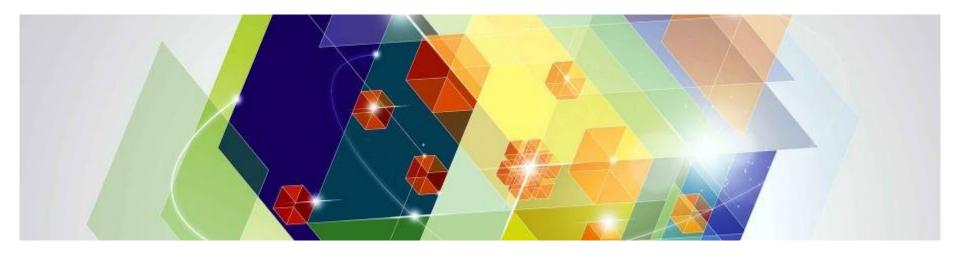


System z Flash Express

Introduction, Uses, and Benefits

SHARE in San Francisco February 5th, 2013 Session 13086

Elpida Tzortzatos: elpida@us.ibm.com Peter Szwed: pkszwed@us.ibm.com





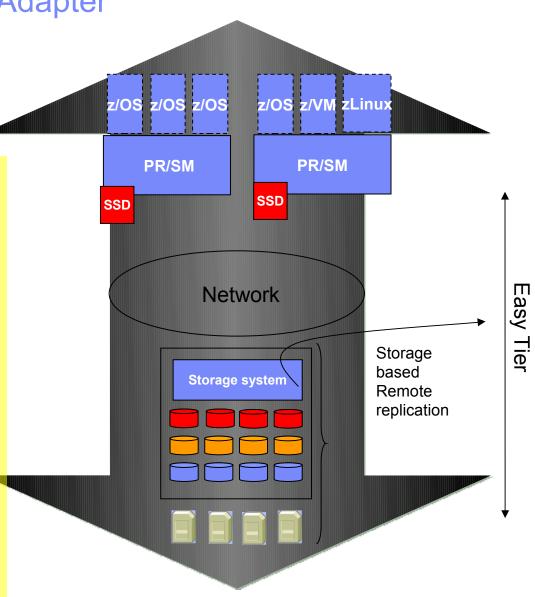
Agenda

- z/OS Customer Value Proposition
- System z Flash Express and z/OS
- Flash Performance Results
- z/OS Flash Roadmap
- Under the Covers Implementation Highlights

System z Flash Express IO Adapter



- Flash Express is a PCIe IO adapter with NAND Flash SSDs (Solid State Drives)
- Flash Express is accessed using the Extended Asynchronous Data Mover
 - Optimized software path for Flash Access based on prior learning with z expanded store
- Flash Express provides continuous availability
 - RAID 10 to cover adapter failure
 - Concurrent Firmware update to cover service
- Flash Express is fully virtualized
 - A single adapter pair can provide Flash to 60 partitions on a CEC
 - Adapter RAS (call home, recovery, etc.) done at system level, not in OS.
 - Transparent migration to new adapter technology





IBM Flash Express – Smarter Availability for Smarter Systems

- Flash Express is an innovative solution designed to help you compete effectively in today's marketplace
 - Automatically improve availability for key workloads at critical processing times
 - Drive availability and performance for workloads that cannot tolerate paging spikes or inconsistent performance
 - Slash latency for critical application processing such as diagnostics collection
- Extends IBM's expertise in memory management introducing a new tier of memory using Flash Express
- Provides a secured, resilient and immediately usable solution
- Planned Flash Express and pageable large page exploiters:

- z/OS V1.13 Language Environment -IMS 12 Common Queue Server -DB2 10 *

- Java SDK601 SR4, and Java SDK7 SR3 and by extension:
 - CICS Transaction Server 5.1
 - WAS Liberty Profile v8.5
 - IMS 12
 - DB2

Traditional WAS 8.0.0x and Traditional WAS 8.5.5 (future) **

4 *DB2 date to be determined. Support for V10 with APARs is planned.

Flash Express Strengthens Availability



- Innovation to drive availability to exceptional levels
 - Extends IBM's expertise in memory management introducing a new tier of memory using Flash Express
 - Is an *industry unique* application of Flash to improve availability
 - Takes the next step in advanced memory management
- Flash Express can improve availability and reduce latency
 - Improves availability during transition periods and spikes
 - Helps accelerate start of day processing batch to online
 - Enables faster snapshots of diagnostics (e.g. SVC dump, standalone dump)
 - With pageable large pages can improve performance of DB2 and Java
 - Ideal for applications with random read access and high read/write ratios
- Helps customers deliver vigorous service levels
 - Designed to help provide *exceptional* availability and fast response time
 - Delivered in tandem with pageable large pages for *superior* performance
- Minimal configuration- no special skills needed
 - Usable immediately; no special training required
 - Easy to set up and dynamically configurable

Representative Use Cases - Flash Express



Flash Express can reduce latency delays from paging to bring system availability to new heights and improve overall service levels

Application related errors will require collection of diagnostics. These diagnostics can be collected faster with Flash Express, reducing paging related delays that can impact your overall availability.

Having your working data resident in Flash can help accelerate start of day processing, and improve service for many industries at the busiest time of their work day- a time when they cannot afford disruptions.

DB2[®] and Java[™] in memory buffer pools work to store and process application data. DB2 and Java can benefit from 1MB pageable large pages with Flash Express, improving overall performance.



Flash Express – What is it?

FLASH Express

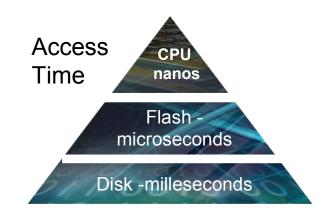
- Flash Express is a PCIe IO adapter with NAND Flash SSDs
- Physically comprised of internal storage on Flash SSDs
- Used to deliver a new tier of memory- storage class memory
- Uses PCIe I/O drawer
- Sized to accommodate all LPAR paging
 - Each <u>card pair</u> provides **1.4 TB** usable storage (2.8 TB total)
 - Maximum 4 card pairs (4 X1.4=5.6 TB)
- Immediately usable
 - Simplifies capacity planning
 - No intelligent data placement needed
 - Full virtualization across partitions
- Robust design
 - Delivered as a RAID10 mirrored pair
 - Designed for long life
 - Designed for concurrent firmware upgrade

Secured

- Flash Express adapter is protected with 128-bit AES encryption.
- Key Management provided based on a Smart Card
- Secure Cryptographic Erase meets audit requirements



One Flash Express Card



Flash memory (SCM) blurs the distinction between memory and storage characteristics

IBM Flash Express – Smarter Availability for Smarter Systems Outstanding Availability and Performance - Innovative Flash Express

- Companies competing for the highest quality of service in today's market must deliver outstanding availability and performance
- Changes in workload processing can impact service levels at critical processing times
- Flash Express is an innovative solution designed to help you improve availability and performance to compete effectively in today's market
 - Automatically improves availability for key workloads at critical processing times
 - Drives availability and performance for workloads that cannot tolerate paging spikes or inconsistent performance
 - Slashes latency for critical application processing such as start of day processing and also collection of diagnostics (SVC dumps, standalone dumps)
 - Delivered as a new adapter card in the PCIe I/O drawer

Benefits

- Improves availability and performance helping companies achieve highest service levels
- Delivers a secured, resilient and immediately usable solution
- Automatic, requires minimal setup, no special training needed





