



# Look What I Found Under The Bar!

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# Agenda

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- Overview of Language Environment storage areas
  - Control blocks, stack, heap
  - Which can you control?
- How to control Language Environment Storage
- Tuning Storage
- More advanced tuning
- Sources of Additional Information



# Language Environment Storage Areas

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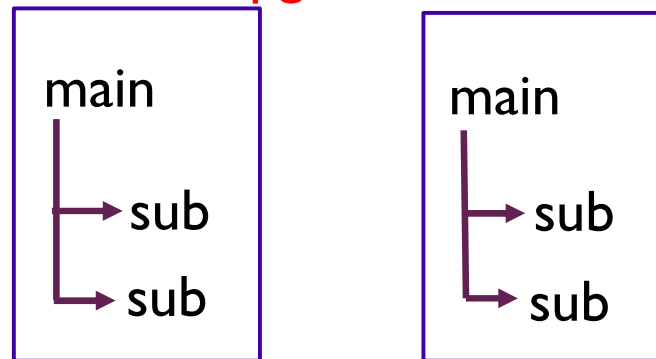
- Language Environment Control Blocks
  - Region level
    - Normally 1 region per address space
  - Process level
    - Normally 1 process per address space
  - Enclave level
    - Potentially many per address space
  - Thread level
    - Potentially very many per address space

