



DFSMS Basics: VSAM VSAM RLS Best Practices

How to leverage VSAM RLS best practices and performance in your environment David LeGendre, dlegendr<u>@us.ibm.com</u> Session : 10967



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Introduction

There are a number of settings, parameters and sizing considerations to be made when configuring VSAM RLS environments. Additionally, there are many performance indicators and that can help to tune those options to a specific installation. The goal of this presentation is to help provide some insights into those topics as well as to supply guidance when problems are encountered.





Agenda

- VSAM RLS I/O Path and Buffer Management Overview
 - RLS sub components
 - GET path example
 - LRU algorithm
- Performance and Tuning Parameters
 - Buffer Pools
 - PARMLIB
 - Cache Structures
 - Lock Structures
 - Data Class
 - Request Level Parameters
 - Data Set Considerations





Agenda

- Performance Measurements
 - SMF 64 records
 - SMF 42 subtypes 15-19
 - RMF Monitor III
 - RMF RLSSC & RLSDS
 - RMF RLSLRU
- SMSVSAM Diagnostics
 - Display commands
 - Hang / wait condition & SMSVSAM dump collection
 - Internal logic errors & required documentation







VSAM RLS I/O Path and Buffer Management Overview





VSAM RLS I/O Path – Sub Components



- VSAM Record Management (VRM)
 - Provides the VSAM interfaces: GET, PUT, POINT, ERASE, etc
 - Parameters passed to VRM are through the RPL control block
- Storage Management Locking Services
 (SMLS)
 - Interfaces with VRM and XCF locking services
 - Obtains, releases, and alters locks in the coupling facility lock structure(s)





VSAM RLS I/O Path – Sub Components

- Sysplex Cache Manager (SCM)
 - Interfaces with Buffer Management Facility and XCF caching services
 - Obtains directory elements and read/write data elements in the coupling facility cache structures
- Buffer Management Facility (BMF)
 - Interfaces with VRM and SCM to locate/add buffers to the local pool
 - Supports buffering past close
 - Data sets reopened for RLS within 10 minutes can reuse valid buffers remaining in the pool
 - Manages the size of the local buffer pool via a Least Recently Used (LRU) manager





VSAM RLS I/O Path – Sub Components

• Performance Goal: Spend the least amount of time in the I/O path!

GET/PUT ↔ VRM ↔ SMLS ↔ XCF Locking services

↔ BMF ↔ SCM ↔ XCF Caching Services

↔ Media Mgr Services (to DASD)





VSAM RLS I/O Path – Example of a GET



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SMSVSAM Address Space



Buffer Management Overview

- 31 Bit Buffer Pool
 - Size: 10M to 1728M
 - Defined using the RLS_MAX_POOL_SIZE(nnnn) parameter
 - Same on all systems
- 64 Bit Buffer Pool
 - Size: 0 or 500M > size > 2,000,000M (2TB)
 - Defined using the RLSAboveTheBarMaxPoolSize(sysname1,nnnn) parameter
 - Can be specified per system





🗉 in Atlanta

Buffer Management Overview – LRU(31 bit)

- 31 bit LRU Comprised of 4 different levels
 - Normal Mode (0 80%)
 - Max_Cycles_Unused = 240
 - Cycles_Unused = Cycles_Unused + 1
 - Cycle iterates every 15 seconds
 - (15 * 240) / 60 = 60 minutes
 - Maintenance Mode (80 120%)
 - Max_Cycles_Unused = Max_Cycles_Unused 1
 - Cycles_Unused = Cycles_Unused + 1
 - Cycle iteration stays at 15 seconds
 - Accelerated Mode (120 200%)
 - Max_Cycles_Unused = Max_Cycles_Unused 4
 - Cycles_Unused = Cycles_Unused + 1
 - Cycle iteration stays at 15 seconds
 - Panic Mode (> 200%)
 - Max_Cycles_Unused = Max_Cycles_Unused 8
 - Cycles_Unused = Cycles_Unused + 1
 - Cycle iteration runs every 5 seconds