System z Flash and z/OS

Allocating z FLASH



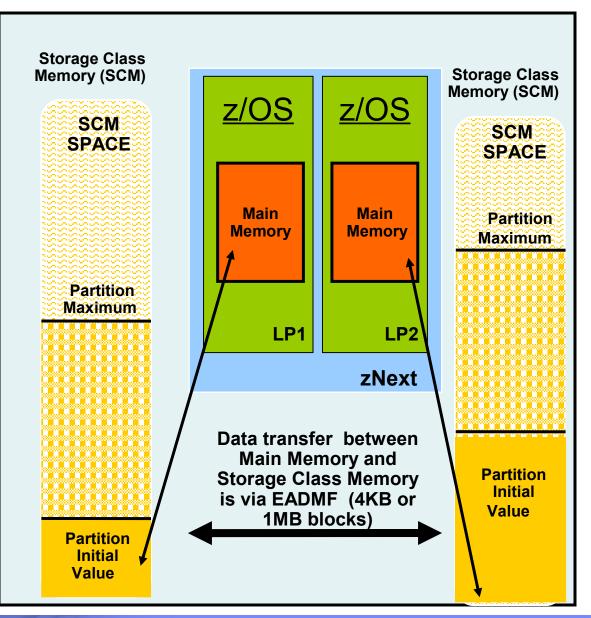
Allocating Flash to a partition

- The initial and maximum amount of Flash Memory available to a particular logical partition is specified at the SE or HMC via a new Flash Memory Allocation panel
- Can dynamically change maximum amount of Flash Memory available to a logical partition
- Additional Flash Memory (up to the maximum allowed) can be configured online to a logical partition dynamically at the SE or HMC
 - For z/OS this can also be done via an operator command
- Can dynamically configure Flash Memory offline to a logical partition at the SE or HMC
 - For z/OS this can also be done via an operator command
- Predefined subchannels, no IOCDS

	.6.164 https://9.12.16					_						
N ÇÎ	lanage Flash A	llocation	1 - P87									
Summa	arv											
	ated: 976 GE	2	Storago in	crement: 16 GB								
	aled. 970 GE able: 1872 GE		Rebuild co									
	tialized: 0 GE		Trobund oo.	inploto. 070								
Unav	ailable: 0 GE	3										
Total:	2848 GE	3										
Partitic	ne											
						_						
*	Select Act	10n	•									
Select	Partition Name	Status	IOCDS	Allocated (GB)	Maximum (GB)							
۲	R70	Active	A0,A1,A2,A3	48	2848							
0	R71	Active	A0,A1,A2,A3	128	2848							
0	R72		A0,A1,A2,A3		2848							
0	R73		A0,A1,A2,A3		2848							
0	R74	Active	A0,A1,A2,A3	80	2848							
0	R75	Active	A0,A1,A2,A3	80	2848							
0	R76	Active	A0,A1,A2,A3	64	2848							
0	R77	Active	A0,A1,A2,A3	64	80							
	D7D											
0	R7B	 ○ R7F Active A0,A1,A2,A3 32 64 										
0		Active	10,11,12,10	Refresh								

z FLASH Virtualization

- Full virtualization of physical Flash PCIe cards across partitions, software sees an Abstracted Flash Storage Space...
 - Allows each logical partition to be configured with its own SCM address space
 - Allocate Flash to partitions by amount, not card size
 - Ability to change underlying technology while preserving API
 - No Hardware Specifics in Software.
 - Error Isolation, Transparent mirroring, Centralized diagnostics, etc.
 - Hardware Logging, FRU Call, Recovery: Independent of software

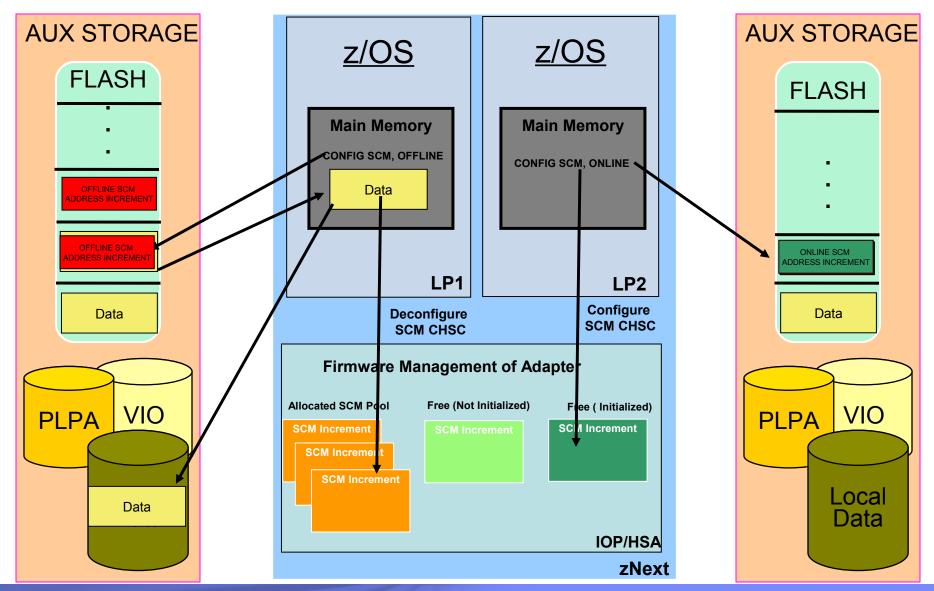


FLASH for z/OS Paging Value

Flash Memory is a faster paging device as compared to HDD

- •The value is NOT in replacing memory with Flash but replacing disk with Flash
- •Flash is suitable for workloads that can tolerate paging and will not benefit workloads that cannot afford to page
- •The z/OS design for Flash Memory does not completely remove the virtual storage constraints created by a paging spike in the system. (Some scalability relief is expected due to faster paging I/O with Flash Memory.)

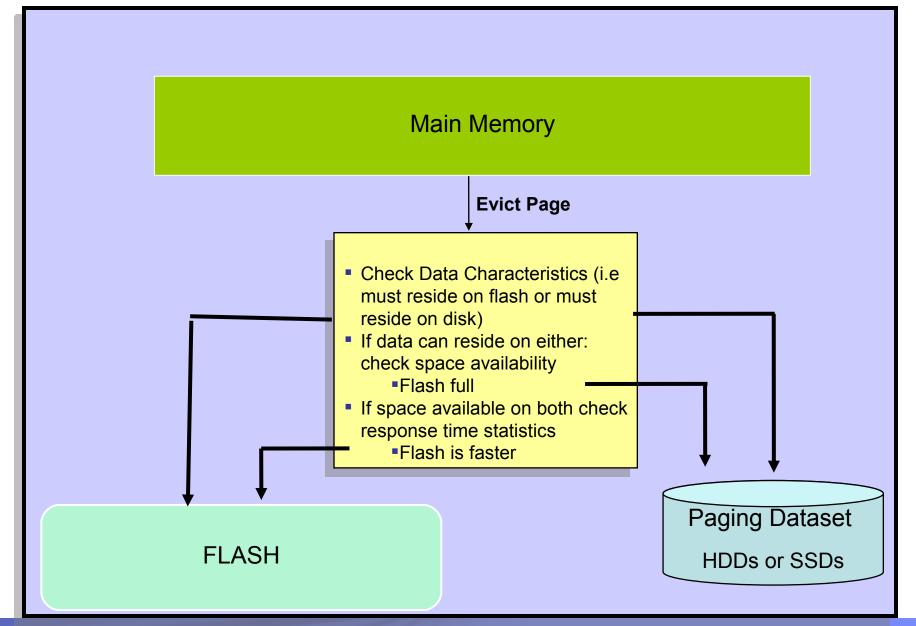
A z/OS Flash Configuration



Typical Customer Configurations for FLASH

- Flash card pair memory size is 1.4TB
 - Min: 1 Card Pair
 - Max: 4 Card Pairs
- •Typical customer configuration is 6 to 8 LPARs per CEC and 40GB 80GB for paging configuration dataset size
- •Even with 10 LPARs per CEC, each LPAR has 140 GB of Flash Memory available for its paging datasets, more than double the current typical customer configuration.
 - All paging data can easily reside on Flash
 - Data will preferably go to Flash and only go to disk (if any) when Flash is full
 - No intelligent placement of data on internal Flash needed

