



Hitachi Dynamic Tiering

Storage Tiering for the Mainframe

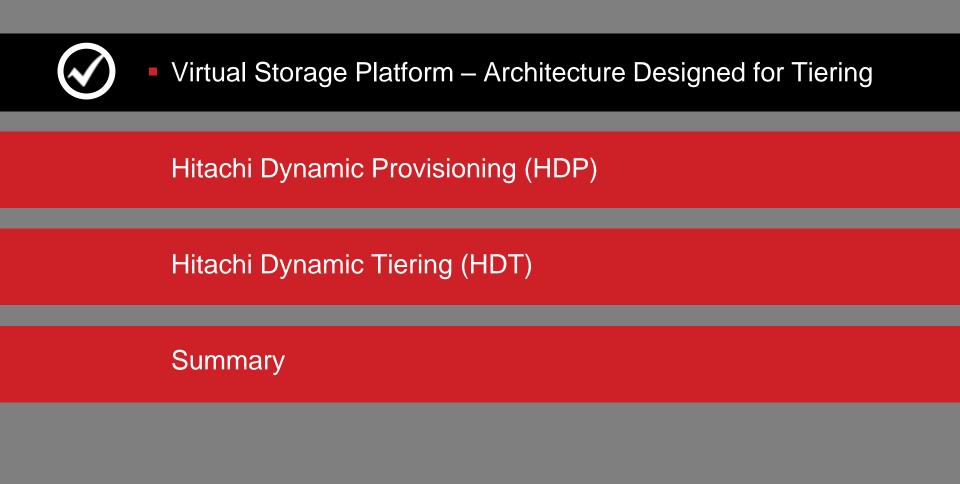
William Smith Hitachi Data Systems

February 5, 2013 Session Number 12680





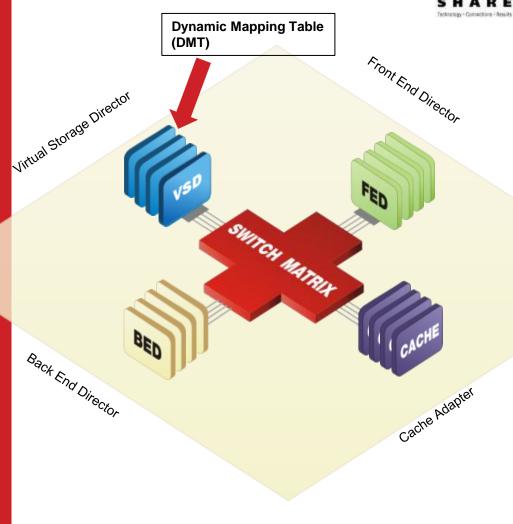




Virtual Storage Platform - 3D Scaling Architecture



- 5th generation multi-dimentional internal crossbar switch
- Native 6G SAS backend
- 192 x FC, 96 x FICON / FCoE
- Independent Scale up, scale out, scale deep growth path
- ITB Global Cache
- 3rd generation imbedded virtualization (Open Systems and Mainframe)
- Granular non-disruptive upgrade paths
- 100% non-disruptive maintenance (7 x 24 x forever)





Drive Types and Capacities

- Current Drive Technology:
 - 2.5" 10K 300GB SAS
 - 2.5" 10K 600GB SAS
 - 2.5" 10K 900GB SAS
 - 2.5" 15K 146GB SAS
 - 2.5" Flash SLC 200GB 3G and 6G SAS
 - 3.5" Flash SLC 400GB 3G and 6G SAS
 - 3.5" 2TB Nearline SAS / 7200 RPM
- New Drive Technology (last quarter):
 - 2.5" 15K 300GB SAS
 - 2.5" Flash MLC 200GB 6G SAS
 - 2.5" Flash MLC 400GB 6G SAS
 - 3.5" Flash MLC 400GB 6G SAS
 - 3.5" 3TB Nearline SAS / 7200 RPM (later this month)
 - New High Density Flash Drives