Percentage of RLS_MAX_POOL_SIZE in use

* 200% of RLS_MAX_POOL_SIZE OR > 1728MB Limit



Buffer Management Overview – LRU(64 bit)

- 64 bit LRU Comprised of 4 different modes
 - Normal Mode (0 80%)
 - Buffers older than 60 minutes are recycled
 - Maintenance Mode (80 90%)
 - Buffers older than 60 minutes are recycled
 - Accelerated Mode (90 100%)
 - Buffers older than 30 minutes are recycled
 - Buffer stealing will be attempted
 - If no available buffers a new one will be generated
 - Panic Mode (> 100%)
 - Buffers older than 5 minutes are recycled
 - Buffer stealing will be attempted
 - If no available buffers a new one will be generated



Percentage of RLSAboveTheBarMaxPoolSize in use







Performance and Tuning Parameters





Performance and Tuning – Buffer Pools

SYS1.PARMLIB(IGDSMSxx)

- RLS_MAX_POOL_SIZE(nnnn MB)
 - Max size of SMSVSAM local buffer pool
 - nnnn = 10 to 9999, anything over 1500 is treated as the max of 1728
 - Recommendation: RLS_MAX_POOL_SIZE(<= 850)</p>
 - Must have adequate cache structure sizes
 - 680M (80%) of buffers to reside in the 31 bit pool for one hour via LRU normal mode
 - Allows for a doubling of the pool to 1700M before entering panic mode (after exceeding the 1728M limit)





Performance and Tuning – Buffer Pools

SYS1.PARMLIB(IGDSMSxx)

- RLSAboveTheBarMaxPoolSize([sys1,val1;sys2,val2...] | [ALL,value])
 - Amount of virtual storage allocated above 2G bar for RLS buffering
 - nnnn = 0, or 500M > size > 2,000,000M (2 TB)
 - Recommendation: RLSAboveTheBarMaxPoolSize(<= 32768)
 - *Must have matching amount of real storage*
 - *Must have adequate cache structure sizes*
 - Recommended when LRU for 31 bit pool regularly in accelerated / panic mode
 - Great for applications needing to traverse large amounts of data repeatedly





Performance and Tuning – Buffer Pools



- Pool Size values are a goal for which the LRU tries to maintain.
- If more buffers are required at any given time, the pool may temporarily exceed the values set.
- Total size of buffer pools should not exceed amount of real storage.
- A paged out buffer is immediately freed by the LRU.





Performance and Tuning – PARMLIB

- SYS1.PARMLIB(IGDSMSxx)
 - RLSINIT(YES | <u>NO</u>)
 - SMSVSAM started during system init
 - Set this to YES or issue V SMS,SMSVSAM,ACTIVE at IPL
 - CF_TIME(xx seconds | <u>3600</u>)
 - Interval for recording of SMF 42 subtype 15 19
 - The default is once / hour
 - The standard is generally once every 15 minutes
 - SMF_TIME(YES | <u>NO</u>)
 - Synchronizes creation of SMF 42 subtypes 2, 15 19 with RMF intervals
 - "YES" helps when correlating data between SMF and RMF during performance diagnosis





Performance and Tuning – PARMLIB

- SYS1.PARMLIB(IGDSMSxx)
 - RLS_MaxCFFeatureLevel(<u>Z</u> | A)
 - Dictates size of data placed in CF cache structures
 - Z: only CIs less than 4K cached
 - Saves space in the cache structure
 - Advantageous if data is primarily read-only and remains valid in local buffers
 - A: CIs up to 32K cached
 - Requires more space in the cache structure
 - Advantageous when shared data is updated across the sysplex
 - Allows use of RLS CF Cache Feature in data class
 - CA_RECLAIM (<u>NONE</u> | {DATACLAS | DATACLASS})
 - Enabling CA reclaim reduces KSDS fragmentation and the need for reorgs
 - None: CA reclaim is not used, regardless of data class setting
 - DATACLAS | DATACLASS: Enables CA reclaim for all eligible data sets





- Cache Structures and Cache Sets
 - Allows for partitioning of workloads and data sets
 - Structures and sets defined in the base config (share control data set)
 - Many cache structures mapped to one cache set
 - One cache set assigned to each storage class
 - Structure allocation details defined in CFRM policy
- Proper cache structure sizing is required to ensure maximum buffer effectiveness!