Flash vs Disk Placement Criteria



Flash vs Disk Placement Criteria

Data Type	Data Page Placement
PLPA	At IPL/NIP time PLPA pages will be placed both on Flash and disk.
VIO	VIO data will always be placed on disk (First to VIO accepting datasets with any spillover flowing to non-vio datasets)
Pageable Large Pages	If contiguous Flash space is available, pageable large page will be written to Flash. If Flash is not available in the system configuration pageable large pages will be backed with 4k page frames.
All other data	If available space exists on both Flash and disk then make a selection based on response time.

z/OS FLASH Use Cases

Paging

z/OS paging subsystem will work with mix of internal Flash and External Disk

-Self Tuning based on measured performance

-Improved Paging Performance, Simplified Configuration

•Begin Paging 1 MB Large Pages only on Flash

-Exploit Flash's random IO read rate to get CPU performance by enabling additional use of Large Pages. Currently large pages are not pagable.

•Begin Speculative Page-In of 4K Pages, 1MB Pages only on Flash

 Exploit Flash's random IO read rate to get Improved Resilience over Disruptions.

-Market Open, Workload Failover

Flash Memory Usage and Invocation

- New PAGESCM= keyword in IEASYSxx defines the amount of flash to be reserved for paging
 - Value may be specified in units of M, G, or T
 - NONE indicates do not use flash for paging
 - ALL (default) indicates all flash defined to the partition is available for paging

New messages issued during IPL indicate the status of SCM

• IAR031I USE OF STORAGE-CLASS MEMORY FOR PAGING IS ENABLED -PAGESCM=ALL, ONLINE=00065536M

OR

• IAR032I USE OF STORAGE-CLASS MEMORY FOR PAGING IS NOT ENABLED -PAGESCM=NONE

- The D ASM and D M commands are enhanced to display flash-related information/status
 - •D ASM lists SCM status along with paging data set status
 - •D ASM,SCM displays summary of SCM usage
 - •D M=SCM display SCM online/offline and increment information
 - •D M=SCM(DETAIL) displays detailed increment-level information

The CONFIG ONLINE command is enhanced to allow bringing additional SCM online

CF SCM(amount),ONLINE

P87: Operating System Message	ges						• X	
2012200 17.27.02 R71	IEE200I 17.27.	02 DISPLAY	ASM 143				P87:R71	
	TYPE FULL	STAT DEV	DATASET NAME					
	PLPA 38%	OK 2002	SYS1.R71.PLPA					
	COMMON 68	OK 2002	SYS1.R71.COMMON					
	LOCAL 0%	OK 2003	SYS1.R71.LOCAL					
	LOCAL 0%	OK 2021	SYS1.R71.LOCAL1					
	LOCAL 0%	OK 2261	SYS1.R71.LOCAL4					
	LOCAL 0%	OK 2269	SYS1.R71.LOCAL5					
	SCM 0%	OK N/A	N/A					
	PAGEDEL COMMAN							
2012200 17.27.45 R71	IEE207I 17.27.	45 DISPLAY	ASM 148					
		LL	SIZE	USED	IN-ERROR			
		0%	33,554,432	13,865	0			
2012200 17.28.04 R71	IEE174I 17.28.							
	STORAGE-CLASS	MEMORY STAT	TUS					
	2848G DEFINED							
	ONLINE							
	0G-128G					=		
	1872G OFFLINE-	AVAILABLE						
	0% IN USE							
	SCM INCREMENT	SIZE IS 160	;			•		
Command:						-		
Priority (select this when re	esponding to priority	(red) messag	jes)					
	1.4.							
<u>Send</u> <u>R</u> espond <u>D</u> e	elete							
Close Help								

A P87: Operating System Messages	Flash Mele							
2012200 17.30.15 R71 IEE174I 17.30.15 DISPLAY M 163		P87:R71						
STORAGE-CLASS MEMORY STATUS - INCREMENT DETAIL								
2848G DEFINED								
ADDRESS IN USE STATUS								
0G 0% ONLINE								
16G 0% ONLINE								
32G 0% ONLINE								
48G O% ONLINE								
64G O% ONLINE								
80G 0% ONLINE								
96G O% ONLINE								
112G O% ONLINE								
ONLINE: 128G OFFLINE-AVAILABLE: 1872G PENDING OFFLINE: 0G								
0% IN USE	=							
SCM INCREMENT SIZE IS 16G	•							
Command: d m=scm(detail)	•							
Priority (select this when responding to priority (red) messages)								
Send Respond Delete								
Close Help								

Display ASM Command

d asm				
IEE200I	17.17.	.46 DI	SPLAY	ASM 944
TYPE	FULL	STAT	DEV	DATASET NAME
PLPA	100%	FULL	02E6	SYS1.PLPA.PAGCOM
COMMON	61%	OK	02E6	SYS1.COMMON.PAGCOM
LOCAL	08	OK	098E	SYS1.LOCAL.PAGEP2
LOCAL	08	OK	0987	SYS1.LOCAL.PAGEP3
LOCAL	08	OK	098F	SYS1.LOCAL.PAGEP4
SCM	118	OK	N/A	N/A

d asm, scm

IEE207I 17.35.02 DISPLAY ASM 947

 STATUS
 FULL
 SIZE
 USED
 IN-ERROR

 IN-USE
 11%
 16,777,216
 2,096,144
 0

Flash Related Commands

D M=SCM

IEE174I 17.57.26 DISPLAY M 230 STORAGE-CLASS MEMORY STATUS 80G DEFINED ONLINE 0G-64G 16G OFFLINE-AVAILABLE 14% IN USE SCM INCREMENT SIZE IS 16G

D M=SCM(DETAIL)

IEE174I 17.57.30 DISPLAY M 232 STORAGE-CLASS MEMORY STATUS - INCREMENT DETAIL 80G DEFINED ADDRESS IN USE STATUS 0G 55% ONLINE 16G 0% ONLINE 32G 0% ONLINE 48G 0% ONLINE 0NLINE: 64G OFFLINE-AVAILABLE: 16G PENDING OFFLINE: 0G 14% IN USE SCM INCREMENT SIZE IS 16G