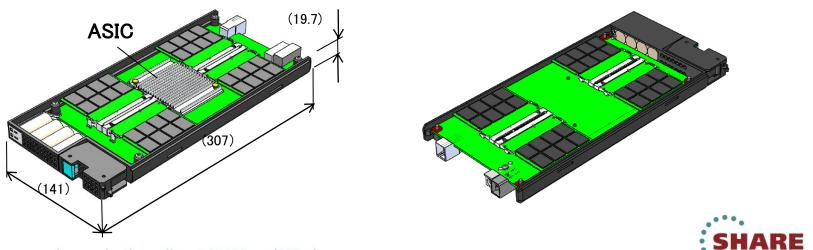
New Hitachi Flash Technology



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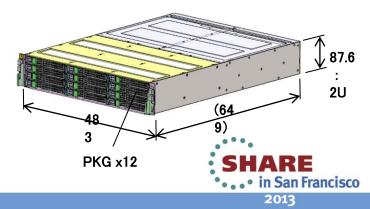
- The Flash Memory Drive (FMD) is Hitachi built specifically for Hitachi arrays.
- Each FMD is as follows:
 - Raw capacity option of 1.6TB
 - eMLC based
 - Each FMD uses an Hitachi designed and built System-On-Chip (SOC) ASIC
 - Multi-Core ASIC performs flash memory management, data transfer, RAS etc.
 - High bandwidth PCIe Gen2 data paths to eMLC memory
 - Much improved performance over industry SSD



New Hitachi Flash Technology



- VSP will offer granularity of 4 Flash Memory Drives (FMD) minimum, in 4 separate FMU (Flash Memory Units)
- VSP can support a total of 192 FMD's
 - 12 Flash FMD's per FMU
 - 4 FMU per Flash Box (FBX similar to current Disk Chassis)
 - 4 FBX per VSP (2 per Control Chassis)
 - Total raw capacity: 307TB (1.6TB FM)
- Target performance is >1 Million Random Read IOPS per FMU
- Available for all VSP's and non-disruptive install on live VSP's
- Game changer in flash technology
 - Cost effectiveness per TB improved considerably
 - Expect market adoption to be very high







Virtual Storage Platform – Architecture Designed for Tiering



Hitachi Dynamic Provisioning (HDP)

Hitachi Dynamic Tiering (HDT)

Summary

Dynamic Provisioning for Mainframe



- Dynamic Provisioning
 - OPEN Systems 2nd Generation
 - Mainframe August 2011
 - Internal and external (virtualized) storage
 - Managed in pages (42MB for OPEN, 38MB for Mainframe)
- Cost savings through:
 - Improved storage utilization
 - Improved performance through ultra-wide striping
 - Non-disruptive addition of storage into the pool
 - Integration with SMS



Hitachi Dynamic Provisioning (HDP) for the Mainframe



- So, just what is a page?
 - For the VSP, there are two page sizes
 - OPEN 42MB
 - Mainframe 38MB
 - They are different to fit the best use of storage space for the underlying OS structures they support
 - For Mainframe, 38MB fits the 3390 architecture with minimal wasted space
- Each storage supplier has a different implementation
 - Some are in MB
 - Some are sub-LUN
 - Some are partial volumes
 - Some are Chunks / Chunklets
- Each is designed to best fit their architecture



Complete your sessions evaluation online at SHARE.org/SFEval

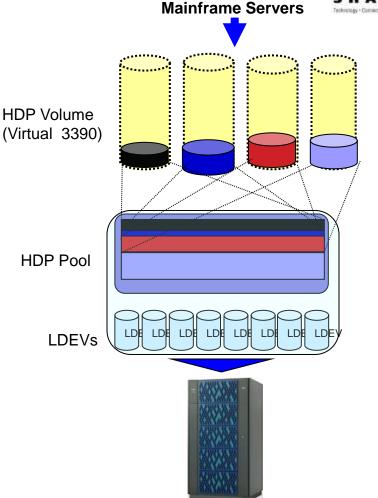
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Hitachi Dynamic Provisioning for Mainframe

Simplifies Capacity Expansion; no IOGEN, no documentation updates,

no changes to GDPS Parms, XRC Sessions, Flashcopy configs etc.