# Language Environment Storage Areas

region - address space

process - application

enclave - pgm - enclave



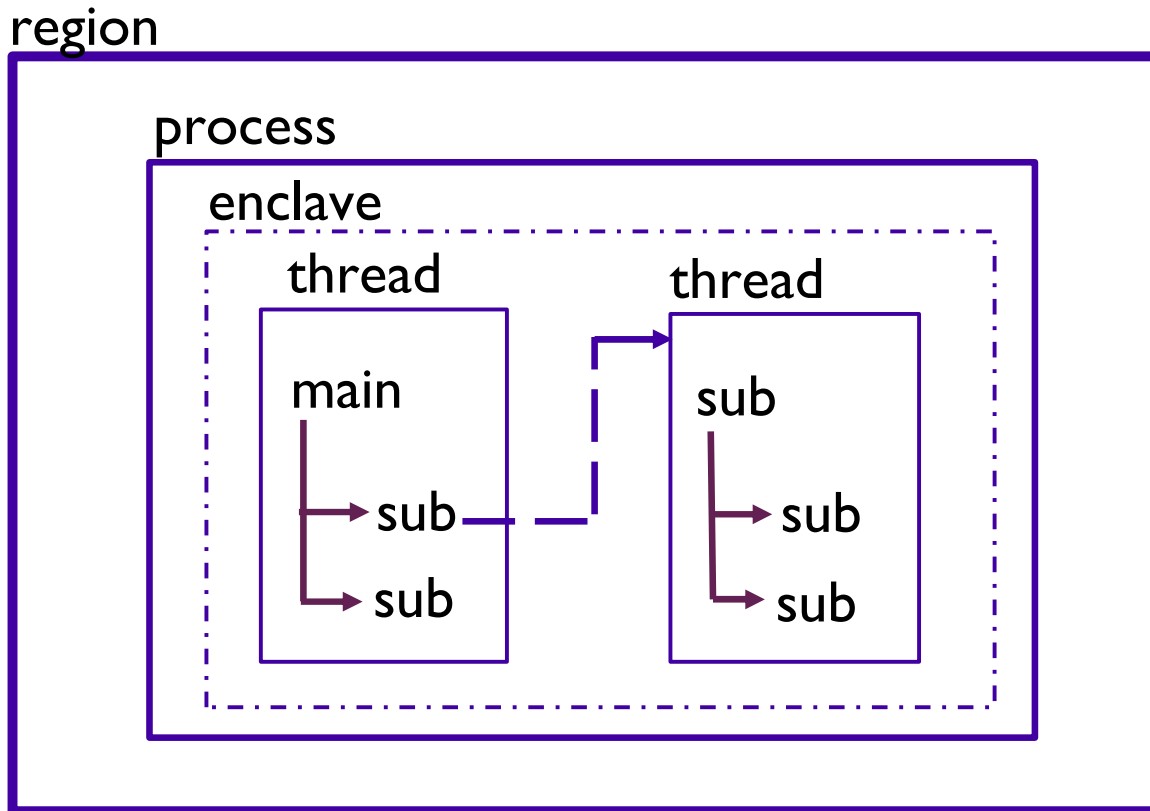


# Language Environment Storage Areas

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- Language Environment Enclave
  - Every “main” program is a new enclave
  - Every “link” is a new enclave
  - Contains
    - CEECAA, CEEEDB, CEEOCB, stacks, heaps, environment variables, and much more!
  - “Extra” enclaves are expensive both to initialize and in storage usage.

# Language Environment Storage Areas



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# Language Environment Storage Areas

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- Language Environment Thread
  - Only created by `pthread_create()`
  - Contains
    - CEECAA, stacks, and a little more
  - Threads are not nearly as expensive as enclaves.
  - Used mostly by C/C++
  - Enterprise PL/I multitasking uses threads





# Language Environment Storage Areas

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- Language Environment Stacks
  - Stacks
    - Last In, First Out structures
    - Allow programs to be reentrant
    - Thread level structures
      - “Main” programs have separate stacks
      - “Linked” programs have separate stacks
      - Pthreads have separate stacks



# Language Environment Storage Areas

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- Language Environment supports 2 independent stacks
  - User stack – (poorly named)
    - Used by user programs and Language Environment
  - Library stack
    - Used “rarely” by Language Environment
    - Always below the 16M line



# Language Environment Storage Areas

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- DATA in stacks
  - “Chunks” are called stack segments
    - Made up of 1 or more DSAs
  - DSA – Dynamic Save Area
    - Also called a “stack frame”
  - DSAs contain
    - Register Save Area (RSA)
    - NAB – Next Available Byte
    - Automatic (local) variables
      - C – int i;
      - PL/I – declare i fixed;
      - NOT COBOL WORKING-STORAGE
        - COBOL LOCAL-STORAGE in stack



# Language Environment Storage Areas

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- Language Environment Heaps
  - Heaps
    - Completely random access
    - Allows storage to be dynamically allocated at runtime
  - Enclave level control structures
    - Each 'main' has a separate stack and heap
    - Each 'link' causes a separate stack and heap
    - pthreads share a single heap for all threads



# Language Environment Storage Areas

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- Language Environment Heaps
  - Four independently maintained sets of heap segments all with similar layouts:
    - User Heap
      - COBOL WORKING-STORAGE
      - C/C++ (malloc or operator new)
      - PL/I dynamic storage (allocate)
    - LE Anywhere Heap
      - COBOL and LE above the line CBs
    - LE Below Heap
      - COBOL and LE below the line CBs
    - Additional Heap
      - Defined by the user



# Controlling Storage

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- Run-time options dealing with stacks
  - STACK(init,inc,ANY|BELOW,KEEP|FREE,dsInit,dsInc)
    - Init - Initial size of storage "chunk" allocated and managed by LE for user stack
    - Inc - When init is full, size of next storage "chunk" (increment)
    - ANY|BELOW - Location of storage
      - ANY Anywhere in 2G virtual storage
      - Below Always below 16M line
        - Required when all31(OFF)
    - KEEP|FREE - What to do when done with inc
      - KEEP Do not free the storage "chunks"
      - FREE Free the storage "chunks"
    - DsInit - Initial size of storage "chunk" (XPLINK)
    - DsInc - When initial full, size of next "chunk" (XPLINK)



# Controlling Storage

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- Run-time options dealing with stacks
    - LIBSTACK(init,inc,KEEP|FREE)
      - Init - Initial size of storage "chunk" allocated and managed by LE for library stack
      - Inc - When init is full, size of next storage "chunk" (increment)
      - KEEP|FREE - What to do when done with inc
        - KEEP Do not free the storage "chunks"
        - FREE Free the storage "chunks"
- NOTE: No ANY|BELOW, LIBSTACK always below the 16M line



