Buffer-to-Buffer Credits, Exchanges, and Urban Legends

Lou Ricci, IBM
Howard L. Johnson, Brocade

8 August 2011 (3:00pm – 4:00pm)
Session 9931
Room Europe 7
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Abstract

- Performance in a FICON network is influenced by the underlying flow control mechanisms of Fibre Channel. In this session, we examine how Buffer-to-Buffer credits flow from the channel to the control unit. We also look at how exchanges are used in FICON applications and how they change with the introduction of zHPF. During both examinations, we explore the role of the FICON Director in managing Buffer-to-Buffer credits and exchanges over a cascaded network. Throughout the session, we debunk the various FICON “Urban Legends” featuring credits and exchanges. Take the opportunity to learn from two of the FICON industry’s leading experts in channel and fabric development and join our session.
Agenda

• **Buffer Credits**
  • What are they and how do they work?
    • How do you *fill the pipe*?
    • What if you *can’t fill the pipe*?
    • What’s wrong with *multiple senders and one receiver*?
    • What happens when the *pipes are different sizes*?
    • What’s it like in the *real world*?
    • How do *cascades Directors* work?

• **Exchanges**
  • What are they and how do they work?
    • What’s an Exchange?
    • How many exchanges are needed?
    • Can they be "reused"?
    • Can you have too many exchanges?

• **Error Sensitivity**
  • Is FICON more sensitive to errors than FCP?
    • How sensitive are FICON frames to loss or corruption?
    • What recovery actions are taken?
    • What are the differences with FCP?
What are they and how do they work?

BUFFER CREDITS
What is Buffer-to-Buffer Credit?

- The greater the BB Credit….
  - A. The faster frames can be sent
  - B. The farther apart the two ports can be
  - C. The larger the frames can be
  - D. None of the above
What is Buffer-to-Buffer Credit?

• The greater the BB Credit…. 
  • A. 
  • B. The farther apart the two ports can be 
  • C. 
  • D.
Flow Control

- Related to the devices' ability to receive and process frames
- Manages when frames are coming faster than they can be processed
- Dropped frames occur when frames are arriving too fast to be processed

- Frames can only be transmitted when the receiver is ready
- Credit establishment communicates the number of frames a device can receive at a time
- The credit value is exchanged at login
- Transmission stops when credit runs out
- The receiver indicates when it is ready to receive more frames
Buffer Credit

- At initialization, the two ports establish credit
  - Each buffer credit corresponds to a frame (regardless of size)
- Each side can support different values
  - Credit Count
- If a port doesn’t have credit, it can’t send a frame
  - Credit Count has reached zero
- Mechanism limits frame drops

Credit accounting after Fabric Login
**Receiver Ready (R_RDY)**

- **R_RDY**
  - Used for link level flow control
  - Called buffer-to-buffer credit (BB Credit)

- **R_RDY is not a frame**
  - It is a “primitive” so it doesn’t consume a buffer

- **Frame transmission**
  - BB Credit is decremented
    - Once for each frame transmitted
  - When BB Credit = 0
    - Transmission stops

- **Frame received**
  - R_RDY is sent
  - Causes transmitter to increment BB Credit

---

**Credit Accounting during Transmission**

- Host says, “I can send 8 frames without stopping.”
- Switch says, “I can 40 frames that way and 16 frames the other way without stopping.”
- Storage says, “I can send 8 frames without stopping.”
- Credit Count 1 ➔ Credit Count 40 ➔ Credit Count 16 ➔ Credit Count 5 ➔ R_RDY ➔ R_RDY
Urban Legend:
Buffer Credits at Zero are a Problem

- Buffer credit determines DISTANCE
  - The distance two nodes can be apart and still maintain full link frame rate

- Buffer credit is the number of FRAME buffers
  - A port provides for it’s NEAREST neighbor for RECEIVING frames
  - Does NOT have to be symmetrical

- Buffer credit is a FRAME count
  - Not a data SIZE
  - A 1 byte frame consumes 1 buffer credit
  - A 2K byte frame consumes 1 buffer credit

- Number of credits needed is determined by:
  - Raw Link Speed
  - Speed of light thru a fiber
  - Distance between two adjacent nodes
Example: A full pipe

BUFFER CREDITS
Initial Conditions

**Channel**

**“Perfect” Switch**

**Control Unit**

Number of B-B Credits the switch advertised to the channel during link init

NOTE: In these animations, both the frames and R_RDY’s are numbered. This is for illustrative purposes only. In reality, neither the frames nor the R_RDY’s are numbered. The arrival of an R_RDY only informs the receiver that a frame has been forwarded, now WHICH frame has been forwarded.
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Example: A not so full pipe

