM 6.2 and z/VM 6.3
 - -z/VM 5.4 and 6.1: VM64961, UM33440 (5.4)
 - Counters come out in a new Monitor record, D5 R13 MRPRCMFC
- We want volunteers to send us MONWRITE data!
 - -Your contributions will help us to understand customer workloads!



CPU MF Counters and CP Monitor, Details

- Counter sample record is in the Processor domain –Monitor record, D5 R13 MRPRCMFC
- MONITOR SAMPLE command manipulates counter collection
- QUERY MONITOR reveals whether counter collection is on
- Requires enabling on HMC
- The D5 R13 records land in your MONWRITE data



IBM Wants Your CPU MF Counter Data

- Your data will help IBM to build a library of customer workloads
- Collect an hour's worth of MONWRITE data...
 - -From a peak period, <- very important!
 - -With CPU MF counters enabled,
 - -With one-minute sample intervals
- Contact Richard Lewis at rflewis at us.ibm.com
- Richard will send you instructions on how to transmit the data to IBM
- No deliverable will be returned to you
- We will be ever grateful for your contribution



Evolution of z/VM LSPR Workload

- From memory-rich to memory-constrained
- From 16-way to 32-way
- From equally-active to unequally-active
- From workload-indexed to RNI-indexed
 We do want your CPU MF counter data
- Our goal is a lab setup that represents z/VM customers' environments



Summary



Summary

- z/VM 6.3 is a performance release
 - -Large memory: we expect scaling to 1 TB
 - -HiperDispatch: we expect improvements for amenable workloads
 - -Large memory dump: necessary for large memory
 - -Lots of CP Monitor and z/VM Performance Toolkit changes
- Keep that CPU MF data coming Richard wants to hear from you
- See <u>http://www.vm.ibm.com/perf/reports/zvm/html/index.html for full</u> <u>z/VM 6.3</u> Performance Report.