# Controlling Storage

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- Run-time options dealing with stacks
  - THREADSTACK(ON|OFF,init,inc,ANY|BELOW,KEEP|FREE,dsInit,dsInc)
    - ON|OFF – Whether or not to use THREADSTACK for pthreads
    - Init - Initial size of storage “chunk” (like STACK)
    - Inc - Increment size of storage “chunk” (like STACK)
    - ANY|BELOW - Location of storage
      - ANY Anywhere in 2G virtual storage
      - Below Always below 16M line
        - Required when all31(OFF)
    - KEEP|FREE - What to do when done with inc
      - KEEP Do not free the storage “chunks”
      - FREE Free the storage “chunks”
    - DsInit, Dsinc – XPLINK “chunk” sizes





# Controlling Storage

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- Runtime options dealing with the heaps
  - HEAP(init,inc,ANY|BELOW,KEEP|FREE,int24,inc24)
    - User heap - mostly application use
    - init - Initial size of the "chunk" of storage obtained to be managed by LE for user heap
    - Inc - When initial "chunk" is full, size of next "chunk" (minimum)
    - ANY|BELOW - Location of "chunk"
      - Not sensitive to ALL31 setting
    - KEEP | FREE - What to do when done with the increment when empty
      - KEEP - Do not free the storage "chunks"
      - FREE - Free the storage "chunks"
    - int24 - Initial size of the "chunk" of storage obtained
      - (if ANY specified but BELOW requested (minimum))
    - inc24 - Size of next "chunk"
      - (if ANY specified but BELOW requested (minimum))



# Controlling Storage

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- Runtime options dealing with the heaps...
  - ANYHEAP(init,inc,ANY|BELOW,KEEP|FREE)
    - LE use - normally above the line
    - init - Same as HEAP.
    - inc - Same as HEAP. (minimum)
    - ANY | BELOW - Location of storage
    - KEEP | FREE - Same as HEAP
  - BELOWHEAP(init,inc,KEEP|FREE)
    - LE use - always below the line
    - init - Same as HEAP.
    - inc - Same as HEAP. (minimum)
    - KEEP | FREE - Same as HEAP



# Initializing Storage

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- STORAGE(getheap, freeheap, dsa alloc)
  - Getheap – Initialize heap storage
    - NONE – no overhead
    - One byte hex value to initialize storage with when heap element obtained
      - 00 similar to WSCLEAR option
      - Relatively low overhead
  - Freeheap – Initial heap storage
    - NONE – no overhead
    - One byte hex value to initialize storage with when heap element freed
      - Useful for debug purposes or security
      - Relatively low overhead



# Initializing Storage

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- STORAGE(getheap, freeheap,dsa alloc)
  - DSA alloc – Initialize stack storage
    - NONE – No initialization – no overhead
    - CLEAR – Entire unused initial stack segment is cleared just before the main program is given control – low overhead
    - A one byte hex value to initialize storage with when stack frame (DSA) is obtained
      - EXTREMELY HIGH OVERHEAD
      - EXTREMELY HIGH OVERHEAD
      - EXTREMELY HIGH OVERHEAD



# Initializing Storage

- Simple program that makes lots of calls
  - STORAGE(,,none)

```
=====
                                REGION          --- STEP TIMINGS ---
STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME  ELAPSED TIME
GO                STORRTO      00      60K      0:00:01.11  0:00:02.42
```

- STORAGE(,,00)

```
=====
                                REGION          --- STEP TIMINGS ---
STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME  ELAPSED TIME
GO                STORRTO      00      60K      0:00:03.17  0:00:04.63
```



# Initializing Storage

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- Best ways to ensure the proper initial value for your variables
  - Use compiler initialization
  - Set them prior to use in your program



# Tuning storage

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- Objectives
  - Use as little storage as possible
  - Have program run as efficiently as possible
- The above objectives are often at odds with each other.
- One way to make a program run faster is to “throw” more storage at it.
  - Care must be taken to use storage wisely
  - Much of what will talk about can be done without recompiling or reworking the program.