- The ideal theoretical *single* cache structure size:
 - Sum of 31 bit and 64 bit buffer goals
 - 31 bit buffer goal = ((RLS_Max_Pool_Size) * Number of Systems)

• 64 bit buffer goal =

(RLSAboveTheBarMaxPoolSize(system1) + RLSAboveTheBarMaxPoolSize(system2) + ... + RLSAboveTheBarMaxPoolSize(systemN))

- Assumptions:
 - RLS_MaxCFFeaturelevel(A): caching all data
 - No sharing of data across the sysplex





- Example:
 - RLS_Max_Pool_Size(850)
 - Number of Systems = 2
 - RLSAboveTheBarMaxPoolSize(System1,2048)
 - RLSAboveTheBarMaxPoolSize(System2,4096)
 - Cache_Structure_Size = (850*2) + 2048 + 4096 = 7844M
 - Cache structure sizes less than the ideal amount should be closely monitored for directory reclaims.





- Recommendations
 - Allocate different cache structures for different applications.
 - Allocate different structures to match internal structure usages.
- If more than one cache structure to be allocated:
 - Divide total ideal theoretical allocation across all structures
 - Base sizes on data set usage distribution
 - Size
 - Number
 - Amount of data accessed





- False Invalids:
 - Take place when a cache structure(s) too small for current buffer allocations
 - Cache structures invalidate entries to make room for new requests
 - Causes additional requests to DASD!
 - Tracked in
 - RMF panels RLSSC, and RLSDS
 - SMF 42 subtypes 15-16







Performance and Tuning – Cache vs Buffer

- A balance between buffer and cache allocations must be made for optimum performance
- buffer space = large and cache space = small
 - Bottleneck in cache will occur between buffers and DASD
 - Cache will invalidate buffer entries to make room for new requests: thrashing
 - False Invalids occur when the cache structure needs more space than it has to satisfy CI requests.
- buffer space = small and cache space = large
 - Potentially wasted coupling facility space is expensive
 - Cache will try to keep as much data in it as possible
 - Buffer LRU will be the constraining factor



Lock Structures and Lock Sets

- Allows for partitioning workloads and environments (i.e. dev and prod)
- Structures and sets defined in the base config (SCDS)
- One lock structure mapped to one lock set
- One lock set assigned to each storage class
- Structure allocation details defined in CFRM policy







Lock_Structure_Size =

(10M * Number_Of_Systems * Lock_Entry_Size)

- Lock_Entry_Size (depends on the CFRM MAXSYSTEM value)
 - MAXSYSTEM <= 7 Lock_Entry_Size = 2
 - MAXSYSTEM >= 8 & <24
 Lock_Entry_Size = 4
 - MAXSYSTEM >=24 & <=32 Lock_Entry_Size = 8
- Example: MAXSYSTEM = 16 and 8 systems in sysplex

Lock_Structure_Size = 10M * 4 * 8 = 320M

Note: Minimum size of 13M is recommended





- Small lock structures result in increased false contention rates
 - True or false contention results in asynchronous lock requests
- To review false contention
 - RMF CF activity report
 - SMF 42 subtype 17
 - D SMS,CFLS command
- Recommended false contention rate is < 5%







Performance and Tuning – Data Class

- BWO (NO | TYPECICS | TYPEIMS)
 - BWO: Backup While Open
 - NO: BWO not used for CICS or IMS VSAM data sets
 - TYPECICS: BWO used for CICS VSAM file control data sets
 - Requires multiple SVC 26 calls to catalog
 - Component 1 lock held across VRM requests
 - TYPEIMS: BWO used for IMS VSAM data sets with IMS 6.1+
 - Recommendation: To maximize concurrent VRM performance, only use BWO when absolutely necessary.
- RLS CF Cache Value (<u>ALL</u> | NONE | UPDATESONLY)
 - ALL: Cache data and index CIs
 - None: Cache index CIs only
 - Updates: Cache CIs for write requests only
 - This setting only applies with RIsMaxCFFeatureLevel(A)!