- Optimize Storage performance by spreading the I/O across all available Disk Drives
- Optimize Storage Capacity by only allocating capacity which is actually used
- Reduce OPEX by automatic thin moves and copies of fat volumes
- Achieve Flash Copy Space Efficiency for Target Volumes
- Leverage Dynamic Volume Extension (DVE) to seamlessly grow physical and logical capacity
- Leverage Extended Address Volumes (EAV)





HDP Mainframe — Basic Concepts



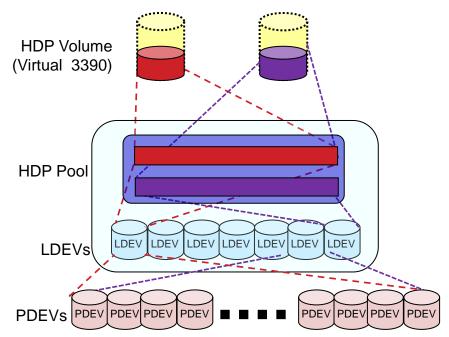
Key Functionalities

- Thin Provisioning
 - Pool Rebalance
 - Pool Volume Reduction
 - Space Reclaim
- Support Mainframe Features
 - PAV, HyperPAV
 - z/HPF
- Support Replication Products
 - TrueCopy and HUR
 - ShadowImage and FlashCopy V2
- Support New Features
 - EAV and DVE
 - FlashCopy SE

 Hitachi Dynamic Tiering Complete your sessions evaluation online at SHARE.org/SFEval

With HDP for Mainframe

 Optimize storage performance by spreading the I/O across all available physical disk drives



Spreading data across more physical volumes reduces the I/Os per volume and improves performance







Virtual Storage Platform – Architecture Designed for Tiering

Hitachi Dynamic Provisioning (HDP)



Hitachi Dynamic Tiering (HDT)

Summary

Hitachi Dynamic Tiering for Mainframe



- Dynamic Tiering
 - Open Systems available since September 2010
 - Mainframe Available since January 2012
 - Automated performance management
 - Cost effective use of storage tiers
 - Internal and external (virtualized) storage
 - Managed in pages (42MB for OPEN, 38MB for Mainframe)
- Cost savings through:
 - Improved storage utilization
 - Reduced storage requirements for Tier 1
 - Less time spent on performance management
 - Re-deploy personnel
 - Integration with SMS



Virtual Storage Platform: Page Level Tiering



Different tiers of storage are in one pool of POOL A pages – up to 128 pools TIER 1 Automatically detects and assigns Tiers based on media type (includes the new Least Referen HAFs SSD). TIER 2 Dynamically: add or remove Tiers Least Referen TIER 3 expand or shrink Tiers expand 3390A volumes Dynamic Tiering Policies at volume level: Default - All Level 1 - Tier 1 only Level 2 - Tier 1 and Tier 2 only Level 3 - Tier 2 only Level 4 - Tier 2 and Tier 3 only Level 5 - Tier 3 only Tier 3 External can have 3 tiers 26 new policies added (% tier)

Virtual Storage Platform: Page Level Tiering



Cycles	
Customer defines strategy, it is executed automatically	
24 hour	
 defined part of 24 hour 	
¹ / ₂ , 1, 2, 4 or 8 hourly	
 All aligned to midnight 	
Weekly or longer	
Jse Hitachi Tier Storage Manager (HTSM)	

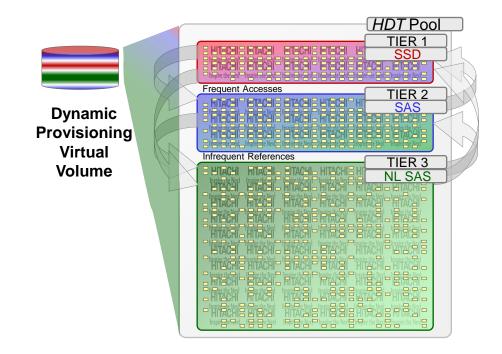
- Set tiering policy schedules to match business cycles and requirements
- Report on data mobility (promotions and demotions)
- Dynamically change or adjust cycle, policies, tiers etc...