BUFFER CREDITS
Suppose the switch is too far away from the channel for the B-B credit it advertised to the channel.
Example: Multiple Senders and One Receiver

BUFFER CREDITS
Suppose there are multiple senders to one receiver. Each sender attempts to send at 100% link speed.
Example: Different sized pipes

BUFFER CREDITS
Fat Pipe / Skinny Pipe

20

8 Gig Link

4 Gig Link
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Example: Real Life?

BUFFER CREDITS
More real life example: Two senders sending at 30% - 50% link rate to one receiver
Example: Cascaded Directors

BUFFER CREDITS
Other Neat Stuff

BUFFER CREDITS
How much credit do I need?

- Good “Rule of thumb”

Number of credits needed = 1 + \( \text{Link speed in Gb/s} \times \text{Distance in km} \)

Frame Size in KB

Example: 20 km at 1 Gb/s
\[
1 + \frac{1 \times 20}{2} = 11
\]

Example: 10 km at 4 Gb/s
\[
1 + \frac{4 \times 10}{2} = 21
\]
How “long” is a frame?

- Traveling at the speed of light, a frame can be very long
  - At 1G, the length of a frame is about 4-kilometers.
  - At 2G, the length of a frame is about 2-kilometers.
  - At 4G, the length of a frame is about 1-kilometer.
  - At 8G, the length of a frame is about 500-meters.
  - At 16G, the length of a frame is about 200-meters.
How “fast” is a frame?

- Speed of light in fibre
  - 200,000 km/second
  - 5-microseconds/km
- Transmission Rates in fibre
  - At 1G, a frame is sent in about 20-microseconds.
  - At 2G, a frame is sent in about 10-microseconds.
  - At 4G, a frame is sent in about 5-microseconds.
  - At 8G, a frame is sent in about 2.5-microseconds.
  - At 16G, a frame is sent in about 1-microsecond.
The parts of a transmission

- **Frame**
  - Building block of an Fibre Channel connection
  - Contains the information to be transmitted
  - The address of the source and destination
  - Control information

- **Sequence**
  - A set of one or more related Frames
  - Transmitted unidirectionally from one port to an other

- **Exchange**
  - An Exchange is one or more nonconcurrent sequences
  - A single operation
  - May be unidirectional or bidirectional
Fibre Channel Frame

The basic building block is the FRAME

<table>
<thead>
<tr>
<th>SOF</th>
<th>Header</th>
<th>Payload</th>
<th>CRC</th>
<th>EOF</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>24</td>
<td>0 – 2048 (2112)</td>
<td>4</td>
<td>4</td>
</tr>
</tbody>
</table>

- Source/Destination link addresses
- Exchange ID
- Protocol Type
- Start or End of Sequence
- Start or End of Exchange

Identifies beginning of a frame

Protects Frame

Identifies end of a frame
Ficon IU Examples

1 Frame IU to transfer a Read CCW

<table>
<thead>
<tr>
<th>SOF</th>
<th>Header</th>
<th>Ficon Header</th>
<th>CRC</th>
<th>EOF</th>
</tr>
</thead>
</table>

3 Frame IU to transfer 4K of data

<table>
<thead>
<tr>
<th>SOF</th>
<th>Header</th>
<th>Ficon Header</th>
<th>CCW Data (2016 bytes)</th>
<th>CRC</th>
<th>EOF</th>
</tr>
</thead>
</table>

<table>
<thead>
<tr>
<th>SOF</th>
<th>Header</th>
<th>CCW Data (2048 bytes)</th>
<th>CRC</th>
<th>EOF</th>
</tr>
</thead>
</table>

<table>
<thead>
<tr>
<th>SOF</th>
<th>Header</th>
<th>CCW Data (32 bytes)</th>
<th>Ficon CRC</th>
<th>CRC</th>
<th>EOF</th>
</tr>
</thead>
</table>
Urban Legend: FICON uses fewer Exchanges Than FCP

- Fibre Channel Architecture defines an Exchange as
  - “A mechanism for identifying and managing an operation between two ports”

- All IUs (a.k.a. Sequences) that make up a single I/O operation are part of an Exchange

- In Ficon, each concurrent I/O operation uses two Exchanges
  - One unidirectional Exchange for IUs from the Channel to the CU
  - A different unidirectional Exchange for IUs from the CU to the Channel

- The PAIR is commonly know as a “Ficon Exchange”
Sequences and IUs

- Each Upper Layer Protocol (ULP) defines the contents and format of its own Information Units (IUs)
  - Commands
  - Data
  - Status
  - Control
  - Etc
- Ficon IUs can be up to 8K (8192) in size
  - 8160 (8K-32) bytes of data
    - 32 bytes contain Ficon Header information
  - 4 frames are needed for the largest IU
- The collection of frame(s) that make up a IU are called a Sequence
  - A Sequence may be as small as a single Frame
How many Exchanges do I need?