# Tuning storage

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- Simple example
  - In a test environment (not production) use the RPTSTG run-time option.
  - A report will be generated describing the storage used by the program.
  - This information can be used to assist with better settings of Language Environment run-time options





# Tuning storage

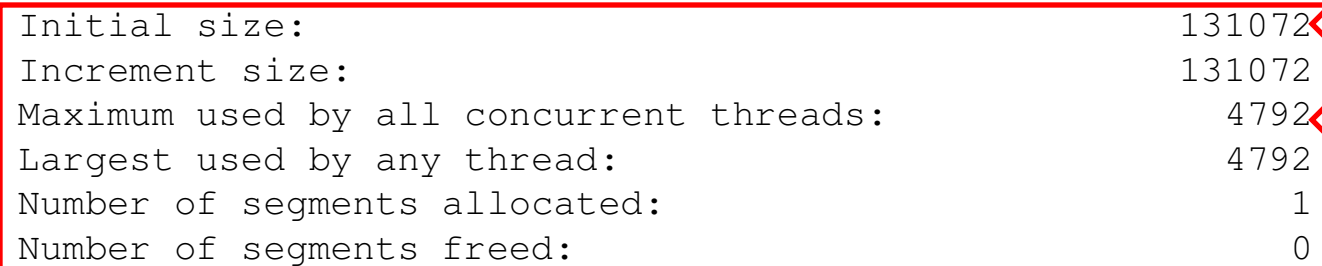
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## ■ Simple example

Storage Report for Enclave main 02/07/11 5:12:26 PM  
Language Environment V01 R12.00

### STACK statistics:

Initial size:	131072
Increment size:	131072
Maximum used by all concurrent threads:	4792
Largest used by any thread:	4792
Number of segments allocated:	1
Number of segments freed:	0



### THREADSTACK statistics:

Initial size:	0
Increment size:	0
Maximum used by all concurrent threads:	0
Largest used by any thread:	0
Number of segments allocated:	0
Number of segments freed:	0

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# Tuning storage

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## ■ Simple example...

LIBSTACK statistics:

Initial size:	4096
Increment size:	4096
Maximum used by all concurrent threads:	0
Largest used by any thread:	0
Number of segments allocated:	0
Number of segments freed:	0

THREADHEAP statistics:

Initial size:	4096
Increment size:	4096
Maximum used by all concurrent threads:	0
Largest used by any thread:	0
Successful Get Heap requests:	0
Successful Free Heap requests:	0
Number of segments allocated:	0
Number of segments freed:	0

# Tuning storage

## ■ Simple example...

HEAP statistics:

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	3328
Successful Get Heap requests:	4
Successful Free Heap requests:	2
Number of segments allocated:	1
Number of segments freed:	0

HEAP24 statistics:

Initial size:	8192
Increment size:	4096
Total heap storage used (sugg. initial size):	0
Successful Get Heap requests:	0
Successful Free Heap requests:	0
Number of segments allocated:	0
Number of segments freed:	0



# Tuning storage

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## ■ Simple example...

ANYHEAP statistics:

Initial size:	16384
Increment size:	8192
Total heap storage used (sugg. initial size):	616
Successful Get Heap requests:	6
Successful Free Heap requests:	2
Number of segments allocated:	1
Number of segments freed:	0

BELOWHEAP statistics:

Initial size:	8192
Increment size:	4096
Total heap storage used (sugg. initial size):	0
Successful Get Heap requests:	0
Successful Free Heap requests:	0
Number of segments allocated:	0
Number of segments freed:	0



# Tuning storage

- Now using stack(8k,4k) heap(4k,4k)

STACK statistics:

Initial size:	8192
Increment size:	4096
Maximum used by all concurrent threads:	4792
Largest used by any thread:	4792
Number of segments allocated:	1
Number of segments freed:	0

...snip...