CF SCM(16G),ONLINE

IEE195I SCM LOCATIONS 64G TO 80G ONLINE IEE712I CONFIG PROCESSING COMPLETE

RMF Paging Activity Report

Session E	[43 x 132]				_			
<u>E</u> ile	<u>Edit Edit_Settings</u>	<u>M</u> enu <u>U</u> ti	lities <u>C</u> ompi					
VIEW Command	D10.FLASH.R70.RM		*	****	** Ton of D	ata ******	*****	Columns 00001 00124 Scroll ===> <u>CSR</u>
								1514 Line(s) not Displayed
001515	1				PAGING	ACTIV	ΙΤΥ	
001516	z/OS V1	R13	SYSTEM	ID R7F		DATE 07/19	/2012	INTERVAL 15.00.000
001518			RPT VE	RSION V1R	13 RMF	TIME 16.00	.00	CYCLE 1.000 SECONDS
	-OPT = IEAOPTRD LF							
001520	OMEMORY OBJECTS							
001522								
001523	MIN	65	Θ	4				
001524 001525	MAX AVG	65 65	0 0	4 4				
	01 MB FRAMES		FIXED			- PAGEABLE		
001527			AVAILABLE	IN-USE		AVAILABLE	IN-USE	
001528	MIN	2,048	1,874		1,045	799	246	
001529 001530	MAX AVG	2,048 2,048	1,874 1,874	174	1,045 1,045	799 799	246 246	
	OHIGH SHARED FRAMES	Z,048 TOTAL		174 STORAGE	1,040	AUX DASD	AUX SCM	
001532								
001533	MIN	136902.1M		1,295		Θ	Θ	
001534	MAX	136902.1M		1,295		0	0	
001535	AVG OHIGH COMMON FRAMES	136902.1M TOTAL	CENTROL	1,295 STORAGE	FIXED 4K	0 AUX DASD	0 AUX SCM	
001537								
001538	MIN	17301504			3,194,870	258,137	438,688	
001539	MAX AVG	17301504			3,194,870	258,137	438,688	
001540	HVG 	17301504		257,190	3,194,870	258,137	438,688	152 Line(s) not Displayed
*****	*****	*****	*****	******	* Bottom of	Data ****	*****	****
MA F						04/015		
Connect	ed to remote server/host xrfmcl.pok.ib	m.com using lu/pool	M04TC079 and port 23					

RMF Paging Activity Report (cont)...

Session	F - [43 x 132]		100	_											
	and the second se	E <u>d</u> it_Sett	ings	Menu	<u>U</u> tilities	<u>C</u> ompil	ers <u>I</u> e	est <u>H</u> el	p						
VIEW	D10	.FLASH.R7	0.RMF	. OUTPUT	(SCM0719)	- 01.00)			*****	******	****		Columns 0000 Scroll == ********	==> <u>CSR</u>
		·						·		·				1545 Line(s) not Di	
<u>0</u> 01546		OF SAMPL	ES =	900)		PAGE	E DATA S	et and s	ICM USAG	iΕ				
001547	7 3 -PAGE									~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~	PAGE			v	
001549		VOLUME	DEV	DEVICE	E SLOTS	S	LOTS US	SED	BAD	ĩN	TRANS	NUMBER			
001550		SERIAL	NUM	TYPE	ALLOC	MIN	MAX	AVG	SLOTS	USE	TIME	IO REQ		0 data set name	
	-PLPA	R7FPG1		33903	35999	13951	13951	13951	Θ	0.00	0.000	Θ	Θ	SYS1.R7F.PLPA	
	2 OCOMMON			33903	35999	1945	1945	1945	0		0.000	0	Θ	SYS1.R7F.COMMON	
	3 OLOCAL 1 OLOCAL	R7FPG2 R7FPG3		33903 33903	597599 597599	17028 15839	$17051 \\ 15855$	$17038 \\ 15846$	0 0		0.000 0.000	43 82	44 111		
	6 OLOCAL	R7FPG3		33903	593999	16231	16236	16233	0 0		0.000	36	46		
	6 OLOCAL	R7FPG5		33903	593999	16255	16264	16257	Ö		0.000	49	72		
001557	7 OLOCAL	R7FPG6	2258	33909	1620K	14622	14630	14625	Θ	0.00	0.000	47	55		
001558	B OSCM	NZA	NZA	NZA	8389K	385374	386033	385673	Θ	0.00	0.000	1,234	1,229	N/A	
								·						134 Line(s) not Di ******	
*****	*****	***	****	****	* * * * * * * * * * * * *	• • • • • • • • • •	****	*** 8011	UM UT Da	119 ****	****	*****	***	*****	****
MA F									6	7/002					
Conne	cted to remote	server/host xrfmo	l.pok.ibm.	com using l	u/pool M04TC079	and port 23									1
		al factoria inte			THE CONTRACT		1000 C.V.V	100	and a second second						

RMF Monitor II – Page Data Set Activity

ℬ ☐ ibmuser						
<u>File Edit View Communi</u>	ication <u>A</u> ctions <u>W</u> ir	ndow <u>H</u> elp				
o re re a se		1	i 🖬 🗎	٠		
	RMF	- PGSP	Page Da	ta Set Act	tivity	Line 1 of 9
		CPU=	6 UI	C= 65K PR=	= Θ	System= 4381 Total
S VOLUME DEV	DEV	SLOTS	PAGE	I/O REQ A	AVG PAGES	08:56:43
T SERIAL NUM			TRAN TIM		PER I/O V	DATA SET NAME
P PAGCOM 02E6 C PAGCOM 02E6 L PAGEP2 098E L PAGEP3 0987 L PAGEP4 098F L PAGE08 0981 L PAGETD 02E5 L PAGE99 0484 S N/A N/A	33903 (33903 33903 33903 33903 33903	$ \begin{array}{r} 100.0 \\ 50.31 \\ 0.00 \\ 0.00 \\ 0.00 \\ 0.00 \\ 0.00 \\ 0.00 \\ 0.00 \\ 4.74 \\ \end{array} $	$\begin{array}{c} 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000\\ 0.000 \end{array}$	0.000 0.000 0.000 0.000 0.000 0.000 0.000 0.000 0.310	0.000 0.000 0.000 Y 0.000 Y 0.000 Y 0.000 Y 0.000 Y 166.667	SYS1.PLPA.PAGCOM SYS1.COMMON.PAGCOM SYS1.LOCAL.PAGEP2 SYS1.LOCAL.PAGEP3 SYS1.LOCAL.PAGEP4 SYS1.COMMON SYS1.LOCAL.PAGETD PAGE2 N/A
Command ===> F1=HELP F7=UP MA e	F2=SPLIT F8=D0WN		SWAP	F4=RETUF F10=LEFT	F11=RI	

SVC Dump Statistics

• VERBX IEAVTSFS

 Shows total dump capture time, system/task nondispatch time, page operations required to dump requested address space (real-to-real copies, page-ins, etc)

SVC Dump Statistics (cont)

Dump start

Dump end

Total dump capture time

00:00:14.357089

10/09/2012 14:30:29.867495

10/09/2012 14:30:44.224584

System nondispatchability	start	10/09/2012	14:30:29.870030

System set nondispatchable 10/09/2012 14:30:29.870048

Time to become nondispatchable 00:00:00.000017

SVC Dump Statistics (cont)

Asid 0071:

Local storage start	10/09/2012 14:30:30.424083
Local storage end	10/09/2012 14:30:43.011936
Local storage capture time	00:00:12.587853
Tasks reset dispatchable	10/09/2012 14:30:43.011944
Tasks were nondispatchable	00:00:12.587861

Defers for frame availability	0
Pages requiring input I/O	170196
Source page copied to target	16987
Source frames re-assigned	566614

Source AUX slot IDs re-assigned 15749





Flash Express Performance Results

- All performance information was determined in a controlled environment.
- Actual results may vary.
- Performance information is provided "AS IS" and no warranties or guarantees are expressed or implied by IBM.





