



Fit For Purpose Platform Architecture Selection and Design

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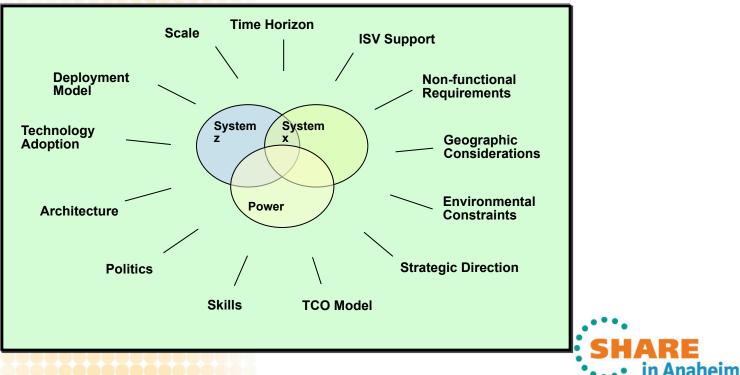
What is Fit for Purpose (F4P)?



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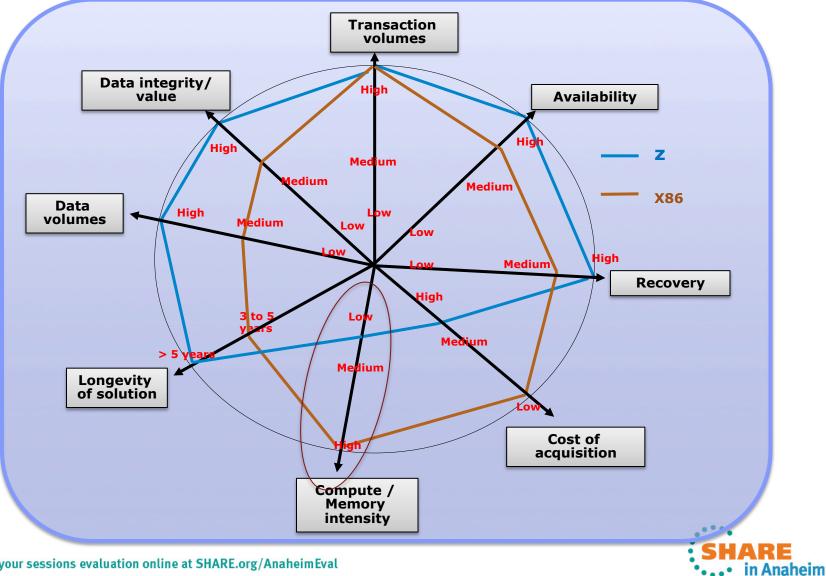
Fit for Purpose is a client centric thought process that when applied, yields infrastructure architecture decisions which are in line with the client's requirements and local conditions.

It is based on the fundamental principles that "one size does not fit all" and that "local factors matter."



A Client's Decision Matrix





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Know the legacy, workload, and costs



Know the current IT Environment



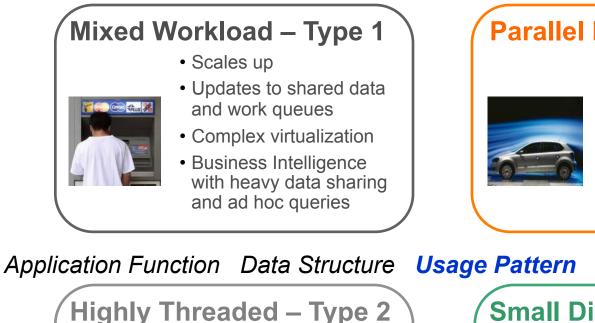
Examine costs

Understand the workload



Fit for Purpose Categorized Workloads





Scales well on large SMP



- Web application servers
- Single instance of an **ERP** system
- Some partitioned databases

Black are design factors

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Parallel Data Structures – Type 3



- Scales well on clusters.
- XML parsing
- Buisness intelligence with Structured Queries
- HPC applications