HEAP statistics:

Initial size:	4096
Increment size:	4096
Total heap storage used (sugg. initial size):	3328
Successful Get Heap requests:	4
Successful Free Heap requests:	2
Number of segments allocated:	1
Number of segments freed:	0



# Tuning storage

---

- A bit more meaty!
  - Testcase requests 50000 random pieces of storage of sizes from 1 to 32K in size
  - Then the storage is freed.
  - We'll run the program without tuning
  - We'll then re-run the program (same random values) after tuning
  - Have we saved storage? Performance?

# Tuning storage

## ■ Untuned example

STACK statistics:

Initial size:	131072
Increment size:	131072
Maximum used by all concurrent threads:	204184
Largest used by any thread:	204184
Number of segments allocated:	3
Number of segments freed:	0

...Snip...

HEAP statistics:

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	819229056
Successful Get Heap requests:	50002
Successful Free Heap requests:	50000
Number of segments allocated:	27949
Number of segments freed:	0

- Note: 27949 segments of 32k each – 915,832,832 bytes

# Tuning storage

- Tuned HEAP(100M,100M) STACK(256K,256K)

STACK statistics:

Initial size:	262144
Increment size:	262144
Maximum used by all concurrent threads:	200944
Largest used by any thread:	200944
Number of segments allocated:	1
Number of segments freed:	0

...Snip...

HEAP statistics:

Initial size:	104857600
Increment size:	104857600
Total heap storage used (sugg. initial size):	818334944
Successful Get Heap requests:	50002
Successful Free Heap requests:	50000
Number of segments allocated:	8
Number of segments freed:	0

- Note: 8 segments of 100M each – 838,860,800 bytes!!!

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# Tuning storage

- Look what else happened!

- Untuned

```
=====
```

					--- STEP TIMINGS ---		
STEPNAME	PROCSTEP	PGMNAME	CC	REGION USED	CPU TIME	ELAPSED TIME	EXCP
CLPG	COMPILE	CBCDRVR	00	72K	0:00:00.06	0:00:02.84	1590
CLPG	PLKED	EDCPRLK	04	60K	0:00:00.01	0:00:00.99	534
CLPG	LKED	HEWL	00	92K	0:00:00.01	0:00:00.63	174
CLPG	GO	PGM=* .DD	00	60K	0:00:15.44	0:00:18.02	505

- Tuned

```
=====
```

					--- STEP TIMINGS ---		
STEPNAME	PROCSTEP	PGMNAME	CC	REGION USED	CPU TIME	ELAPSED TIME	EXCP
CLPG	COMPILE	CBCDRVR	00	72K	0:00:00.06	0:00:03.67	1555
CLPG	PLKED	EDCPRLK	04	60K	0:00:00.01	0:00:01.25	535
CLPG	LKED	HEWL	00	92K	0:00:00.01	0:00:00.50	170
CLPG	GO	PGM=* .DD	00	60K	0:00:00.12	0:00:01.45	501

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# Tuning storage

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- What about KEEP vs FREE
  - Testcase requests 50000 random pieces of storage of sizes from 1 to 32K in size
  - Free 20000 pieces, then get 20000 more
  - Free everything
  - We'll run the program without tuning and FREE
  - We'll run the program without tuning and KEEP
  - What have we done to storage and performance?

# Tuning storage

## ■ Untuned

HEAP statistics: **(Using FREE)**

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	819982896
Successful Get Heap requests:	70002
Successful Free Heap requests:	69999
Number of segments allocated:	39122
Number of segments freed:	39120

HEAP statistics: **(Using KEEP)**

Initial size:	32768
Increment size:	32768
Total heap storage used (sugg. initial size):	819983152
Successful Get Heap requests:	70002
Successful Free Heap requests:	70000
Number of segments allocated:	27952
Number of segments freed:	0

- Note: You can't determine storage used to back segments now

# Tuning storage

- Performance – not a huge difference but KEEP is faster!