Flash Express Performance Test Setup

- z/OS Tests were designed to demonstrate flash performance under paging workloads that are typically encountered in a z/OS enterprise environment
 - SSD performance is not only about the number of IOPS but about steady performance over time and consistent latency
 - **Preconditioned SSDs** with random-write IO engage the device's wear leveling, error handling, and flash management algorithms
 - Comparison DASD Characteristics used current device configurations
 - DS8800 model 2107-951
 - 60 GB cache, cache hit rates of 95-100% were observed during the tests
 - DASD was not shared with any other systems and did not have any I/O traffic other than the paging traffic used for these tests
 - Configured 16 local page datasets spread across 8 LCUs



Flash Express Performance Benefits

Test Results

- FLASH paging benefits
 - Improved availability through faster paging at critical times
 - Faster workload transitions (e.g.; morning startup)
 - meaning less time to reach peak transaction rates
 - Faster SVC dumps (reduced non-dispatchable time
)
 - meaning higher availability more transactions can be run
- Pageable Large Page benefit
 - Java realizes performance benefits from use of large 1MB pageable pages
 - Large pages benefits for JIT Code Cache, 31 bit Java applications
 - No authorization needed to access fixed large pages
 - Approximately 5-8% CPU improvement from PLP





Workload Configuration Block Diagram Building block – A WAS instance accessing CICS and DB2

Each WAS instance has a WAS Control Region and 3 WAS Servant Regions.

Each WAS Control Region has a 0.5GB heap plus a JIT Code cache. Each WAS Servant Region has a 2GB heap plus a JIT Code Cache.

WAS Cor	WAS Control Region (0.5GB heap)						
WAS Servant Region (2GB heap)	WASWASServantServanRegionRegior		 WAS 7 (3 servants each with two GB heap) + 1 control region (.5 GB Heap) CICS V4.2, DB210 on a zEC12 Storage: DS8800 2107-951 with 60GB cache, very fast device 				
	CICS transactions						
	 SVC dump measurements were taken for an 18 GB dump. 						

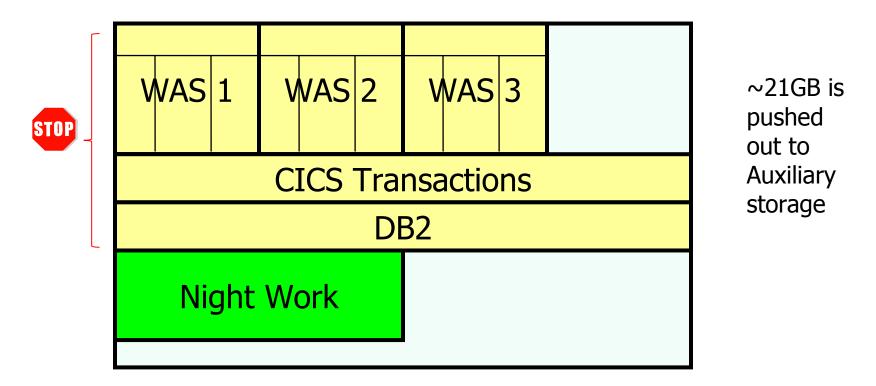


I. Morning Transition

Transition from night batch to OLTP

WAS workload to CICS and DB2 represents OLTP work which is then stopped

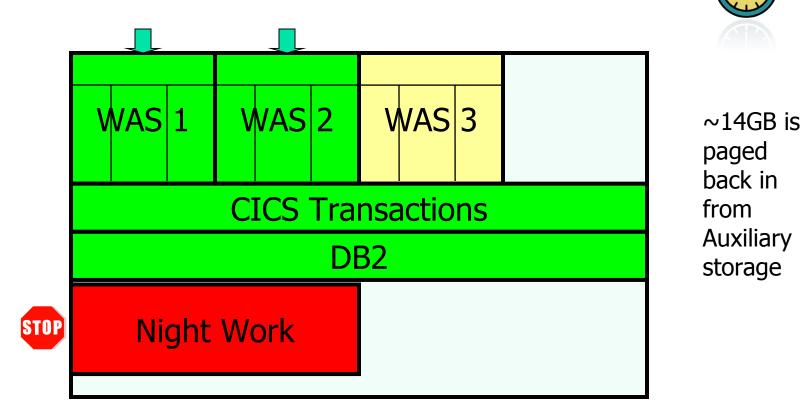
Simulated overnight work consumes real storage pushing other pages out



IBM

Morning Transition - Transition from night batch to OLTP

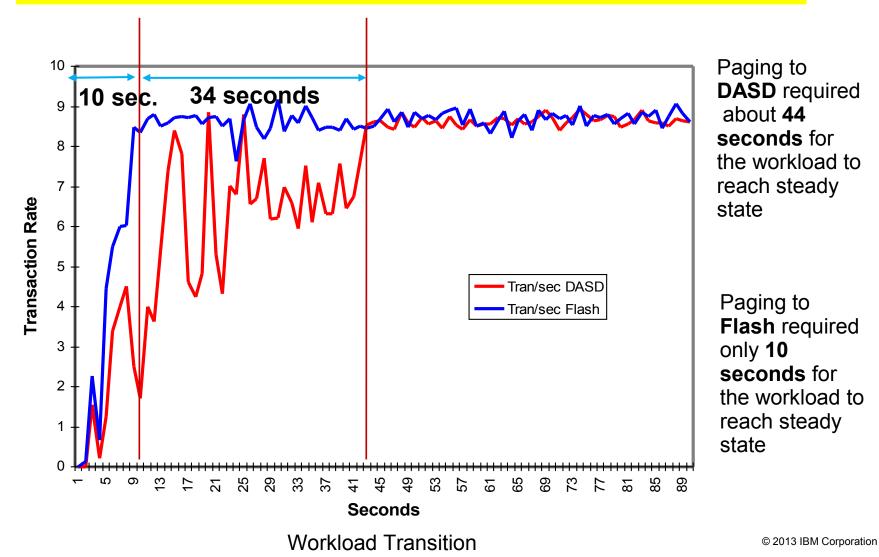
The "Night Work" is then stopped and OLTP work is started (WAS 1 and WAS 2) Measure the time needed to bring the OLTP work to full speed.





Morning Transition - Results

 During morning transition, workloads using Flash Express reached peak throughput in under 1/4th the time





Morning Transition - Results Apparent in First 45 Seconds

Transaction completion & response time	DASD	Flash	Improvement
Total Transactions within first 45 seconds	251	343	37% increase
Average response time within first 45 seconds	0.62	0.06	90% reduction

Units in seconds

Paging to Flash Express during morning transition showed up to a 10 times faster response time and up to a 37% increase in throughput within the first 45 seconds

(1) Test was for the first 45 seconds of morning transition time

Workload Transition

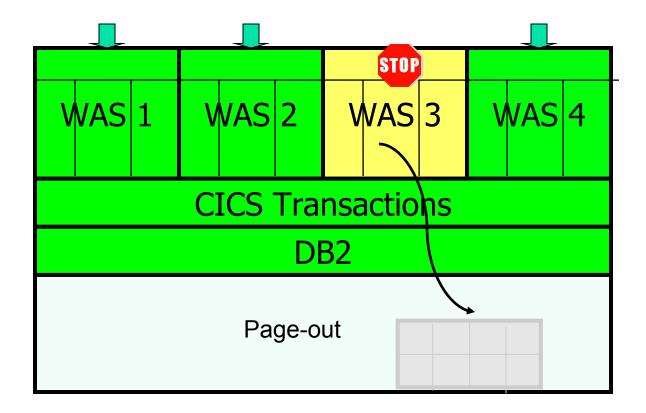


II. SVC Dump

SVC dump with pages out

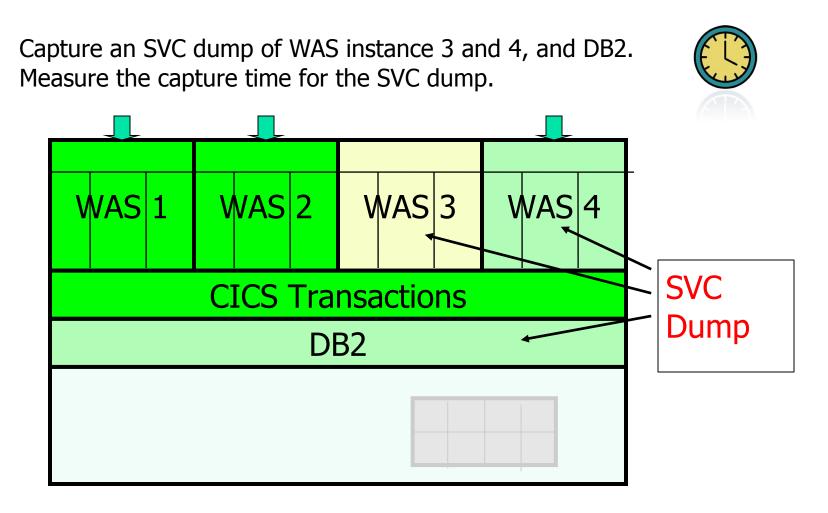
Three of four WAS instances were active.