Usage Pattern SLA Integration Scale

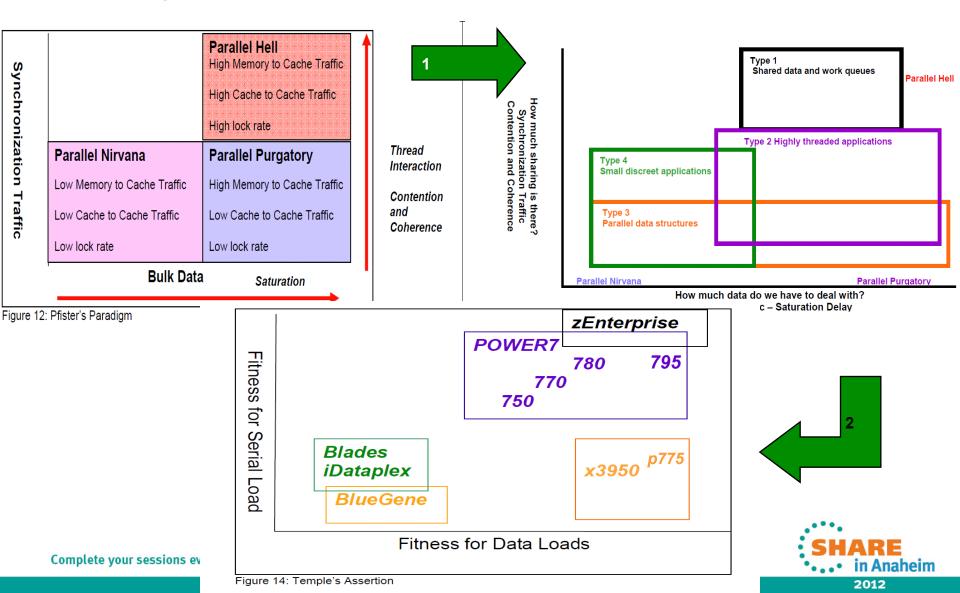


Blue are local factors

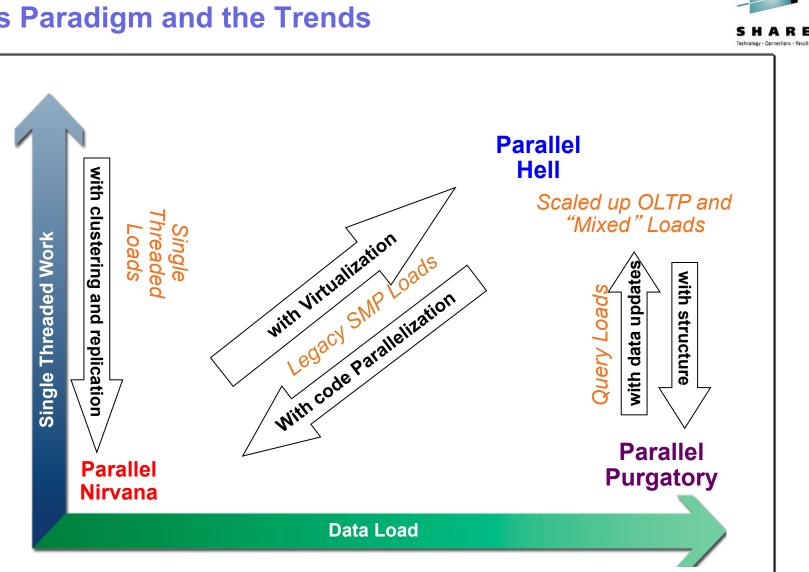


Using Pfister's paradigm we can map the workload types to our existing platforms and new platforms we build as a result of a WOS study