Performance and Tuning – Data Class

- RLS Above the 2-GB Bar (YES | <u>NO</u>)
 - Specifies whether SMSVSAM can use 64-bit virtual buffering for the data set
 - NO:
 - Recommended for heavy insert and update data sets
 - YES:
 - Recommended for heavy read data sets where data is reaccessed within one hour's time frame
 - Must specify a non-zero RLSAboveTheBarMaxPoolSize
 - The current design of 64 bit buffering uses 10-20% more CPU for equivalent 31 bit requests, however, the large 64 bit pool size allows for increased buffer hits over the 31 bit pool. 64 bit buffering will be redesigned in a future release to address this CPU increase.





Performance and Tuning – Data Class

- Extent Constraint Removal (YES | NO)
 - Specifies whether data sets can have > 255 extents
- CA Reclaim (<u>YES</u> | NO)
 - Specifies whether DASD space for empty CAs in KSDS data sets will be reused
 - Value saved in catalog
 - Only used if CA reclaim function enabled in IGDSMSxx or with SETSMS command
 - Recommended to avoid data set fragmentation and to reduce reorg frequency







Performance and Tuning – Data Set

CISIZE

- Currently, both RLS and base VSAM will serialize the *entire* data set when
 - CI / CA split is in progress
 - CI / CA reclaim is in progress
 - Accessing spanned records
- Recommendation: When reasonable, increase CI size for data sets with higher insert activity.
 - Larger is better in terms of CPU
 - Larger also means larger buffers: performance vs. storage tradeoff
 - Larger CIs will contain more records which may lead to contention on DASD flush





Performance and Tuning – ACB / GENCB

- RLSREAD (<u>NRI</u> | CR | CRE)
 - NRI
 - No Read Integrity
 - Will not get record lock
 - CR
 - Consistent Read
 - Will get / release record lock
 - CRE
 - Consistent Read Extended
 - Will get record lock
 - Lock released at commit (recoverable data sets only)





Performance and Tuning – RPL

• OPTCD

- ASY | SYN
 - Asynchronous: SRB
 - Synchronous: TCB
- DIR | SEQ | SKP
 - Direct: Process request starting with specific record
 - Sequential: Process request record by record
 - Skip Sequential: Position on a record and move sequentially
- NRI | CR | CRE
 - No Read Integrity
 - Consistent Read
 - Consistent Read Extended







Performance Measurements





• SMF 64 Records

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SMF64RIO:

SMF64NLR:

- Cut by CLOSE and EOV on a per ACB basis:
 - SMF64DLR: number of logical records
 - SMF64DDE: number of delete requests
 - SMF64DIN: number of insert requests
 - SMF64DUP: number of update requests
 - SMF64DRE: number of retrieve requests
 - SMF64BMH: number of BMF hits in the local buffer pool
 - SMF64CFH: number of CF hits in the RLS cache structure
 - number of requests read from DASD
 - SMF64DEP: total number of requests
 - number of logical records at open





- SMF 42 Subtypes 15, 16, 17, 18, 19
 - **Subytpe 15** RLS statistics by storage class
 - Subtype 16 RLS statistics by data set
 - Must use V SMS,MONDS(spherename),ON to collect subtype 16 statistics
 - **Subtype 17** RLS CF lock structure usage
 - **Subtype 18** RLS CF caching statistics
 - Subtype 19 Buffer Manager LRU statistics
- Note: Only one system in the sysplex collects the SMF 42 records. The system collecting the records is displayed in the D SMS,SMSVSAM operator command.





- RMF Monitor III
 - All statistics gathered from interface call with SMSVSAM address spaces
 - Collected on RMF defined intervals
 - Can be matched to SMF intervals with SMF_TIME(YES)
 - Any SMSVSAM data in RMF monitor reports also found in SMF 42 records





• RLSSC & RLSDS

- RLSSC: VSAM RLS activity by storage class
- RLSDS: VSAM RLS activity by data set
- Two different organizations split into sequential and direct requests
- Provides information on:
 - Local buffer access (BMF)
 - CF cache structure hit %
 - DASD hit %
 - % valid buffer / cache hits
 - % false invalid buffer / cache hits
- Assists in determining sizes
 - CF cache structures
 - Buffer pools
- Helps analyze overall application performance





RMF Sysplex Report Selection Menu Selection ===>

Enter selection number or command for desired report.