Media tier groupings supported by Hitachi VSP	Order of tiers
2.5" SSD (200GB) SLC or eMLC 2.5" SSD (400GB) SLC or eMLC	1
2.5" SAS15Krpm (146GB) 2.5" SAS15Krpm (300GB)	2
2.5" SAS10Krpm (300GB) 2.5" SAS10Krpm (600GB) 2.5" SAS10Krpm (900GB)	3
2.5" SAS 7.2Krpm (2TB, 3TB)	4
3.5" SATA (2TB)	5
External	Lowest
External (June '12)	Any Tier



Hitachi Dynamic Tiering





- What determines if a page moves up or down?
- When does the relocation happen?



Page Relocation



- At the end of a monitor cycle the counters are recalculated
 - Either IOPH (Period) or weighted average (Continuous)
- Page counters with similar IOPH values are grouped together
- IOPH groupings are ordered from highest to lowest
- Tier capacity is overlaid upon the IOPH groupings to decide on values for Tier Ranges
 - Tier Range is the 'break point' in IOPH between tiers
- Relocation processes page by page looking for pages on the 'wrong' side of a Tier Range value
 - i.e. high IOPH in a lower tier
- You can see the IOPH groupings and Tier Range values in HDT reporting "Pool Tier Properties"



Pool Tier Properties

☆ View Tier Properties

View Tier Properties

Tier1	Drive Type/RPM	SSD/-
	Capacity(Used/Total)	2.45 TB / 3.13 TB [78 %]
	Performance Utilization	5%
Tier2	Drive Type/RPM	SAS/15k
	Capacity(Used/Total)	34.96 TB / 43.92 TB [79 %]
	Performance Utilization	62 %
Tier3	Drive Type/RPM	SATA/7.2k
	Capacity(Used/Total)	10.42 TB / 42.99 TB [24 %]
	Performance Utilization	100 %

What is being used now in the pool in terms of capacity and performance

HITACHI

☆ Performance Graph (SPC-1(0)) — 🔿 Tier1 Range 100000 — \land Tier2 Range 10000 Average 1000 pool. b-o------Number 100 of I/O(per Hour) 10 1 to" 0 0 10 20 30 40 Capacity[TB] 2010/09/17 13:39 - 2010/09/17 15:49 Close Help

The I/O distribution across all pages in the pool.

Combined with tier range this is what HDT is using to decide "where the pages should go to"



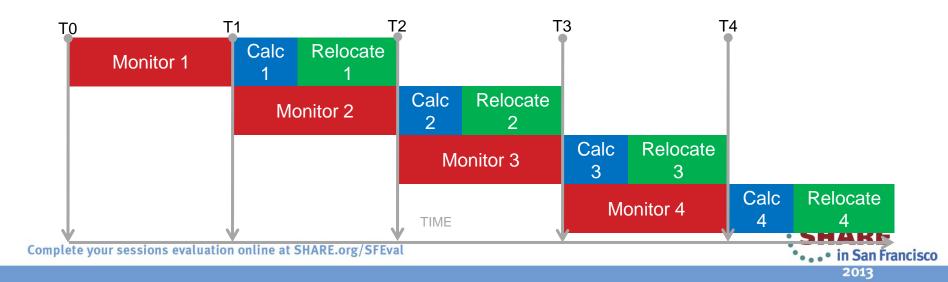


LET TIERING LEARN YOUR WORK LOAD



MF-HDT Cycle

- Cycle Time set at the MF-HDT Pool Level
- Manual Mode
 - User can start and stop performance monitoring using any interval up to 7 days
- Automatic Mode
 - Continuous Monitoring followed by Relocation cycles
 - Monitor Interval from 30min, 1h, 2h, 4h, 8h or 24h(default)