- **FREE**

```
=====
-                                     REGION          --- STEP TIMINGS ---
- STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME    ELAPSED TIME    EXCP
- CLPG      COMPILE   CBCDRVR       00      72K      0:00:00.06   0:00:04.95      1496
- CLPG      PLKED     EDCPRLK       04      60K      0:00:00.01   0:00:02.46      504
- CLPG      LKED      HEWL          00      92K      0:00:00.01   0:00:01.13      171
- CLPG      GO        PGM=* .DD    00      60K      0:00:25.79   0:01:02.34      474
=====
```

- **KEEP**

```
=====
-                                     REGION          --- STEP TIMINGS ---
- STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME    ELAPSED TIME    EXCP
- CLPG      COMPILE   CBCDRVR       00      72K      0:00:00.06   0:00:03.15      1493
- CLPG      PLKED     EDCPRLK       04      60K      0:00:00.01   0:00:00.87      505
- CLPG      LKED      HEWL          00      92K      0:00:00.01   0:00:00.46      171
- CLPG      GO        PGM=* .DD    00      60K      0:00:22.34   0:00:24.85      469
=====
```

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# Tuning storage

- Look what happens when we tune.

HEAP statistics:

```
Initial size: 104857600
Increment size: 104857600
Total heap storage used (sugg. initial size): 819088944
Successful Get Heap requests: 70002
Successful Free Heap requests: 69999
Number of segments allocated: 8
Number of segments freed: 0
```

```
-----
-                                     REGION          --- STEP TIMINGS ---
- STEPNAME  PROCSTEP  PGMNAME      CC      USED      CPU TIME  ELAPSED TIME  EXCP
- CLPG      COMPILE   CBCDRVR      00      72K      0:00:00.06  0:00:02.67    1499
- CLPG      PLKED     EDCPRLK     04      60K      0:00:00.01  0:00:00.81    547
- CLPG      LKED      HEWL        00      92K      0:00:00.01  0:00:00.32    171
- CLPG      GO        PGM=*.DD    00      60K      0:00:00.15  0:00:01.04    496
```

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# More advanced tuning

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- What about those pesky Language Environment control blocks?
  - No externals to help
  - Effort can be made to reduce the number of enclaves
  - Use dynamic calls rather than linking to next program
  - Hard to see the results without using system tools... but let's try



# More advanced tuning

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- Simple program does a LINK to another program
  - A new enclave is created
  - This 2<sup>nd</sup> program continues to get storage until it runs out
  - It is able to obtain 21568K of storage



# More advanced tuning

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- Add to program to call down through 5 nested enclaves
  - Last enclave is able to obtain 20576K of storage
- Add to program to call down through 10 nested enclaves
  - Last enclave is able to obtain 19808K of storage
- Storage being consumed is to:
  - Load programs
  - Create enclave control blocks
    - This includes stacks and heaps
  - 1760K of storage usage (21568K-19808K)





# More advanced tuning

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- Change programs to use dynamic call rather than LINK
  - One call case – 21728K of storage available
  - Five call case – 21664K of storage available
  - Ten call case – 21600K of storage available
- Note how much less storage is consumed.
  - Basically just the amount to load the programs
  - 128K for 10 calls deep (21728K-21600K)



# Summary

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- Storage run-time option has high overhead for initializing the stack
- Use RPTSTG to tune your stack, heap and other storage sizes
- KEEP is faster than FREE
- Use dynamic call versus LINK
  - Requires program update or recompile



# Sources of Additional Info

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- All Language Environment documentation available on the Language Environment Web site
  - Language Environment Debugging Guide
  - Language Environment Programming Reference
  - Language Environment Programming Guide
- Language Environment Web site
  - [http://www.ibm.com/systems/z/os/zos/features/lang\\_environment](http://www.ibm.com/systems/z/os/zos/features/lang_environment)