Sysplex Reports

1 SYSSUM	Sysplex performance summary	(SUM)
2 SYSRTD	Response time distribution	(RTD)
3 SYSWKM	Work Manager delays	(WKM)
4 SYSENQ	Sysplex-wide Enqueue delays	(ES)
10 RLSSC	VSAM RLS activity by storage class	(RLS)
11 RLSDS	VSAM RLS activity by data set	(RLD)
12 RLSLRU	VSAM LRU overview	(RLL)





RLSSC - Sysplex Totals View

Command ===>	RMF V1R	8 VSA	M RLS	Activi	ty - S	YSPLEX	Line 1 Scroll ==	of 12 => HALF
Samples: 59	Systems:	2 Da	te: 10	/31/06	Time:	13.16.0	0 Range: 60) sec
< 2	2GB / > 2	GB						
LRU Status : Go	od / Acc	el						
Contention % : 0	0.0/0.	0						
False Cont % : 0	0.0/0.	0						
Stor Class Access	s Resp		Rea	nd			BMF	Write
	Time	Rate	BMF%	CF%	DASD%	Valid%	False Inv%	Rate
RLS_SC1								
Below 2GB DIR	0.004	665.6	88.2	0.5	11.3	100	0.01	0.00
SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB DIR	0.004	665.6	88.2	0.5	11.3	100	0.01	0.00
SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
RLS_SC2								
Below 2GB DIR	0.005	200.0	90.5	0.0	9.5	100	0.00	0.00
SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00





RLSSC - System/CF Structure View

Command>		RMF V1	R8 VSAM	RLS A	ctivit	y - S	YSPLEX	Line 1	of 23
									-> TIALI
Samples: 120	Sy	/stems:	2 Da	te: 10,	/31/06	lime:	13.25.0	0 Range: 12	0 Sec
	< 20	GB / > 2	GB						
LRU Status	: Goo	od / Acc	el Sto	rage C	lass :	RLS_SC	1		
Contention %	: 0.	0 / 0.	0 Cac	he Set	:	PUBLIC	1		
False Cont %	: 0.	0 / 0.	0						
System/CF	Access	Resp		Read	d b			BMF	Write
		Time	Rate	BMF%	CF%	DASD%	Valid%	False Inv%	Rate
*ALL									
Below 2GB	DIR	0.000	14.98	83.0	0.0	17.0	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB	DIR	0.000	14.98	83.0	0.0	17.0	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
SYS1									
CACHE01									
Below 2GB	DIR	0.000	7.49	83.0	0.0	17.0	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB	DIR	0.000	7.49	83.0	0.0	17.0	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00





RLSDS - Sysplex Totals View

Command>	RMF V1R8 V	SAM RLS Ac	tivity	- SYS	SPLEX	Line 1	of 20 -> HALE
Samploc: 120	Systoms: 2	Data: 10	/31/06	Timor	12 25 00	Pango: 12	
	$\gamma \gamma $	Date. 10	/ 51/00	TIME.	13.23.00	Range. 12	0 Sec
< 2	GB / > ZGB						
LRU Status : Go	od / Accel						
Contention % : 0	0.0 / 0.0						
False Cont % : 0	0.0 / 0.0						
Sphere/DS Access	Resp	Rea	d		BM	1F	Write
	Time R	ate BMF%	CF%	DASD%	Valid% F	alse Inv%	Rate
BMAI.VSAMIN.MEGA							
BMAI.VSAMIN.MEGA.	AIX.DATA						
Below 2GB DIR	0.003 0	.01 0.0	0.0	100	0.0	0.00	0.00
SEQ	0.000 0	.00 0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB DIR	0.003 0	.01 0.0	0.0	100	0.0	0.00	0.00
SEQ	0.000 0	.00 0.0	0.0	0.0	0.0	0.00	0.00
BMAI.VSAMIN.MEGA.	AIX.INDEX						
Below 2GB DIR	0.003 0	.03 50.0	0.0	50.0	100	0.00	0.00
SEQ	0.000 0	.00 0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB DIR	0.003 0	.03 50.0	0.0	50.0	100	0.00	0.00
SEQ	0.000 0	.00 0.0	0.0	0.0	0.0	0.00	0.00