• Little’s Law states:
  • The number of “things” in a system can be determined by multiplying the average arrival rate of those “things” by the average time each “thing” stays in the system.

• Applied to Ficon:
  • The average number of Exchanges active at any given time = Average I/O rate * Average response time

  • Example: 5000 Ficon I/Os / Second on a given channel with .4ms service time\(^1\) needs 2 Active Exchanges (pairs) at any given time

\(^1\) The amount of time the I/O is active in the channel
ERROR SENSITIVITY
Urban Legend: FICON is more sensitive than FCP

- Is Ficon More Sensitive to Errors than FCP?
  - Reasons for Link Errors are the same

- Is a Ficon frame more likely to get lost, damaged or corrupted than FCP?
  - The probability is the same
  - Frames are frames

- When a Ficon frame gets lost, damaged or corrupted, is the recovery action different from FCP?
  - Both protocols retry when errors occur
    - FCP by the Device Driver
    - Ficon by IOS/ERP
So What are the Differences?

- z Operating Systems tend to provide more detailed messages
- Ficon does provide additional debug data and actions
  - RNID
  - Link Error Status Blocks
  - Extensive State Change Processing

Table 89 - Link Error Status Block format for RLS command

<table>
<thead>
<tr>
<th>Bits</th>
<th>31</th>
<th>..</th>
<th>00</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>Link Failure Count</td>
<td></td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>Loss-of-Synchronization Count</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>Loss-of-Signal Count</td>
<td></td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>Primitive Sequence Protocol Error</td>
<td></td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>Invalid Transmission Word</td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>Invalid CRC Count</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Source: FC-FS-3 INCITS/T11 Draft Standard v0.92
See www.t11.org
Thank you

SUMMARY
Summary

• Buffer Credits
  • Distance
  • Flow Control

• Exchanges
  • Unidirectional
  • Bidirectional

• Error Sensitivity
  • Recovery
  • Reporting
SHARE, Orlando, August 2011
Buffer-to-Buffer Credits, Exchanges, and Urban Legends
Session 9931
THANK YOU!
Speaker Biography

• Lou Ricci
  • IBM
    • 32-years
    • 24-years in channel development
    • An inventor of FICON
  • FICON Firmware Team Leader

• Contact Information
  • lricci@us.ibm.com
Speaker Biography

• Patty Driever
  • IBM
  • System z I/O and Networking Technologist

• Contact Information
  • pgd@us.ibm.com
Speaker Biography

• Howard L. Johnson
  • BROCADE
  • Technology Architect, FICON
    • 27 years technical development and management

• Contact Information
  • howard.johnson@brocade.com
BONUS SLIDES
Speed of Light

- **299 792 458 meters per second**
  - In a vacuum
  - National Institute of Standards and Technology
- **300 000 000 meters per second**
  - Common approximation
  - Wikipedia
- **1.44 – 1.46**
  - Refractive index of 1300 nanometer fiber
  - Encyclopedia of Laser Physics and Technology
    - http://www.rp-photonics.com/refractive_index.html
- **1.5**
  - Common approximation
  - Wikipedia
- **200 000 000 meters per second**
  - Calculated speed of light in single-mode fiber
  - \( V = \frac{C}{N} \)
  - \( 200 000 000 = \frac{300 000 000}{1.5} \)
End to End Credit

- Device to Device Flow Control
  - Between source and destination
    - Not the links
  - Similar to buffer-to-buffer flow control
    - At N_Port Login
      - Report available receive buffers (EE_Credit)
      - Transmitter counts buffers transmitted (EE_Credit_CNT)
    - Receiver acknowledges frame (ACK)
      - ACK 1 (a single data frame in a sequence) – most common
      - ACK n  (several (N) consecutive data frames in a sequence)
      - ACK 0  (all data frames in a sequence) – not used
Virtual Channels

- Technology to allocate BB_Credits to particular data flows
  - Class F traffic has one data flow
  - Assigned with Zoning by using special Zone names
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