One WAS instance was stopped and most pages were paged out.





SVC Dump- Diagnostics capture





SVC Dump - Results

*Flash Express SVC dump elapsed time was up to 25% shorter

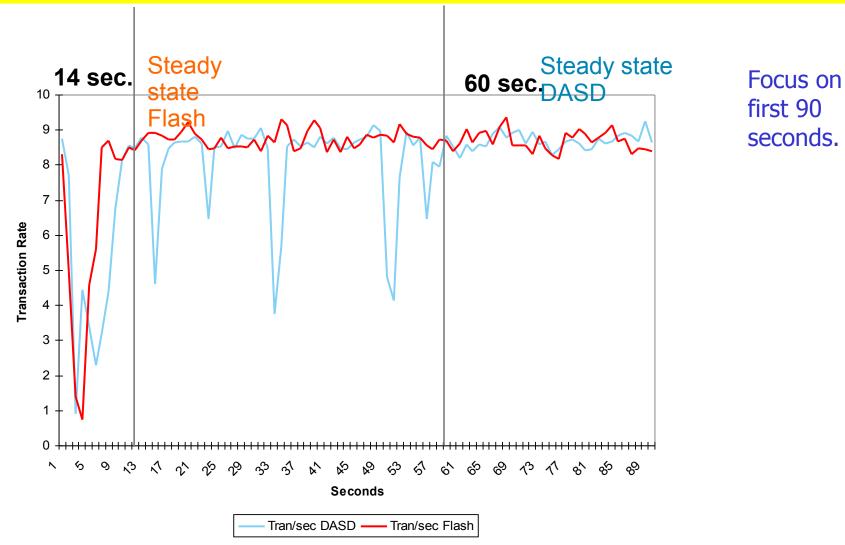
SVC Dump Metrics	DASD	Flash
SVC Dump size (in bytes):	18GB	18GB
% of pages from Aux storage:	50%	53%
DUMP Elapsed time:	189	143
Max address space non-dispatchable seconds	58.89	13.74
System non-dispatchable seconds	1.34	0.55

Let's graph these results....



SVC Dump - Results

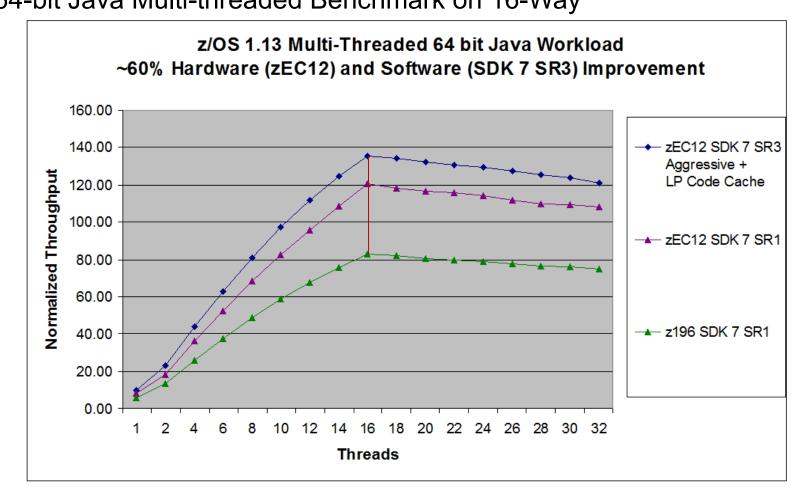
In SVC dump test, steady state performance was achieved up to 4 times faster *



* Transaction steady state was reached in 14 seconds with Flash Express, vs. 60 seconds DASD

IBM System z III. WAS and Java

z/OS Java SDK 7:16-Way Performance Shows up to 60% Improvement 64-bit Java Multi-threaded Benchmark on 16-Way



Aggregate 60% improvement from zEC12 and Java7SR3

- zEC12 offers a ~45% improvement over z196 running the Java Multi-Threaded Benchmark
- Java7SR3 offers an additional ~13% improvement (-Xaggressive + Flash Express pageable 1Meg large pages)





WAS benchmark: z/OS Performance for Pageable Large Pages

The WAS Day Trader benchmarks showed up to an 8% performance improvement using Flash Express.

Java 7 SR3	JIT	Java Heap	Multi Threaded	WAS Day Trader 2.0	
31 bit	yes	yes	4%		
64 bit	yes		1%	3%	
64 bit		yes	4%	5%	
* WAS Day Trader 64-bit Java 7 SR3 with JIT code cache & Java Heap					

DETAILS

- **64-bit Java heap (1**M fixed large pages (FLPs) or 1M Pageable (PLPs)) versus 4k pages Java heap 1M PLPs improve performance by about
 - 4% for Multi-Threaded workload
 - 5% for WAS Day Trader 2.0
- 64-bit Java 7 SR3 with JIT code cache 1M PLPs vs without Flash
 - 3 % improvement for traditional WAS Day Trader 2.0*
 - 1 % improvement for Java Multi-Threaded workload
- 31-bit Java 7 SR3 with JIT code cache and Java heap 1M PLPs vs without Flash
 - 4 % improvement for Java Multi-Threaded workload

* Note: This test used 64-bit Java 7 SR3 with JIT code cache & Java Heap leveraging Flash and pageable large pages. Also , tests used WAS Day Trader app that supports PLP; earlier version of 31-bit Java did not allocate 1M large pages



Performance Summary for Flash Express⁽¹⁾

WORKLOAD TRANSITION

During morning transition, workloads using Flash Express reached peak throughput in under 1/4th the time
 Paging to Flash Express during morning transition showed up to a 10 times faster response time and up to a 37% increase in throughput within the first 45 seconds

WAS JAVA PERFORMANCE BENCHMARKS

***** The WAS Day Trader benchmarks showed up to an 8% performance improvement using Flash Express.(2)

* This test used 64-bit Java 7 SR3 with JIT code cache & Java Heap leveraging Flash and pageable large pages.

Improved Availability During Diagnostics

- In SVC dumps, availability was up to 4 times higher for workloads and up to twice as high for systems*
- In SVC dump tests, steady state performance was achieved up to 4 times faster *
- Flash Express SVC dump elapsed time was up to 25% shorter

* Transaction steady state was reached in 14 seconds with Flash Express, vs. 60 seconds DASD.

DB2

◆ Up to 28% improvement in DB2™ throughput due to faster CPU and leveraging Flash Express with Pageable Large Pages (PLP)*

• Workloads leveraging Flash Express with PLP can see up to a 8%** price performance improvement over the z196.