RLSDS - System View

	F	RMF V1R8	VSAM R	LS Act:	ivity	- SYS	PLEX	Line 1	of 9
Command ===>								Scroll ==	=> HALF
Samples: 120) Sy	/stems:	2 Da	te: 10,	/31/06	Time:	13.25.0	0 Range: 12	0 Sec
	< 20	GB / > 2	GB						
LRU Status	: Good	d / Acc	el				Storage	Class : RLS	SC1
Contention %	s : 0.	0 / 0.	0		C	ache Se	t :	RLSCSET	-
False Cont %	s : 0.	0 / 0.	0	[Data S	et	: RLSAD	SW.VF01D.INV	'ENTOR.IND
System/CF	Access	Resp		Read	d			BMF	Write
-		Time	Rate	BMF%	CF%	DASD%	Valid%	False Inv%	Rate
*ALL									
Below 2GB	DIR	0.000	24.27	97.7	2.2	0.1	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
Above 2GB	DIR	0.000	24.27	97.7	2.2	0.1	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
SYS1									
RLSCACHE01									
Below 2GB	DIR	0.000	0.20	100	0.0	0.0	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00
SYS2									
RLSCACHE01									
Above 2GB	DIR	0.000	24.07	97.7	2.2	0.1	100	0.00	0.00
	SEQ	0.000	0.00	0.0	0.0	0.0	0.0	0.00	0.00





- RLSLRU
 - Provides buffer manager LRU statistics for each system
 - Helps to size 32 and 64 bit pools
 - Properly sized pools help to tune cache structures





• RLSLRU

An example of a RLSMAXPOOLSIZE = 500 and RLSABOVETHEBARMAXPOOLSIZE = 2000 specification:

	RMF	V1R8 VS	AM LRU	Overview	- SYS	PLEX		Line 1 of	f 2
Command ===>	>						Scro	oll == =>	HALF
Samples: 120	9 System	s: 2	Date:	10/31/06	Time:	13.25.00	Rar	nge: 120	Sec
MVS	Avg CPU -	Buffer S	Size -	Accel	Reclaim		Read		
System	Time	Goal I	High	<u>%</u>	0/0	BMF%	CF%	DASD%	
SYS1									
Below 2GB	0.023	500	352	0.0	0.0	82.0	8.0	10.0	
Above 2GB	3.543	2000	1552	0.0	0.0	95.0	2.0	3.0	
SYS2									
Below 2GB	4.457	500	612	100.0	0.0	31.2	10.5	58.3	



RLSLRU - Buffer Counts by Pool View

RMF VSAM LRU Overview - Buffer Counts by Pool The following details are available for MVS System: SYS2 Press Enter to return to the Report panel. Fixed Storage Below 2 GB : 0 Above 2 GB : Real Storage % Below 2 GB : 0 Above 2 GB : ----- Above 2 GB ---- Above 2 GB ----High Low Avg Low High Avg Fixed Pages 0 0 Buffer Counts by Pool: More: + 2K 4K 6K 8K 10K 12K 14K 16K . . . 30K 32K



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SMSVSAM Diagnostics





- D SMS,CFLS(lock_structure)
 - Displays lock structure statistics to the console
 - Contention values for IGWLOCK00 (default) or specified structure
 - Tracks
 - LockRate: Number of lock requests / sec
 - ContRate: Percentage of requests that encountered contention
 - FContRate: Percentage of requests encountering false contention
 - WaitQLen: Average number of requests waiting behind locks





D SMS,CFLS(lock_structure)