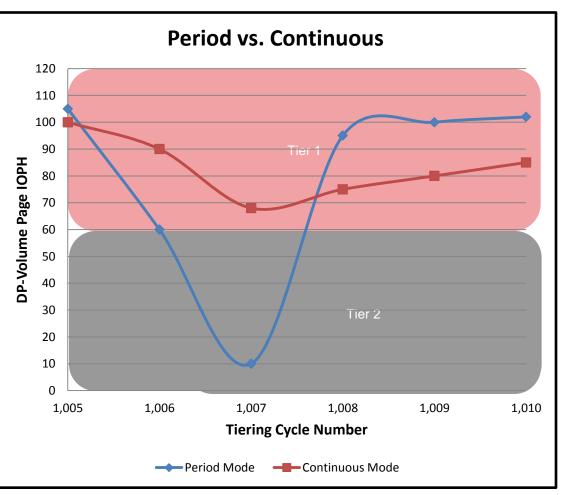


LET TIERING LEARN YOUR WORK LOAD



MF-HDT Monitoring Modes

- Monitoring Modes set at the MF-HDP Pool Level
- Period
 - The value used in the Calculation cycle is the Actual IO Load on DP-Volume Page from Previous Monitoring Cycle
- Continuous
 - The value used in the Calculation cycle is the Weighted Average of Multiple Previous Monitoring Cycles for DP-Volume Page
 - Reduces Page Thrashing
 - May Slow Migration to upper Tiers

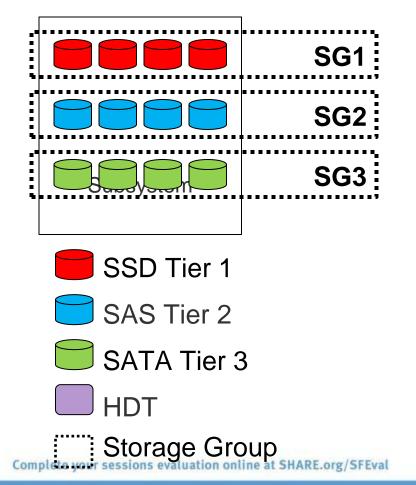




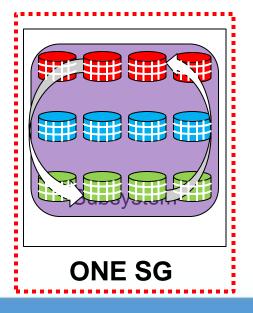
DFSMS Storage Groups and MF HDP and HDT



- Tiered Storage can be implemented with DFSMS Storage Groups
- DFSMS Attributes such as Data Class and Storage Class are used to direct allocation to the right Tier



 With MF HDT may combine all SG's into one and use HDT to move data to the right Tier automatically based on Workload performance





HDP Volumes are managed at page-level

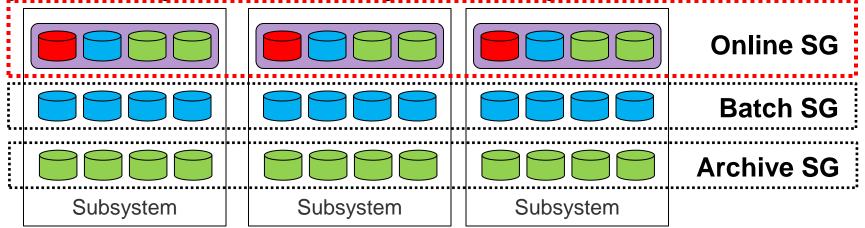
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DFSMS Storage Groups and MF HDT