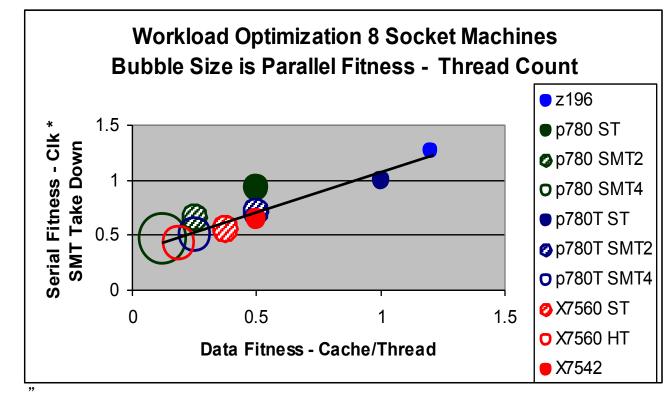


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Workload Optimization

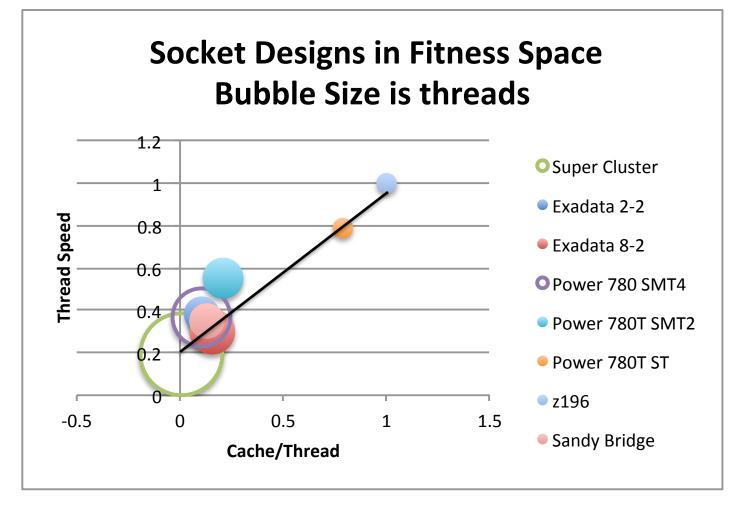




There is a clear trade off at work here:

- To have more threads you must give up thread speed and cache/thread
- Machine capacity metrics govern how that tradeoff is made. In turn the metrics are designed for the "style" of computing used by each machine's base market.







Throughput and Capacity *Relating Fitness to workloads*



• We observe:

Throughput ~ Thread Count x Thread Speed

Also:

Thread Capacity ~ Cache / Thread x Thread Speed

• We Assert:

Performance ~ Thread Capacity and Throughput

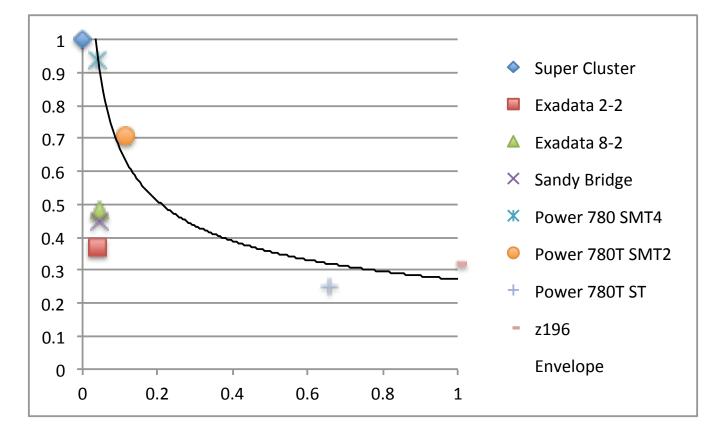
• We Define:

Performance = w(Thread Capacity) +(1-w)(Throughput)



Machines have different Throughput and Thread Capacity



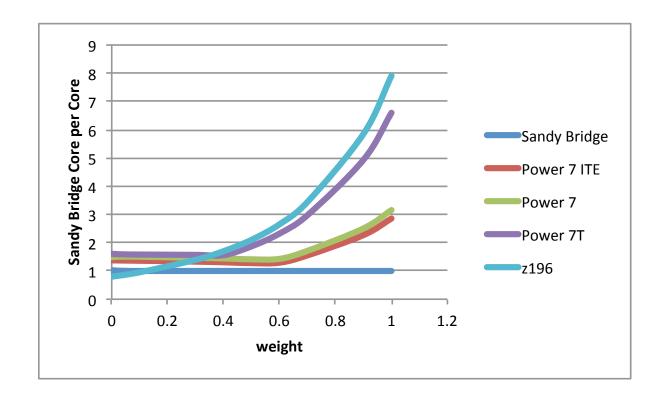


Note that Very High Throughput \rightarrow Very Low Thread Capacity (and Vice Versa)