IGW320I 00:01:48 Display SMS,CFLS(IGWL0CK00) PRIMARY STRUCTURE: TGWI OCK00 VERSION: C57C859902B9264F STZE: 95232K RECORD TABLE ENTRIES: 34438 USED: 747 SECONDARY STRUCTURE: IGWLOCK00 VERSION: C58548B866A7576E SIZE: 95232K RECORD TABLE ENTRIES: 34438 USED: 747 LOCK STRUCTURE MODE: DUPLEXED STATUS: ENABLE LockRate System Interval ContRate FContRate Wait0Len 190 1 Minute 15.10.000 0.331 0.00 J90 1 Hour 11.9 0.000 0.047 0.01 8 Hour 190 8.2 0.001 0.039 0.00 0.001 1 Dav 17.3 0.01 J90 0.042 (09) 1 Minute 7.5 0.000 0.161 0.00 7.0 0.000 0.00 (09) 1 Hour 0.020 (09) 8 Hour 6.6 0.001 0.033 0.00 (09) 1 Day 20.9 0.001 0.029 0.00

LockRate = number of lock requests per second CONTRATE = % of lock requests globally managed FCONTRATE = % of lock requests falsely globally managed WaitQLen = Average number of requests waiting for locks





- D SMS,SMSVSAM,DIAG(Contention)
 - Displays latch contention on SMSVSAM resources
 - Latches are 8 bytes of storage used to serialize system resources
 - Two states: Holder & Waiter
 - Latches are not shared
 - The contention display shows
 - Latch address
 - Holder and any waiters if the latch is in contention
 - Elapsed time of contention period
 - System specific (will not show latches around the plex)
 - Elapsed time of more than a few seconds is likely a problem
 - Should be used in conjunction with D GRS,C





D SMS, SMSVSAM, DIAG(Contention)

09.55.29	SYSTEM1		IGW34	43I VSAM F	RLS DI	LAG STATUS	5 (V.01)
RESOU	RCE	W	AITER		HOL	_DER E	ELAPSED
TYPE	ID	JOB NAME	ASID	TASK	ASID	TASK	TIME
LATCH	7F158C70	SMSVSAM	003A	008DA250	003A	008D7218	00:00:06
DESC	RIPTION:	IGWLYSPH	- SHM	OBJECT PO	00L		
LATCH	7F151E78	SMSVSAM	003A	008D7218	003A	008DC1C8	00:00:21
DESC	RIPTION:	IGWLYDTS	- SHM	OBJECT PO	00L		
LATCH	7BAD43B8	SMSVSAM	003A	008DC1C8	002D	007F3000	00:19:09
LATCH	7BAD43B8	SMSVSAM	003A	008D5A48	002D	007F3000	00:22:09
LATCH	7BAD43B8	SMSVSAM	003A	008D6938	002D	007F3000	00:33:23
LATCH	07F1B1D0	SMSVSAM	003A	008D64F8	003A	008D6CF0	01:47:20
LATCH	07F1D3B8	SMSVSAM	003A	008D6CF0	0000	00000000	11:23:30





• D SMS, SMSVSAM, QUIESCE

- Displays outstanding quiesce activity for SMSVSAM
- A Quiesce event is a decision to shutdown / allow CICS access to an RLS dataset
- Any CICS regions participating in the event will be in the display
- If there is not an event in progress, IGW540I rejects the command
- Invaluable to determine CICS regions holding up a quiesce request
- CICS regions with 00.00.00 in the "completed" section are likely in trouble
 - Dump these regions with SMSVSAM prior to termination
- Often the region can be terminated instead of SMSVSAM
 - Allows the quiesce event to finish
 - Saves a plex-wide SMSVSAM outage





• D SMS, SMSVSAM, QUIESCE

With quiesce activity

IGW540I 13.30.45 DISPLAY SMS, SMSVSAM, QUIESCE

MVS1

SPHERE NAME: DLLEHR.TEST1

 SYSTEM NAME: MVS1
 START TIME:
 .27.50 TOTAL ELAPSE TIME: 57.02.55

 PARTICIPATING SUB-SYSTEM STATUS:
 SCHEDULED:
 COMPLETED:
 ELAPSE:

 SUB-SYSTEM NAME:
 CI1AORP1
 .27.50
 00.00.00
 57.02.55

 SUB-SYSTEM NAME:
 CI1AORP2
 .27.50
 .27.50
 00.00.00

Without quiesce activity

IGW540I 07.54.28 DISPLAY SMS, SMSVSAM, QUIESCE DISPLAY SMSVSAM QUIESCE SPHERE IS REJECTED. NO QUIESCE EVENTS ARE ACTIVE ON THIS SYSTEM.