- HDT can be applied to selected Storage Groups only
- DFSMS Horizontal Storage Groups and HDT Pools with Dynamic Page based Tiering can be aligned



SSD Tier 1







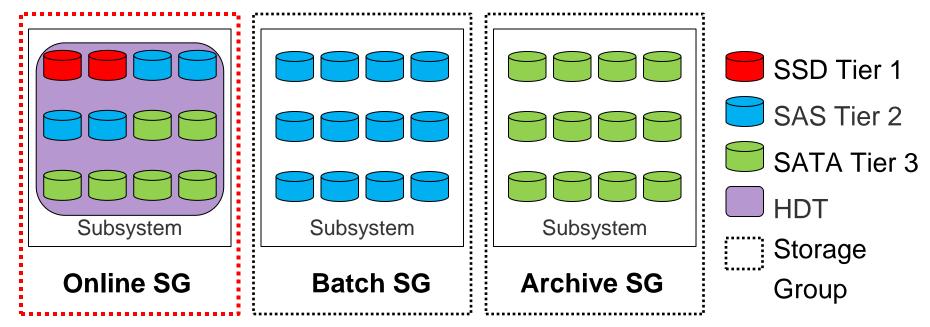
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DFSMS Storage Groups and MF HDT



- HDT can be applied to selected Storage Subsystems only
- DFSMS Vertical Storage Groups and HDP Pools with Dynamic Page based Tiering can be aligned





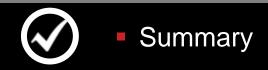
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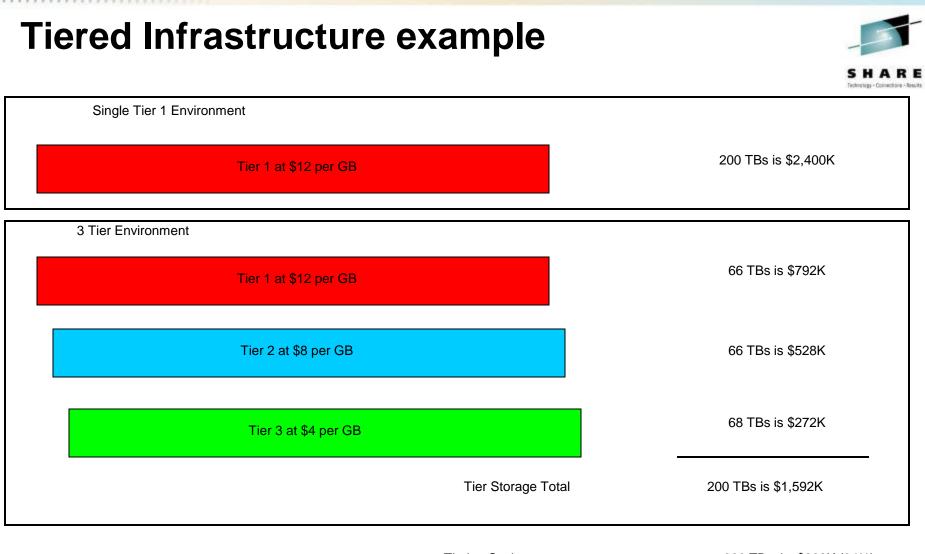


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Tiering Savings

200 TBs is \$808K (34%)

Tiering allows "Buy the same capacity of storage for less"

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Summary

Let Tiering Learn your work Load

- Two different Page Sizes
 - 38MB Page Size for Mainframe
 - 42MB Page Size for OPEN
- Up to 128 unique pools
- All features and functions the same for OPEN and Mainframe
- Tiering Policies to better control data placement and movement
 - Default All
 - Level 1 Tier 1 only
 - Level 2 Tier 1 and Tier 2 only
 - Level 3 Tier 2 only
 - Level 4 Tier 2 and Tier 3 only
 - Level 5 Tier 3 only
 - 26 new policies recently added.
 - Full support of Mainframe standard features including DFSMS



Francisco



Thank You.