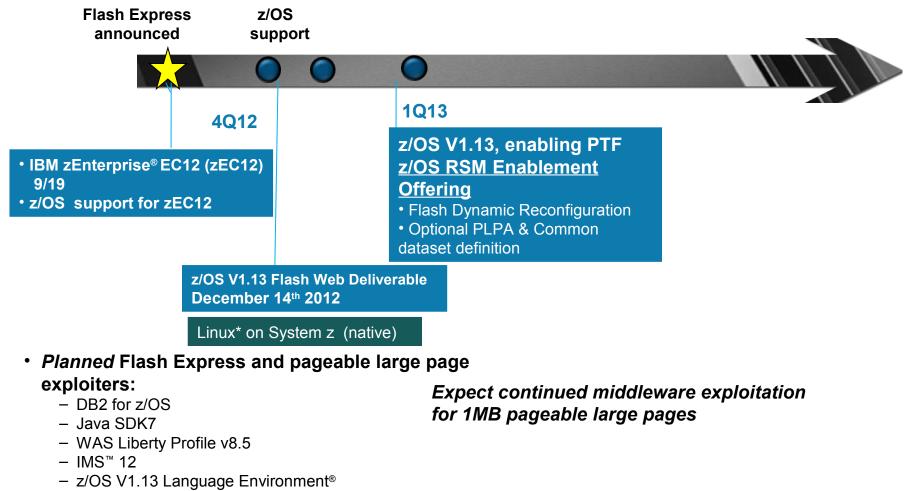
* PLP for DB2 helps DB2 to achieve "additional" up to 3% additional performance on top of zEC12 CPU expected throughput improvements of 25%.

- ** based on average 5% discount for zEC12 workloads under the AWLC pricing plus up to 3% more performance per MSU with Flash Express.
 - (1) All tests are comparing the use of Flash Express as compared to using DASD (DS8800)
 - (2) System non dispatchability and address space non dispatchability time were dramatically reduced enabling work to be processed that would otherwise have been stopped

z/OS Flash Roadmap



Flash Express Exploitation Flash support in z/OS sets the stage for further use



- Other (CICS®)