Therefore to achieve <u>Very High Throughput</u> workload must have <u>Very Low Weight</u> Complete your sessions evaluation online at SHARE.org/AnaheimEval



Single Core Relative Capacity



Note that relative capacity *is not linear with weight.*



Local factors take the form of Operational Trade offs

- Operations governed by "Normalized Headroom"
 - HR = $(1-u)/u = c^2 N t_0 / t_{wait}$
 - HR(avg) = kcN² => (SLA)(Variability)(Scaling)
 - U = 1/(1+HR)
 - $t = t_0 + t_{wait}$
 - t_{wait} = (t₀)(c²N)(u/(1-u)) = (t₀)(c²N)/HR = (1/weighted capacity)(variability)(scaling)/HR M/G/1 system



What are you Optimizing?



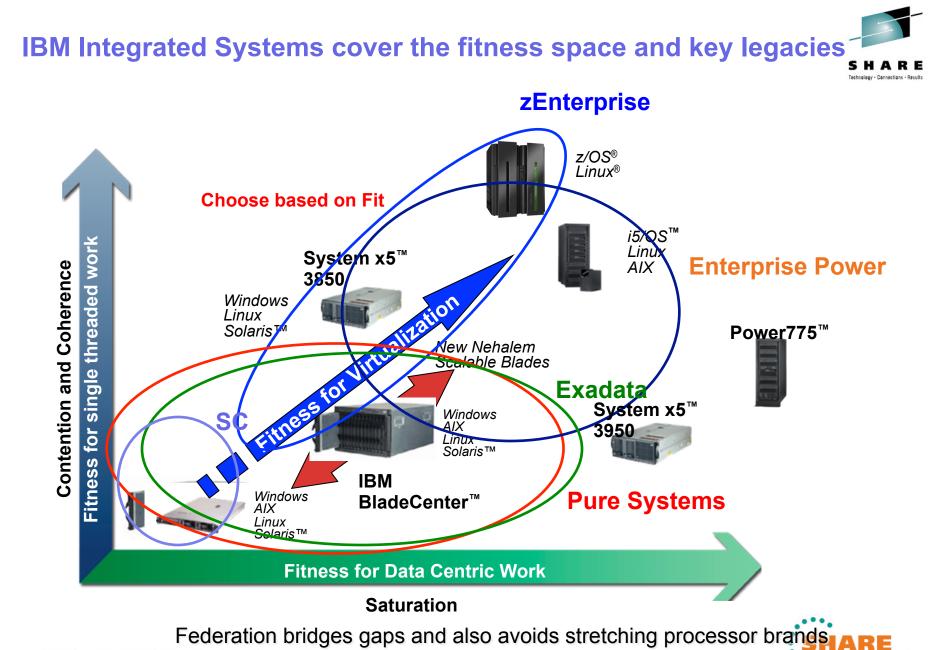
Is t measured at completion of each user thread?Look at Thread Capacity v Throughput

Want efficiency with good enough response time? Look for Thread Capacity

Or is t measured at completion of many threads? Look for Throughput

Can "wait time" tolerance be bought by reduction of network latency? Trade off wait time for efficiency; look for Capacity



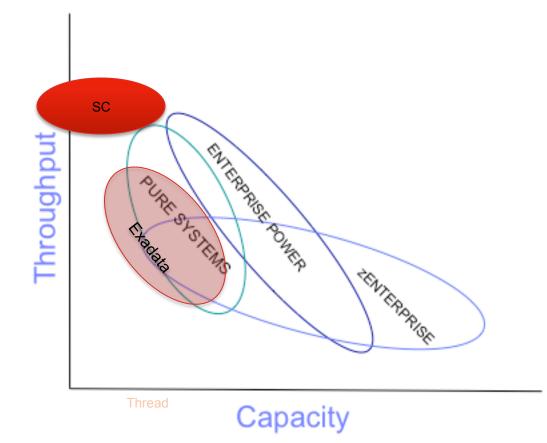


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Positioning with Throughput and Thread Capacity



Today Pure Systems is a match for workloads with lower weight.