SMSVSAM Diag – Hang / Wait in RLS

- Hang / wait conditions in RLS
 - May be related to activities such as
 - Accessing data sets for RLS
 - Starting / stopping of RLS client spaces (e.g. CICS)
 - Starting / stopping of SMSVSAM
 - SMSVSAM commands not responding
 - Detecting / pinpointing this condition requires
 - Display command output
 - Plex-wide dumps of SMSVSAM and related client spaces





SMSVSAM Diag – Hang / Wait in RLS

- Display commands to run
 - D SMS, SMSVSAM, ALL
 - D SMS, TRANVSAM, ALL (if TVS is in use)
 - D SMS, SMSVSAM, DIAG (CONTENTION) (system scope)
 - D SMS, SMSVSAM, QUIESCE (system scope)
 - D GRS,C (system scope)
 - D XCF,STR,STRNM=[IGWLOCK00 | lock_structure]
 - D XCF,STR,STRNM=[cachestructurename]





SMSVSAM Diag – Hang / Wait in RLS

- Plex-wide dumps
 - Most required info is in SMSVSAM address and data spaces
 - Both SMSVSAM and MMFSTUFF dataspaces are needed
 - SMSVSAM is thoroughly plex-wide so the plex must be dumped
 - Creating an IEADMCxx parmlib member is recommended
 - Example of a plex-wide parmlib definition:

```
COMM=(SMSVSAM hang)
JOBNAME=(*MASTER*,SMSVSAM),
DSPNAME=('SMSVSAM'.*),
SDATA=(ALLNUC,COUPLE,CSA,GRSQ,LPA,LSQA,PSA,RGN,SUM,SQA,SWA,TRT,XESDATA),
REMOTE=(SYSLIST=(*('SMSVSAM')),DSPNAME,SDATA)
```





SMSVSAM Diag – Internal Logic Errors

- Internal logic errors
 - Surfaced as an ABEND0F4 with unique RC and RSN
 - SVC dump taken (typically plex-wide)
 - Depending on severity, processing may be stopped
 - Dumps may be accompanied by syslog flower box
 - IGW400I / IGW405I
 - VRM request type and key
 - Jobname / data set name
 - ASID / TCB
 - User and data space RPL
 - VRM footsteps
 - When encountered, necessary doc includes
 - Plex-wide 0F4 dumps
 - OPERLOG
 - Plex-wide software detail EREP reports (LOGREC)







RLS Maintenance and References





Maintenance and References – PSP Bucket

- The RLS PSP bucket is on the web at:
 - http://www14.software.ibm.com/webapp/set2/psearch/sea rch?domain=psp
- 'Search for: SMSVSAM'
- Maintenance organized by release
 - Upgrade ZOSV1R11, Subset DFSMS
 - Upgrade ZOSV1R12, Subset DFSMS
 - Upgrade ZOSV1R13, Subset DFSMS





References

- VSAM Demystified
 - SG24-6105-01
- z/OS V1R13 DFSMSdfp Storage Administration
 - SC26-7402-16
- z/OS V1R13.0 DFSMS: Using Data Sets
 - SC26-7410-11
- z/OS MVS System Management Facilities
 - SA22-7630-24





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Conclusion

- RLS performance increases as a result of shortening the I/O path.
- Reducing time in I/O can be best accomplished when
 - Data is in the buffers
 - Data is in the cache structures on the coupling facility
- Reducing CPU utilization at the data set level is another important means of performance improvement.
 - Better data set tuning
 - Avoid serialization intensive situations
 - Take advantage of offerings in the SMS configuration







DFSMS Basics: VSAM VSAM RLS Best Practices

How to leverage VSAM RLS best practices and performance in your environment David LeGendre, dlegendr<u>@us.ibm.com</u> Session : 10967