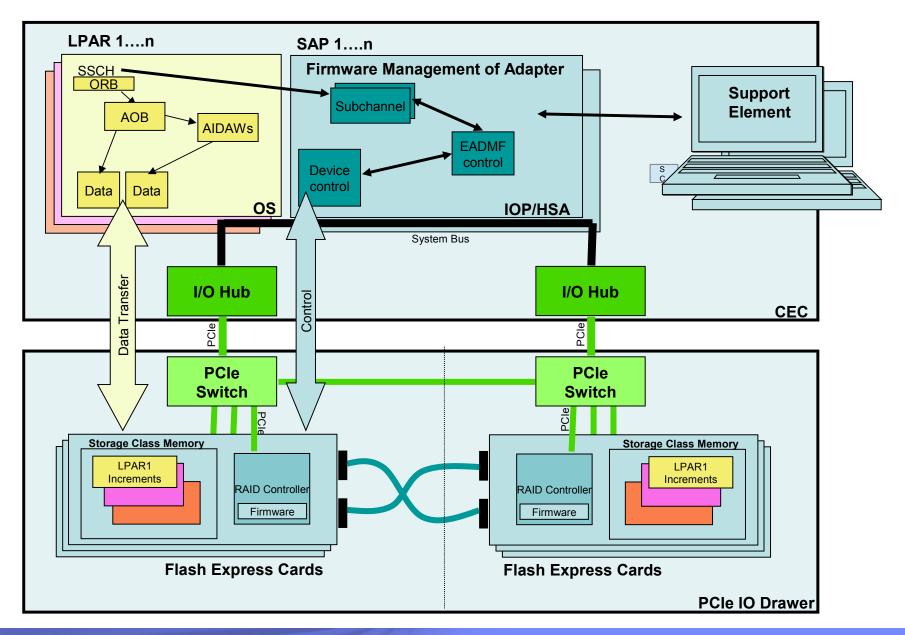
47

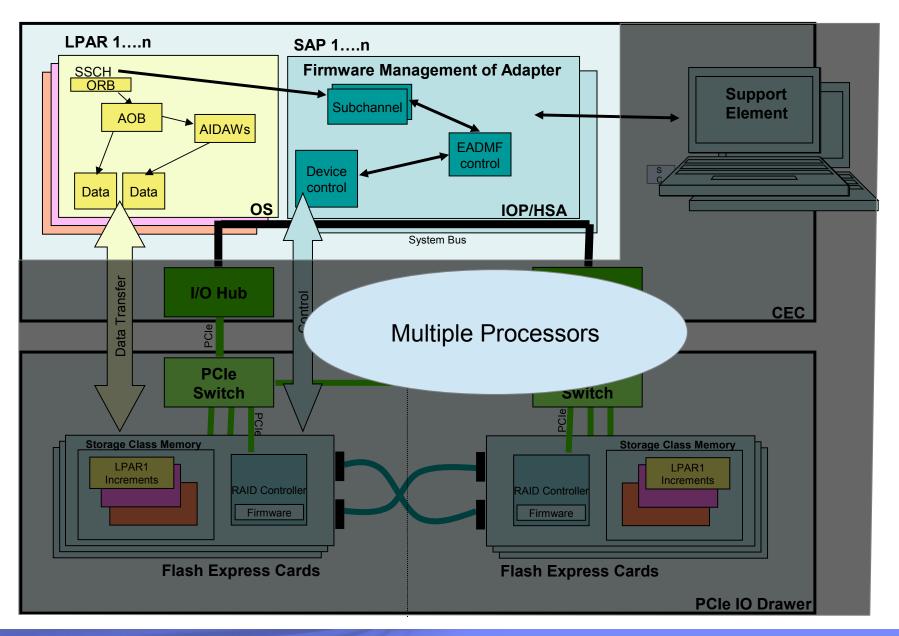


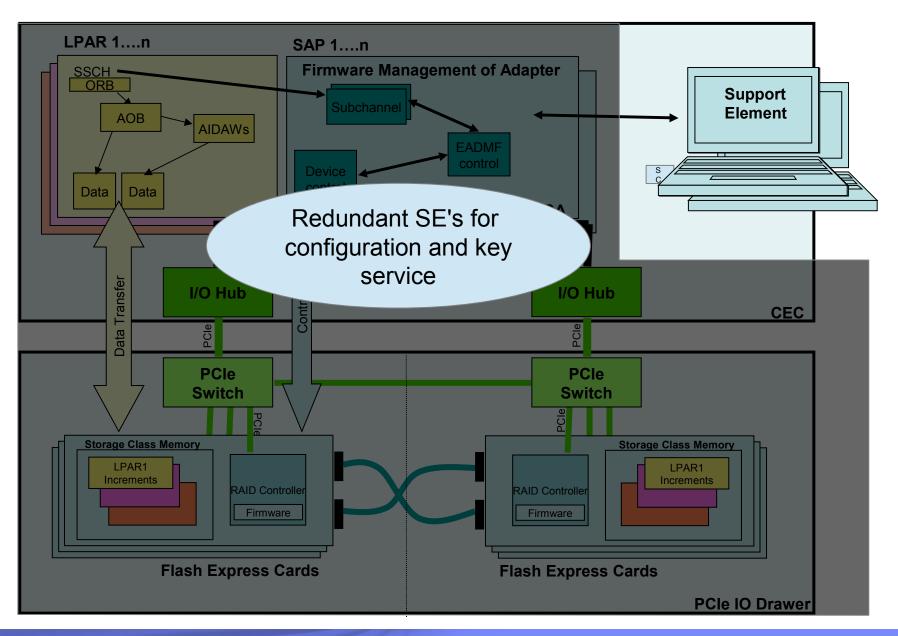
Flash Express Implementation

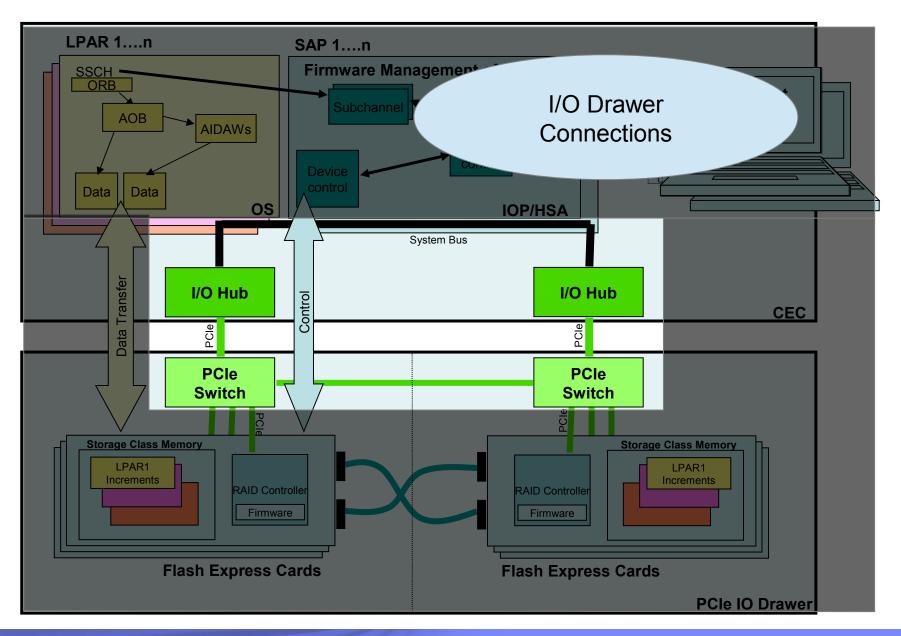
- ✓ System Overview
- ✓ Redundant Physical Structures
- ✓ Data Protection Mechanisms
- ✓ Data and Key Encryption
- ✓ Non-Disruptive Service Techniques

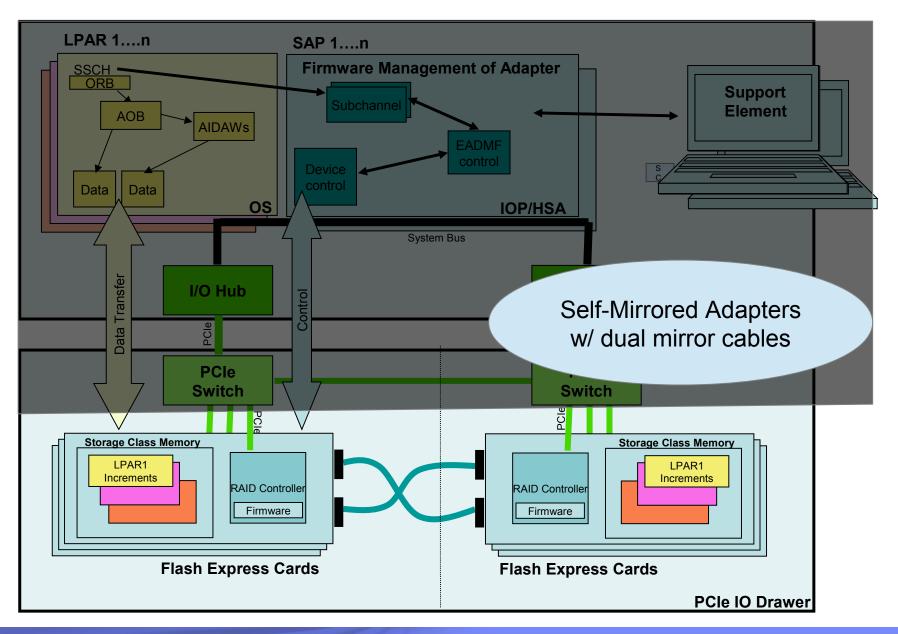
FLASH Express System Overview



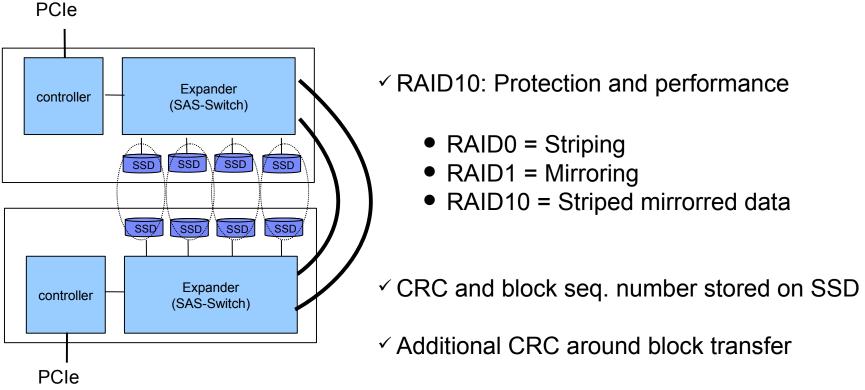








Data Protection Mechanisms



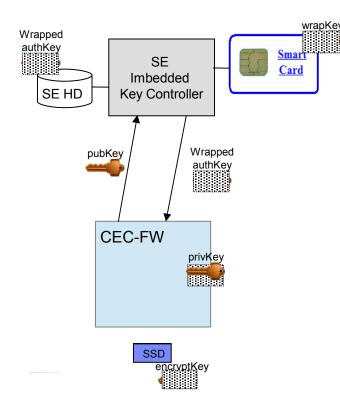
Comm links (SAS, PCIe) provide embedded protection & recovery

CEC-based hardware address protection on communication from adapter

✓ ECC on internal system memory

Data and Key Encryption

- ✓ On SSD, data is protected with inline encryption (hidden encryptKey)
- ✓ Access to SSD is via authentication key (authKey) served from SE

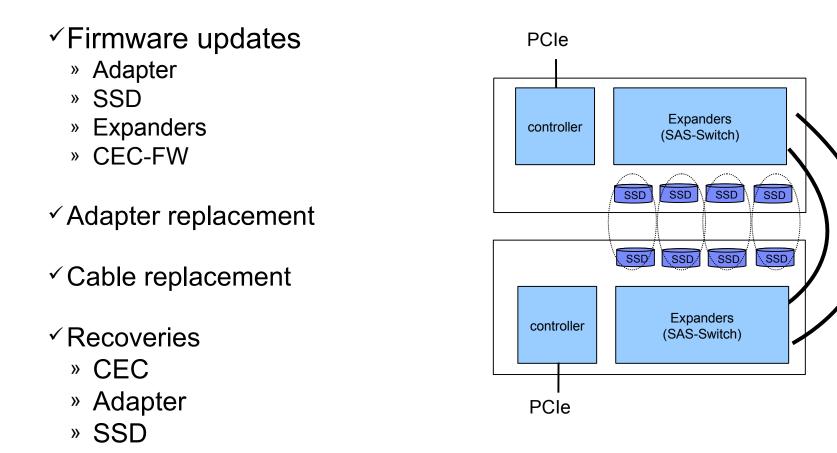


During Flash install, in smart card on SE:

- Create authKey (aka PIN)
- Wrap authKey in an encrypted file
- wrapKey stored in smart card
- Wrapped key file stored on SE
- $SE \rightarrow CEC$ -FW authkey service:
 - asymetric protocol pub/private
 - IOP sends public key to SE
 - In smart card, Key file unwrapped then encrypted with CEC pubKey
 - Encrypted authKey sent to CEC
 - CEC 'unwraps' authKey using its privKey

✓ AuthKey used during SSD format and subsequent power cycles

Non-disruptive Service Strategy



THANK YOU



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